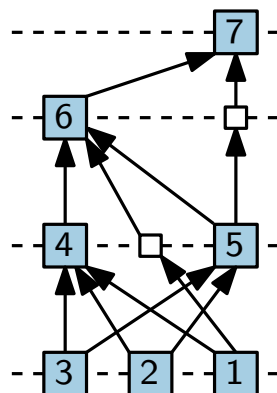
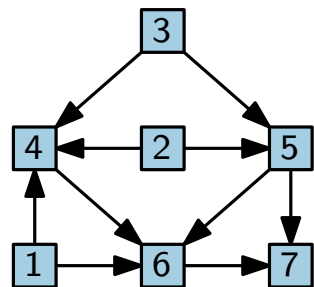
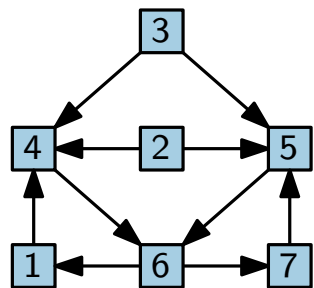


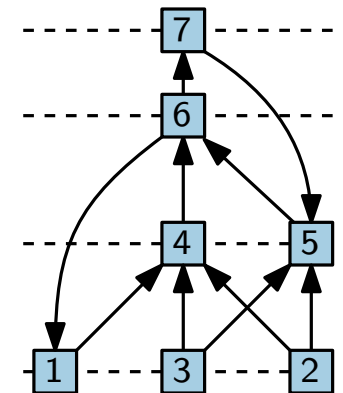
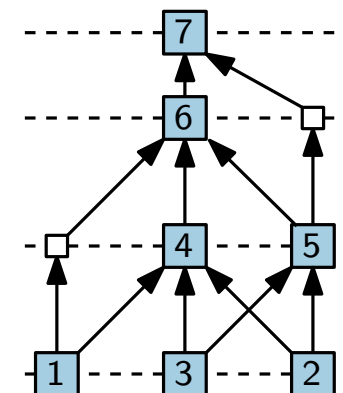
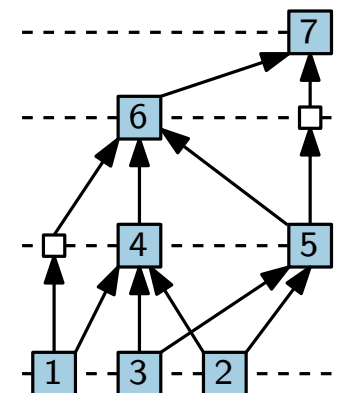
Visualization of Graphs

Lecture 8: Hierarchical Layouts: Sugiyama Framework

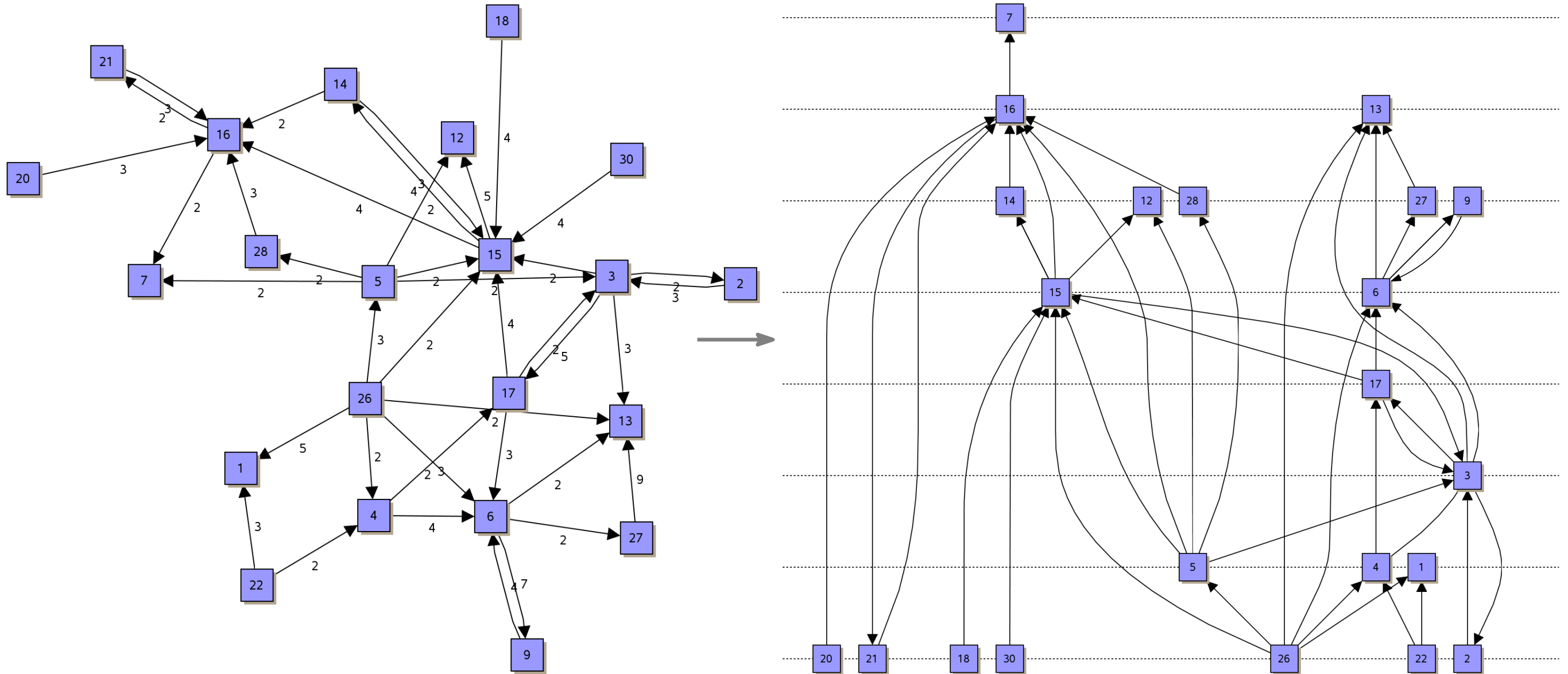


Alexander Wolff

Summer term 2026



Hierarchical Drawings – Motivation



Hierarchical Drawing

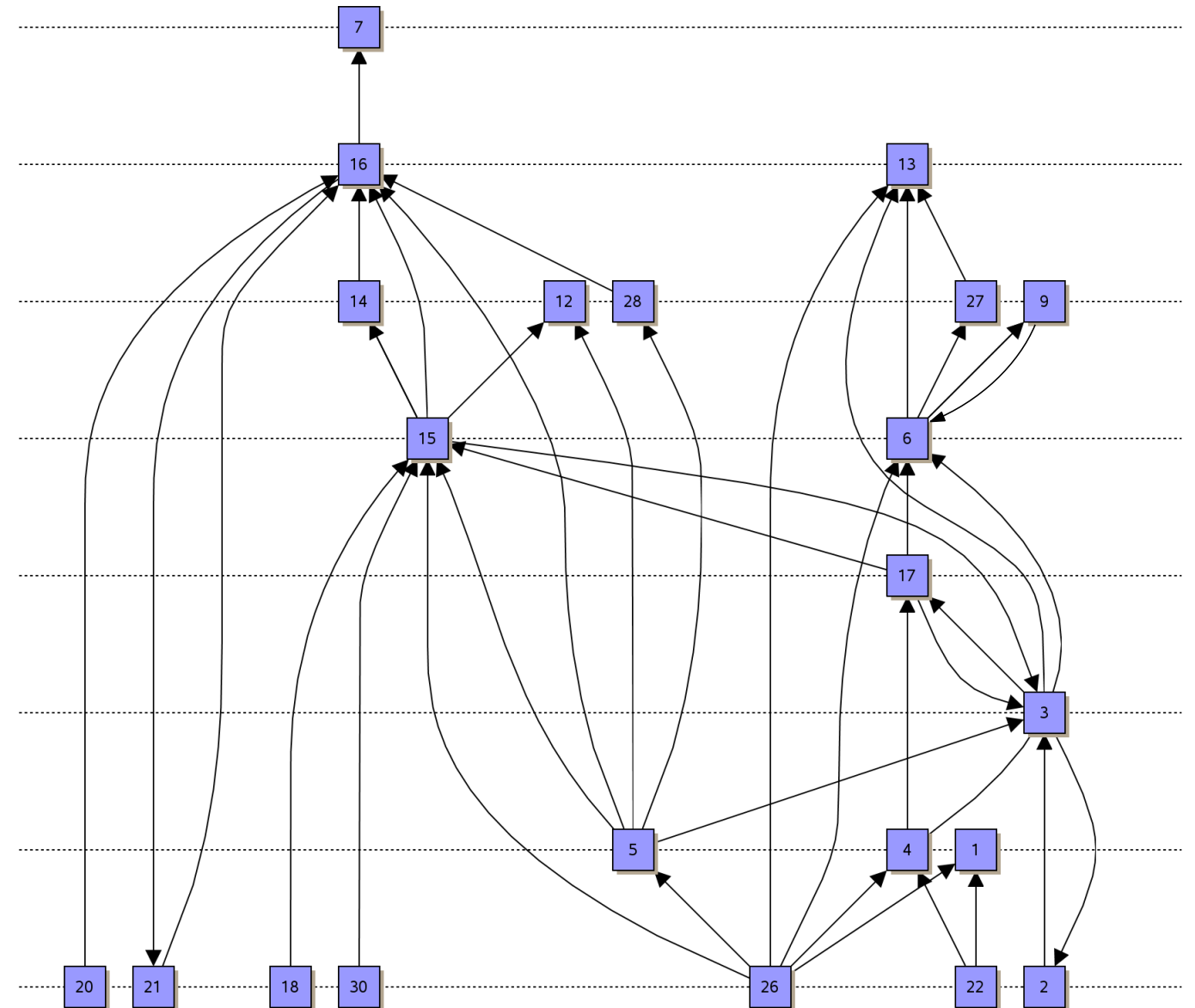
Problem Statement:

- Input: digraph G
- Output: drawing of G that “closely” reproduces the hierarchical properties of G

Desirable Properties:

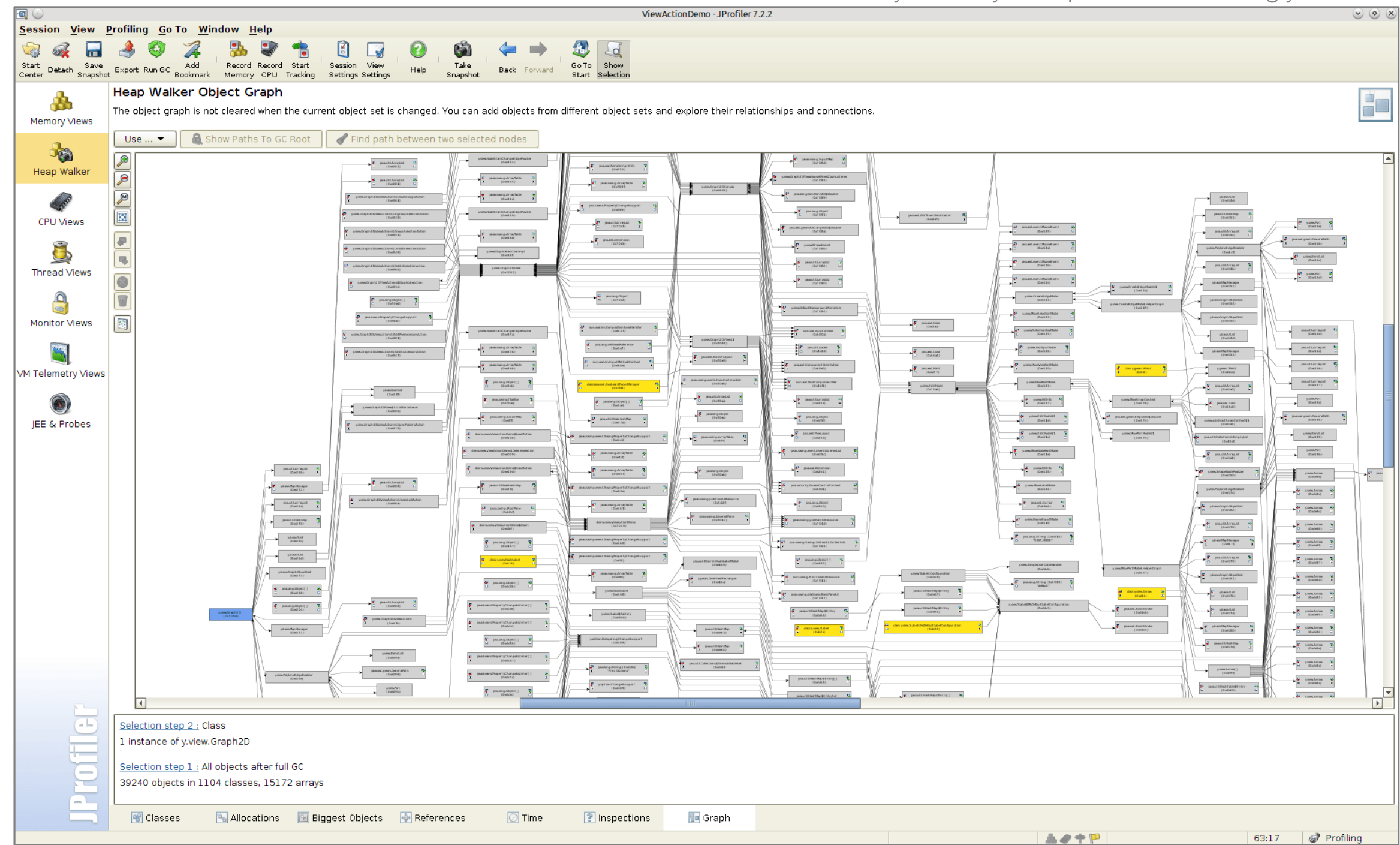
- edges are directed upwards,
- vertices lie on (few) horizontal lines,
- few pairs of edges cross,
- edges are short,
- vertices are evenly spaced.

Criteria can be contradictory!

