



Introduction to Informatics for Students from all Faculties

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Lecture 05 Algorithmic Thinking

November 19, 2024

· Lecture L05: Algorithmic Thinking

- 19.11.2024
- **Educational objective:** We introduce algorithms for basic problems like binary search and sorting. We discuss the runtime of algorithms and cover basic data structures.
 - Python Data Structures
 - A Simple Algorithm
 - Binary Search Algorithm
 - Basic Sorting Algorithms

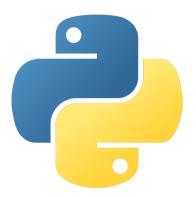
• Exercise Sheet 4 due 26.11.2024

Motivation

- we introduced high-level programming languages that are translated to machine code
- we wrote a first "Hello World" program in C and inspected the machine code generated by the compiler
- we distinguished between compiled languages like C/C++ and interpreted languages like python
- we introduced basics of the **python syntax**

open issues

- how to write programs that solve actual problems?
- what are algorithms and how we can we implement them in high-level languages like python?
- need to develop algorithmic thinking, which is key to understand how a computer (scientist) works



Recap: python **Syntax**

- python programs are stored in text files (typically with extension .py)
- one line in text file = one instruction

key python statements

- assignment (=) used to assign value to a variable
- def used to define a function
- import statement used to import functions from modules
- ► if and else used to conditionally execute instructions
- ► for and while used to repeatedly execute instructions in a loop
- ▶ "blocks" of instructions grouped by indentation level
- python is whitespace-sensitive, i.e. placement of newline, space or tab characters changes semantics
- python enforces meaningful formatting of code, making programs easy to read for humans

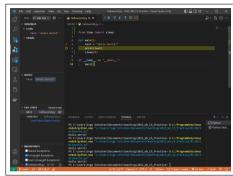
```
import time

def main():
    for i in range(5):
        text = "Hello World!"
        print(text)
        text = 42
        print(text)
        sleep(5)
```

Note that the "grammatical structure" of a (programming) language is called "syntax", which contains the Greek
words "syn" (together) and "taxis" (ordering/composition). The syntax of a language defined keywords and
determines the ordering of characters that constitutes a valid sentence or (program) in a (programming)
language.

Integrated Development Environment (IDE)

- all we need to write python program is text editor and python interpreter (i.e. executable python.exe)
- sufficient for small single-file programs
- what about complex software with hundreds of files and millions of lines in code?
- integrated development environments (IDEs) are specialized tools to support and simplify development of complex software
- IDEs provide advanced functions to edit and format code, semantically highlight/color keywords, compile and/or execute program, and find errors



Open Source IDE Visual Studio Code

definition

An Integrated Development Environment (IDE) is a software that simplifies the programming of computers. It minimally provides functions to edit source code files, compile and/or execute programs, and find errors at compile- and run-time.

Python Data Structures

- all programming languages support basic data types
 - ightharpoonup integer numbers, i.e. 42, -55, 0
 - floating point numbers, i.e. 4.52, 1.567e2, 2.0e 2
 - character types, i.e. "c", "t"
 - string types, i.e. "Lecture"
- python is a dynamically-typed language, i.e. we can assign any type to a variable
- what if we need more complex structures to store data?
 - list of numbers
 - all sentences of a book
 - mapping from numbers to text
 - queue of jobs to be executed in sequence
- python standard library provides complex data types that can hold list, sequences, dictionaries of values

```
import time

def main():
    for i in range(5):
        text = "Hello World!"
    print(text)
    text = 42
    print(text)
    sleep(5)
```

Python Lists

python lists can hold ordered sequence of elements of any type

adding / removing eleemnts

- append allows to append additional values at end of list
- pop can be used to remove and return element at a given index (or at the end)
- remove deletes first occurrence of a value
- we can use zero-based integer indexing to read/write elements at specific position
- slicing operator [start:end:step] can be used to return new list with selected elements
- using append and pop(0) we can use list as queue, where elements are returned in

```
# create list
l = [2.42.120.18.420]
# adding/removing elements
l.append('hello')
l.remove(120)
print(l.pop())
print(l.pop(0))
# index-based access
[[1] = 43]
# elements up to index 2
# (excluding 2)
print([:2])
# elements starting from index 1
# (including 1)
print([[1:])
# initialize list with 42 zero entries
```

