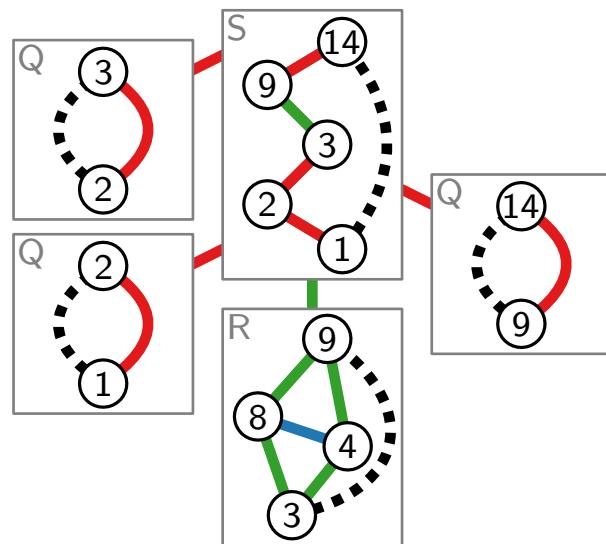


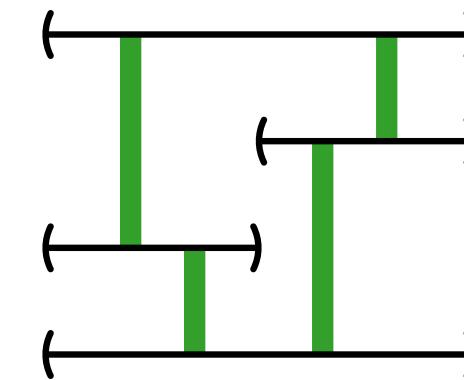
# Visualization of Graphs

## Lecture 9: Partial Visibility Representation Extension



### Part I: Problem Definition

Alexander Wolff



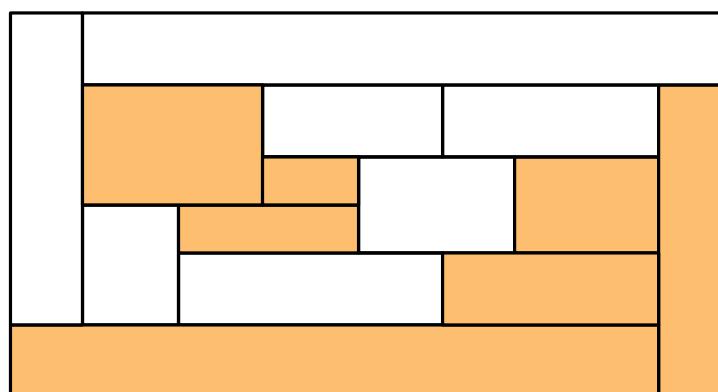
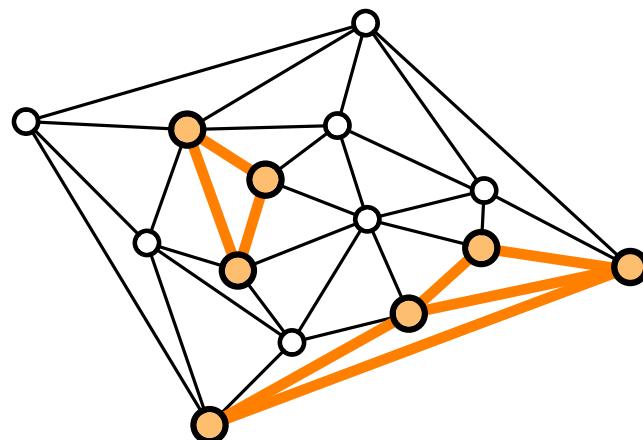
# Partial Representation Extension Problem

Let  $G = (V, E)$  be a graph.

Let  $V' \subseteq V$  and  $H = G[V']$

Let  $\Gamma_H$  be a representation of  $H$ .

Find a representation  $\Gamma_G$  of  $G$  that *extends*  $\Gamma_H$

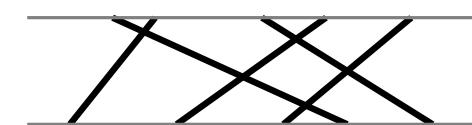


Polytime for:

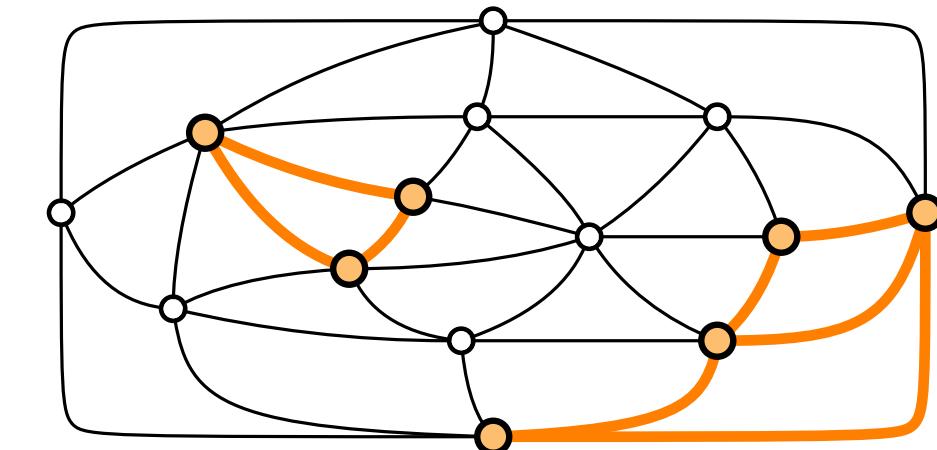
- (unit) interval graphs



- permutation graphs



- circle graphs

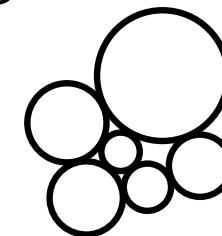


NP-hard for:

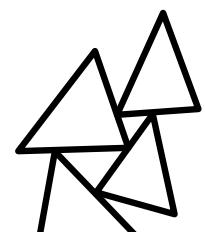
- planar straight-line drawings

- contacts of

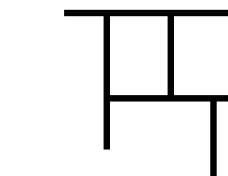
- disks



- triangles

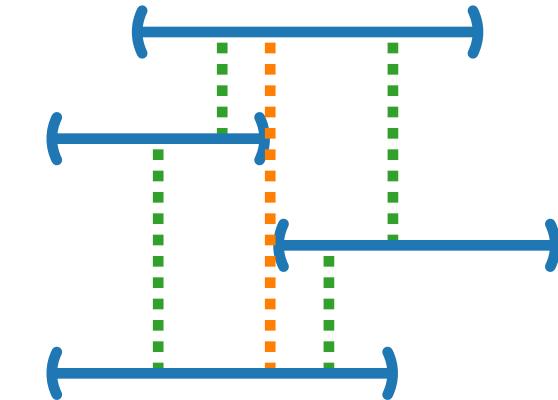
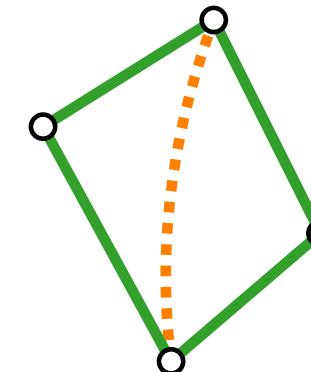


- orthogonal segments



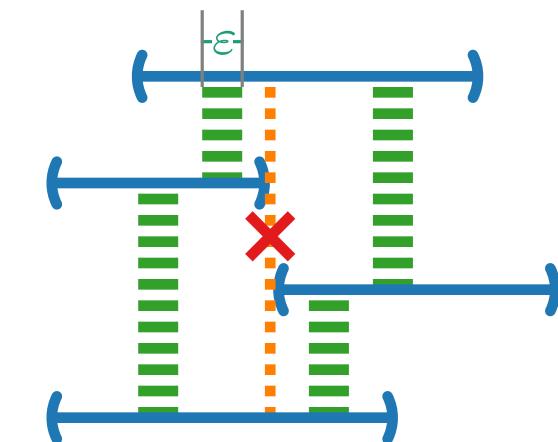
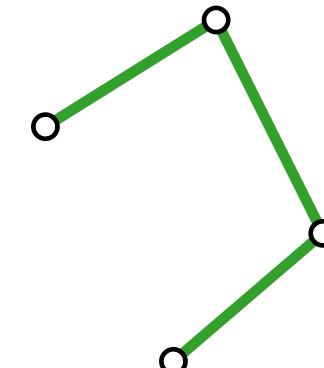
# Bar Visibility Representation

- Vertices correspond to horizontal open line segments called **bars**.
- **Edges** correspond to unobstructed vertical lines of sight.
- What about unobstructed **0-width** vertical lines of sight? Do all visibilities induce edges?

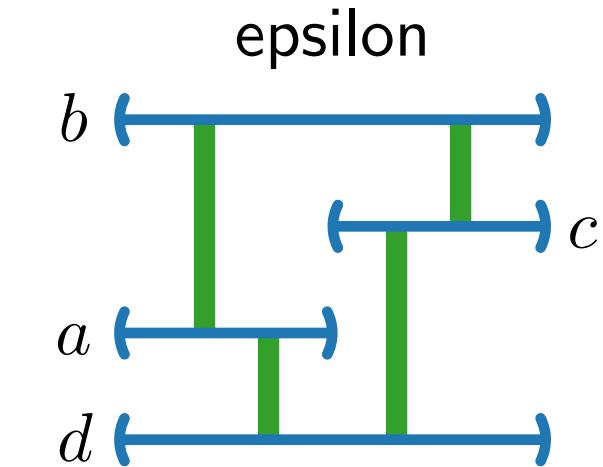
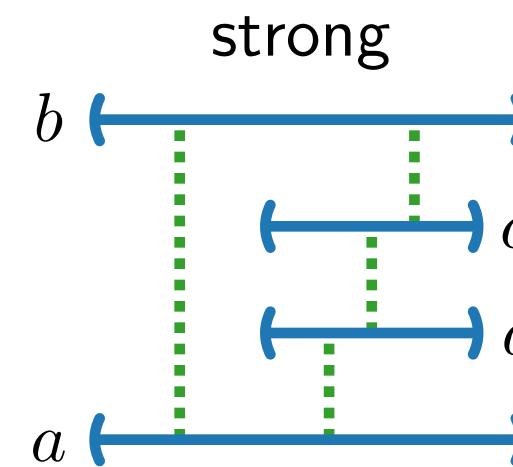
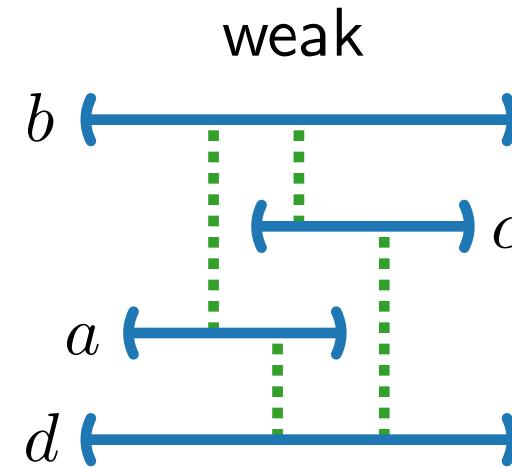
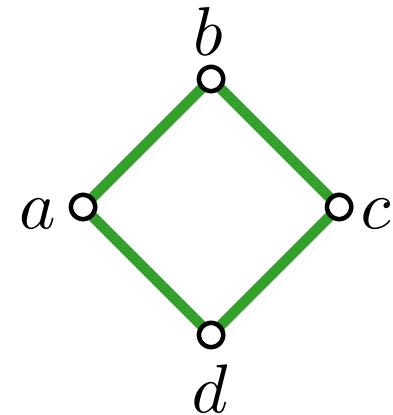


## Models.

- **Strong:**  
Edge  $uv \Leftrightarrow$  unobstructed **0-width** vertical lines of sight.
- **Epsilon:**  
Edge  $uv \Leftrightarrow \varepsilon$ -wide vertical lines of sight for  $\varepsilon > 0$ .
- **Weak:**  
Edge  $uv \Rightarrow$  unobstructed vertical sightlines exists, i.e., any subset of *visible* pairs



# Problems



## Recognition Problem.

Given a graph  $G$ , **decide** whether there exists a weak/strong/ $\varepsilon$  bar visibility representation  $\psi$  of  $G$ .

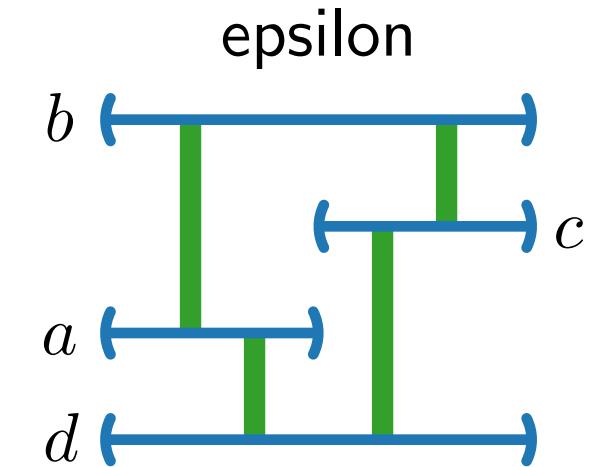
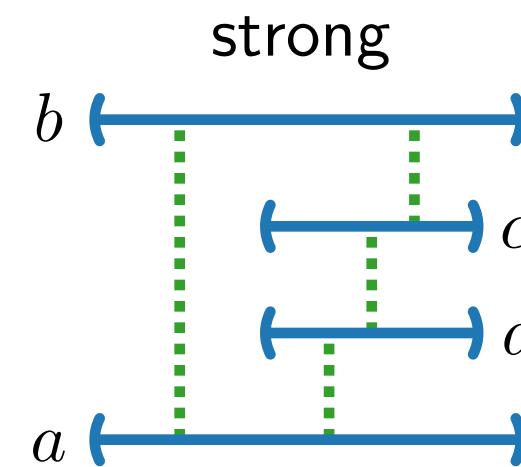
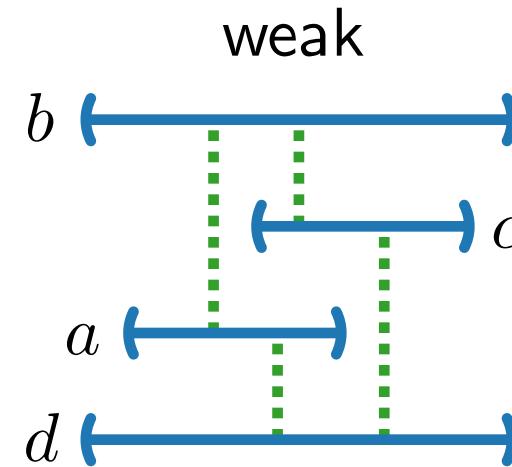
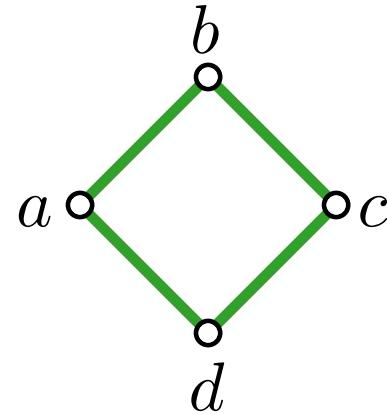
## Construction Problem.

