

Exercise Sheet #4

Advanced Algorithms (WS 2023/24)

Exercise 1 – Flow problem with vertex capacities

Let $G = (V, E)$ be a flow network where the capacity function c is not only defined on the edges but also on the vertices. More precisely, for an s - t flow f in G there is a capacity constraint $\sum_{u \in V} f(u, v) = \sum_{u \in V} f(v, u) \leq c(v)$ for each $v \in V \setminus \{s, t\}$.

Show that a maximum flow problem on a flow network with vertex and edge capacities can be reduced to a maximum flow problem in a normal flow network G' , i.e., without vertex capacities. How many more vertices and edges does G' require precisely?

3 Points

Exercise 2 – Implementing the generic push-relabel algorithm

In the lecture we have only outlined how the methods PUSH and RELABEL work and not discussed how we actually find the next applicable operation at all.

- a) Show how to implement PUSH in $\mathcal{O}(1)$ time, RELABEL in $\mathcal{O}(|V|)$, and how to select an applicable operation in $\mathcal{O}(1)$ time. In particular, what data structure would you use to maintain overflowing vertices?
6 Points
- b) Using these implementations, what is the resulting runtime of the push-relabel algorithm?
1 Point

Exercise 3 – Push–relabel for integer flow network

Show that if each edge in a flow network has integral capacity, then the push–relabel algorithm finds an integral maximum flow. **2 Points**

Exercise 4 – Maximum bipartite matching

A *matching* in a graph $G = (V, E)$ is a subset of edges $M \subseteq E$ such that for all $v \in V$, at most one edge of M is incident to v . A *maximum matching* is a matching with maximum cardinality, that is, a matching M such that for any matching M' , we have $|M| \geq |M'|$. A *bipartite* graph $G = (L, R, E)$ has two vertex sets L and R and each edge in E connects a vertex in L to a vertex in R .

- a) Explain how we can transform the problem of finding a maximum matching in a bipartite graph $G = (L, R, E)$ into a maximum flow problem on a flow network N . **4 Points**
- b) Prove that the cardinality of a maximum matching M in G equals the value of a maximum integer flow f in your constructed flow network N . **4 Points**

Please hand in your solutions on Wuecampus until the beginning of the next lecture, that is 14:15 on Wednesday, November 29.