Python Tuples

- by appending, assigning or removing elements,
 python lists can dynamically change their size and
 elements can change during lifetime of list
- requires complex implementation that makes some operations relatively slow
- for fixed-size ordered sequences that cannot change, we can use python tuples
- indexing and slicing works the same as for lists
- elements cannot change and size of tuple cannot grow or shrink
- we can use + operator to concatenate two tuples, returning a new tuple

```
# create tuple
t = (2, 42, 120, 189, 420)
# index-hased access
print(t[1])
# slicina
print(t[:2])
print(t[1:])
# NOT VALID
t[1] = 43
# returns new tuple with additional
# elements
t2 = t + (4.5.6)
print(t2)
```

Python Sets

- lists and tuples are ordered sequences
- checking whether an element is in a list/tuple requires to test all elements → naive linear search
- for unordered collection of objects without duplicates we can use python set
- useful to eliminate duplicate elements and quickly test for membership
- python sets are unordered and thus do not support indexing or slicing

```
# create set
s = {2, 'hello', 120, 42, 189, 420}
print(s)
# check membership
print('hello' in s)
# add element
s.add(32)
# remove element
s.remove('hello')
# NOT VALID
s[1]
s[:4]
```

Python Dictionaries

- we often need an associative mapping that maps unique keys to values
 - string/number as unique identifier (key)
 - arbitrary data of record (value)
- useful to quickly find data that are stored under a given key
- we can read/write entries using index syntax similar to lists
- but: index does not need to be an integer
- reverse lookup (i.e. find key(s) for a given value) not supported

```
# create dictionary
d = { 'hello': 'Hallo'.
      'world': 'Welt'.
      'teacher': ['Ingo', 'Scholtes'
# check membership of key
print('hello' in d)
# access value of given kev
print(d['hello'], d['teacher'])
# assign value to (new) key
d['audience'] = 'Studierende'
```

Practice Session

- we show how to install the Open Source python distribution Anaconda
- we use the **integrated development environment** (IDE) Visual Studio Code to write and execute a simple python program
- we use VS Code to rename variables and refactor code
- we use the **debugger of Visual Studio Code** for a step-wise execution of python statements
- we demonstrate lists, tuples, sets, and dictionaries in python

```
import time
def main():
   text = "Hello World!"
   print(text)
   text = 42
   print(text)
   sleep(5)
```

practice session

see directory 05-01 in gitlab repository at

What is an algorithm?

definition → LO1 - Motivation

An algorithm is a sequence of precisely defined (mathematical) instructions that must be executed to solve a given problem.

- algorithm takes a (possibly empty) input and produces after a finite number of steps – a desired output
- expressing an algorithm in terms of a programming language allows us to implement it on a computer

Example: pecil-and-paper algorithm to add two numbers

```
step 1 start at right-most position
```

step 4
$$\,$$
 for sums ≥ 10 additionally carry over 1 to position on the left

step 5 move one position to left and go to step 2

Group Exercise 05-01

Assume that we want to implement the pen-and-pencil algorithm to add two decimal numbers with an arbitrary number of digits. Specify a reasonable input and output of this algorithm.
Develop a python function add that implements the pen-and-pencil algorithm, using control structures like while, for, if as well as a python list

- In lecture LO2 we have seen how we can use digital logics to implement the addition of two binary numbers in terms of hardware. Thanks to the fact that this operation is implemented in the ALU of the CPU, we can directly add two 32 or 64 bit numbers by a single machine instruction (e.g. ADD). This implies that we can use the ADD operator + in high-level languages like python, which is directly mapped to this machine instruction.
- As an exercise, we pretend that there was no such operation that allows to add numbers with more than one
 digit. Let us implement the pen-and-pencil algorithm to add decimal numbers in python.
- Note that such an algorithm can still be useful if we want to add two numbers that cannot be represented by 64 bits or less, which may not be supported by the ALU of a common CPUs.