Given a graph  $G$ , **construct** a weak/strong/ $\varepsilon$  bar visibility representation  $\psi$  of  $G$  – if one exists.

## Partial Representation Extension Problem.

Given a graph  $G$  and a **set of bars**  $\psi'$  of  $V' \subset V(G)$ , **decide** whether there exists a weak/strong/ $\varepsilon$  bar visibility representation  $\psi$  of  $G$  **where**  $\psi|_{V'} = \psi'$  (and **construct**  $\psi$  if a representation exists).

# Background



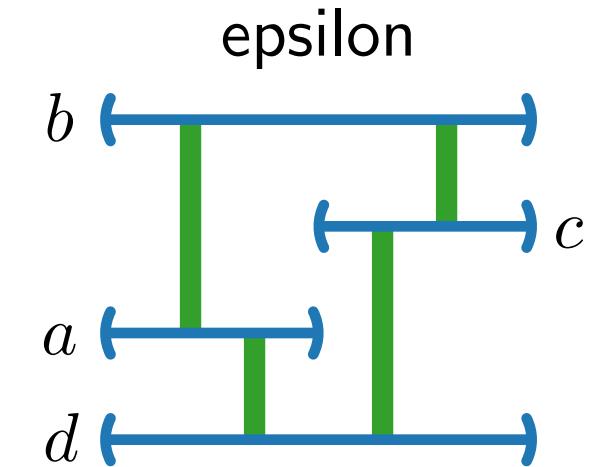
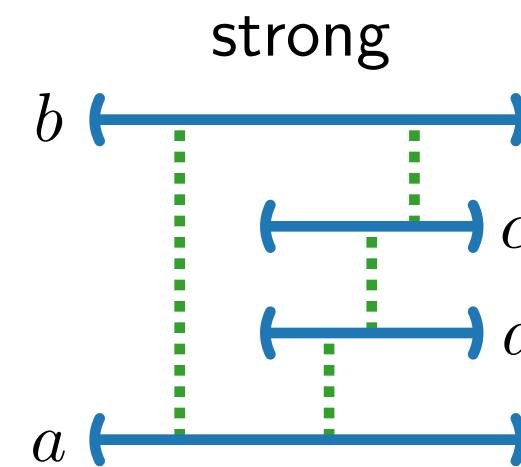
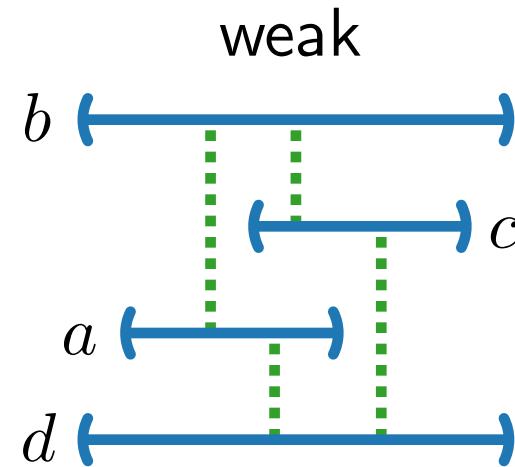
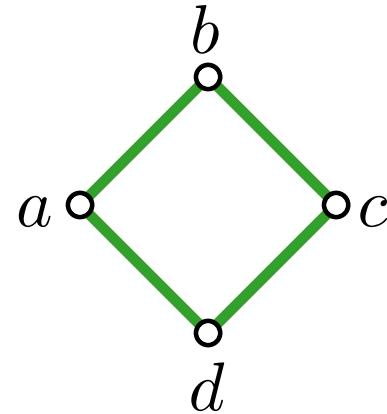
## Weak Bar Visibility.

- All planar graphs. [Tamassia & Tollis '86; Wismath '85]
- Linear time recognition and construction [T&T '86]
- Representation Extension is NP-complete [Chaplick et al. '14]

## Strong Bar Visibility.

- NP-complete to recognize [Andreae '92]

# Background

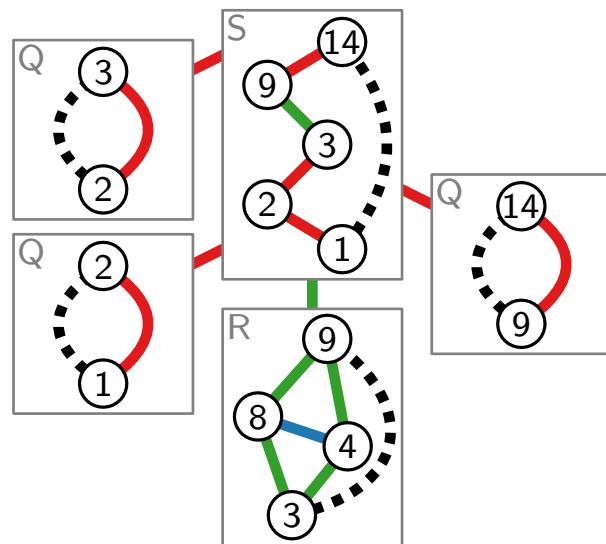


## $\varepsilon$ -Bar Visibility.

- Planar graphs that can be embedded with all **cut vertices** on the outerface. [T&T '86, Wismath '85]
- Linear-time recognition and construction [T&T '86]
- Representation extension? **This Lecture!**

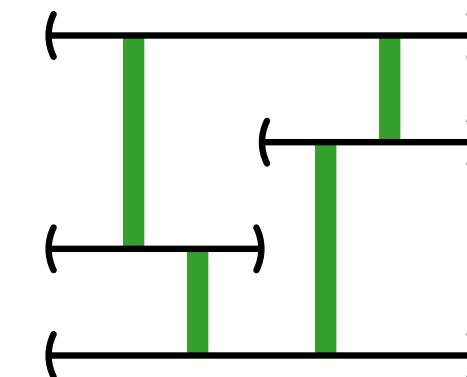
# Visualization of Graphs

## Lecture 9: Partial Visibility Representation Extension



Part II:  
Recognition & Construction

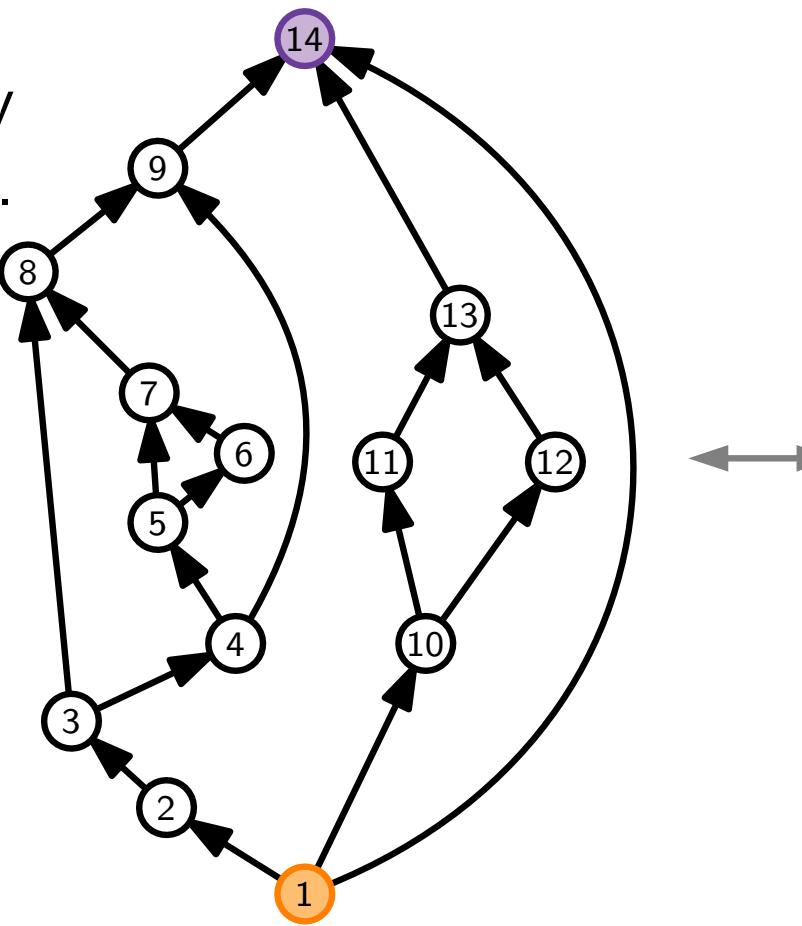
Alexander Wolff



# $\varepsilon$ -bar Visibility and $st$ -Graphs

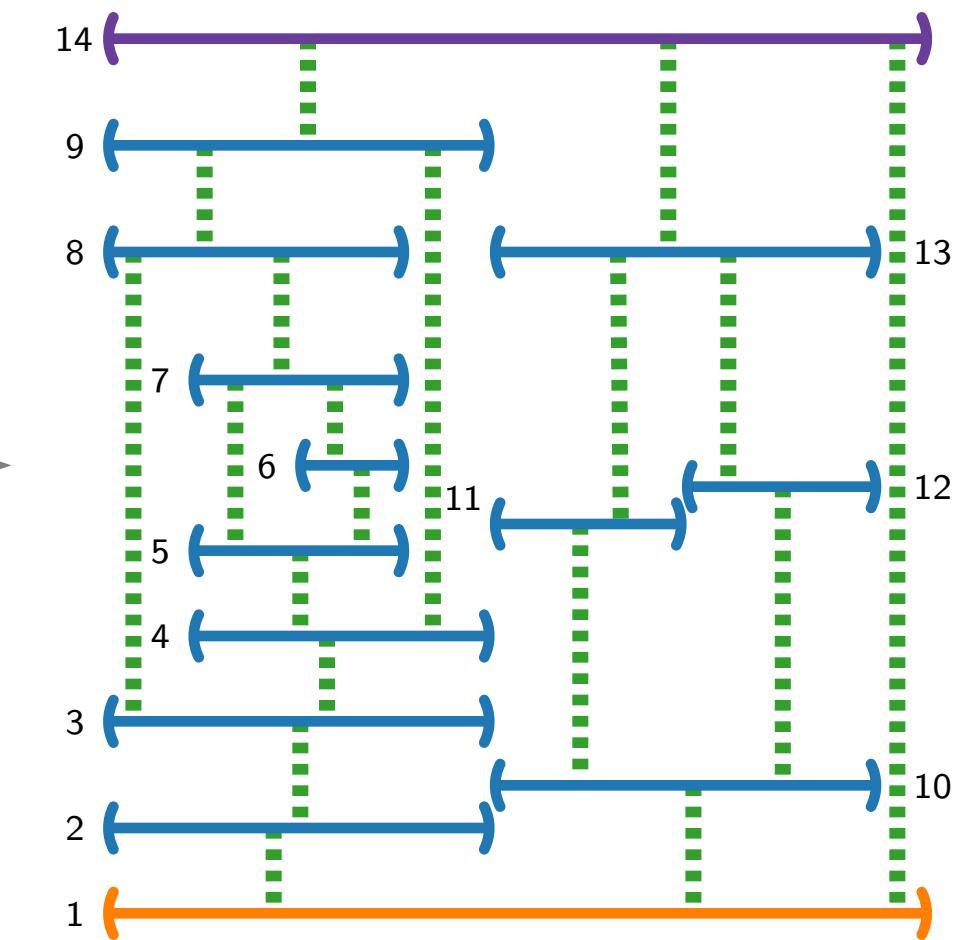
Recall that an  **$st$ -graph** is a planar digraph  $G$  with exactly one **source  $s$**  and one **sink  $t$**  where  $s$  and  $t$  occur on the outer face of an embedding of  $G$ .

- $\varepsilon$ -bar visibility testing is easily done via  $st$ -graph recognition.
- Strong bar visibility recognition... open!
- In a **rectangular** bar visibility representation  $\psi(s)$  and  $\psi(t)$  span an enclosing rectangle.



## Observation.

$st$ -orientations correspond to  $\varepsilon$ -bar visibility representations.



# Results and Outline

## Theorem 1.

[Chaplick et al. '18]

**Rectangular**  $\varepsilon$ -Bar Visibility Representation Extension can be solved in  $\mathcal{O}(n \log^2 n)$  time for *st*-graphs.

- Dynamic program via SPQR-trees
- Easier version:  $\mathcal{O}(n^2)$

## Theorem 2.

[Chaplick et al. '18]

$\varepsilon$ -Bar Visibility Representation Extension is NP-complete.

- Reduction from PLANAR MONOTONE 3-SAT

## Theorem 3.

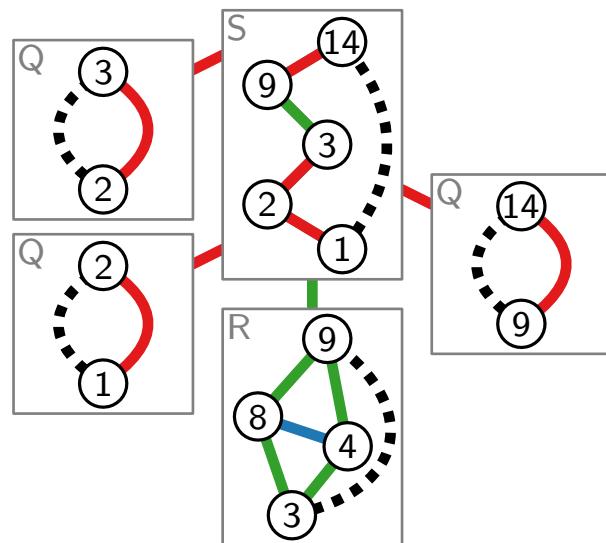
[Chaplick et al. '18]

$\varepsilon$ -Bar Visibility Representation Extension is NP-complete for (series-parallel) *st*-graphs when restricted to the **integer grid** (or if any fixed  $\varepsilon > 0$  is specified).