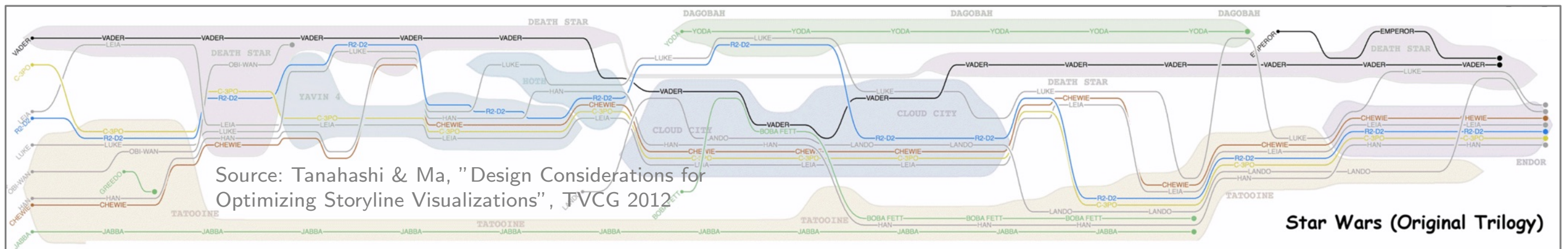
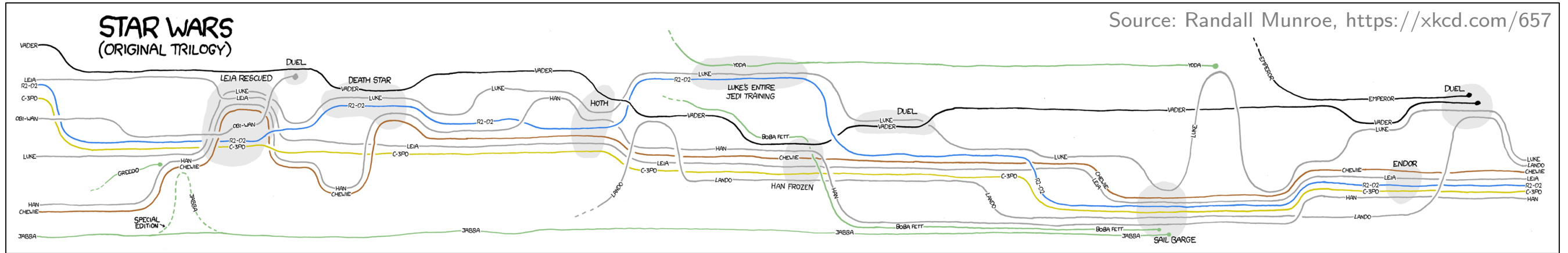


Hierarchical Drawing – Applications

yEd Gallery: Java profiler JProfiler using yFiles

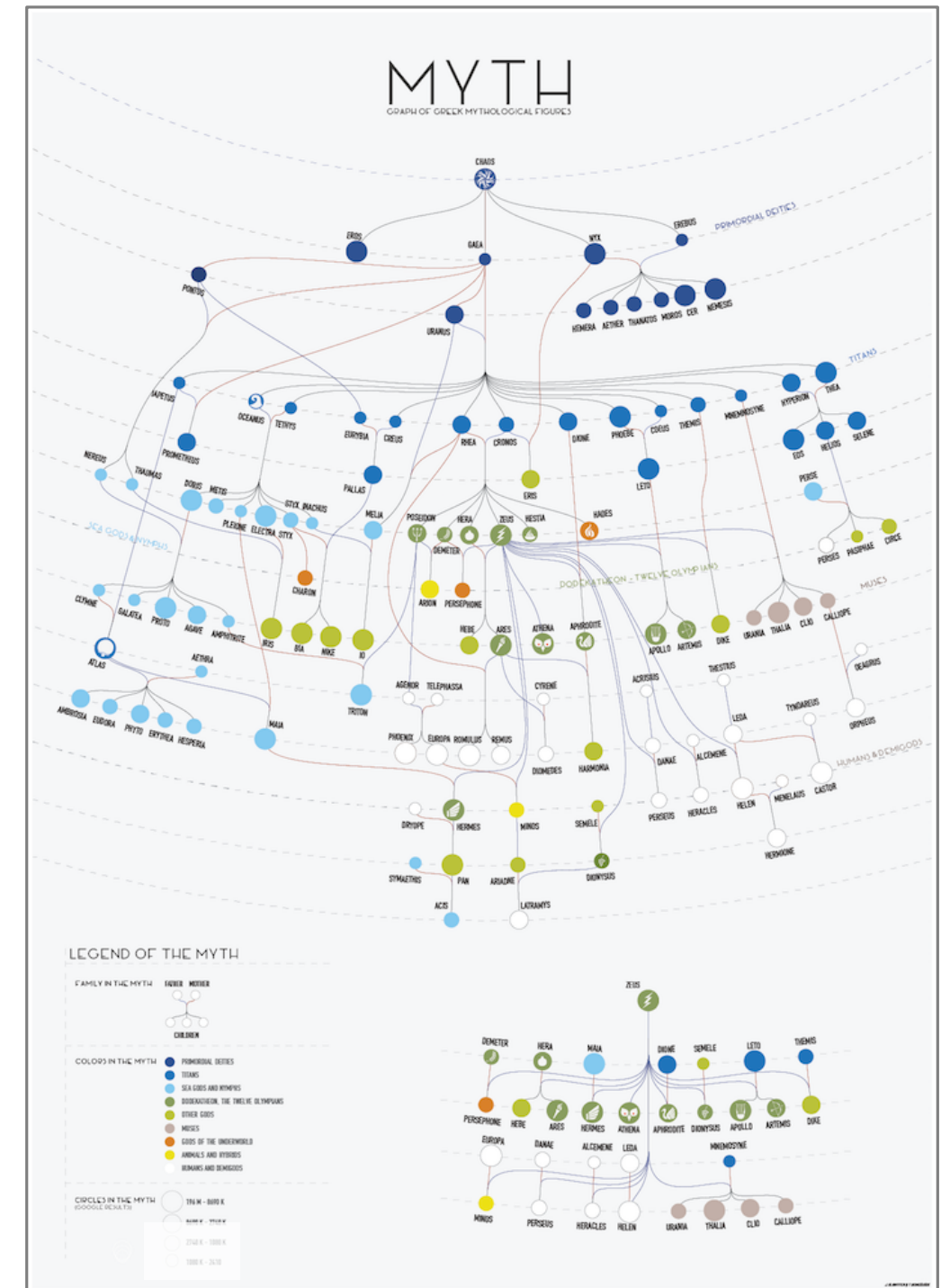


Hierarchical Drawing – Applications

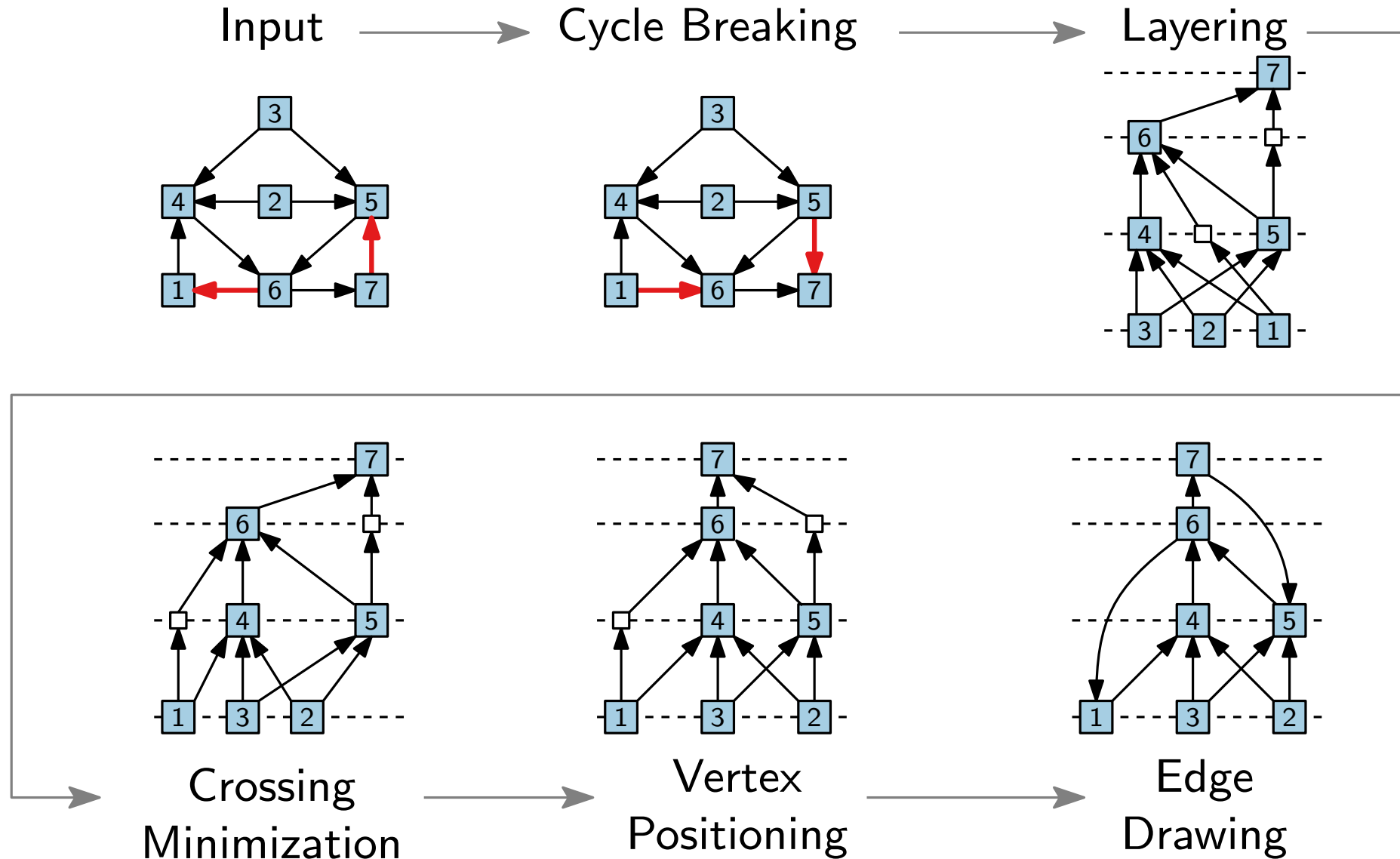


Hierarchical Drawing – Applications

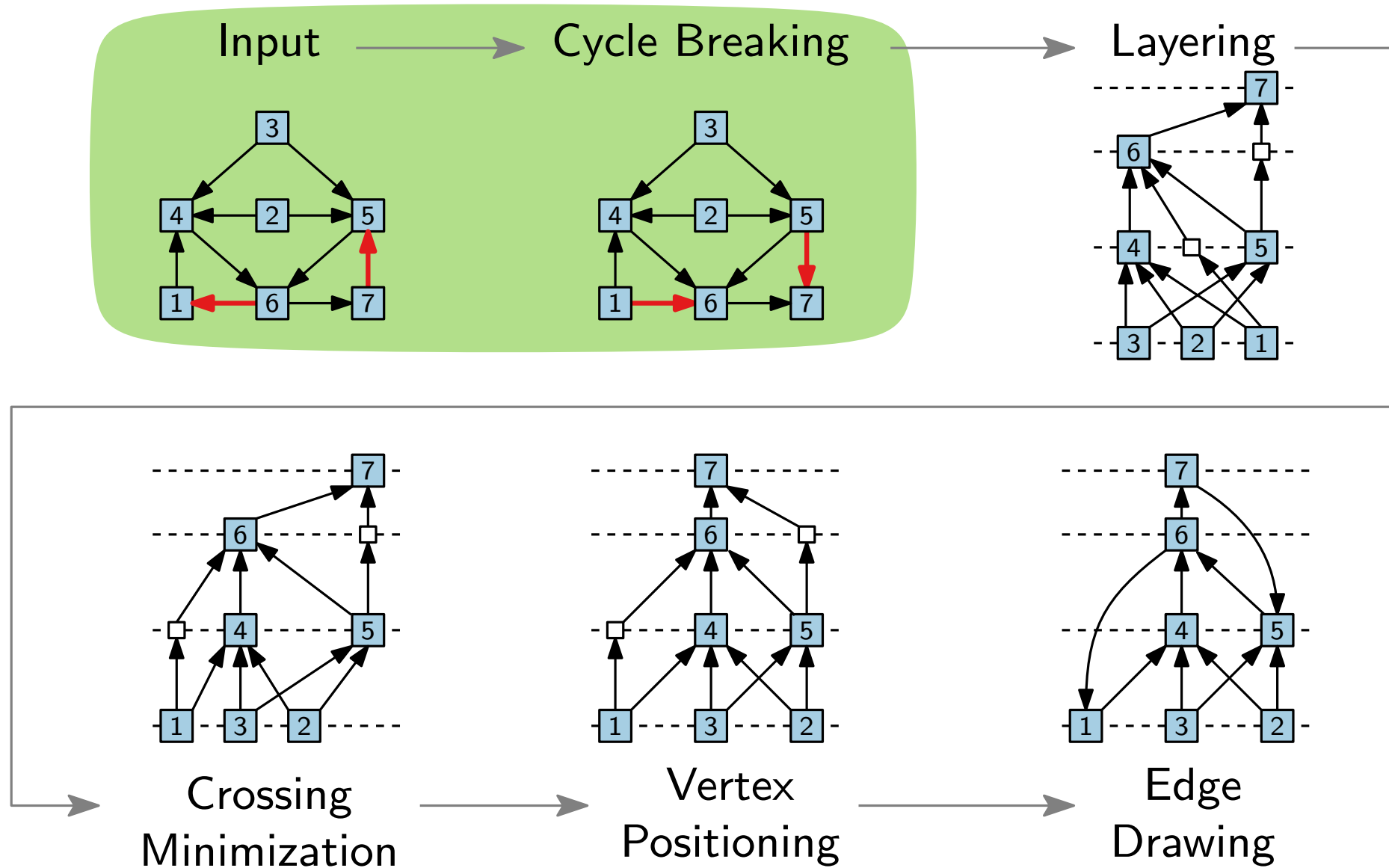
Source: Visualization that won the Graph Drawing Contest, Creative Track, 2016. Klawitter & Mchedlidze



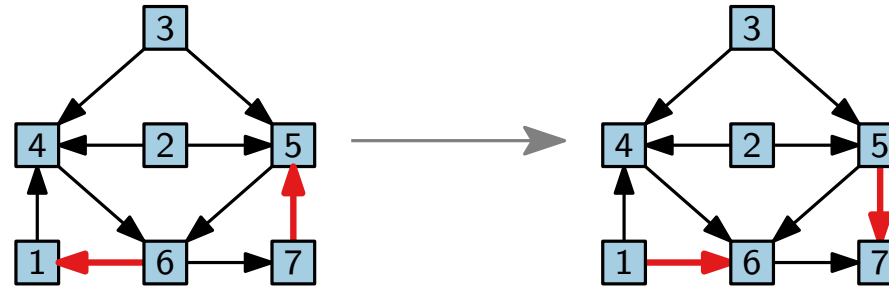
Classical Approach: Sugiyama Framework [Sugiyama, Tagawa, Toda '81]



Step 1: Cycle Breaking



Step 1: Cycle Breaking



Approach.

- Find minimum-size set E^* of edges that are not upward.
- Remove E^* and insert reversed edges.

Problem MINIMUM FEEDBACK ~~ARC~~ SET (~~F~~AS).

- Input: directed graph G
- Output: minimum-size set $E^* \subseteq E(G)$ such that $G^* = (V(G), \cancel{E(G)} \setminus \cancel{E^*})$ acyclic

... NP-hard 😞

edges in E^* but reversed

$(E(G) \setminus E^*) \cup E_{\text{rev}}^*$

Heuristic 1

[Berger, Shor '90]

GreedyMakeAcyclic(Digraph G):

$E' \leftarrow \emptyset$

foreach $v \in V(G)$ **do**

if $|E^{\rightarrow}(v)| \geq |E^{\leftarrow}(v)|$ **then**

$E' \leftarrow E' \cup E^{\rightarrow}(v)$

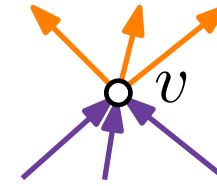
else

$E' \leftarrow E' \cup E^{\leftarrow}(v)$

 remove v and $E(v)$ from G .

return $G' = (V(G), E')$

■ G' is a DAG.

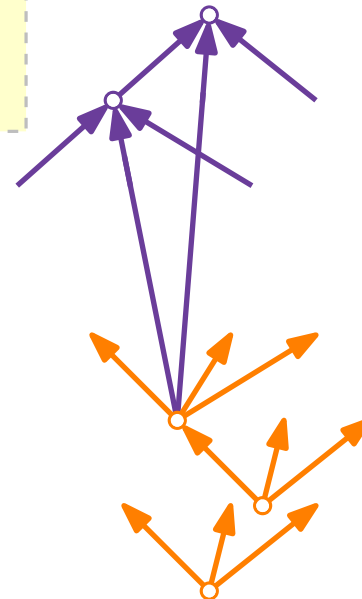


$$E^{\rightarrow}(v) := \{(v, u) : (v, u) \in E(G)\}$$

$$E^{\leftarrow}(v) := \{(u, v) : (u, v) \in E(G)\}$$

$$E(v) := E^{\rightarrow}(v) \cup E^{\leftarrow}(v)$$

Proof Idea.