Practice Session

- we use lists to implement the algorithm developed in the previous group exercise in python
- we test our algorithm with different inputs



practice session

see directory 05-02 in gitlab repository at

 $\rightarrow \texttt{https://gitlab2.informatik.uni-wuerzburg.de/ml4nets_notebooks/2024_wise_infhaf_notebooks/2024_wi$

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Search problems

- we are frequently confronted with standard problems that can be solved by well-understood standard algorithms
- exemplary standard problem: search for an object

search problems

Search problems are a class of problems that seek to quickly find a given object within a certain data structure.

- examples
 - problem 1: search name "Turing" in arbitrary list of 10,000 names
 - problem 2: search name "Turing" in a phonebook
- optimal solution to the search problem depends on prior knowledge on the data structure



 $\underline{\mathsf{image}\;\mathsf{credit:}}\;\mathsf{DALL\text{-}E}\;\mathsf{generated}\;\mathsf{image},\mathsf{prompt}\;\mathsf{``needle}\;\mathsf{in}\;\mathsf{a}\;\mathsf{haystack''}$

Searching in sorted data

- we can sort a list of objects whenever for pair of objects a and b we can determine whether a ≥ b
- we assume that we search object x in a list sorted in ascending order, i.e. list / where for index i we have

$$I[i+1] \geq I[i]$$

▶ naive algorithm checks x == I[i] for index i = 0, 1, ...

binary search algorithm

- test if x is larger/smaller/equal than middle element c
- ightharpoonup if x == c return object
- if x > c repeat search in elements right of c
- if x < c repeat search in elements left of c
- binary search is example for divide-and-conquer algorithm
- both algorithms give correct result, but which one is "better"?

searching for x = 17

sorted list

2	3	5	11	17	23	47
---	---	---	----	----	----	----

step one

2	3	5	11	17	23	47
-	-	-	-	17	23	47

step two

				-		
-	-	-	-	17	23	47
-	-	-	-	17	-	-

step three



found x = 17!

Complexity of binary search

we can evaluate algorithms in terms of **computational complexity**, i.e. we count how many steps they **maximally require** to produce the correct output for a given input?

example 1: how many steps do we need in list / with 32 objects

naive (lir	near) search algorithm	binary s	earch algorithm
step	tested element	step	tested element
1	0	1	16
2	1	2	8 or 24
3	2	3	4, 12, 20, or 28
		4	2, 6, 10, 14, 18, 22, 26, or 30
32	31	5	0, 1, 3, 5, 7, 8, 9, 11, 13, 15, 16, 17, or 31

requires 32 steps for 32 objects

requires 5 steps for 32 objects

Complexity of binary search

we can evaluate algorithms in terms of **computational complexity**, i.e. we count how many steps they **maximally require** to produce the correct output for a given input?

example 2: how many steps do we need in list / with 64 objects

naive (li	near) search algori
step	tested element
1	0
2	1
3	2
4	3
63	64

binary s	earch algorithm
step	tested element
1	32
2	16 or 48
3	8, 24, 40, or 56
4	4, 12, 20, 28, or 60
5	2, 6, 10, 14, 18, 22, or 62
6	0, 1, 3, 5, 7, 9, 11, 13, 15 or 63

requires 64 steps for 64 objects

requires 6 steps for 64 objects

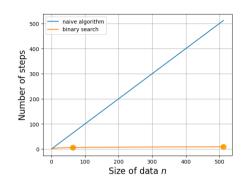
Linear vs. logarithmic complexity

- naive search algorithm requires one step for each each entry in the input list, i.e. runtime is proportional to the input size
- we say an algorithm has linear complexity if for input with size n it requires at most

$$c + x \cdot n$$

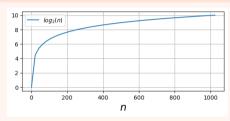
steps for some numbers c and x

- thanks to sorted input, binary search algorithm requires less than linear number of steps
- how does number of steps grow as we increase input size n?