- Reduction from 3-PARTITION

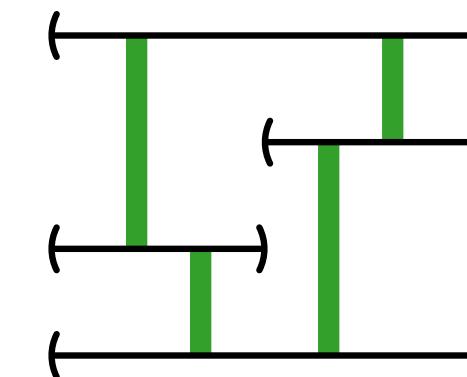
# Visualization of Graphs

## Lecture 9: Partial Visibility Representation Extension



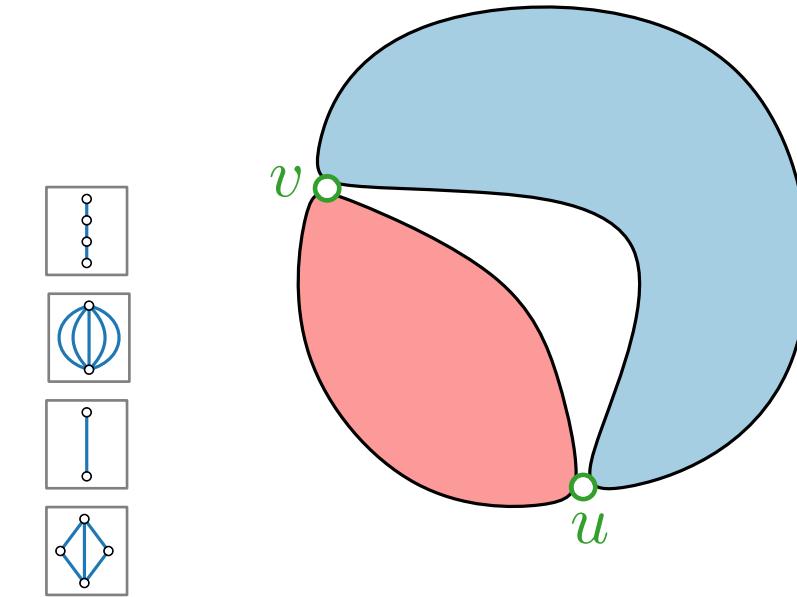
Part III:  
SPQR-Trees

Alexander Wolff

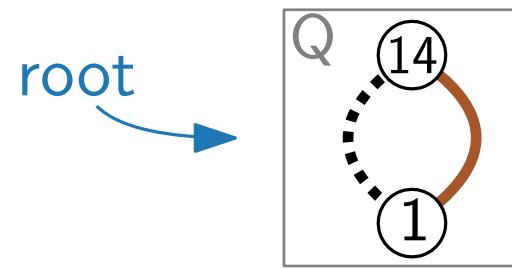
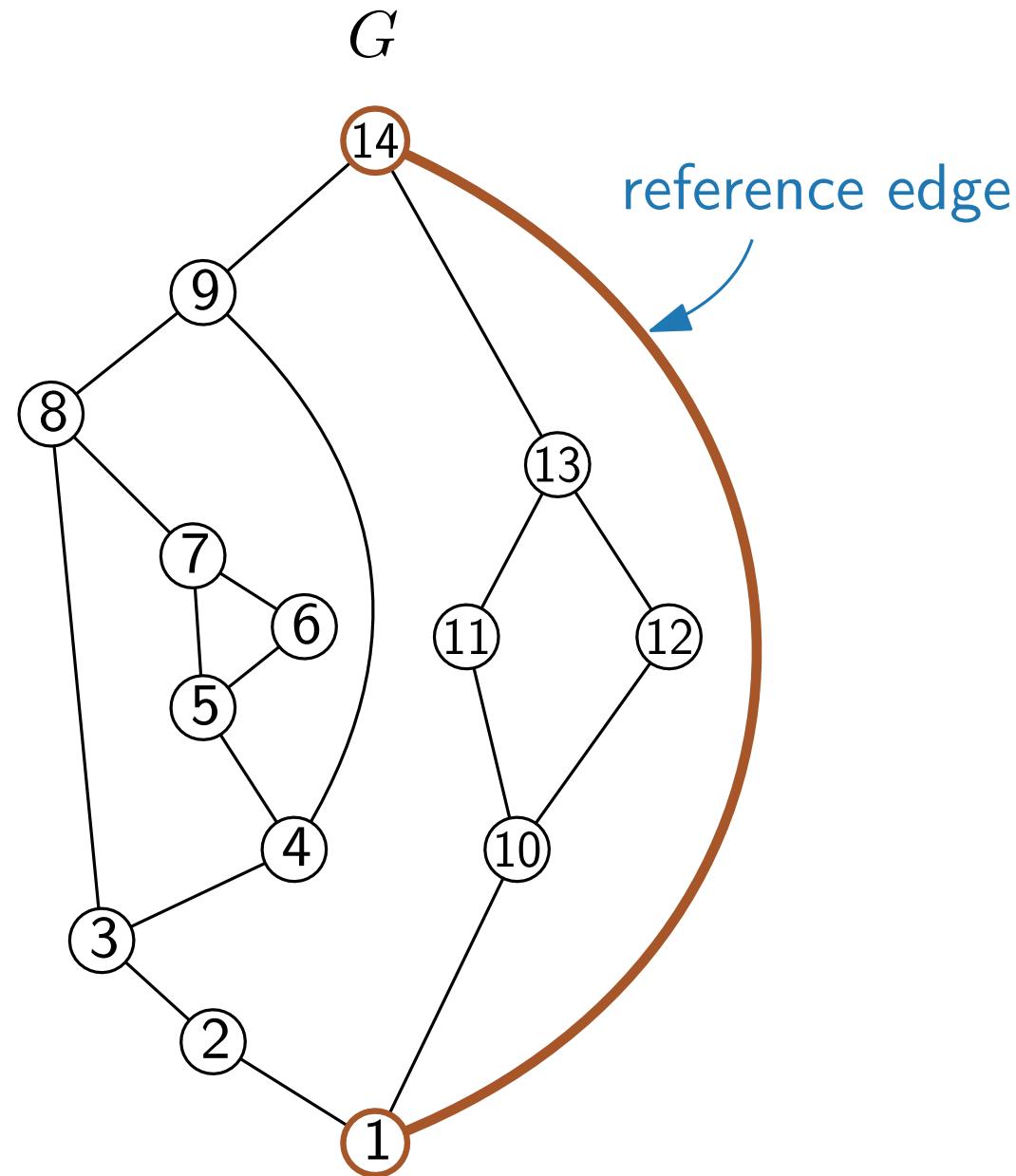


# SPQR-Tree

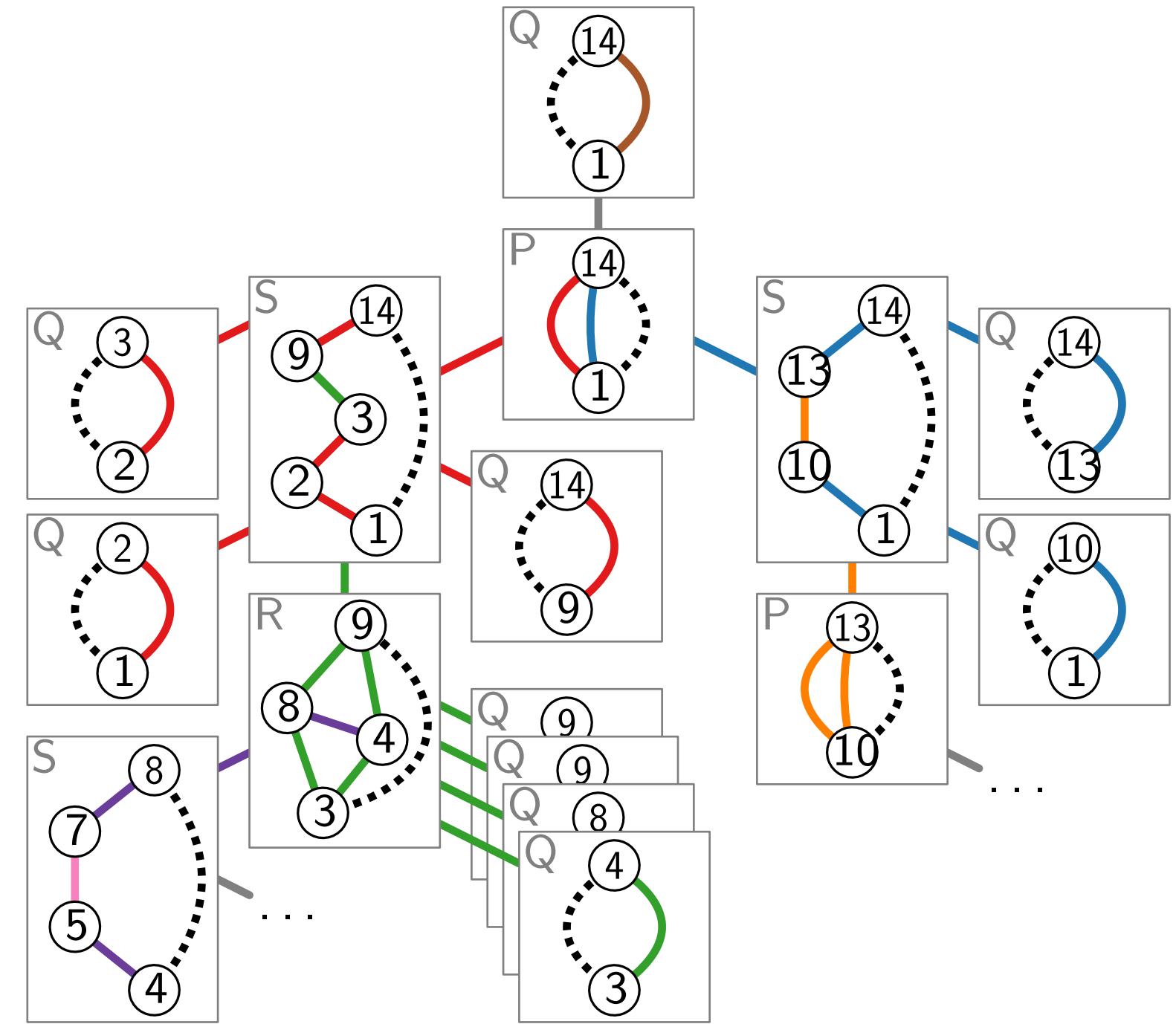
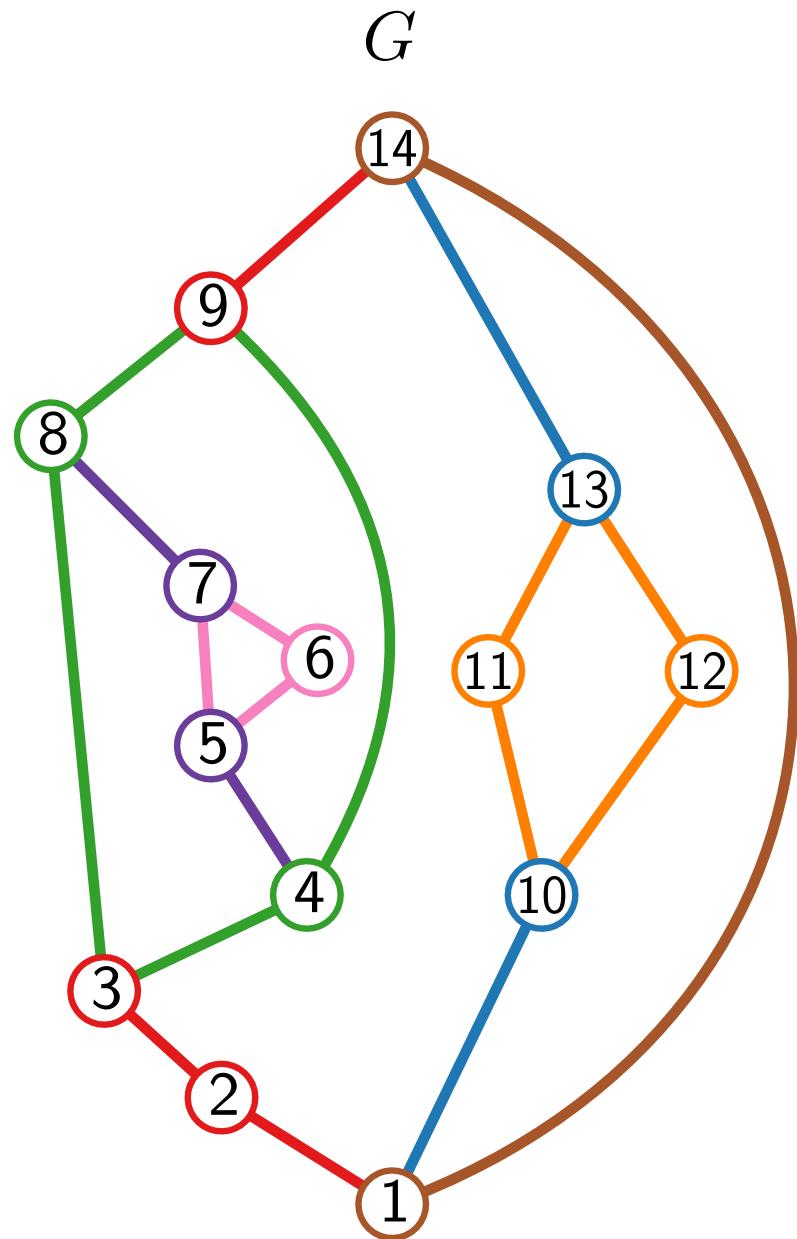
- An **SPQR-tree**  $T$  is a decomposition of a planar graph  $G$  by **separation pairs**.
- The nodes of  $T$  are of four types:
  - **S**-nodes represent a series composition
  - **P**-nodes represent a parallel composition
  - **Q**-nodes represent a single edge
  - **R**-nodes represent 3-connected (*rigid*) subgraphs
- A decomposition tree of a series-parallel graph is an SPQR-tree without **R**-nodes.
- $T$  represents all planar embeddings of  $G$ .
- $T$  can be computed in  $\mathcal{O}(n)$  time. [Gutwenger, Mutzel '01]



# SPQR-Tree Example



# SPQR-Tree Example

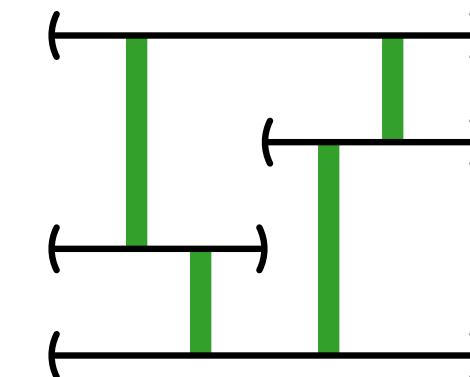
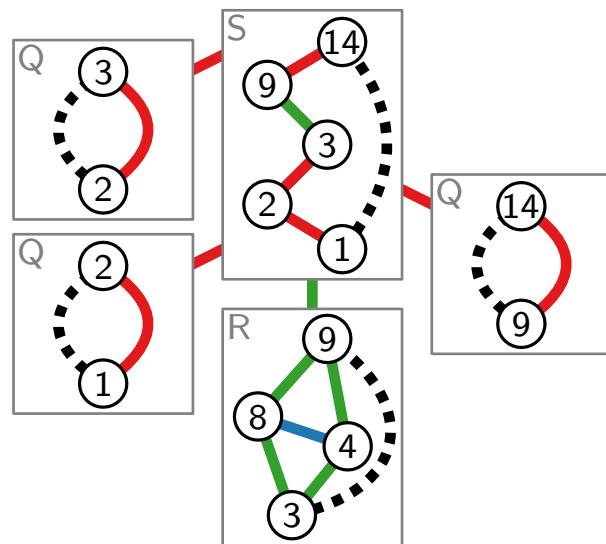