Place the vertices on distinct y-coordinates.
y-coordinates increase/decrease towards the middle.
All edges point upwards.

Heuristic 1

[Berger, Shor '90]

GreedyMakeAcyclic(Digraph G):

$E' \leftarrow \emptyset$

foreach $v \in V(G)$ **do**

if $|E^{\rightarrow}(v)| \geq |E^{\leftarrow}(v)|$ **then**

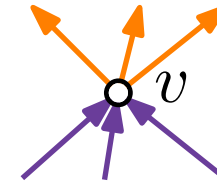
$E' \leftarrow E' \cup E^{\rightarrow}(v)$

else

$E' \leftarrow E' \cup E^{\leftarrow}(v)$

 remove v and $E(v)$ from G .

return $G' = (V(G), E')$



$$E^{\rightarrow}(v) := \{(v, u) : (v, u) \in E(G)\}$$

$$E^{\leftarrow}(v) := \{(u, v) : (u, v) \in E(G)\}$$

$$E(v) := E^{\rightarrow}(v) \cup E^{\leftarrow}(v)$$

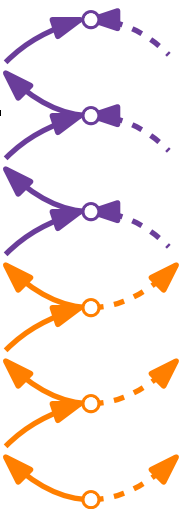
Proof Idea.

Use the vertex order from before (edges in E' upwards).

In this order, add the edges of $E(G) \setminus E'$ in rev. direction.

Added edges have other endpoint more in the middle.

→ All edges point upwards.



- G' is a DAG.
- $E(G) \setminus E'$ is a feedback set.

Heuristic 1

[Berger, Shor '90]

GreedyMakeAcyclic(Digraph G):

$E' \leftarrow \emptyset$

foreach $v \in V(G)$ **do**

if $|E^{\rightarrow}(v)| \geq |E^{\leftarrow}(v)|$ **then**

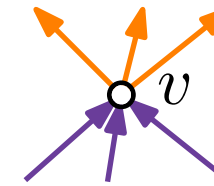
$E' \leftarrow E' \cup E^{\rightarrow}(v)$

else

$E' \leftarrow E' \cup E^{\leftarrow}(v)$

 remove v and $E(v)$ from G .

return $G' = (V(G), E')$



$$E^{\rightarrow}(v) := \{(v, u) : (v, u) \in E(G)\}$$

$$E^{\leftarrow}(v) := \{(u, v) : (u, v) \in E(G)\}$$

$$E(v) := E^{\rightarrow}(v) \cup E^{\leftarrow}(v)$$

■ G' is a DAG.

■ $E(G) \setminus E'$ is a feedback set.

■ Runtime: $\mathcal{O}(|V(G)| + |E(G)|)$

■ Quality guarantee: $|E'| \geq |E(G)|/2$

Heuristic 2

[Eades, Lin, Smyth '93]

GreedyMakeAcyclic2(Digraph G):

$E' \leftarrow \emptyset$

while $V(G) \neq \emptyset$ **do**

while $V(G)$ contains a **sink** v **do**

$E' \leftarrow E' \cup E^{\leftarrow}(v)$

 Remove v and $E^{\leftarrow}(v)$.

 Remove all **isolated vertices** from $V(G)$.

while $V(G)$ contains a **source** v **do**

$E' \leftarrow E' \cup E^{\rightarrow}(v)$

 Remove v and $E^{\rightarrow}(v)$.

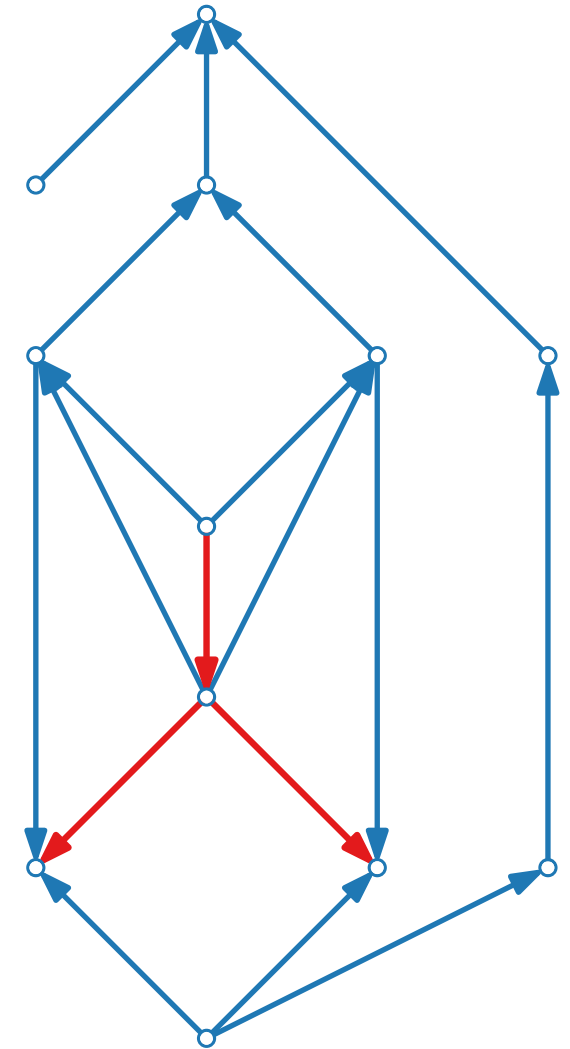
if $V(G) \neq \emptyset$ **then**

 Let $v \in V(G)$ such that $|E^{\rightarrow}(v)| - |E^{\leftarrow}(v)|$ maximal.

$E' \leftarrow E' \cup E^{\rightarrow}(v)$

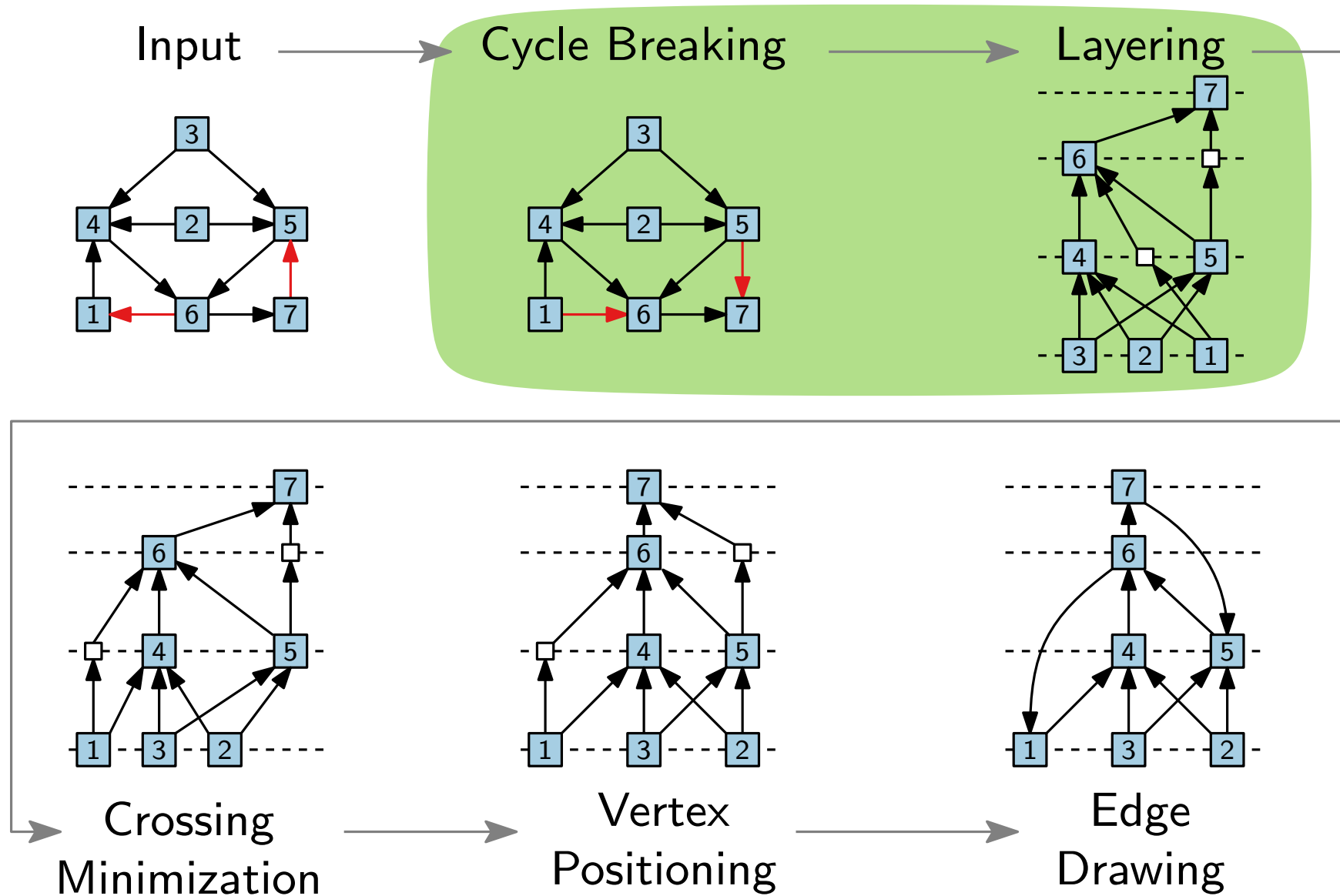
 Remove v and $E(v)$ from G .

return $G' = (V(G), E')$



- Time: $\mathcal{O}(|V(G)| + |E(G)|)$ [Main idea: Use bins for sinks and sources, and a bin for each $|E^{\rightarrow}(v)| - |E^{\leftarrow}(v)|$]
- Quality guarantee: $|E'| \geq |E(G)|/2 + |V(G)|/6$

Step 2: Layering



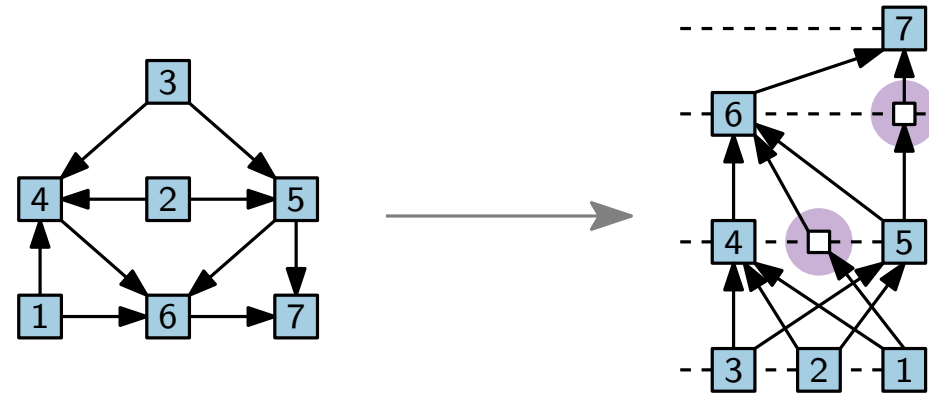
Step 2: Layering

Problem.

- Input: Acyclic digraph G .
- Output: Layering $y: V(G) \rightarrow \{1, \dots, n\}$,
such that, for every $(u, v) \in E(G)$, $y(u) < y(v)$.

Objective is to *minimize* ...

- number of layers, i.e., $\max_{v \in V(G)} y(v)$
- length of the longest edge, i.e., $\max_{(u,v) \in E(G)} y(v) - y(u)$
- width, i.e., $\max_{i \in \{1, \dots, n\}} |\{v \in V(G) : y(v) = i\}|$
- total edge length, i.e., number of **dummy vertices**.

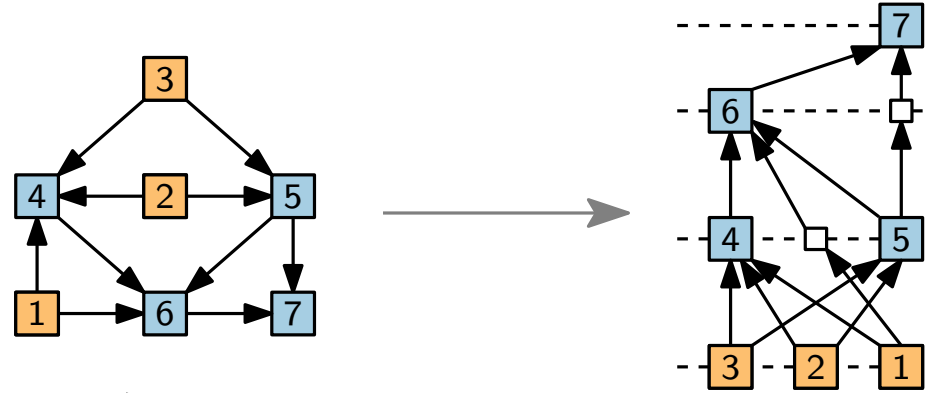


Whenever an edge spans across a layer, we insert a *dummy vertex*.

Minimize Number of Layers

Algorithm.

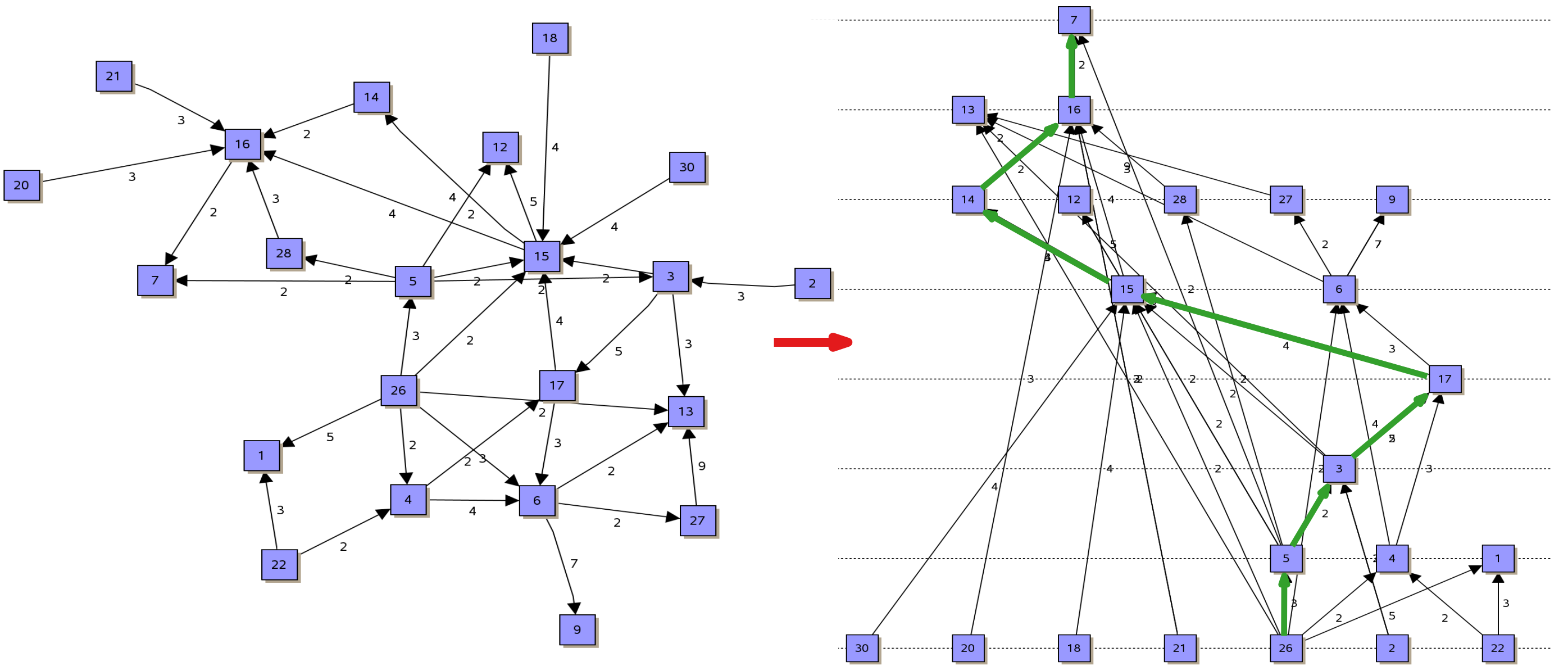
- for each **source** q ,
set $y(q) := 1$
- for each **non-source** v ,
set $y(v) := \max \{y(u) \mid (u, v) \in E(G)\} + 1$



Observation.