Logarithms

Logarithm

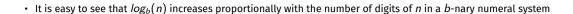


For a number n, the logarithm $log_b(n)$ with base b of n is the number x such that $b^x = n$.

logarithms base b=2

logarithms base
$$b=10$$

n	binary number	$log_2(n)$	n	$log_{10}(n)$
$2^0 = 1$	1	0	$10^{0} = 1$	0
$2^1 = 2$	10	1	$10^1=10$	1
$2^2 = 4$	100	2	$10^2 = 100$	2
$2^3 = 8$	1000	3	$10^3 = 1000$	3
$2^4 = 16$	10000	4	$10^4 = 10000$	4



Complexity of search algorithms

we can evaluate algorithms in terms of **computational complexity**, i.e. we count how many steps they **maximally require** to produce the correct output for a given **input of size** *n*?

how many steps do we need in list I with n = 64 objects

naive (li	near) search algori
step	tested element
1	0
2	1
3	2
4	3
64	63

binary search algorithm					
step	tested element				
1	32				
2	16 or 48				
3	8, 24, 40, or 56				
4	4, 12, 20, 28, or 60				
5	2, 6, 10, 14, 18, 22, or 62				
6	0, 1, 3, 5, 7, 9, 11, 13, 15 or 63				

requires *n* steps for *n* objects

requires $log_2 n$ steps for n objects

Sorting problem

- for binary search, we assumed that the list of objects is sorted
- to sort objects we must be able to compare **them,** i.e. for each pair a, b we must be able to determine a > b
- how can we compare pairs of
 - numbers.
 - words.
 - books.
 - emoiis?

sorting problem

The sorting problem refers to the problem of sorting a list of pairwise comparable objects in ascending or descending order.

in the following, we consider the sorting problem for a list of integer numbers

input:

7	2	47	23	5	11
---	---	----	----	---	----

desired output:

2	5	7	11	23	47
---	---	---	----	----	----

BubbleSort algorithm

- simple idea: repeatedly compare pairs of numbers and swap them if they are in the wrong order
- with each swap ...
 - larger numbers progressively move to right
 - smaller numbers progressively move to left
- in each pass of the algorithm, we must compare all subsequent pairs of numbers in the list
- if we have zero swaps during a pass, we know that the list is sorted!
- in the example, we needed
 - ▶ $4 \cdot 5 = 20$ comparisons
 - ightharpoonup 4 + 2 + 1 = 7 swaps
- how many comparisons do we need in best/worst case?

third pass

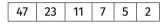
2	7	5	11	23	47
2	7	5	11	23	47
2	7	5	11	23	47
2	5	7	11	23	47
2	5	7	11	23	47
2	5	7	11	23	47
2	5	7	11	23	47
2	5	7	11	23	47
2	5	7	11	23	47
2	5	7	11	23	47

5 comparisons, 1 swap

Worst-case complexity of BubbleSort

worst-case runtime

For an input list sorted in reverse order BubbleSort algorithm requires n passes with n-1 comparisons each.



$$n = 6$$

$$n \cdot (n-1) = 6 \cdot 5 = 30$$
 comparisons

best-case runtime

For an input list that is already sorted ${\tt BubbleSort}$ algorithm requires a single pass with n-1 comparisons.

$$n = 6$$

$$n-1=5$$
 comparisons

Linear vs. polynomial complexity

for input list with n elements, BubbleSort has worst-case runtime of

$$n\cdot (n-1)=n^2-n$$

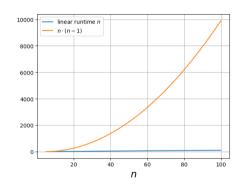
i.e. number of required steps grows as **second power** (i.e. square) of input size *n*

we call expressions of the form

$$a_k \cdot n^k + a_{k-1} \cdot n^{k-1} + \dots a_0 \cdot n^0$$

polynomial

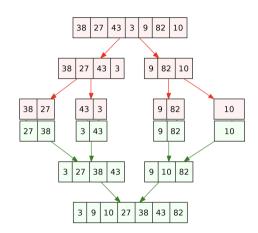
for polynomials with power larger than one, runtime grows over-proportionally with input size



linear vs. polynomial growth of complexity

MergeSort

- can we sort a list faster than BubbleSort?
- **assume that we have two already sorted lists** l_1 and l_2
- ▶ in $n = n_1 + n_2$ steps we can **merge** l_1 **and** l_2 into a new sorted list l
- we can apply divide-and-conquer idea behind binary search to sorting
- phase 1: repeatedly split input until we are left with lists with one element (which are already sorted)
- phase 2: repeatedly merge increasingly large (sorted) lists until full list is sorted