# Visualization of Graphs

## Lecture 9: Partial Visibility Representation Extension

Part IV:  
Rectangular  
Representation Extension

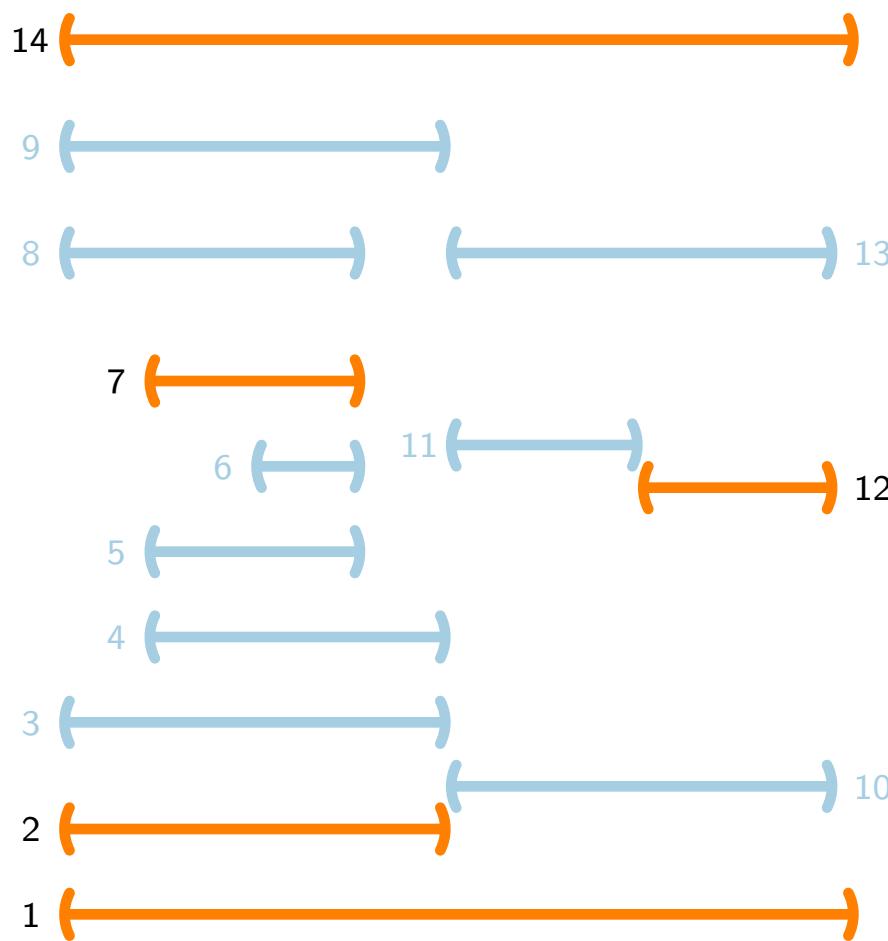
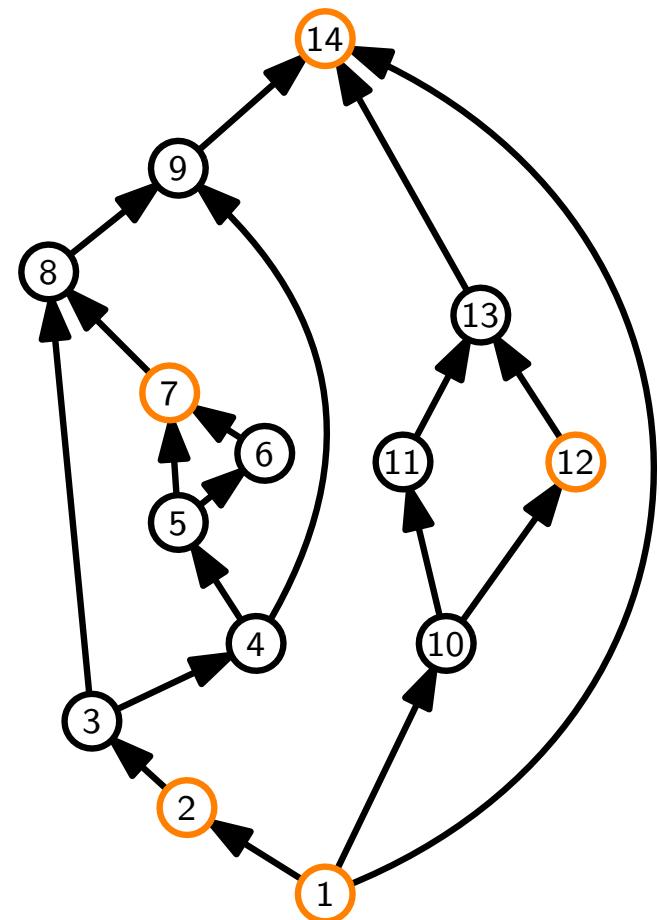
Alexander Wolff



# Representation Extension for st-Graphs

**Theorem 1'.**

**Rectangular**  $\varepsilon$ -Bar Visibility Representation Extension can be solved in  $\mathcal{O}(n^2)$  time for *st*-graphs.

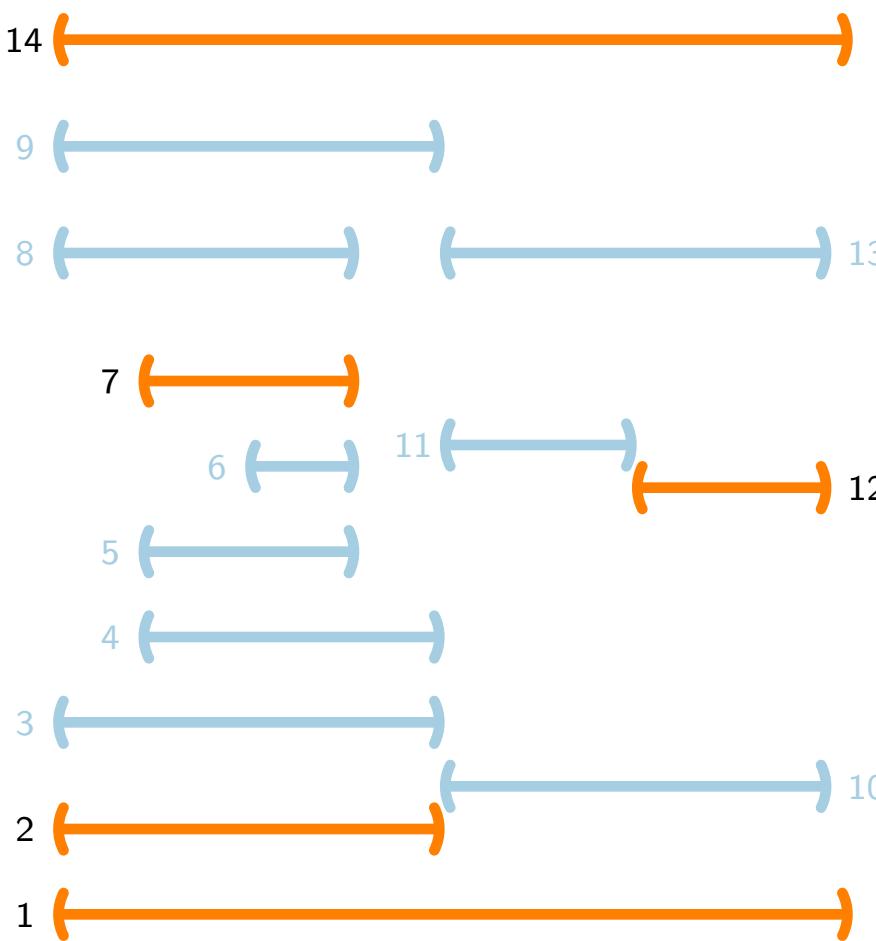
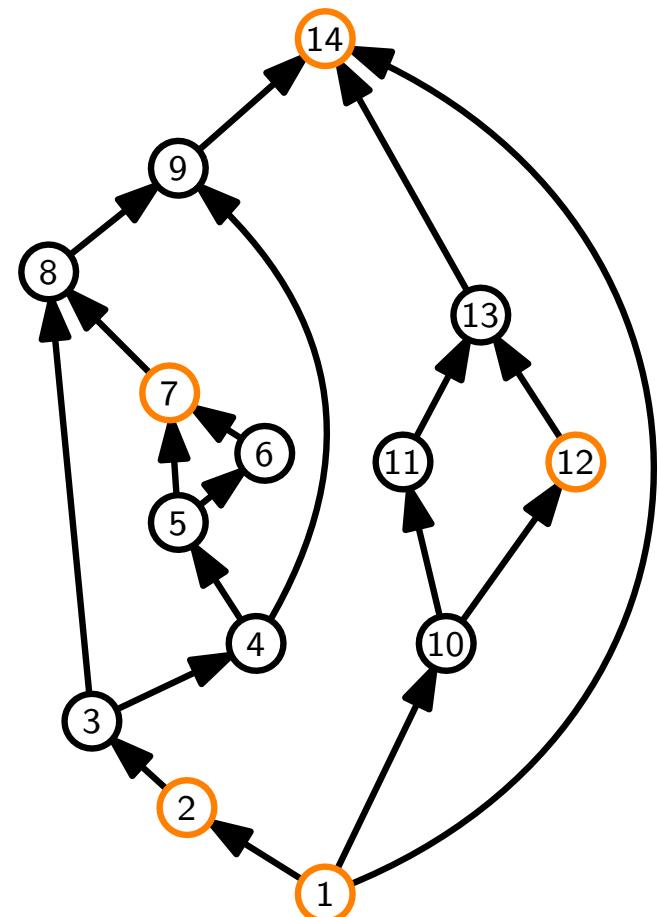


■ Simplify with assumption on y-coordinates

# Representation Extension for st-Graphs

**Theorem 1'.**

**Rectangular**  $\varepsilon$ -Bar Visibility Representation Extension can be solved in  $\mathcal{O}(n^2)$  time for *st*-graphs.



- Simplify with assumption on y-coordinates
- Look at connection to SPQR-trees – tiling
- Solve problems for **S**-, **P**-, and **R**-nodes
- Dynamic program via SPQR-tree

# y-Coordinate Invariant

- Let  $G = (V, E)$  be an  $st$ -graph, and let  $\psi'$  be a representation of  $V' \subseteq V$ .
- Let  $y: V \rightarrow \mathbb{R}$  such that
  - for each  $v \in V'$ ,  $y(v) =$  the y-coordinate of  $\psi'(v)$ .
  - for each edge  $(u, v)$ ,  $y(u) < y(v)$ .

## Lemma 1.

$G$  has a representation extending  $\psi' \Leftrightarrow$

$G$  has a representation extending  $\psi'$

where the y-coordinates of the bars are as in  $y$ .

We can now assume that all y-coordinates are given!

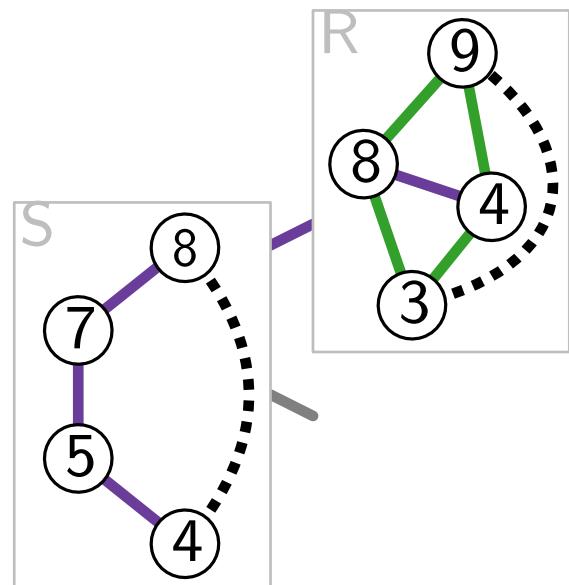
**Proof idea.** The relative positions of **adjacent** bars must match the order given by  $y$ .

So, we can adjust the y-coordinates of any solution to be as in  $y$  by sweeping from bottom to top.

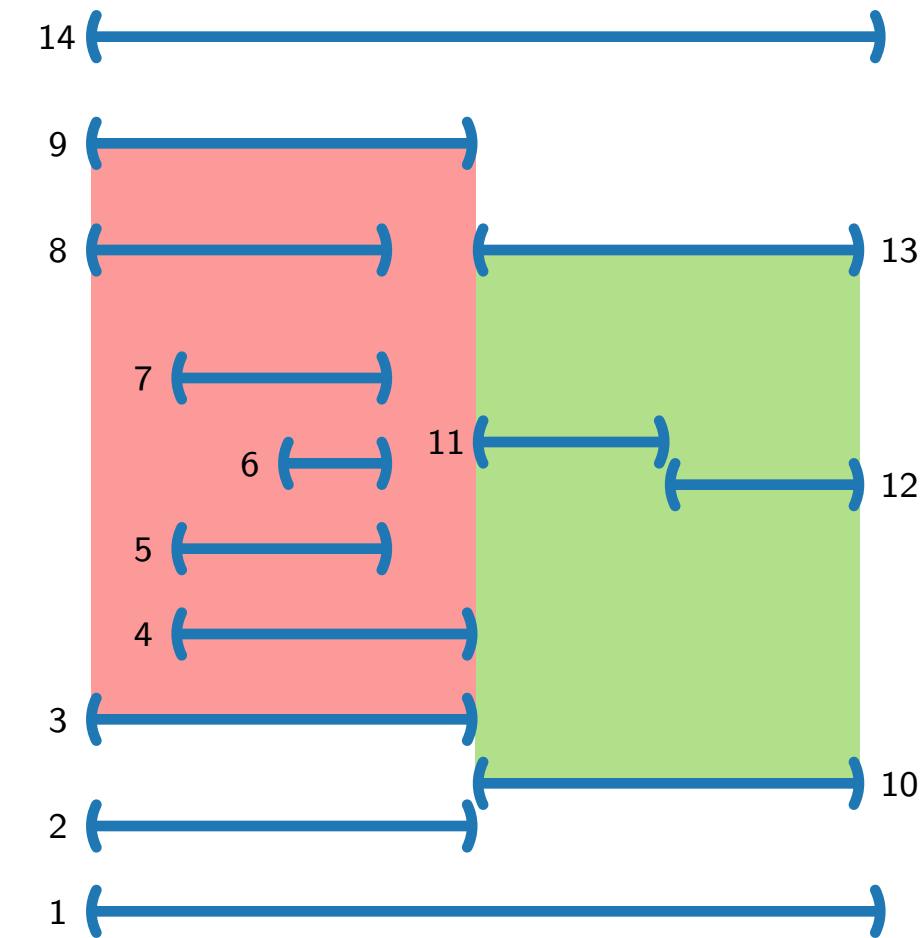
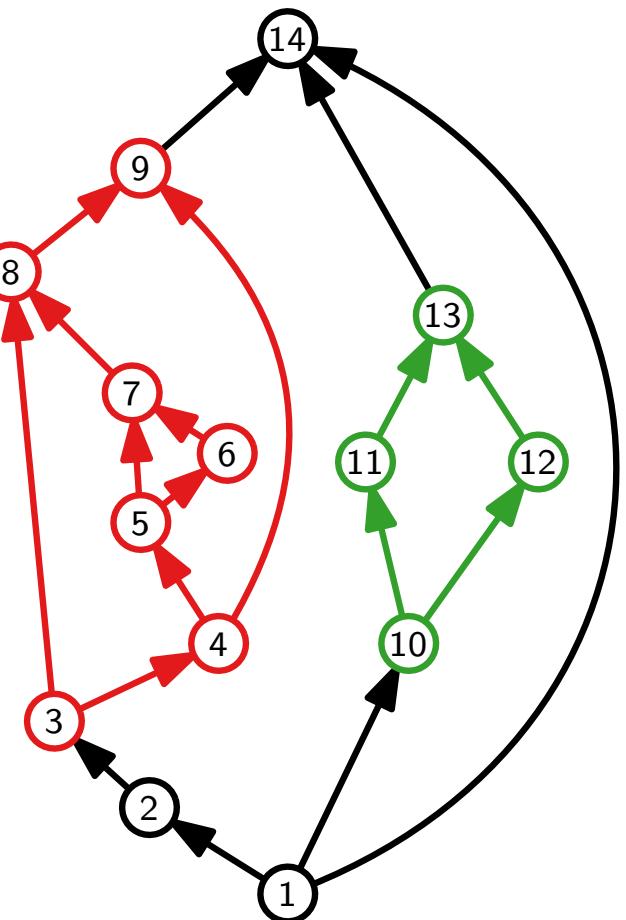
# But why do SPQR-Trees help?