- $y(v)$ is the length of the longest path from a **source** to v plus 1.
... which is optimal!
- Can be implemented in linear time, for example, using a recursive algorithm.
- Closely related to topological sorting.

Example



Minimize Total Edge Length – ILP

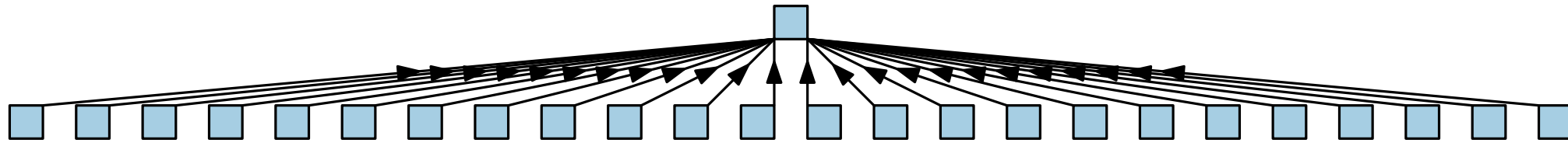
Can be formulated as an integer linear program:

$$\begin{array}{ll}
 \text{Minimize} & \sum_{(u,v) \in E(G)} (y(v) - y(u)) \\
 \text{subject to} & y(v) - y(u) \geq 1 \quad \forall (u,v) \in E(G) \\
 & y(v) \geq 1 \quad \forall v \in V(G) \\
 & y(v) \in \mathbb{Z} \quad \forall v \in V(G)
 \end{array}$$

One can show that:

- Constraint matrix is **totally unimodular** (every square submatrix has $\det$ in $\{-1, 0, 1\}$).
 $\Rightarrow$ Extreme point solutions of the LP relaxation (the ILP without $y(v) \in \mathbb{Z}$) are integer.
- The total edge length can be minimized in polynomial time.

Width



Drawings can be very wide.

Narrower Layer Assignment

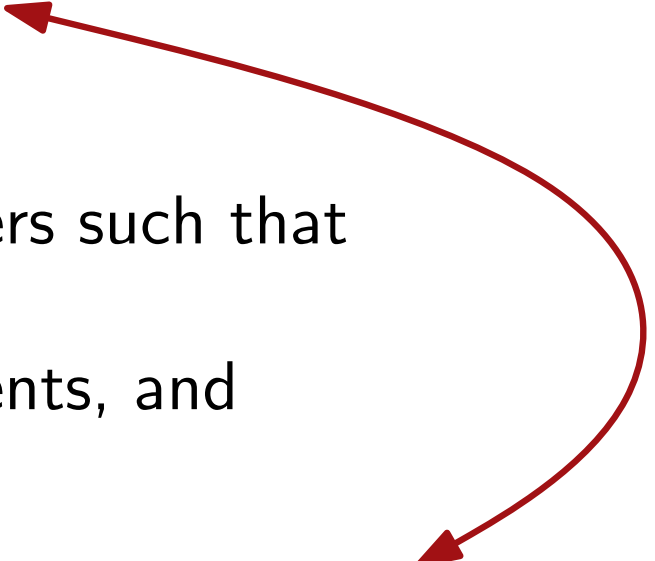
Problem: Layering with a given maximum width.

- Input: Acyclic digraph G , width $W > 0$
- Output: Assignment of the vertices of G to layers such that
 - the assignment is a layering,
 - each layer contains at most W elements, and
 - the number of layers is minimized.

Problem: Precedence-Constrained Multi-Processor Scheduling.

- Input: n jobs with unit processing time, W identical machines, partial ordering $<$ on the jobs.
- Output: Schedule respecting $<$ such that completion time (known as *makespan*) is minimized.
- NP-hard, $(2 - 1/W)$ -approximation, no $(4/3 - \varepsilon)$ -approximation ($W \geq 3$)

same!

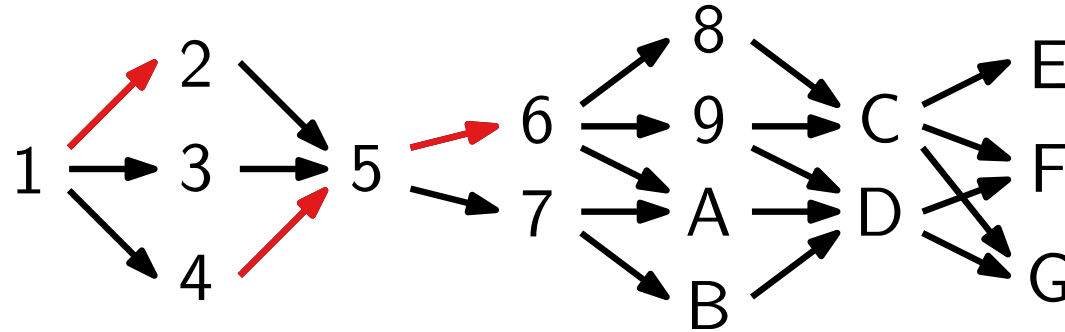


Approximating Precedence-Const. Multi-Processor Scheduling

- Jobs stored in a list L , which is topologically sorted.
- A job in L is *available* if all its predecessors have been scheduled.
- For each point in time $t = 1, 2, \dots$, we can schedule $\leq W$ available jobs.
- As long as there are free machines and available jobs, take the first available job and assign it to a free machine.

Approximating Precedence-Const. Multi-Processor Scheduling

Input: Precedence graph (divided into layers of arbitrary width)



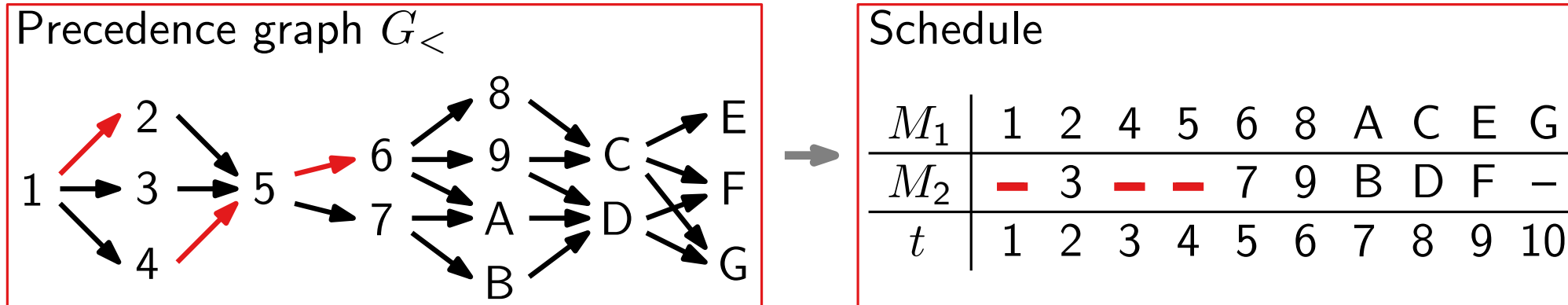
Number of machines is $W = 2$.

Output: Schedule

M_1	1	2	4	5	6	8	A	C	E	G
M_2	–	3	–	–	7	9	B	D	F	–
t	1	2	3	4	5	6	7	8	9	10

Question: Good approximation factor?

Approximating PCMPS – Analysis for $W = 2$



„The art of the lower bound“

$\text{OPT} \geq \lceil n/2 \rceil$ and $\text{OPT} \geq \ell := \text{Number of layers of } G_{<} (= \text{length of longest path in } G_{<})$

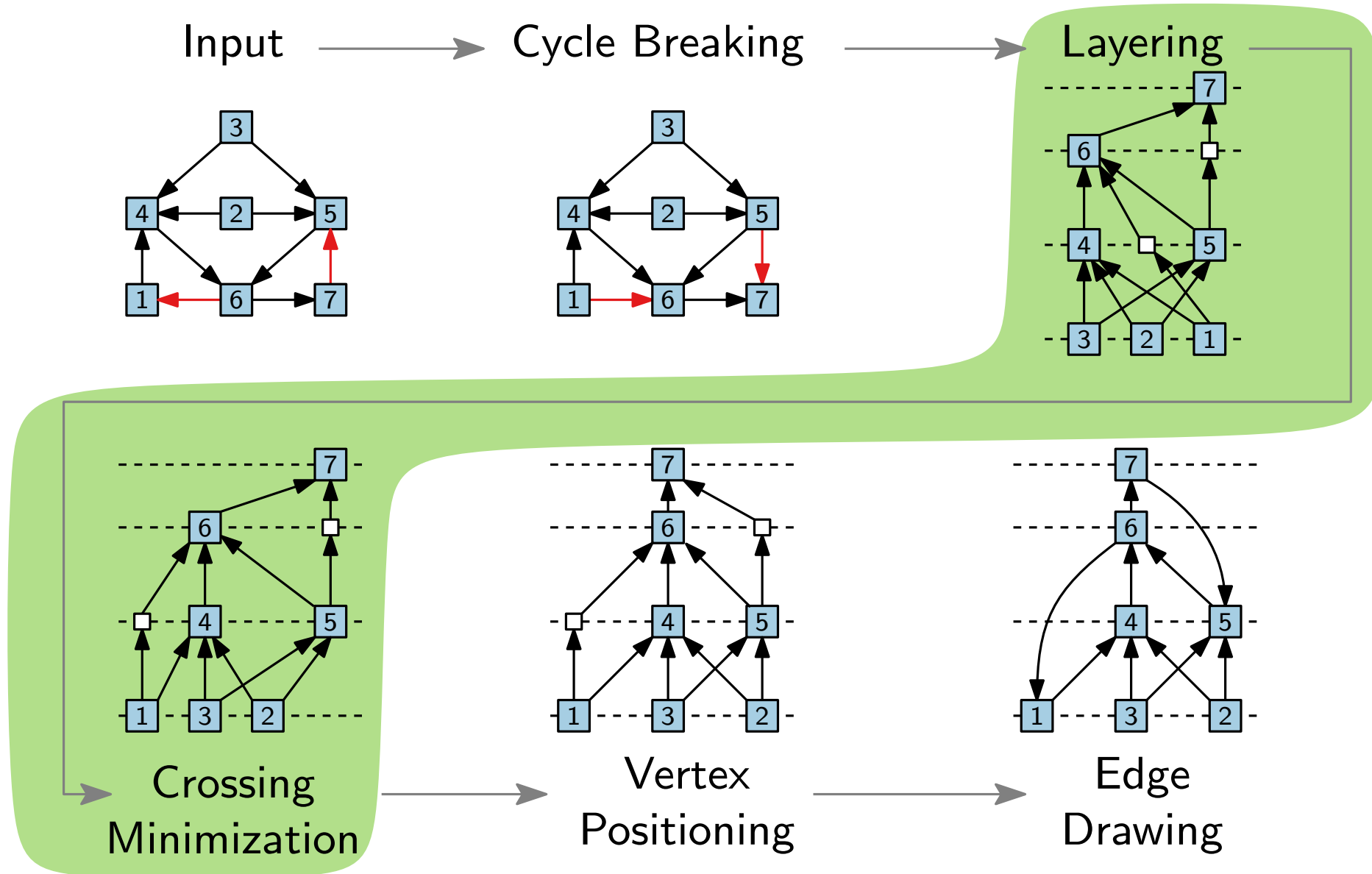
Goal: measure the quality of our algorithm using the lower bounds

$\leq (2 - 1/W) \cdot \text{OPT}$ in general case

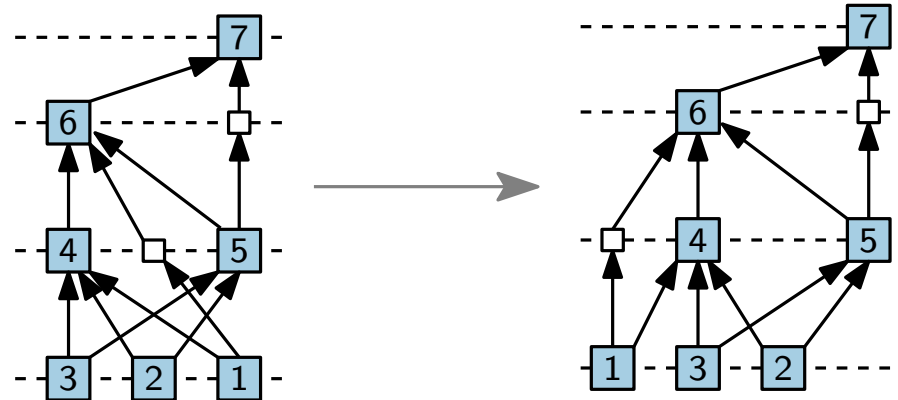
Bound. $\text{ALG} \leq \lceil \frac{n+\ell}{2} \rceil \approx \lceil n/2 \rceil + \ell/2 \leq 3/2 \cdot \text{OPT}$

↑ insertion of pauses (—) in the schedule
(except the last) maps to layers of $G_{<}$

Step 3: Crossing Minimization



Step 3: Crossing Minimization



Problem.

- Input: Graph G , layering $y: V(G) \rightarrow \{1, \dots, n\}$
- Output: (Re-)ordering of vertices in each layer such that the number of crossings is minimized.
- NP-hard, even for two layers. [Garey & Johnson '83]
- Hardly any approaches optimize over multiple layers. 😞

Iterative Crossing Reduction

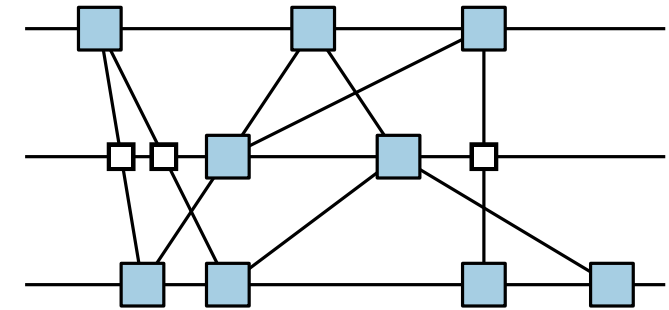
Observation. The number of crossings depends only on permutations of adjacent layers.

Idea.

- Permute one layer after the other.
- Treat dummy vertices as “regular” vertices.

Algorithm scheme.

- (1) choose a random permutation of L_1
- (2) iteratively consider pairs of adjacent layers (L_i, L_{i+1})
- (3) minimize crossings by permuting L_{i+1} while keeping L_i fixed
- (4) repeat steps (2)–(3) in the reverse order (starting from topmost layer L_h)
- (5) repeat steps (2)–(4) until no further improvement is achieved
- (6) repeat steps (1)–(5) with different starting permutations on L_1



one-sided crossing minimization

One-Sided Crossing Minimization

Problem.

