Chair of Computer Science I
Algorithms and Complexity
University of Würzburg

Würzburg, December 23, 2022

Prof. Dr. Alexander Wolff
M.Sc. Marie Diana Sieper

Homework Assignment #8
Approximation Algorithms (Winter Semester 2022/23)

Exercise 1 – Greedy for the KNAPSACK problem

Consider the following greedy approach for the KNAPSACK problem: Iterate over the objects in descending order according to the ratio between profit and size, i. e. $\frac{\text{profit}(a_1)}{\text{size}(a_1)} \geq \ldots \geq \frac{\text{profit}(a_n)}{\text{size}(a_n)}$. Add the current object to the knapsack if there is enough space left.

Show that this approach has an arbitrarily bad approximation ratio.

[5 points]

Exercise 2 – 1/2-Approximation for the KNAPSACK problem

Consider the following algorithm for the KNAPSACK problem: Let $a_1, \ldots, a_n$ be the objects in descending order according to the ratio between profit and size. Assume that each individual object fits in the knapsack. Now determine a smallest index k such that the total size of the first k objects exceeds the capacity of the knapsack. Finally, return the better of the two solutions $\{a_1, \ldots, a_{k-1}\}$ and $\{a_k\}$.

Show that this algorithm is a 1/2-approximation.

[7 points]

Exercise 3 – MINIMUM MAKESPAN SCHEDULING on Two Machines

Consider the following problem: Given n jobs with runtimes $t_1, \ldots, t_n \in \mathbb{N}$, distribute them over two machines such that the total makespan is minimal. More formally, we search for a partitioning of $T := \{t_1, \ldots, t_n\}$ into two sets M_1 and M_2 such that $t(M_1, M_2) := \max\{\sum M_1, \sum M_2\}$ is minimized.

Develop an FPTAS for this problem.

[8 points]

Suggestion: First, develop an exact algorithm with pseudo-polynomial runtime. Then, scale and round accordingly.

Please hand in your solutions via WueCampus until 12:00 on Friday, January 13.