## Lemma 2.

The SPQR-tree of an  $st$ -graph  $G$  induces a recursive **tiling** of any  $\varepsilon$ -bar visibility representation of  $G$ .

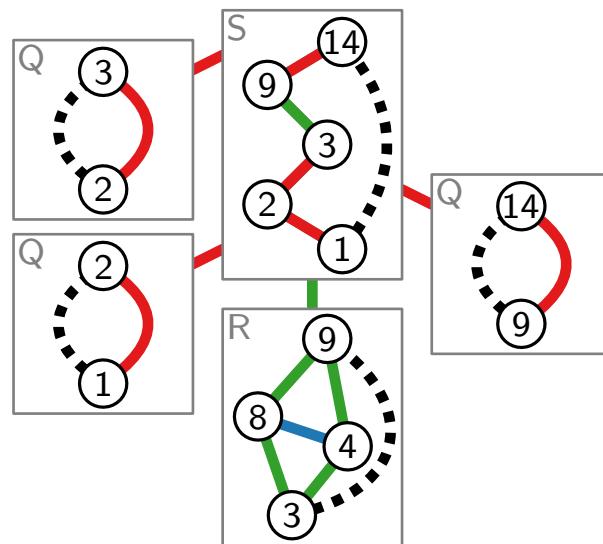


Solve tiles  
bottom-up



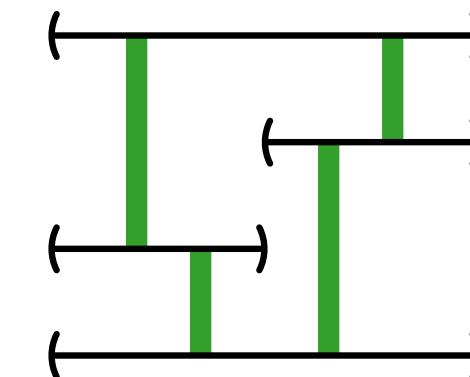
# Visualization of Graphs

## Lecture 9: Partial Visibility Representation Extension



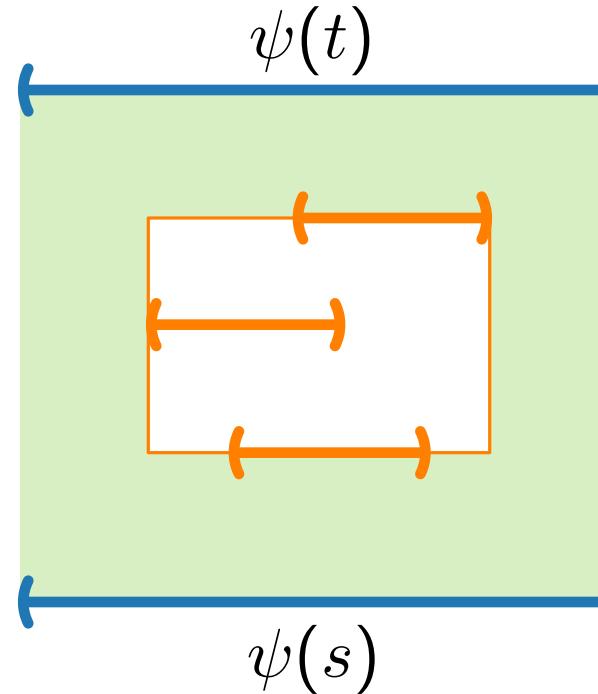
Part V:  
Dynamic Program

Alexander Wolff



# Tiles

**Convention.** **Orange** bars are from the partial representation

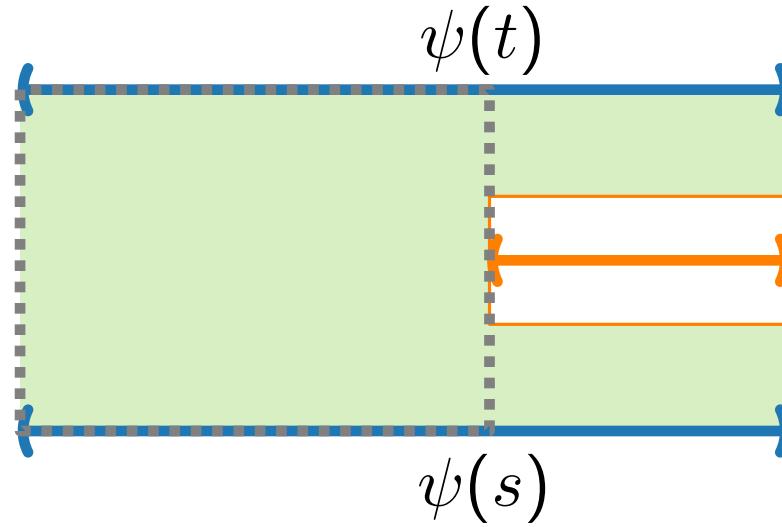


## Observation.

The bounding box (tile) of any solution  $\psi$  **contains** the bounding box of the partial representation.

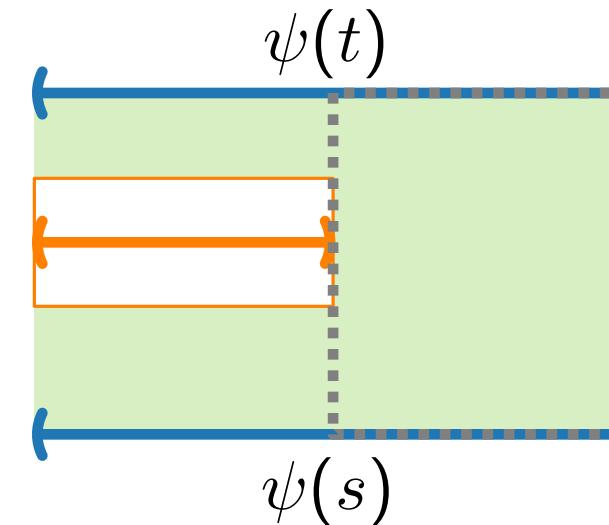
How many different **types** of tiles are there?

# Types of Tiles



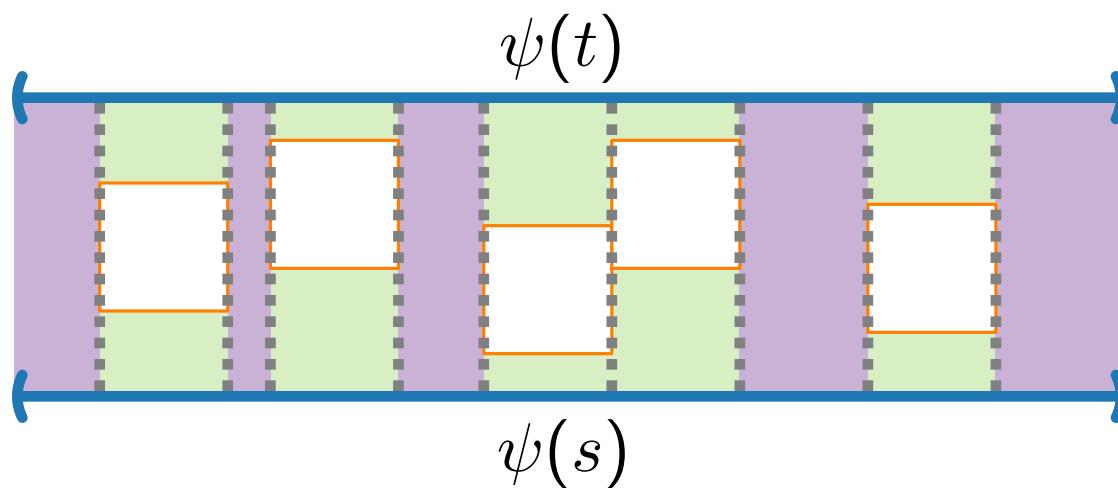
- Left **Fixed** – due to the orange bar
- Right **Loose** – due to the orange bar

- Right **Fixed** – due to the orange bar
- Left **Loose** – due to the orange bar



Four different types: **FF**, **FL**, **LF**, **LL**

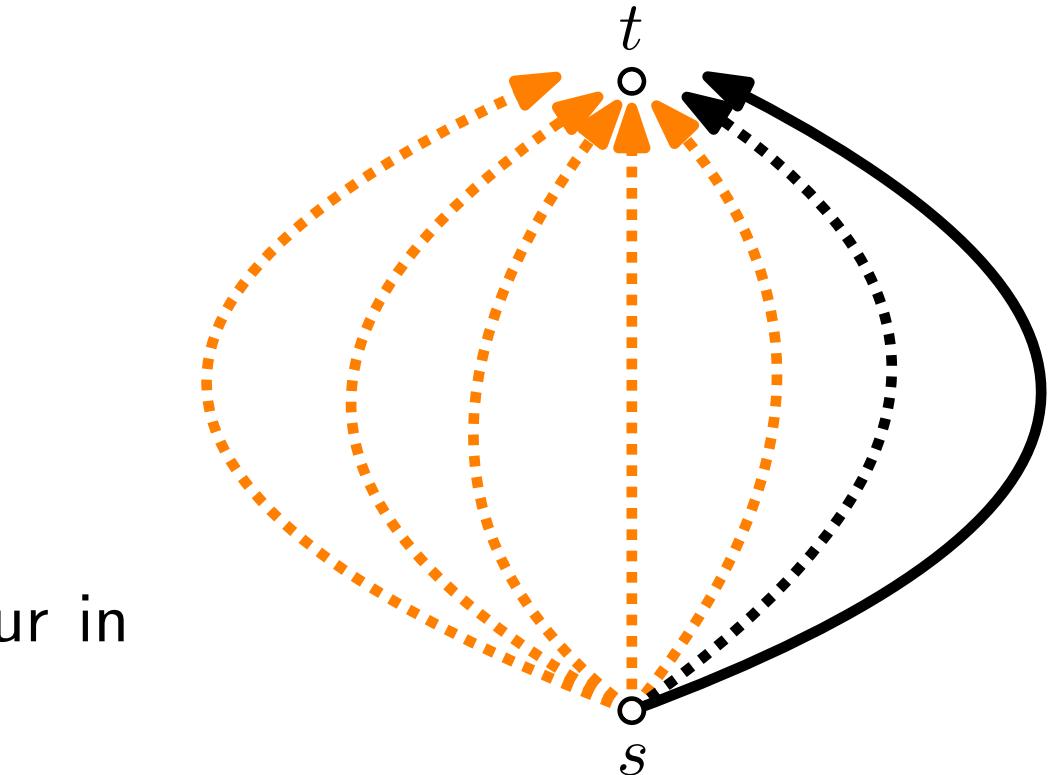
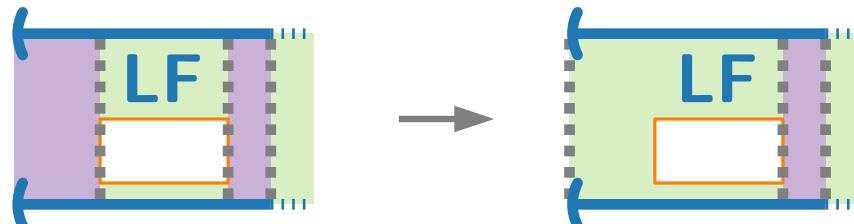
# P-Nodes



- Children of P-node with **prescribed bars** occur in given left-to-right order
- But there might be some **gaps**...

## Idea.

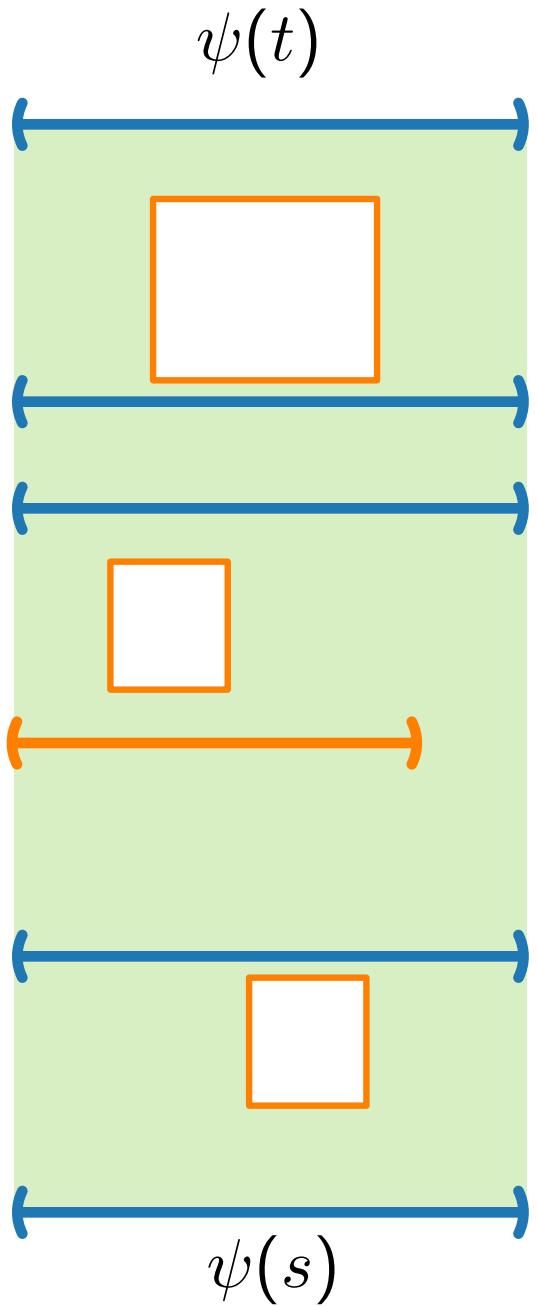
Greedily *fill* the **gaps** by preferring to “stretch” the children with prescribed bars.



## Outcome.

After processing, we must know the valid types for the corresponding subgraphs.