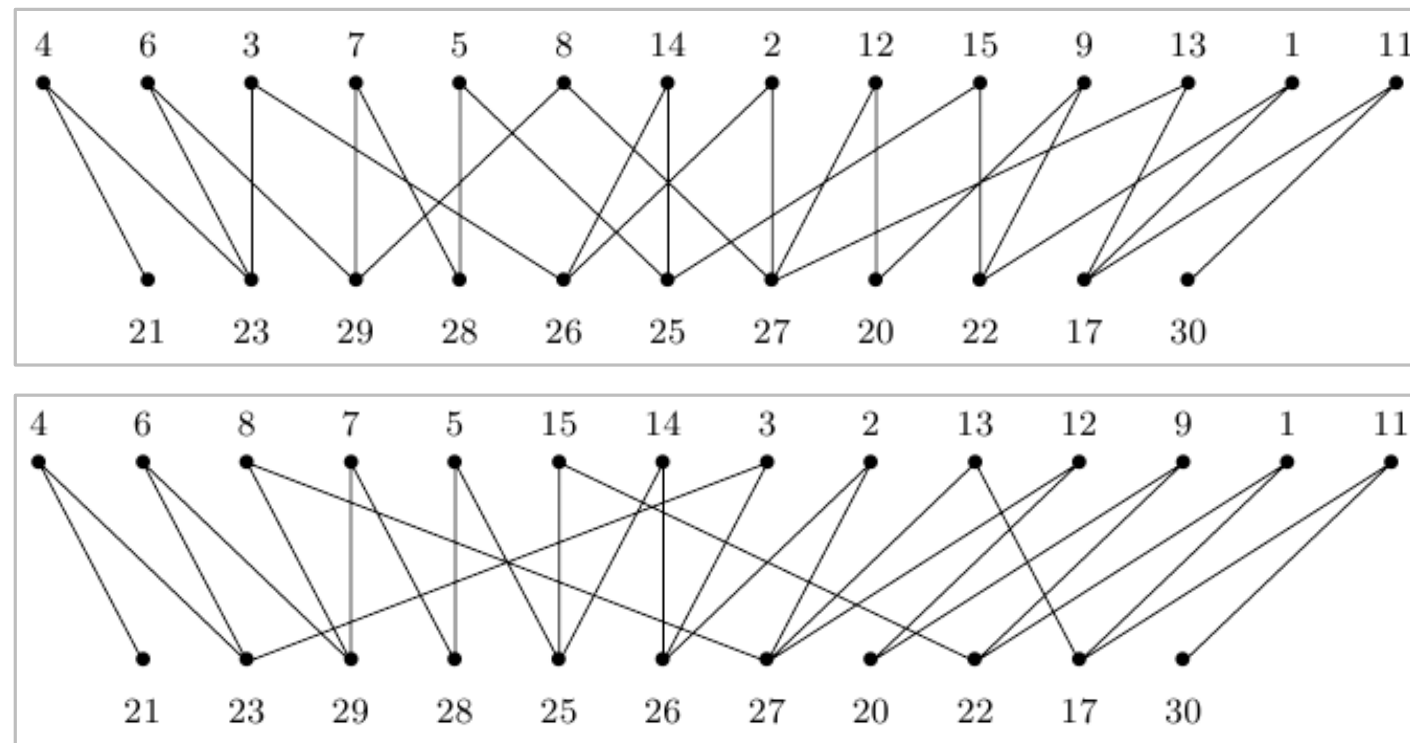
- Input:
 - bipartite graph G with $V(G) = L_1 \cup L_2$,
 - permutation π_1 on L_1
- Output: permutation π_2 of L_2 minimizing the number of edge crossings.

One-sided crossing minimization is NP-hard.

[Eades & Whitesides '94]

Algorithms.

- barycenter heuristic
- median heuristic
- Greedy-Switch
- ILP
- FPT



[Kaufmann & Wagner: Drawing Graphs]

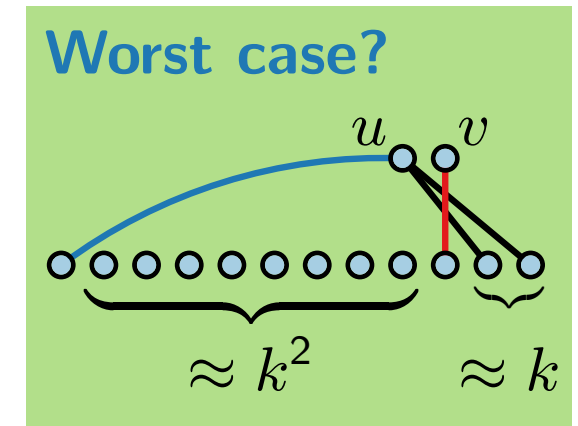
Barycenter Heuristic

[Sugiyama et al. '81]

- Intuition: There are few crossing if vertices are “close” to their neighbors.
- The barycenter of $u \in L_2$ is the mean rank of u 's neighbors on layer L_1 :

$$\text{bary}(u) := \frac{1}{\deg(u)} \sum_{v \in N(u)} \pi_1(v).$$

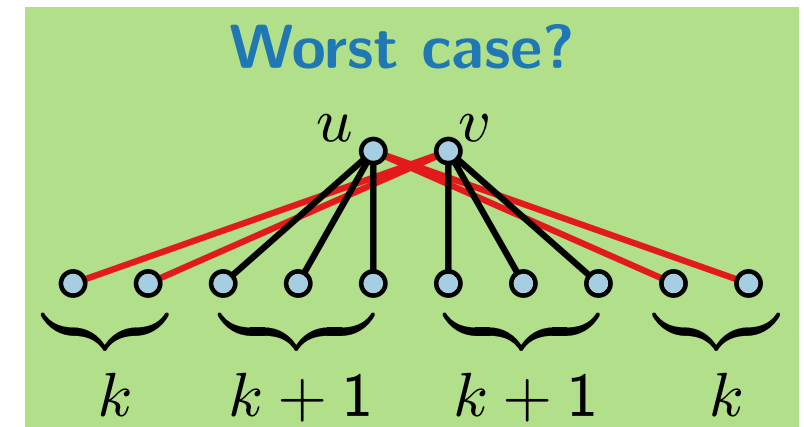
- To get π_2 , sort L_2 ascendingly according to $\text{bary}(\cdot)$.
- Vertices with the same barycenter keep their old relative ranks.
- Linear runtime (in the number of vertices and edges).
- Relatively good results in practice.
- Finds crossing-free solutions if they exist. $\leftarrow$ Exercise!
- Factor- $O(\sqrt{|V(G)|})$ approximation.



Median Heuristic

[Eades & Wormald '94]

- $\{v_1, \dots, v_k\} := N(u)$ with $\pi_1(v_1) < \pi_1(v_2) < \dots < \pi_1(v_k)$
- $$\text{med}(u) := \begin{cases} 0 & \text{if } N(u) = \emptyset, \\ \pi_1(v_{\lceil k/2 \rceil}) & \text{otherwise.} \end{cases}$$
- To get π_2 , sort L_2 ascendingly according to $\text{med}(\cdot)$.
- For vertices with the same median, place vertices of odd degree to the left of vertices of even degree (and keep the old relative ranks among the odd/even-degree vertices).
- Linear runtime (in the number of vertices and edges).
- Relatively good results in practice.
- Finds crossing-free solutions if they exist. ← Exercise!
- Factor-3 approximation. Proof in [GD Ch 11]



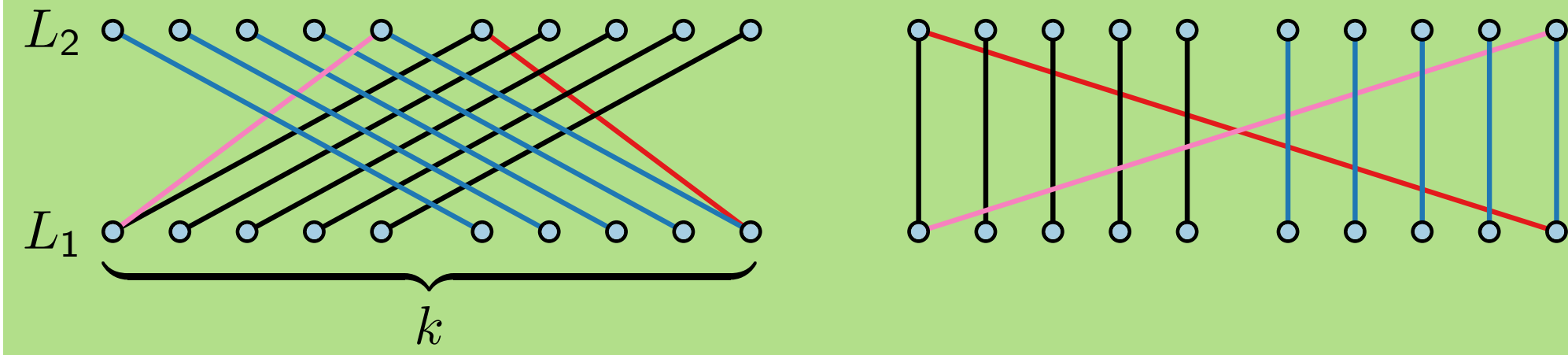
crossings: $2k(k+1) + k^2$ vs. $(k+1)^2$

Greedy-Switch Heuristic

[Eades & Kelly '86]

- Iteratively swap pairs of neighboring vertices on L_2 as long as the number of crossings decreases.
- Runtime: $O(|L_2|)$ per iteration; at most $|L_2|$ iterations $\Rightarrow O(|L_2|^2)$ time.
- Suitable as post-processing for other heuristics.

Worst case?



crossings: $\approx k^2/4$

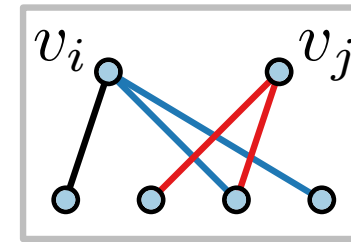
$\approx 2k$

Integer Linear Program (ILP)

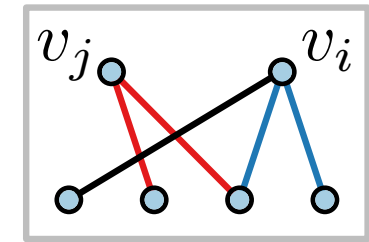
[Jünger & Mutzel, '97]

- Constant $c_{ij} := \#$ crossings between edges incident to v_i and v_j if $\pi_2(v_i) < \pi_2(v_j)$
- Variable x_{ij} for each $1 \leq i < j \leq n_2 := |L_2|$

$$x_{ij} = \begin{cases} 1 & \text{if } \pi_2(v_i) < \pi_2(v_j), \\ 0 & \text{otherwise.} \end{cases}$$



$$c_{ij} = 3$$



$$c_{ji} = 2$$

- Number of crossings of a permutations π_2 :

$$\text{cross}(\pi_2) = \sum_{i=1}^{n_2-1} \sum_{j=i+1}^{n_2} (c_{ij} - c_{ji})x_{ij} + \underbrace{\sum_{i=1}^{n_2-1} \sum_{j=i+1}^{n_2} c_{ji}}_{\text{constant}}$$

Integer Linear Program (ILP)

- Objective (minimize the number of crossings):

$$\text{minimize } \sum_{i=1}^{n_2-1} \sum_{j=i+1}^{n_2} (c_{ij} - c_{ji}) x_{ij}$$

- Transitivity constraints:

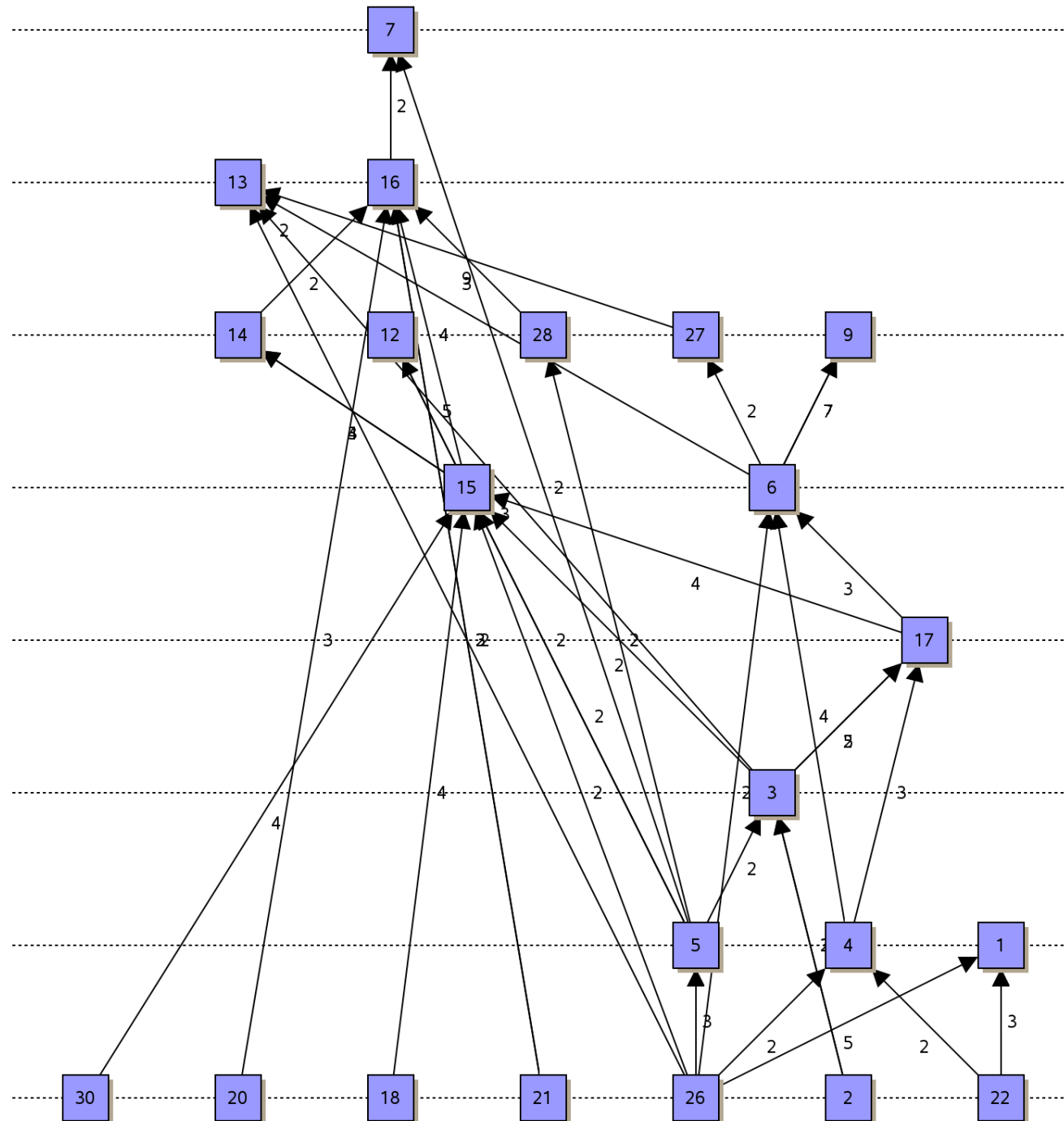
$$0 \leq x_{ij} + x_{jk} - x_{ik} \leq 1 \quad \text{for } 1 \leq i < j < k \leq n_2$$

i.e., if $x_{ij} = \underset{0}{1}$ and $x_{jk} = \underset{0}{1}$, then $x_{ik} = \underset{0}{1}$