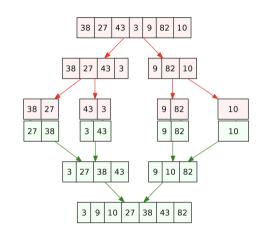
Complexity of MergeSort

- assume that for a list with n elements MergeSort takes T(n) steps
- each split/merge then requires

$$2 \cdot T(\frac{n}{2}) + n$$

- ightharpoonup starting with T(1)=0 we have
 - $T(2) = 2 \cdot T(1) + 2 = 2$
 - $T(4) = 2 \cdot T(2) + 4 = 2 \cdot 2 + 4 = 8$
 - $T(8) = 2 \cdot T(4) + 8 = 2 \cdot 8 + 8 = 24$
 - $T(16) = 2 \cdot T(8) + 16 = 2 \cdot 24 + 16 = 64$
 - ▶ ..
- we can calculate runtime of MergeSort as

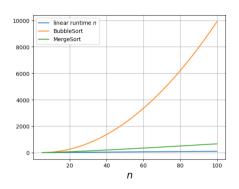
$$T(n) \approx n \cdot \log_2(n)$$



Complexity of sorting?

- with BubbleSort we can sort n numbers in n-1 steps in best case and $n \cdot (n-1)$ in worst case
- ► MergeSort improves worst-case complexity of BubbleSort from n^2 to $n \log_2(n)$
- **on average** MergeSort requires $n \log_2(n)$ steps
- ▶ to sort n objects based on pairwise comparisons, there is no algorithm exist that requires less than $n \log_2(n)$ steps on average
- <u>but:</u> there are specialized algorithms to **sort** *n* **integer numbers in a fixed range** with linear runtime → self-study

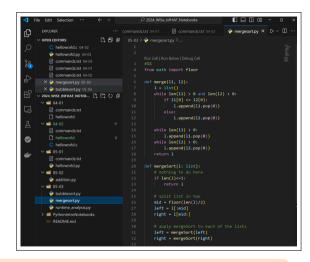
questions



worst-case complexity of MergeSort vs.
BubbleSort

Practice Session

- we implement BubbleSort in python
- we implement the divide-and-conquer method MergeSort
- we study the runtime of both algorithms in increasingly large input lists



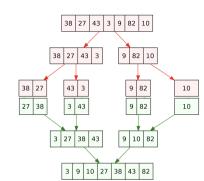
practice session

see directory 05-03 in gitlab repository at

 $\rightarrow \texttt{https://gitlab2.informatik.uni-wuerzburg.de/ml4nets_notebooks/2024_wise_infhaf_notebooks/2024_wi$

In summary

- we covered basic python data structures like sets, tuples, lists, and dictionaries
- we introduced basic algorithms for standard computational problems like searching and sorting
- we evaluated the computational complexity of sort and search algorithms
- we highlighted the difference between logarithmic, linear, and polynomial runtime



Self-study questions

- 1. Explain the differences between a set, a tuple, a dictionary and a list in python.
- 2. Give a formulation of the Pen-And-Pencil algorithm to add two numbers in python and explain it in your own words.
- 3. Extend the Pen-And-Pencil algorithm from the group exercise such that it can add numbers given as sequences of digits in an arbitrary *k*-nary numeral system.
- 4. Give a formulation of the Binary Search algorithm in python and explain it in your own words.
- 5. Could we generalize the Binary Search algorithm such that in each step we split the list into three equally large parts, which would lead to a runtime $log_3(n)$?
- 6. Give a formulation of the BubbleSort algorithm in python and explain it in your own words.
- 7. Give an example for an input for which the BubbleSort algorithm performs the maximum/minimum number of comparisons.
- 8. Count the number of swaps in an input list with *n* elements, where BubbleSort performs the maximum number of comparisons.
- 9. Give a formulation of the MergeSort algorithm in python and explain it in your own words.
- 10. Investigate the BucketSort algorithm for integers in a fixed range and explain why it takes less than $n \log_2 n$ steps on average.

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Literature

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