# S-Nodes

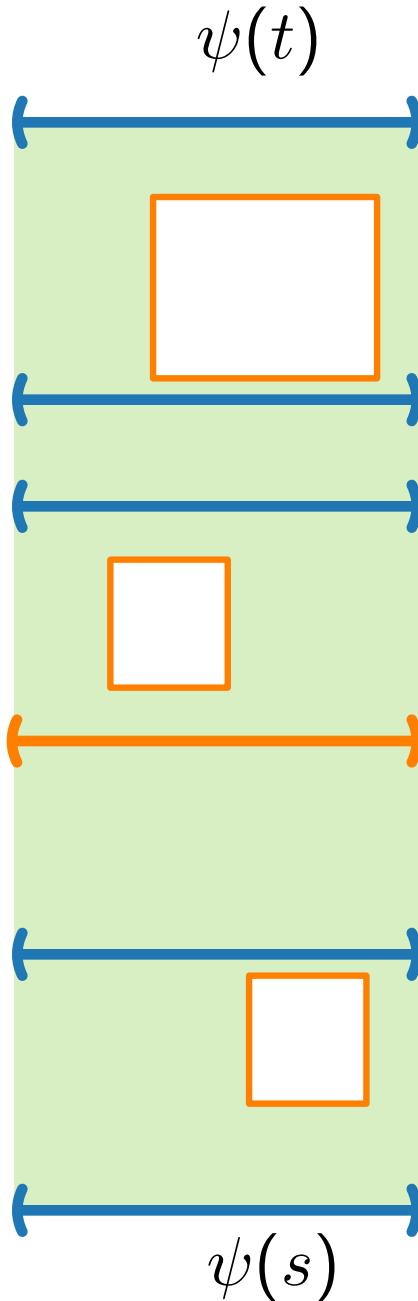


$t$

$s$

This **fixed vertex**  
means we can only  
make a **Fixed-Fixed**  
representation!

# S-Nodes

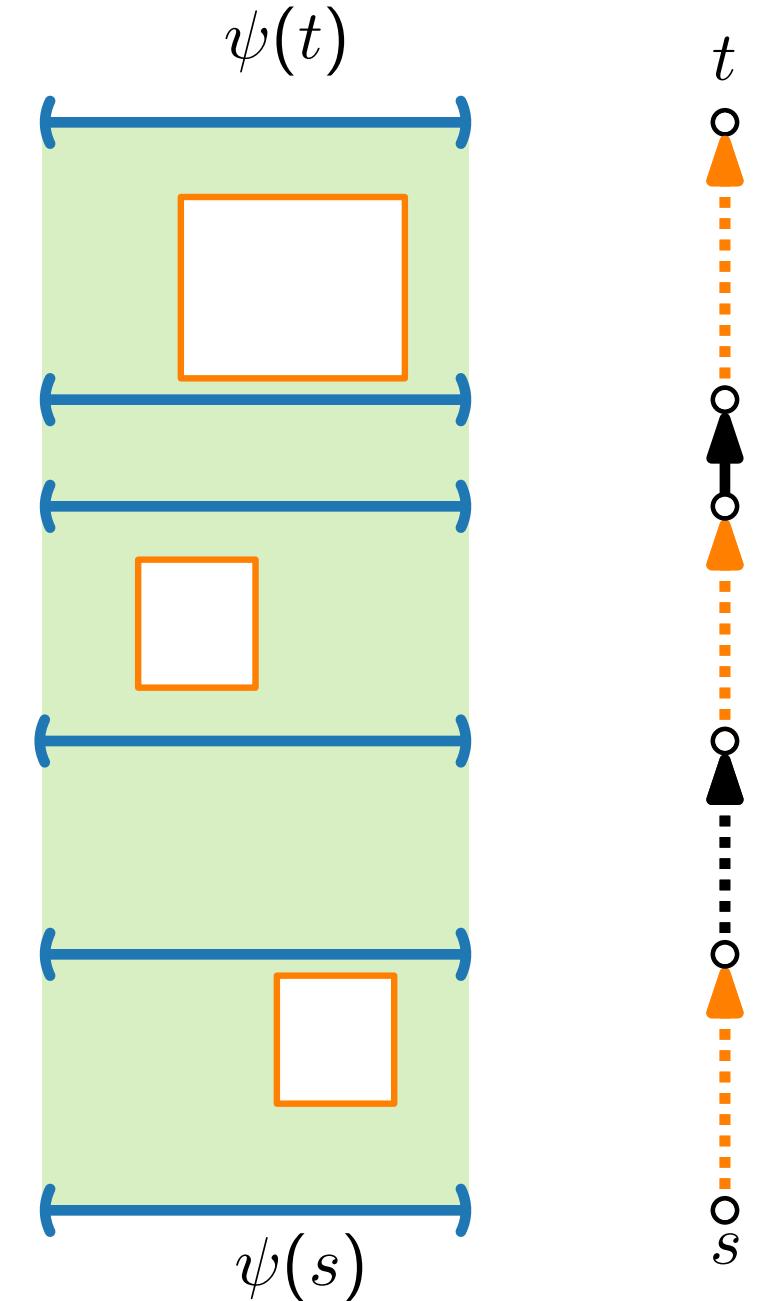


$t$

$s$

This **fixed vertex**  
means we can only  
make a Fixed-Fixed  
representation!

Here we have a  
chance to make all  
(**LL**, **FL**, **LF**, **FF**)  
types.

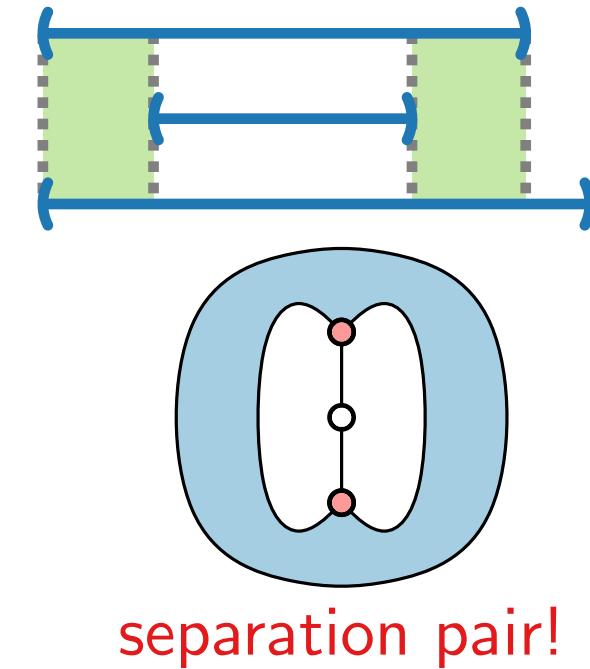
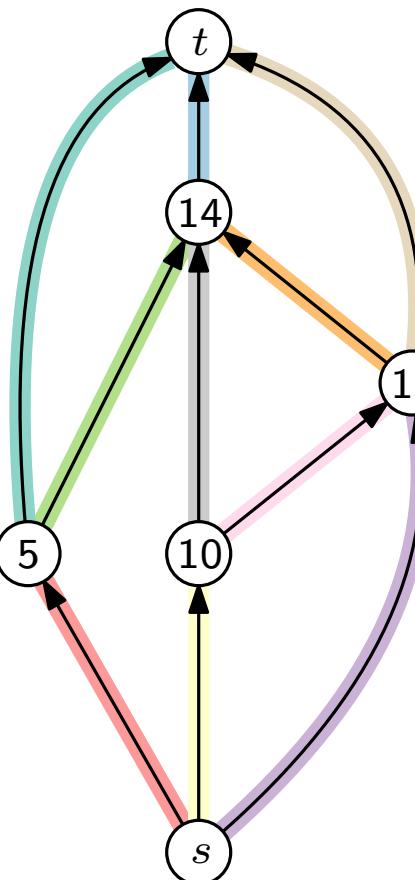
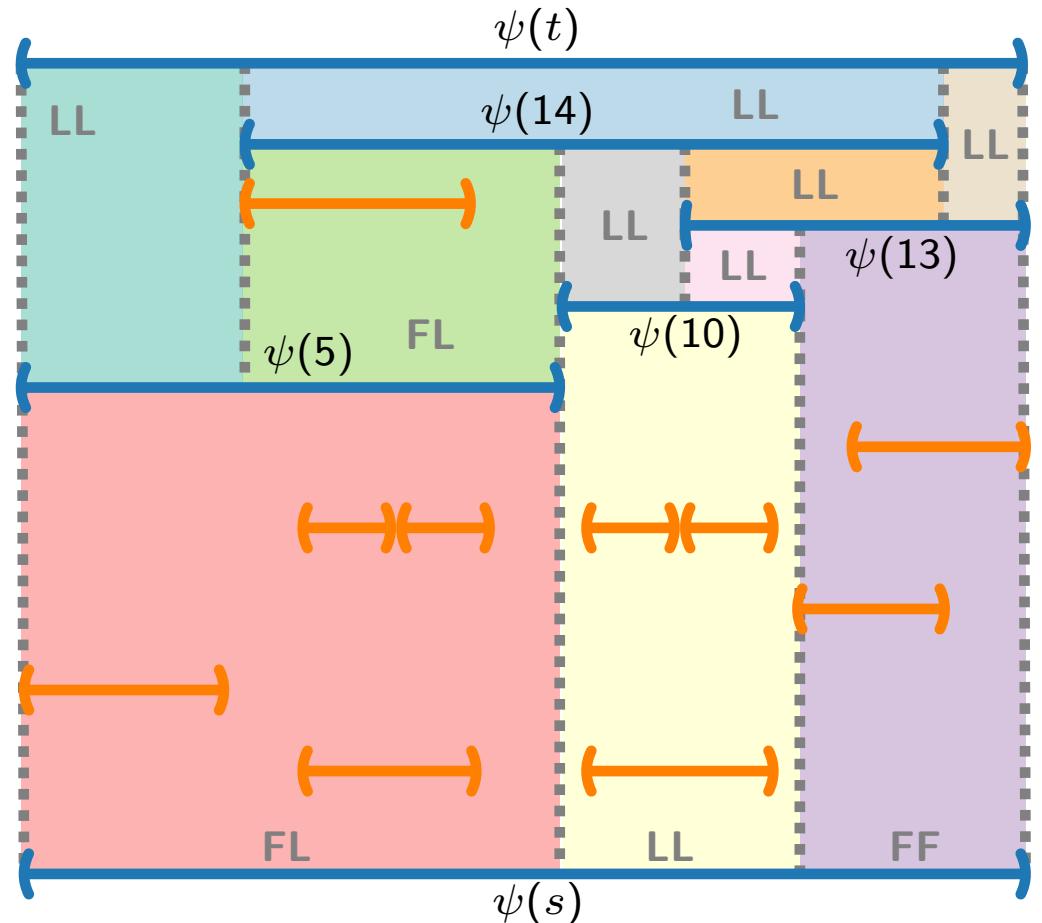


$t$

$s$

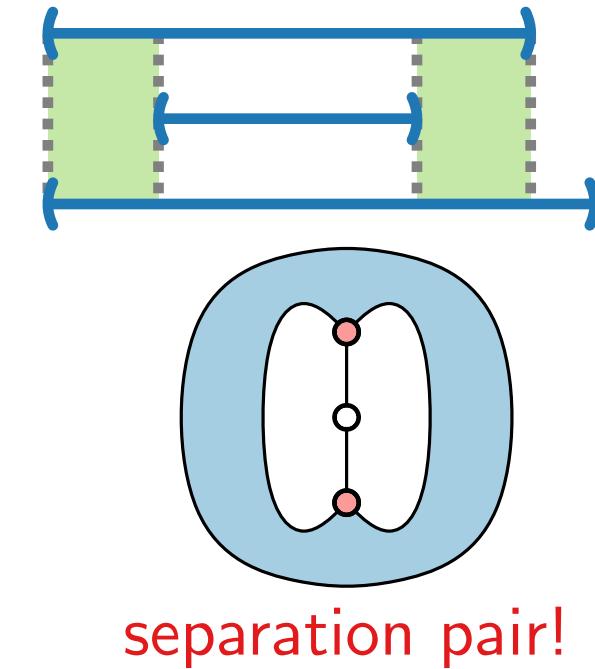
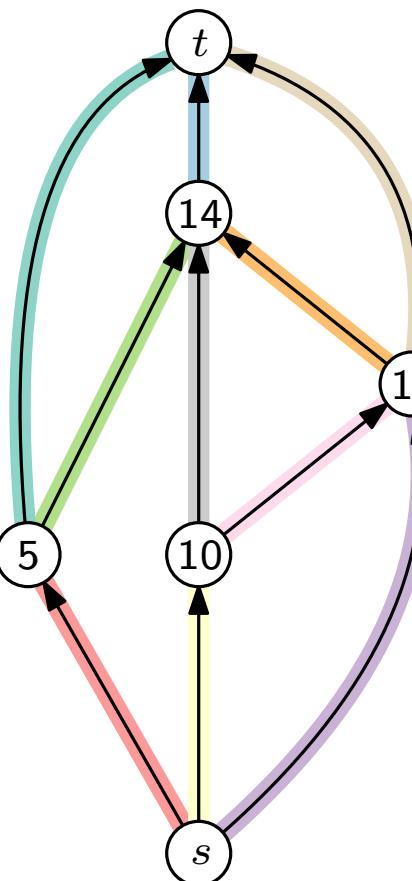
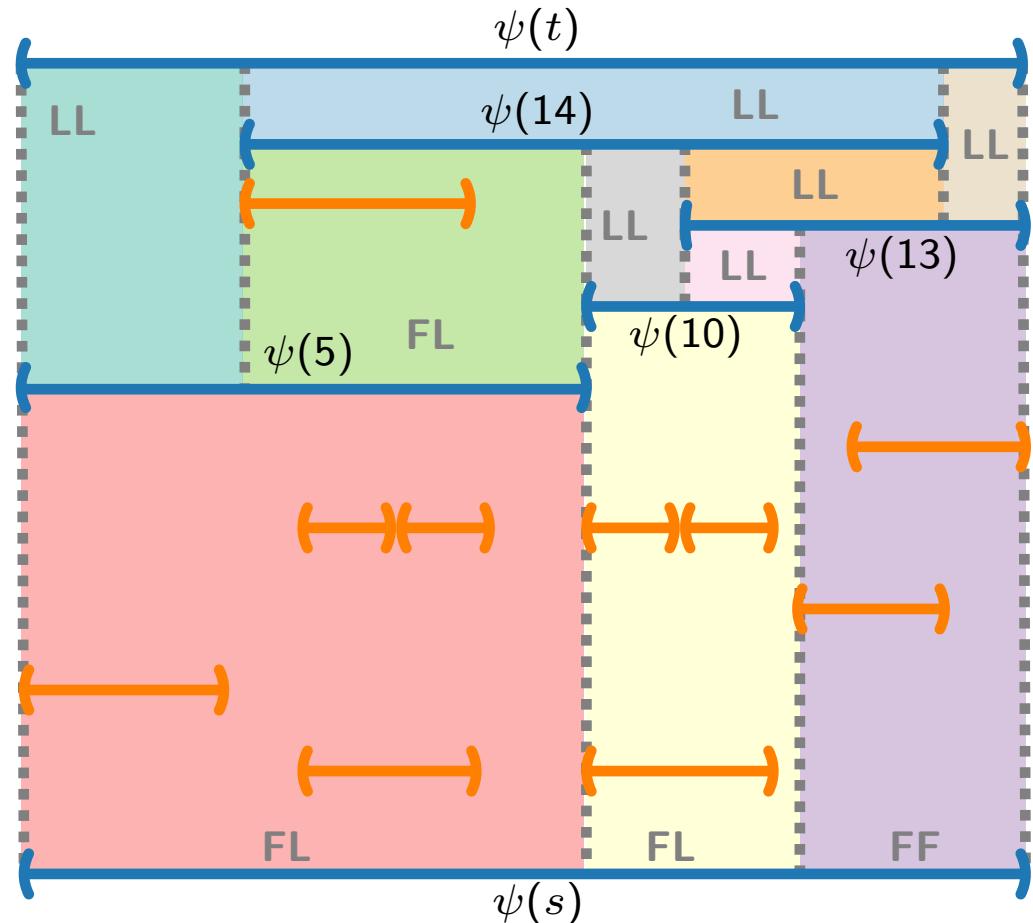
# R-Nodes with 2-SAT Formulation