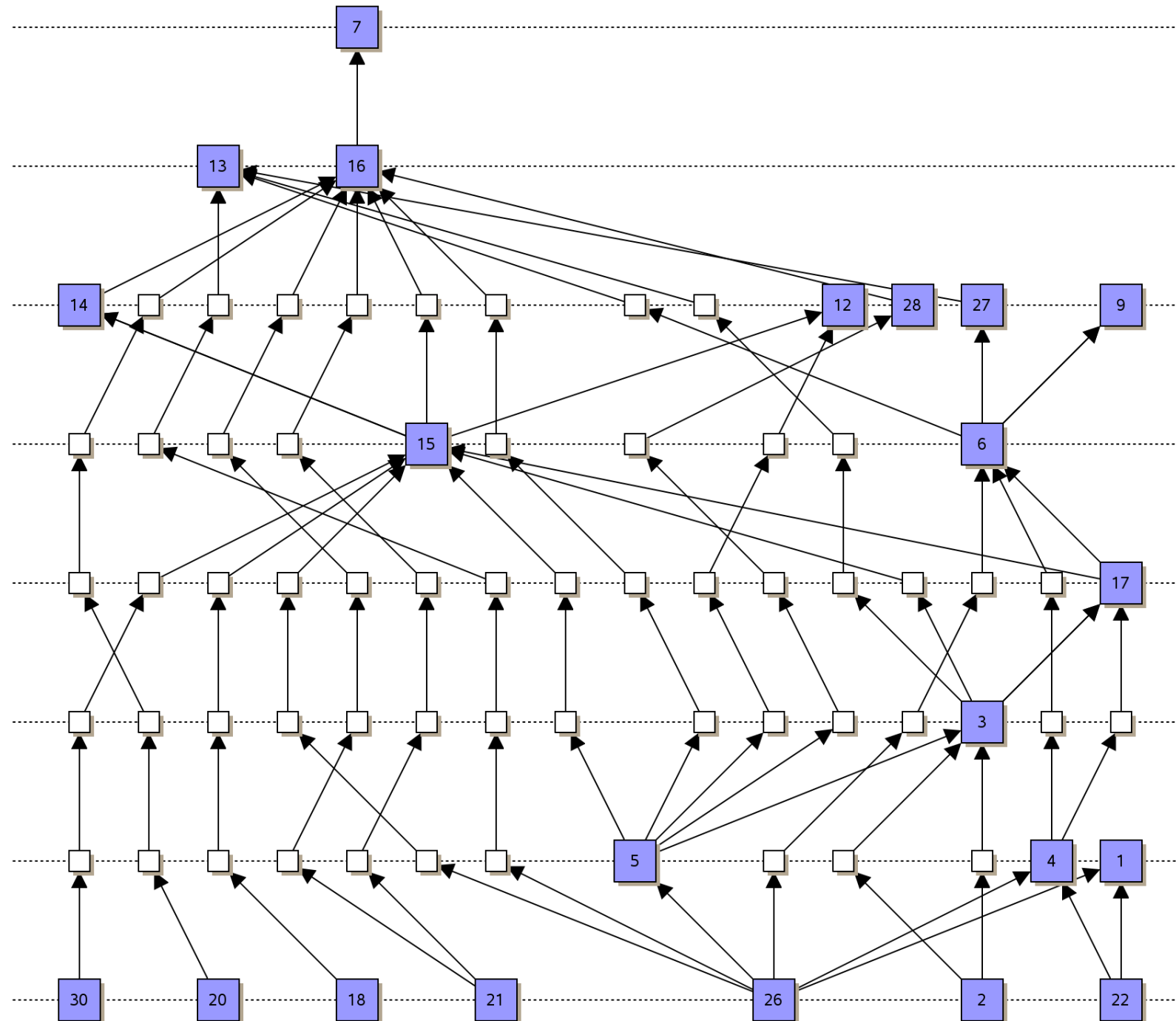
Properties.

- branch-and-cut technique applicable for this ILP
- useful for graphs of small to medium size
- finds optimal solution
- solution in polynomial time is not guaranteed

