



ALGORITHMS IN AI & DATA SCIENCE 1 (AKIDS 1)

Adversarial Search (aka Game Playing) Prof. Dr. Goran Glavaš

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#### Content

- Introduction
- Minimax
- Heuristic Minimax
- Alpha-Beta Pruning

Based on the materials from Prof. Dr. Jan Šnajder: <u>https://www.fer.unizg.hr/ download/repository/AI-4-GamePlaying.pdf</u>

## Game Playing: Adversarial Search

- Games also represent a state space search problem, but with the difference that there is an adversary
- In each game state one must make an **optimal decision** about which move to make next, i.e., one must find an **optimal strategy**



## Game Playing: Adversarial Search

- We will focus on games that are
  - Two-player: that is, one adversary
  - Deterministic: no stochasticity/randomness like dice rolling
  - Complete information: both players have complete knowledge of the game (e.g., know the rules, see the full board, etc.)
  - Zero sum: one player winning (losing) means the other player losing (winning)



## Adversarial Search (Game): Formalization

Game (Adversarial Search Problem)

**Game** is a state space search problem quadruple ( $s_0$ , *succ*, *terminal*, *utility*) where  $s_0$  is the **initial game state**, *succ* :  $S \rightarrow \wp(S)$  defines legal game moves (i.e., **transitions** between states), *terminal*:  $S \rightarrow \{True, False\}$  is a predicate that tests if a state is a **game end state** and *utility*:  $S \rightarrow \mathbb{R}$  is a **payoff function** that assigns numeric values to terminal (end) states of the game for the player(s).

- For example, in chess there are three possible outcomes
  - Win, loss, draw
  - utility(s) ∈ {+1, 0, −1}
- Initial state s<sub>0</sub> and successor function *succ* implicitly define the **game tree**

#### Game Tree



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- Let's call the players MAX (*"us"*, the computer) and MIN (the opponent)
- MAX tries to maximize its win, MIN tries to minimize MAX's win
- Players take turn: nodes at even depths MAX; odd depth MIN



## Minimax method: Optimal Strategy

- Q: What is the optimal game-playing strategy for MAX?
  - Q: Is it to take a step in the direction of its own maximal utility?



## Minimax method: Optimal Strategy

- MAX's optimal strategy is the one that ensures the highest win for MAX, but assuming that MIN uses the same strategy
- Each player chooses a strategy so as to minimize the maximum loss

 To determine the optimal strategy of a player who's turn is next, we compute the minimax value of the root node



## Minimax value



 $m(s_0) = max(min(3, 2, 1),$ min(1, 0, 2),min(-5, 3, 1))



## Minimax value: algorithm

```
max_val(s)
if terminal(s)
return utility(s)
m = -inf
for t in succ(s)
m = max(m, min_val(t))
return m
```

```
min_val(s)
if terminal(s)
return utility(s)
m = +inf
for t in succ(s)
m = min(m, max_val(t))
return m
```

 Note that this is essentially a <u>DFS</u> implemented via two mutually recursive functions

m(s) -   

$$\begin{bmatrix}
utility(s) & if terminal(s) \\
max_{t \in succ(s)} & m(t) & if s is a MAX node \\
min_{t \in succ(s)} & m(t) & if s is a MIN node
\end{bmatrix}$$



## Minimax: remarks

- In practice, the opponent's strategy is unknown
  - Most probably different from that of MAX player
  - The opponent's moves therefore cannot be predicted perfectly (otherwise the game would be boring anyways)
- Because of this in order to make the optimal move, in each turn the players need to re-compute their optimal strategy, starting from the current position as the root of the game tree
- Minimax is a **depth-first search** 
  - Q: Time and space complexity?

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## Imperfect decisions

- Minimax is a **DFS** that makes its time complexity **O(b<sup>m</sup>)** 
  - b: the branching factor of the game
  - m: total number of moves needed for the game to end (reach terminal state)
- For most games **b** and **m** are **big enough** that we **don't have the time** to search through the complete game tree all the way to terminals
- We need to cut the search off at a certain depth d (we can afford at most runtime of O(b<sup>d</sup>))
  - For states at depth d, we need to **estimate** the utility with a **heuristic**
  - E.g., in chess, a simple heuristic can be the sum of "worth" of remaining pieces (vs. the value of remaining pieces of the opponent)

## Heuristic minimax

```
max val(s, d, h)
  if terminal(s)
    return utility(S)
  if d = 0
    return h(S)
  m = -inf
  for t in succ(s)
    m = max(m, min val(t, d-1, h))
  return m
min val(s, d, h)
  if terminal(s)
    return utility(s)
  if d = 0
    return h(S)
  m = +inf
  for t in succ(s)
    m = max(m, max_val(t, d-1, h))
  return m
```

- Game heuristics typically combinations of different features h(s) = w<sub>1</sub>\*x<sub>1</sub>(s) + w<sub>2</sub>\*x<sub>2</sub>(s) + ... + w<sub>n</sub>\*x<sub>n</sub>(s)
- Example features for chess:
  - x<sub>1</sub> = value of all my remaining pieces
  - x<sub>2</sub> = value of all opponent's remaining pieces
  - $x_3 = do I have a queen left$
  - ...
  - Weights w<sub>1</sub>, w<sub>2</sub>, ..., w<sub>n</sub> indicate how important each feature is



- Q: What is the outcome of the game if both players can explore all nodes (until game terminal nodels)?
- Q: What is the outcome if each player can search only two levels deep?





- The two players are **unlikely** to use the same heuristic
- Q: what is the end state (outcome) if each player searches two levels deep but use different heuristics (MAX blue, MIN red)?
  - Important: both heuristics are estimates of utility for MAX!
- Depends on whose turn it is <sup>(C)</sup>



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- Number of states increases exponentially with the number of turns
- We can effectively cut this number in half with alpha-beta pruning
- Q: Can we compute the minimax value of a node without traversing the whole game (sub)tree?
- Yes!

```
m(s) = max( min(3, 2, 1),
min(1, X, X),
min(-5, X, X)) = 1
```



# Alpha-beta pruning

- We prune whenever we're certain that unexplored moves cannot change the result of the max/min operation
  - Cannot lead to a move that is better than the best found so far
- If pruning below the MIN node (but when root is MAX node): alpha pruning
- If pruning below the MAX node (but when root is MIN node): beta pruning



## Minimax with alpha-beta pruning

#### beta pruning

```
max_val(s, α, β)
if terminal(s)
    return utility(s)

m = α
for t in succ(s)
    m = max(m, min_val(t, m, β))
    if m ≥ β
    return β
return m
```

#### alpha pruning

```
min_val(s, α, β)
if terminal(s)
    return utility(s)

m = β
for t in succ(s)
    m = min(m, max_val(t, α, m))
    if m ≤ α
    return α
return m
```

**Initially:** max\_val(s<sub>0</sub>,-inf, +inf)

## Alpha-beta pruning example



## Alpha-beta pruning example



## Alpha-beta pruning example



- Game playing is a search problem in which opposing players take turns (i.e., we have an adversary)
- Minimax algorithm finds an optimal strategy that minimizes the maximum expected loss that an opponent can inflict
- In real games it is impossible to search through the complete game tree, thus we cut the search off at a certain depth and use a heuristic function to estimate the values of game states
  - Different players use different heuristic functions
  - The opponent's heuristic is in principle unknown
- Alpha-beta pruning reduces the number of nodes to traverse

### Questions?