- for each child (edge)  $e$ :
  - find all types of  $\{\text{FF,FL,LF,LL}\}$  that admit a drawing
  - 2 variables  $l_e, r_e$  encoding fixed/loose type of its tile
  - consistency clauses



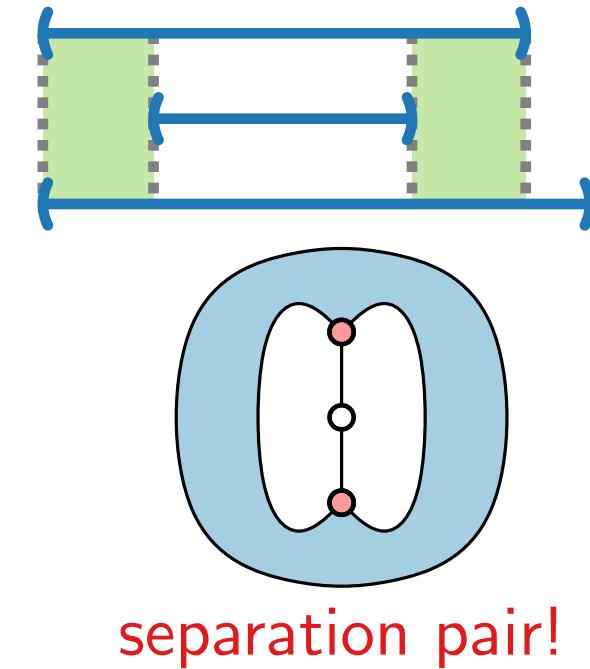
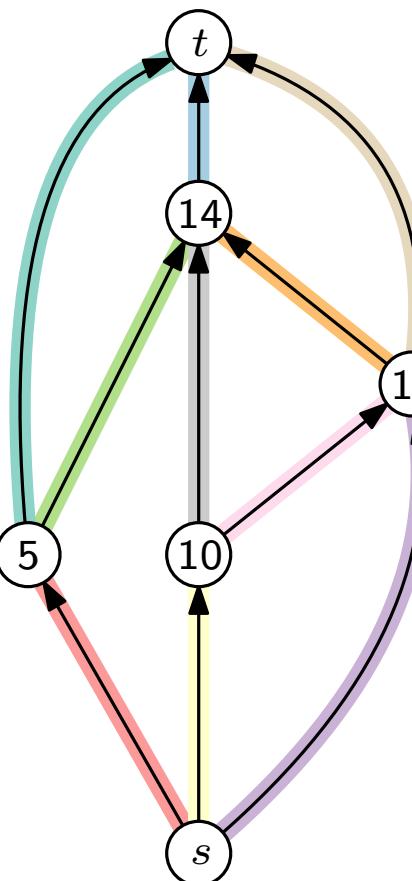
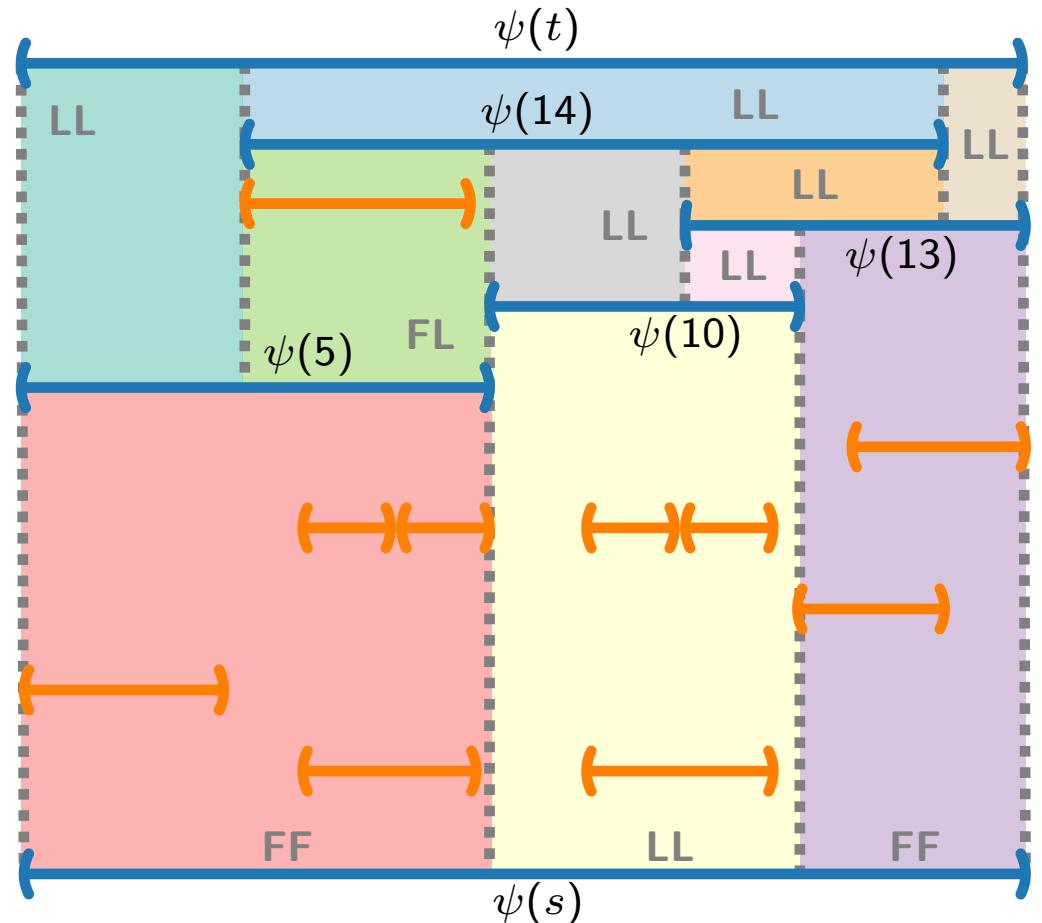
# R-Nodes with 2-SAT Formulation

- for each child (edge)  $e$ :
  - find all types of  $\{\text{FF,FL,LF,LL}\}$  that admit a drawing
  - 2 variables  $l_e, r_e$  encoding fixed/loose type of its tile
  - consistency clauses



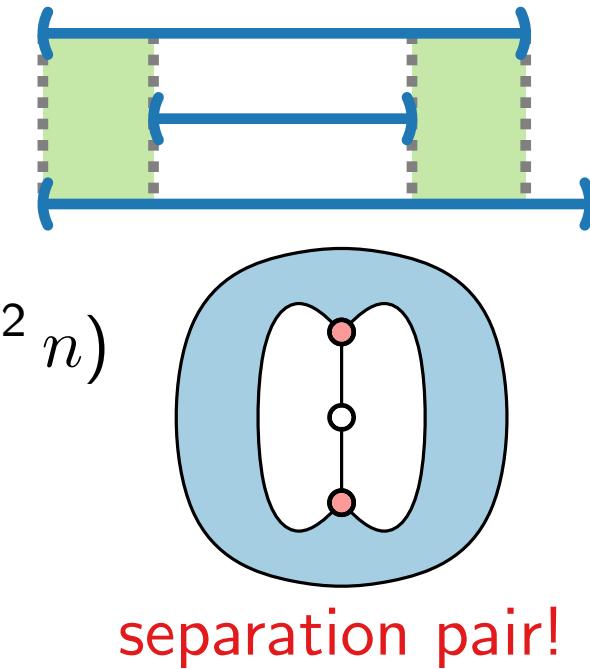
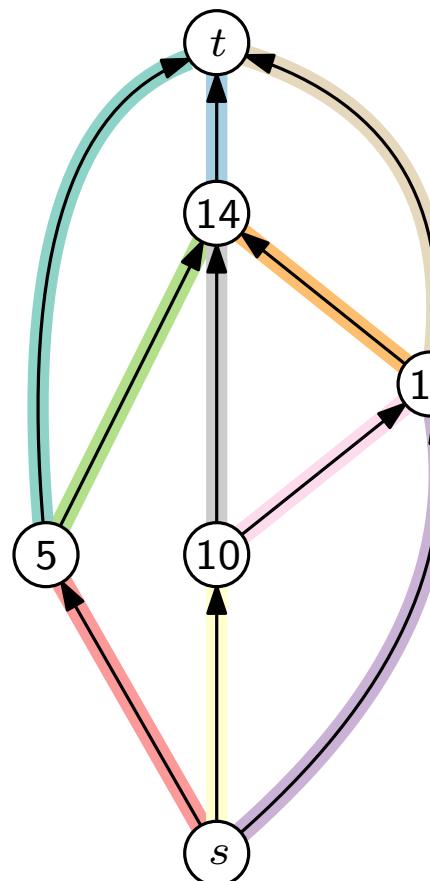
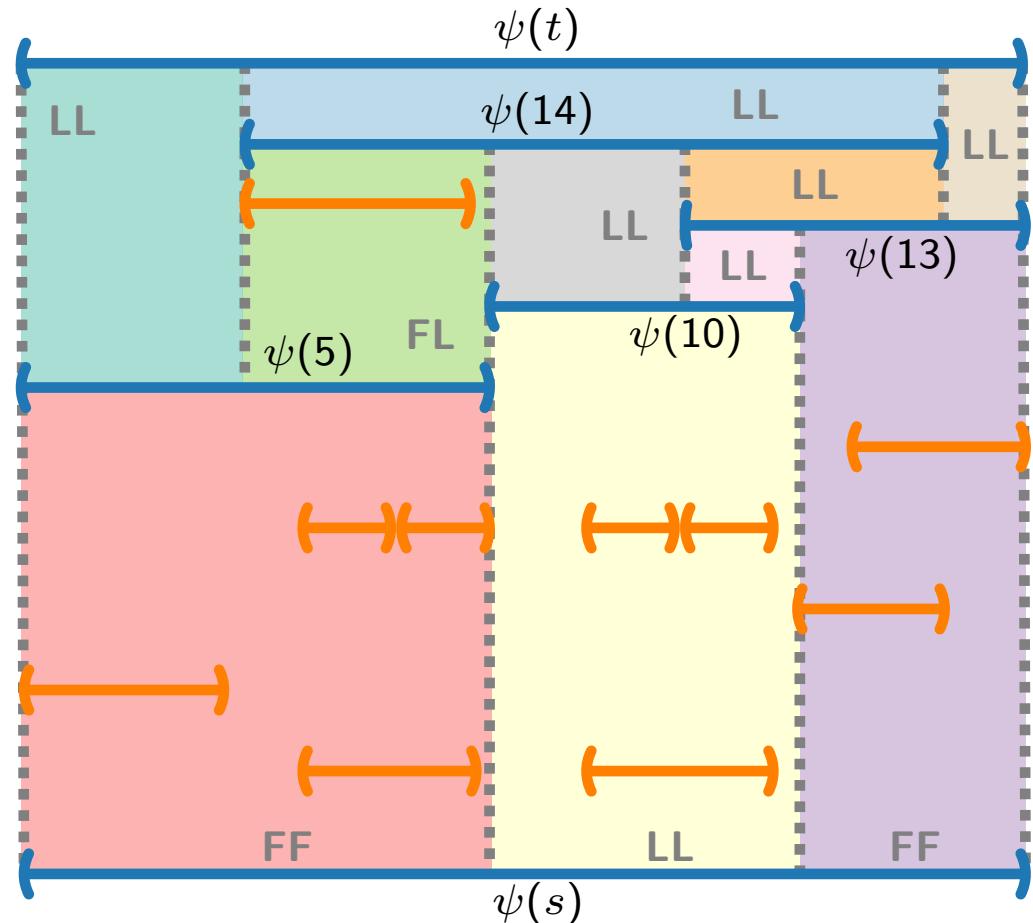
# R-Nodes with 2-SAT Formulation

- for each child (edge)  $e$ :
  - find all types of  $\{\text{FF,FL,LF,LL}\}$  that admit a drawing
  - 2 variables  $l_e, r_e$  encoding fixed/loose type of its tile
  - consistency clauses



# R-Nodes with 2-SAT Formulation

- for each child (edge)  $e$ :
  - find all types of  $\{\text{FF,FL,LF,LL}\}$  that admit a drawing
  - 2 variables  $l_e, r_e$  encoding fixed/loose type of its tile
  - consistency clauses –  $O(n^2)$  many, but can be reduced to  $O(n \log^2 n)$

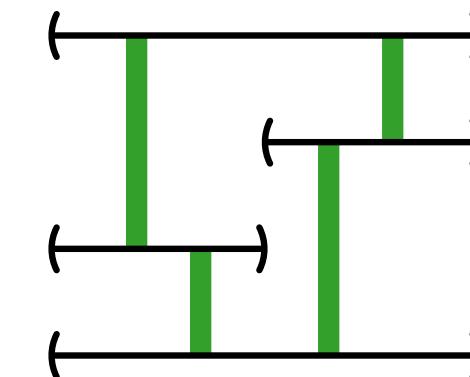
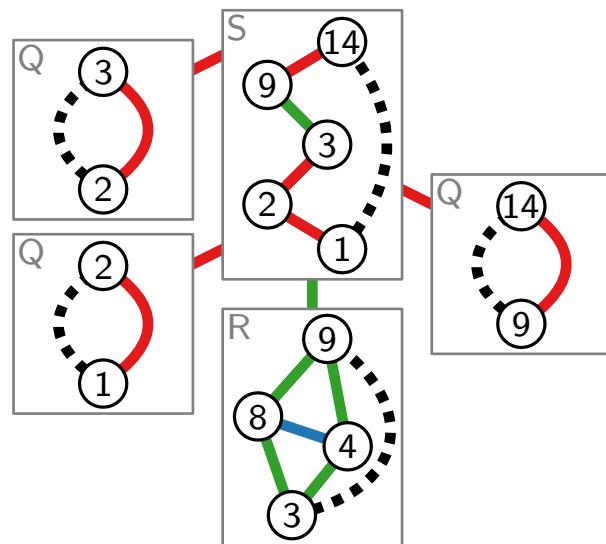


# Visualization of Graphs

## Lecture 9: Partial Visibility Representation Extension

Part VI:  
NP-Hardness  
of the General Case

Alexander Wolff



# NP-Hardness of RepExt in the General Case

## Theorem 2.

$\varepsilon$ -Bar Visibility Representation Extension is NP-complete.