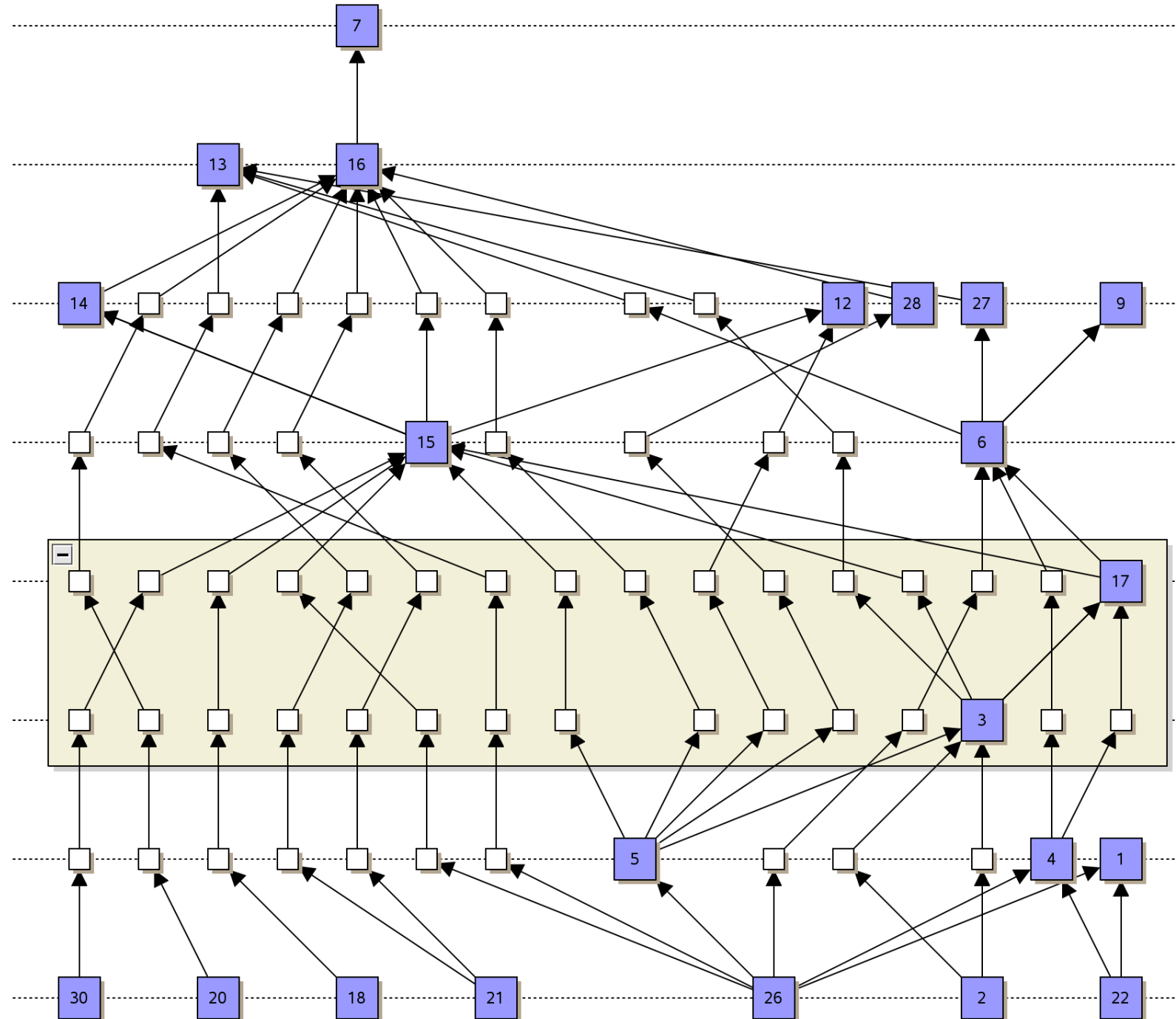
Iterations on Example



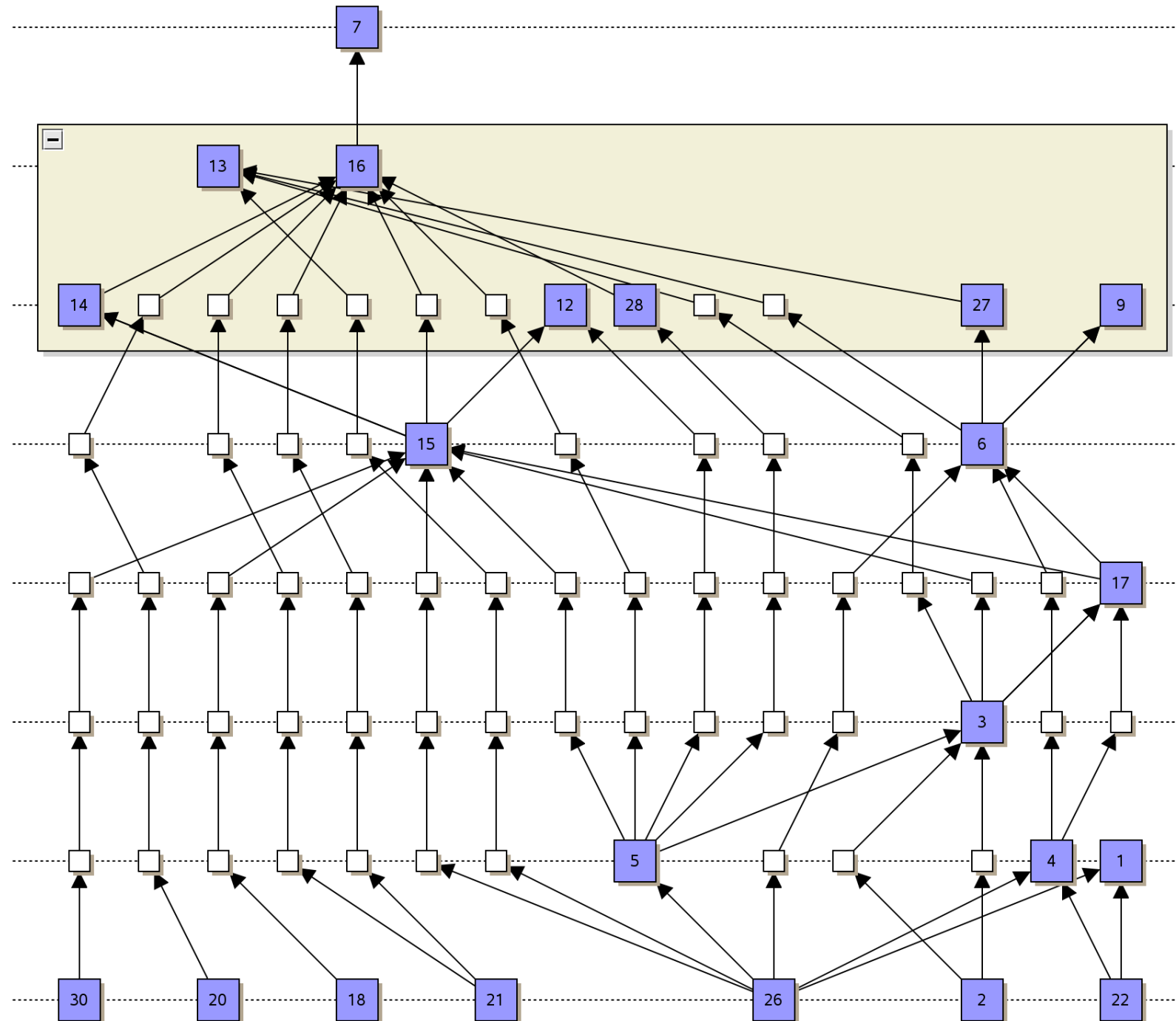
Iterations on Example



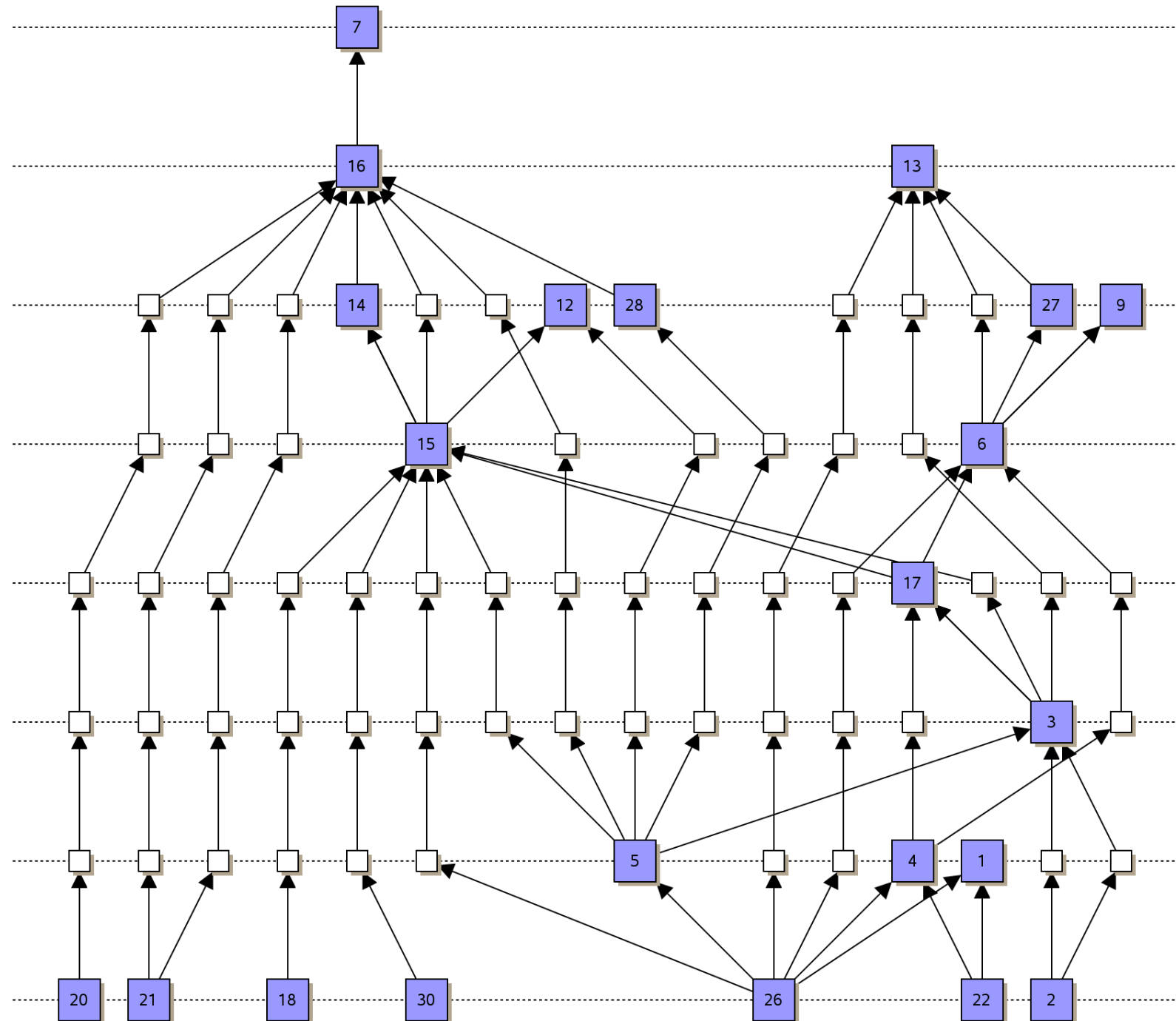
Iterations on Example



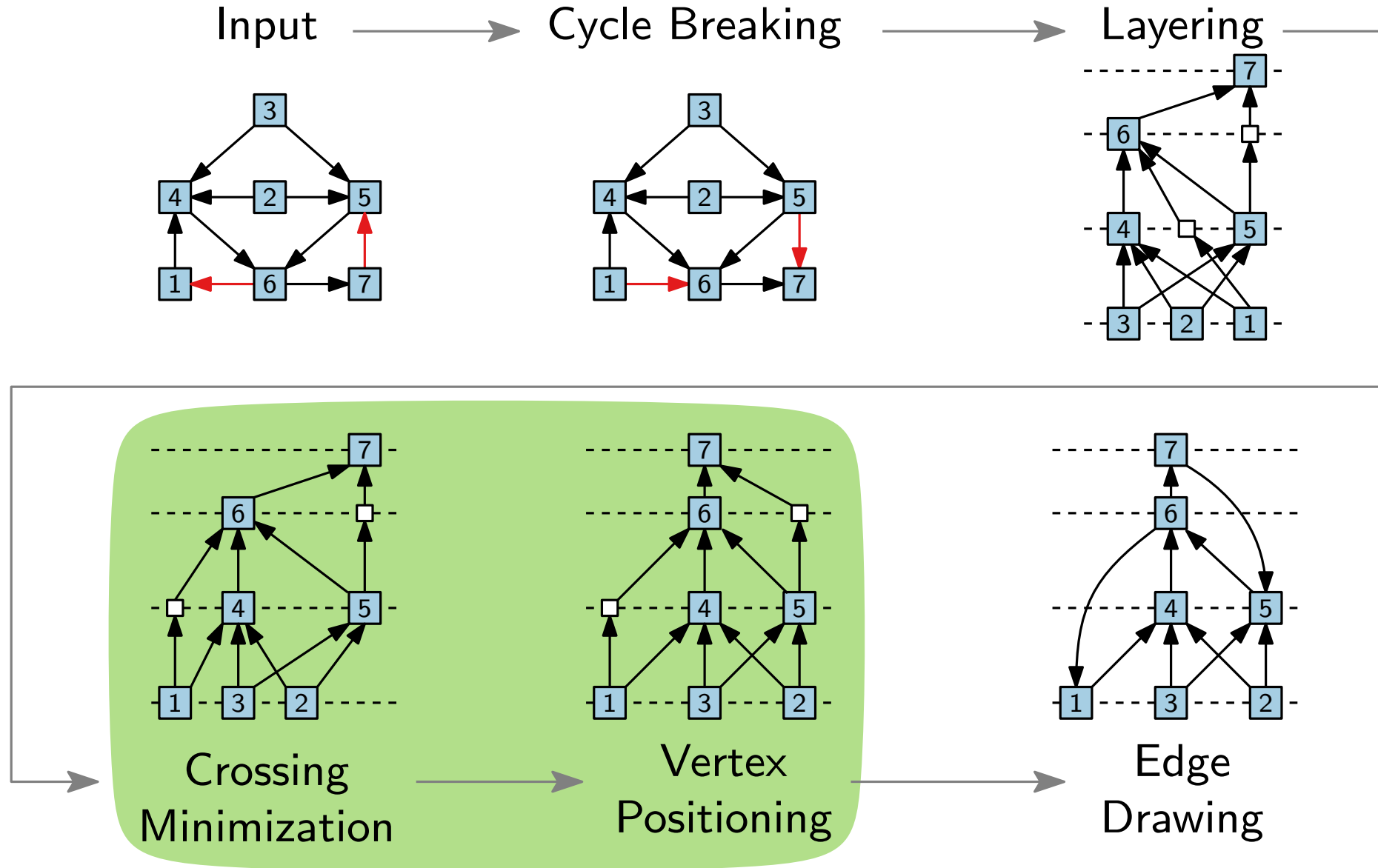
Iterations on Example



Iterations on Example



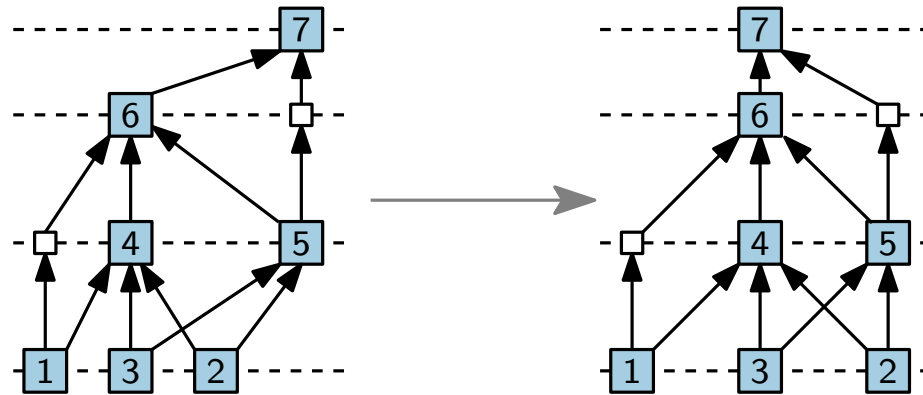
Step 4: Vertex Positioning



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Goals.

- paths of a single edge should be (close to) straight
- vertices on a layer evenly spaced
- prefer vertical edges



- **Exact:** Quadratic Program (QP)
- **Heuristic:** Iterative approach

Quadratic Program

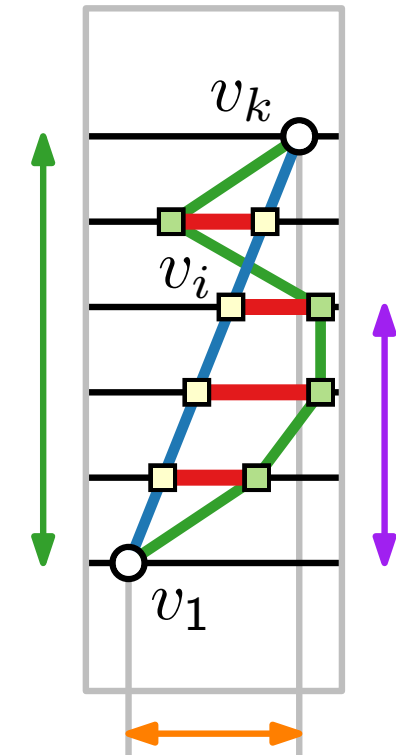
- Let $e = (v_1, v_k)$ be an edge of G , and let $p_e = (v_1, \dots, v_k)$ be the corresponding path with dummy vertices $v_2, \dots, v_{k-1}$.
- x -coordinate of v_i according to the line segment $\overline{v_1 v_k}$ (with equal spacing of the layers):

$$\overline{x(v_i)} = x(v_1) + \frac{i-1}{k-1} (x(v_k) - x(v_1))$$

- Define the deviation from the line

$$\text{dev}(p_e) := \sum_{i=2}^{k-1} \left(x(v_i) - \overline{x(v_i)} \right)^2$$

- Objective function: $\min \sum_{e \in E} \text{dev}(p_e)$
- Constraints for all vertices v, w in the same layer with w to the right of v :
 $x(w) - x(v) \geq \rho$ $\leftarrow$ min. horizontal distance

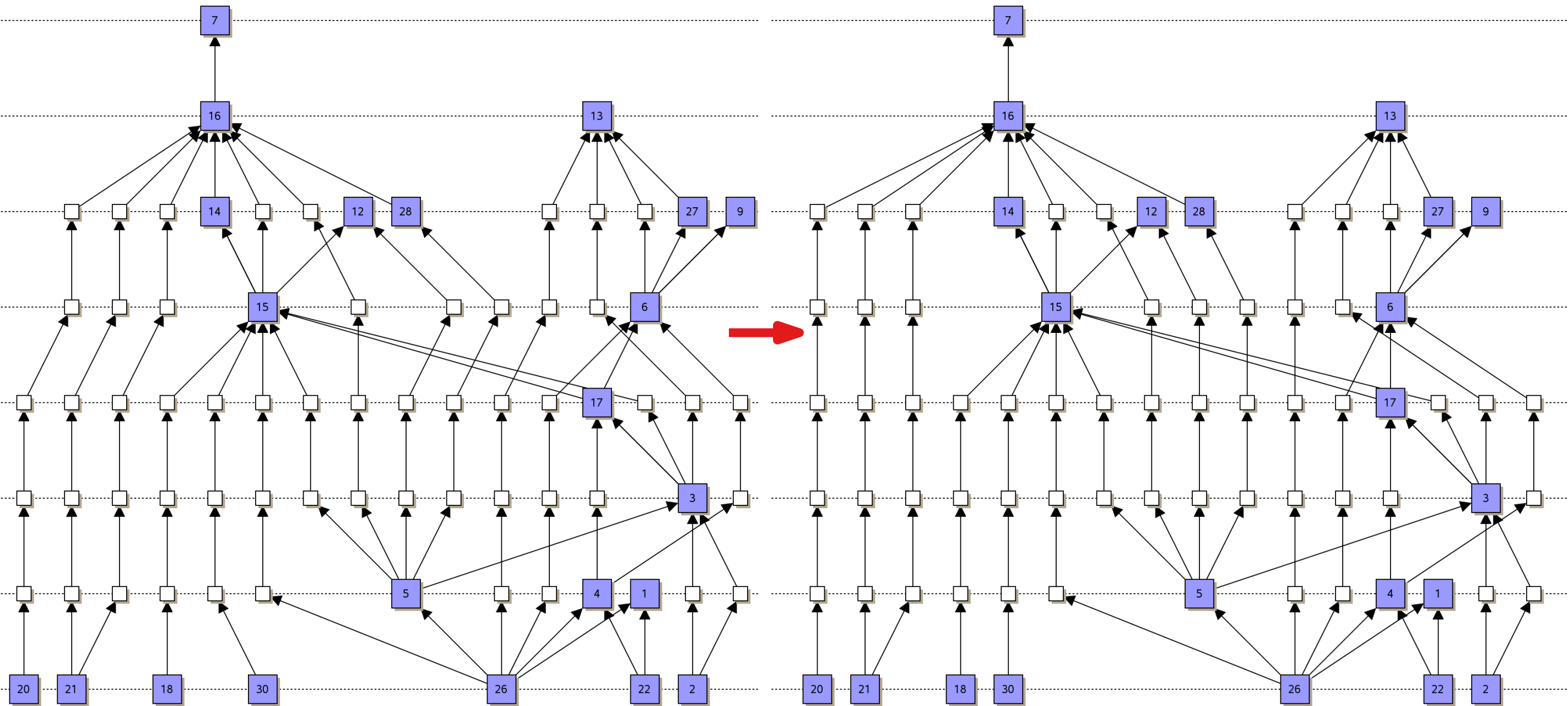


- QP is time-expensive.
- Width can be exponential.

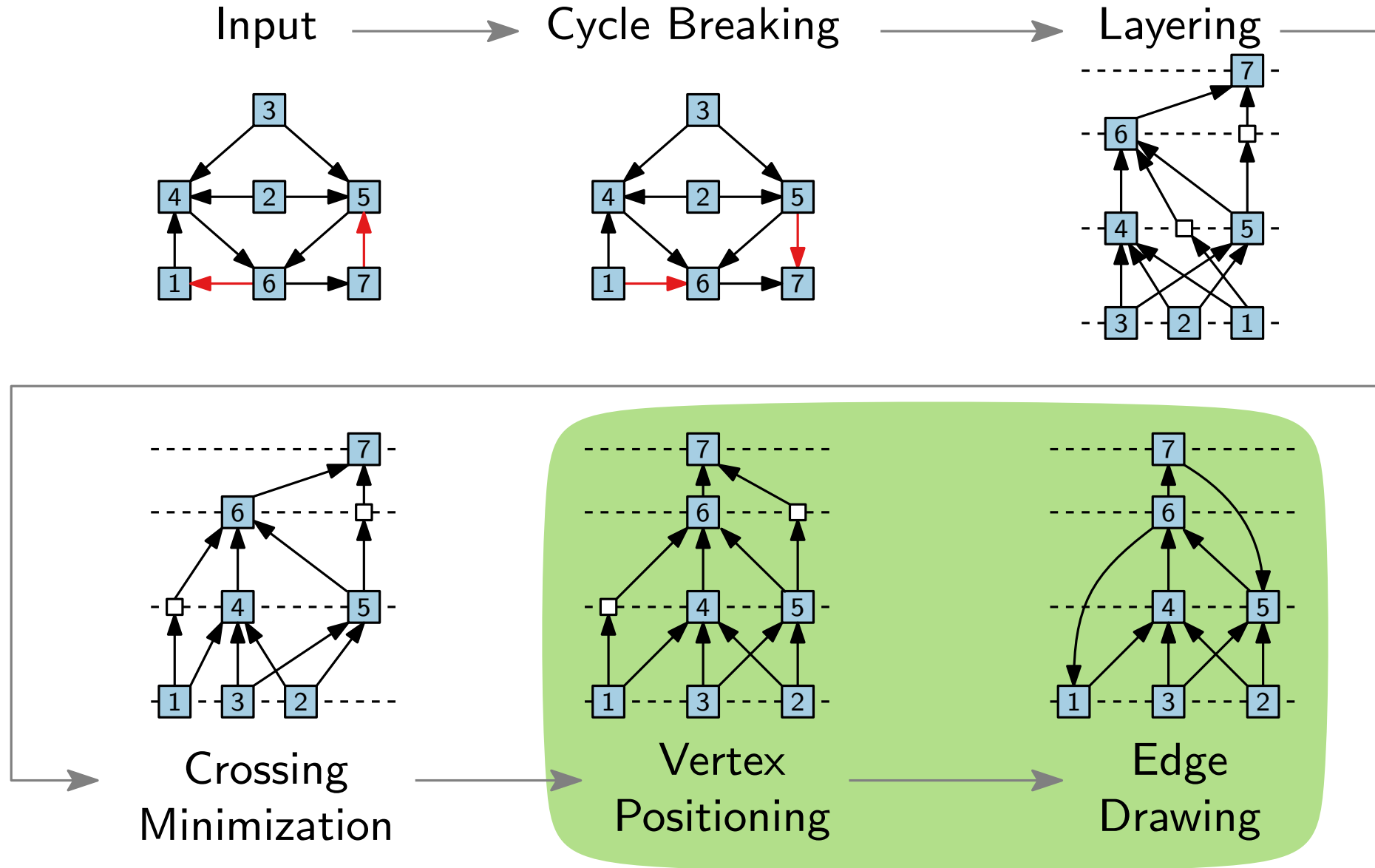
Iterative Heuristic

- Compute an initial layout
- Apply the following steps as long as improvements can be made:
 1. vertex positioning
 2. edge straightening
 3. compactifying the layout (to reduce the width)
- Other algorithms, e.g., the one of Brandes and Köpf
[GD 2002, see also Brandes, Walter, Zink: arXiv 2020]:
 - tries to align vertices vertically
 - does horizontal compaction afterwards
 - linear running time

Example



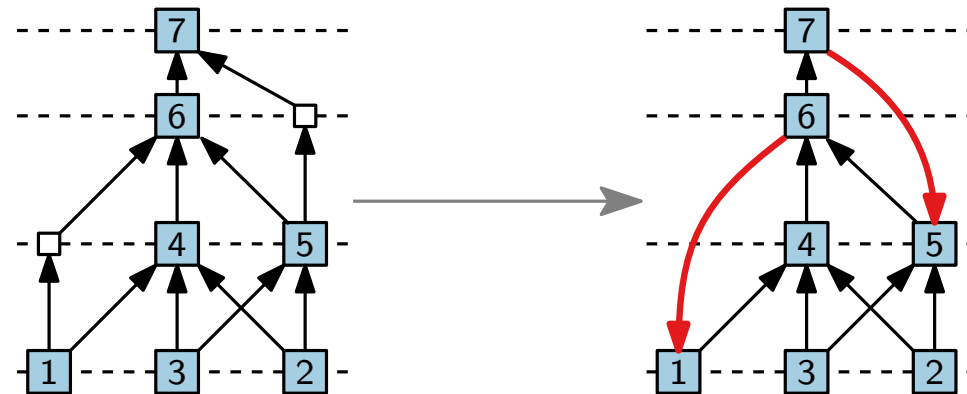
Step 5: Drawing Edges



Step 5: Drawing Edges

Possibility.

