

## Homework Assignment #6

### Approximation Algorithms (Winter Semester 2022/23)

#### Exercise 1 – METRIC-k-CLUSTER

Let  $G = (V, E)$  be a complete graph with edge weights  $c: E \rightarrow \mathbb{Q}_{\geq 0}$  that satisfy the triangle inequality. Let  $k$  be a positive integer. We want to find a partition of  $V$  into  $k$  sets of vertices  $V_1, \dots, V_k$ , called *clusters*, such that the weight of the most expensive intra-cluster edge is minimized. In other words, we have to minimize

$$\max_{1 \leq i \leq k, u, v \in V_i} c(u, v).$$

- Devise a factor-2 approximation algorithm for this problem. [7 points]
- Show that under the assumption  $P \neq NP$ , there exists no factor- $(2 - \varepsilon)$  approximation algorithm for this problem, where  $\varepsilon > 0$ .

*Suggestion:* Use the hardness of the graph coloring problem: Given a graph  $G = (V, E)$  and a parameter  $k$ , can the vertices in  $V$  be colored with at most  $k$  colors such that no two adjacent vertices share a color? [6 points]

#### Exercise 2 – Greedy for METRIC-k-CENTER

Consider the following greedy algorithm for METRIC-k-CENTER.

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**Algorithm 1:** Greedy-Metric-k-Center( $G = (V, E; c), k$ )

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- 1 Pick an arbitrary vertex  $v \in V$
- 2  $S \leftarrow \{v\}$
- 3 **while**  $|S| < k$  **do**
- 4     $u \leftarrow$  Vertex with  $c(u, S) = \max_{v \in V} c(v, S)$
- 5     $S \leftarrow S \cup \{u\}$

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Show that this algorithm is a factor-2 approximation for METRIC-k-CENTER.

[7 points]