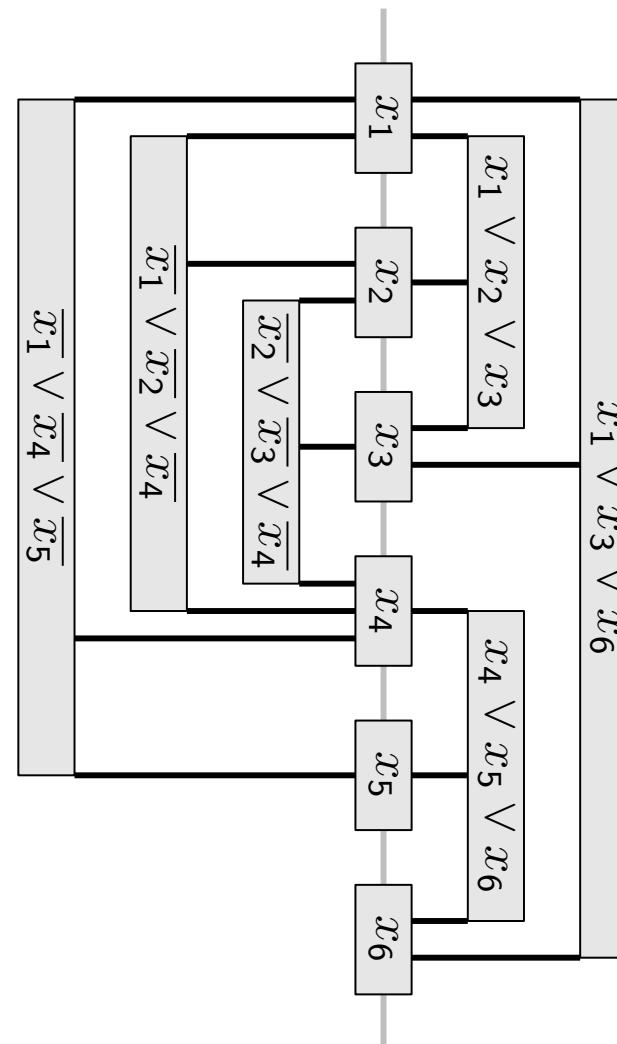
- Reduction from Planar Monotone 3-SAT

# NP-Hardness of RepExt in the General Case

## Theorem 2.

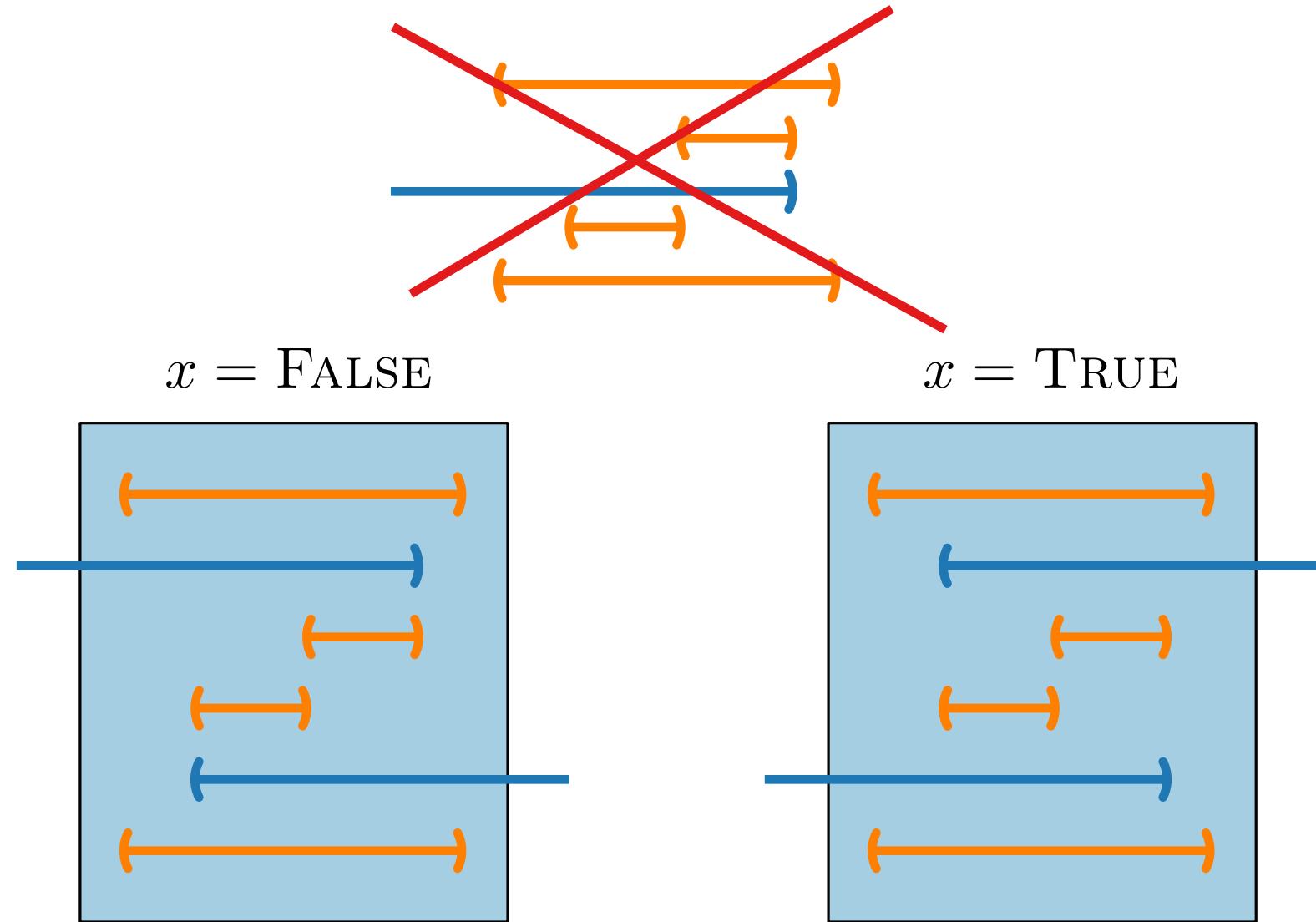
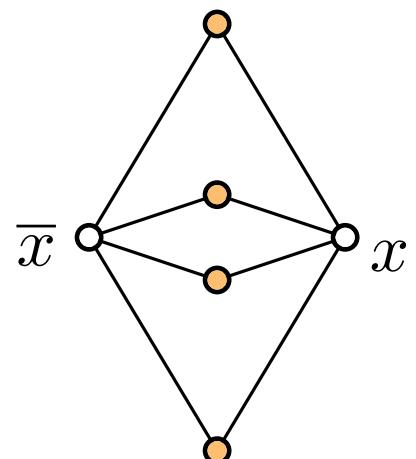
$\varepsilon$ -Bar Visibility Representation Extension is NP-complete.

- Reduction from Planar Monotone 3-SAT



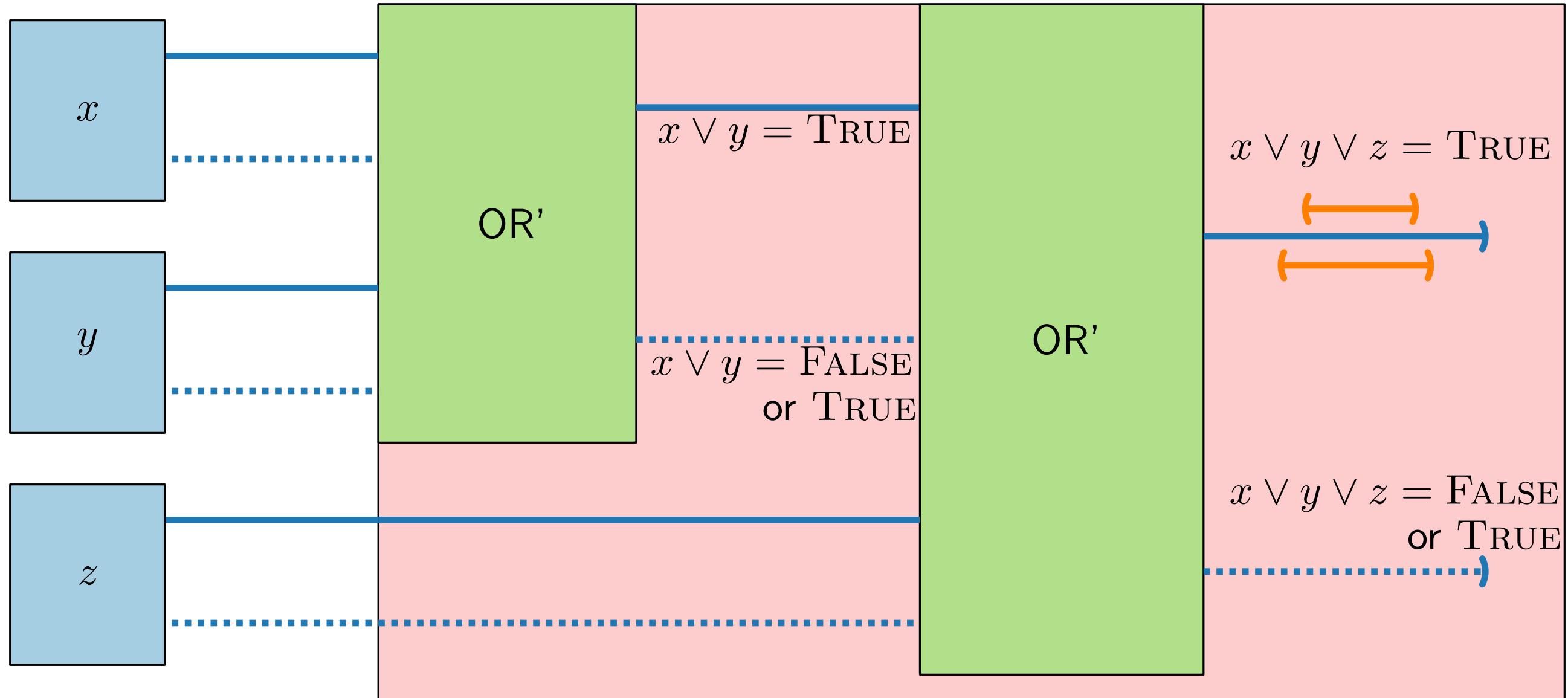
- NP-complete  
[Berg & Khosravi '10]

# Variable Gadget

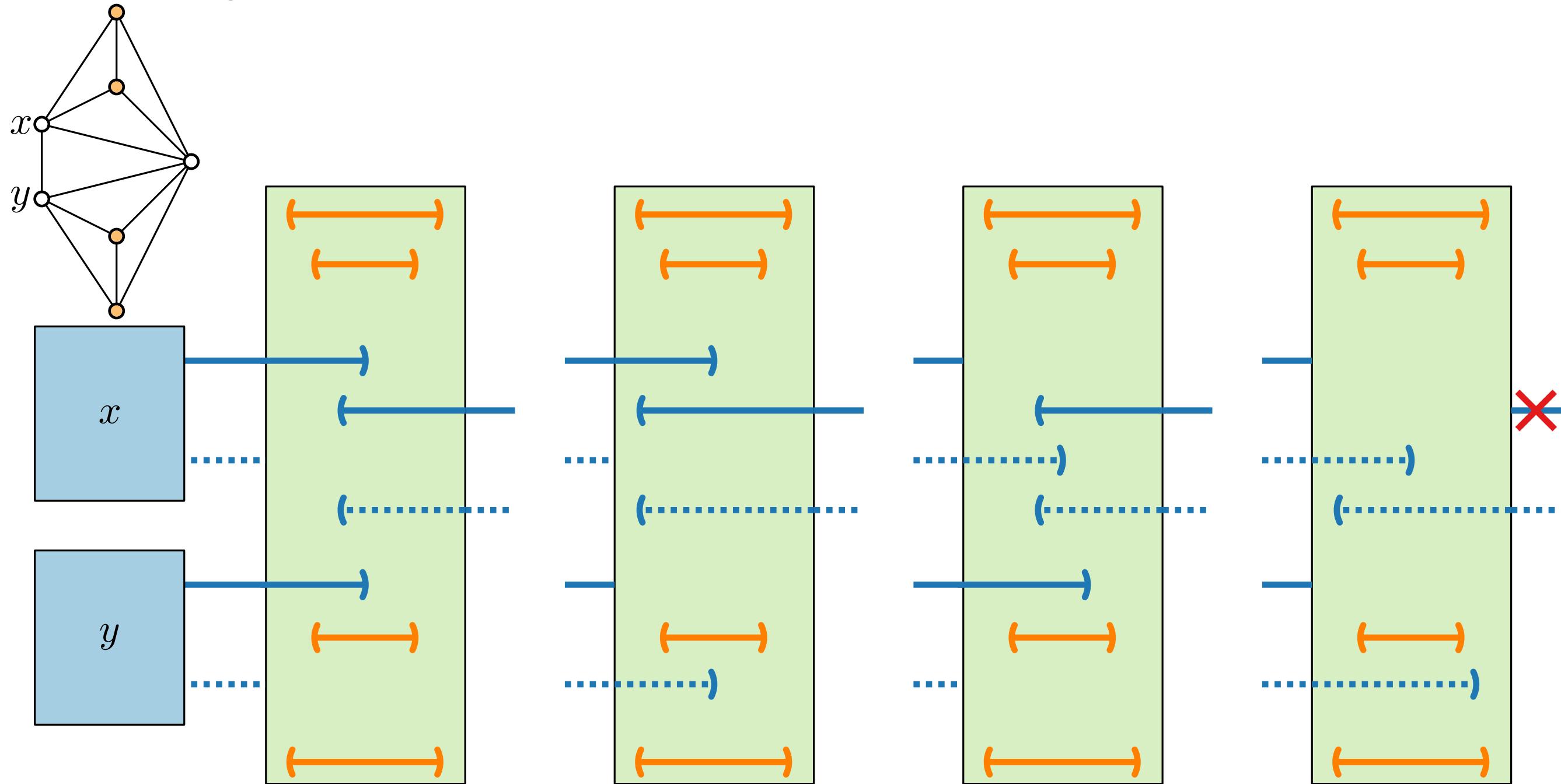


# Clause Gadget

$$x \vee y \vee z$$



# OR' Gadget



# Discussion

- *Rectangular*  $\varepsilon$ -Bar Visibility Representation Extension can be solved in  $O(n \log^2 n)$  time for *st*-graphs.
- $\varepsilon$ -Bar Visibility Representation Extension is NP-complete.
- $\varepsilon$ -Bar Visibility Representation Extension is NP-complete for (series-parallel) *st*-graphs when restricted to the *Integer Grid* (or if any fixed  $\varepsilon > 0$  is specified).

## Open Problems:

- Can ~~rectangular~~  $\varepsilon$ -Bar Visibility Representation Extension be solved in polynomial time for *st*-graphs? For DAGs?
- Can **Strong** Bar Visibility Recognition / Representation Extension can be solved in polynomial time for *st*-graphs?

# Literature

Main source:

- [Chaplick, Guśpiel, Gutowski, Krawczyk, Liotta '18]  
The Partial Visibility Representation Extension Problem

Referenced papers:

- [Gutwenger, Mutzel '01] A Linear Time Implementation of SPQR-Trees
- [Wismath '85] Characterizing bar line-of-sight graphs
- [Tamassia, Tollis '86] Algorithms for visibility representations of planar graphs
- [Andreea '92] Some results on visibility graphs
- [Chaplick, Dorbec, Kratchovíl, Montassier, Stacho '14]  
Contact representations of planar graphs: Extending a partial representation is hard