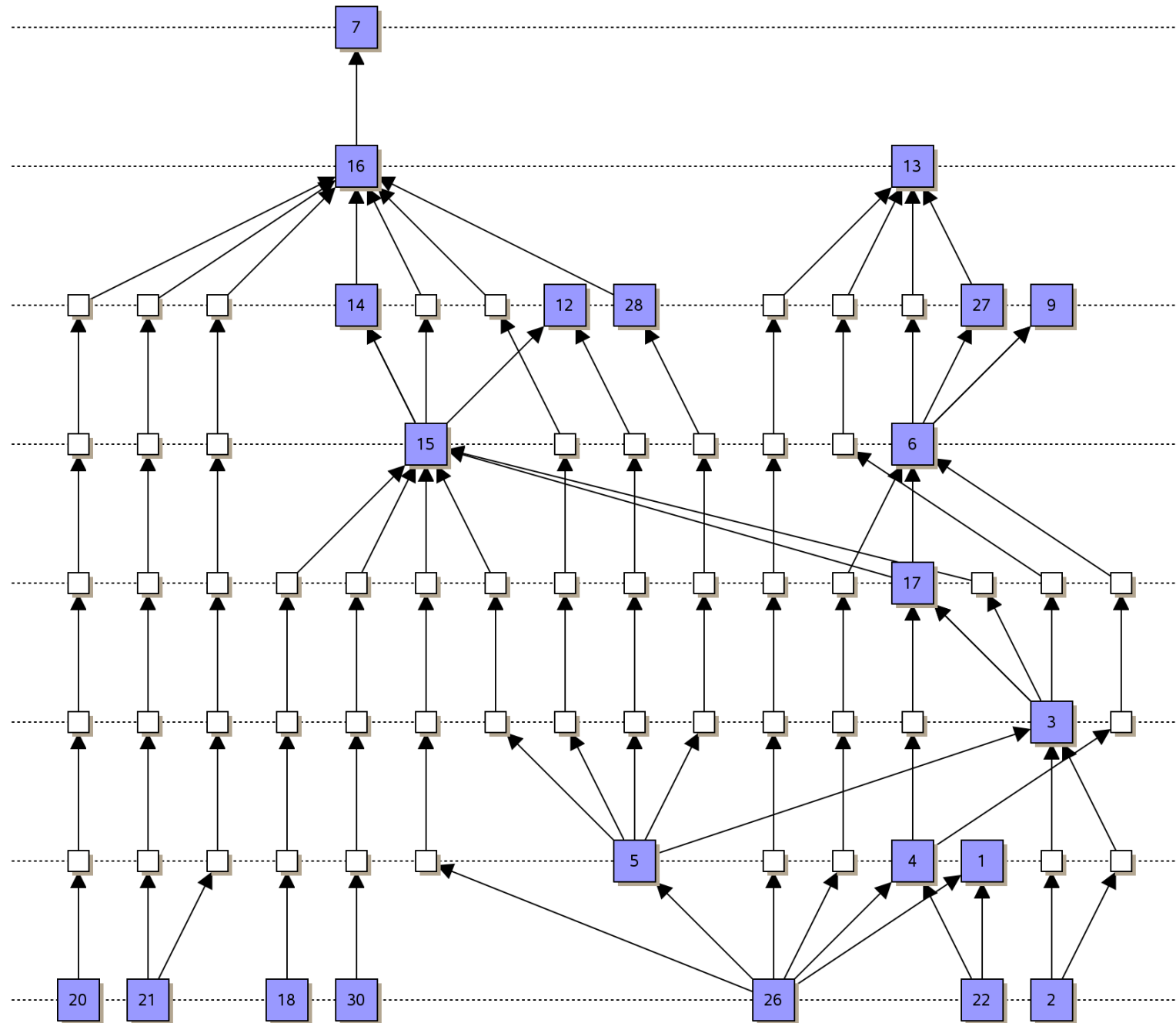
Substitute polylines by Bézier curves.



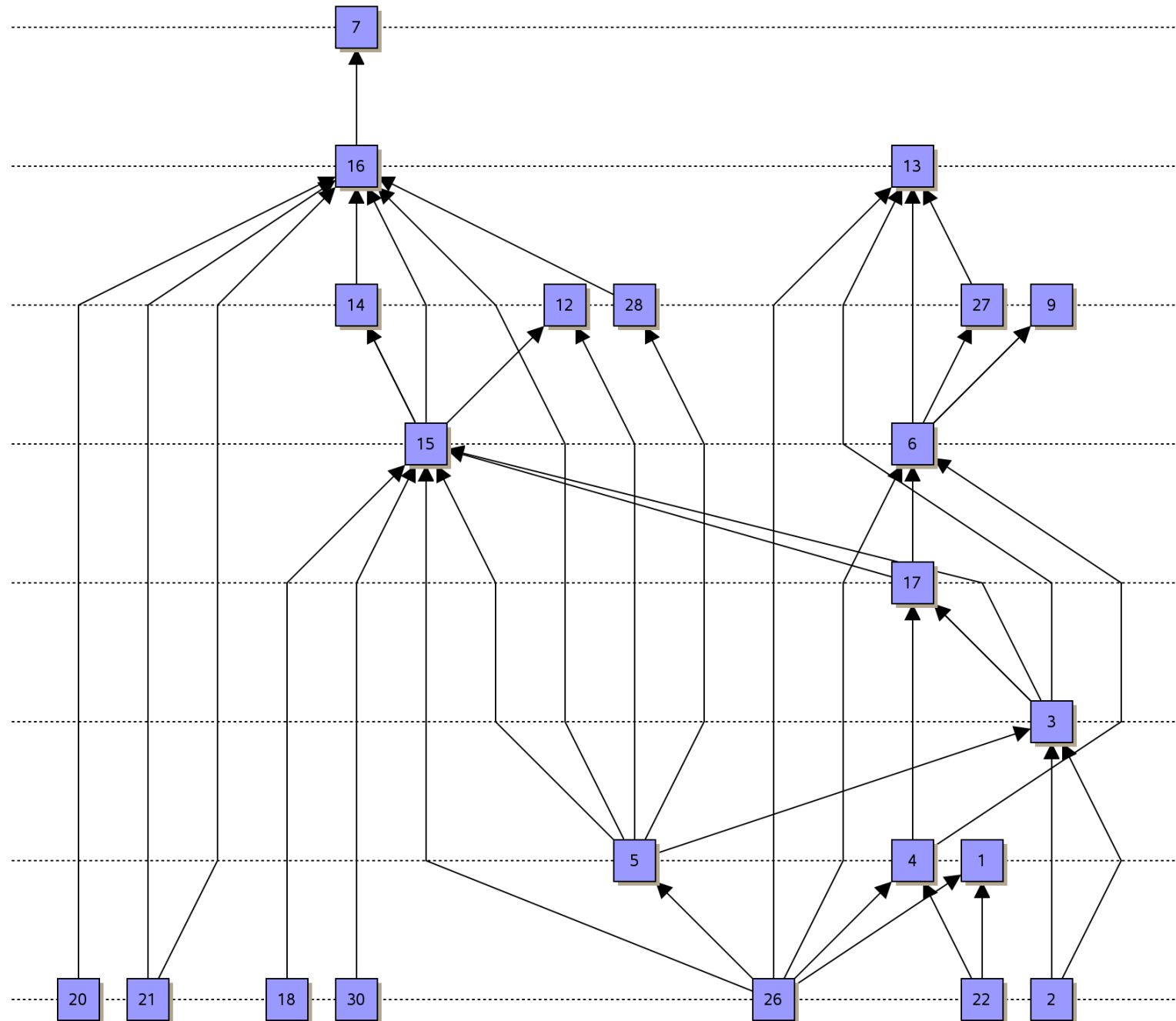
Remark.

Draw reversed edges downwards.

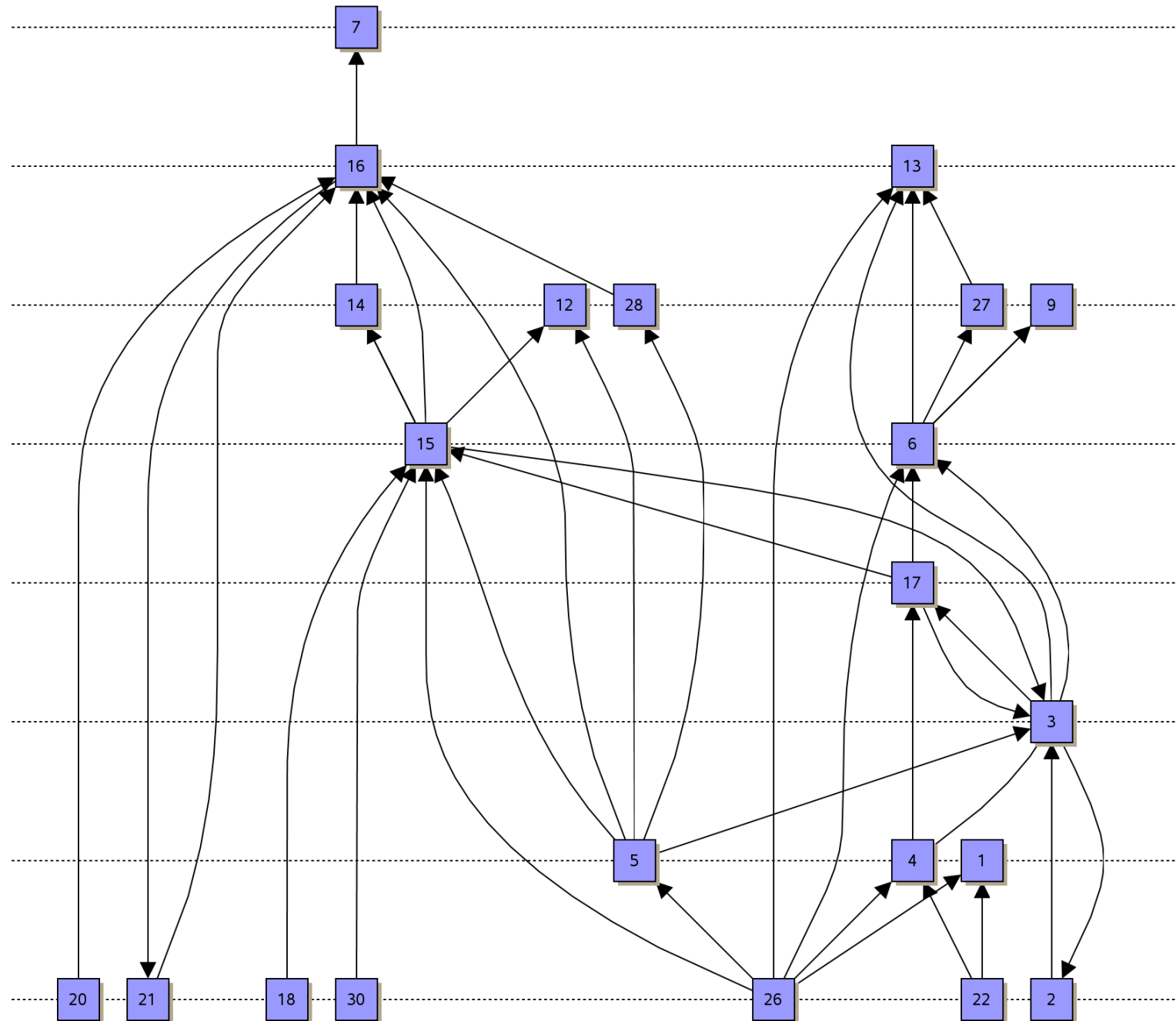
Example



Example

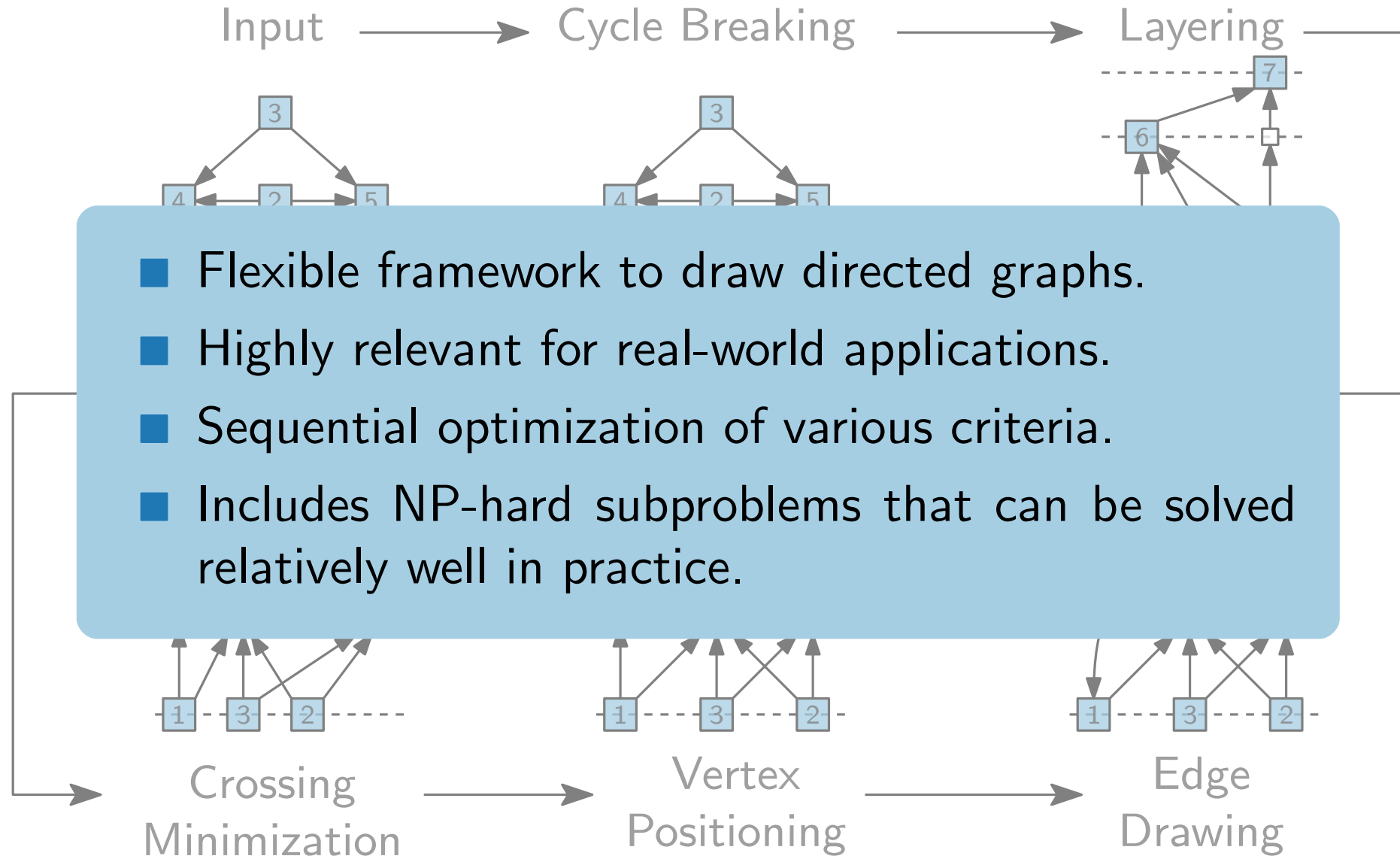


Example



Classical Approach – Sugiyama Framework

[Sugiyama, Tagawa, Toda '81]



Literature

Detailed explanations of steps and proofs in

- [GD Ch. 11] and [DG Ch. 5]

based on

- [Sugiyama, Tagawa, Toda '81]

Methods for visual understanding of hierarchical system structures

and refined with results from

- [Berger, Shor '90] Approximation algorithms for the maximum acyclic subgraph problem
- [Eades, Lin, Smith '93] A fast and effective heuristic for the feedback arc set problem
- [Garey, Johnson '83] Crossing number is NP-complete
- [Eades, Kelly '86] Heuristics for reducing crossings in 2-layered networks.
- [Eades, Whiteside '94] Drawing graphs in two layers
- [Eades, Wormland '94] Edge crossings in drawings of bipartite graphs
- [Jünger, Mutzel '97]

2-Layer Straightline Crossing Minimization: Performance of Exact and Heuristic Algorithms