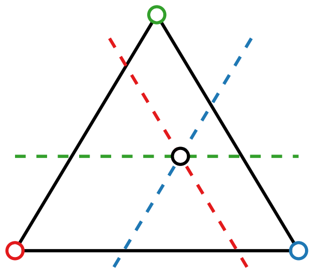
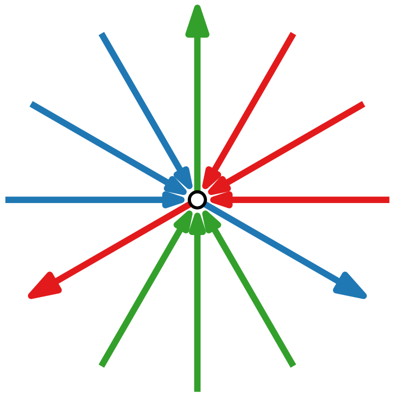
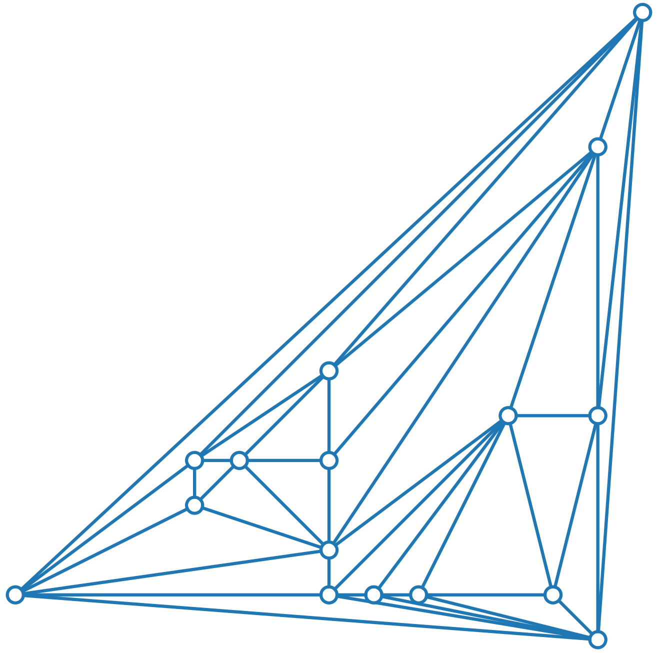
Julius-Maximilians-
UNIVERSITÄT
WÜRZBURG

Visualization of Graphs

Lecture 4:
Straight-Line Drawings of Planar Graphs II:
Schnyder Woods







Johannes Zink

Summer semester 2024

Planar Straight-Line Drawings

Theorem. [De Fraysseix, Pach, Pollack '90]
Every n-vertex planar graph has a planar straight-line drawing of size $(2n - 4) \times (n - 2)$.

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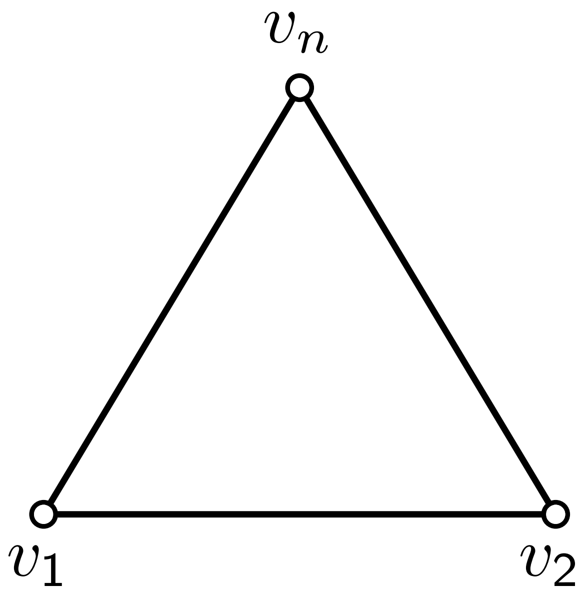
Idea.

- Fix outer triangle.



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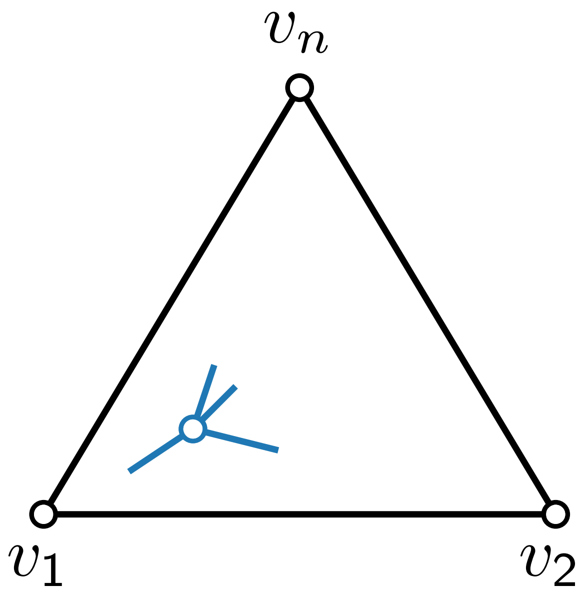

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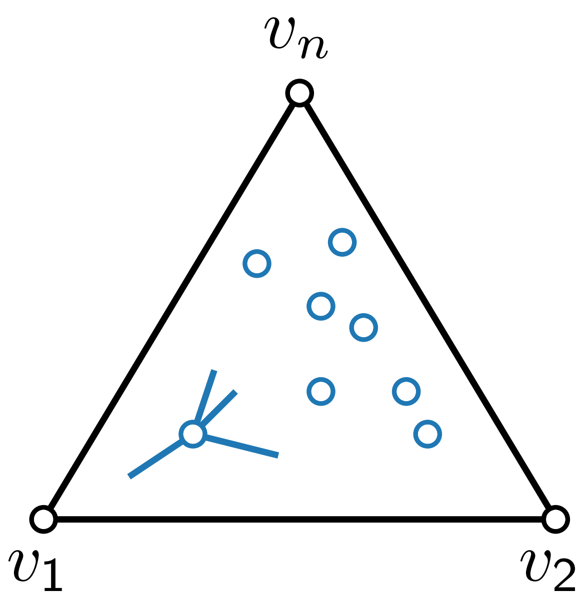

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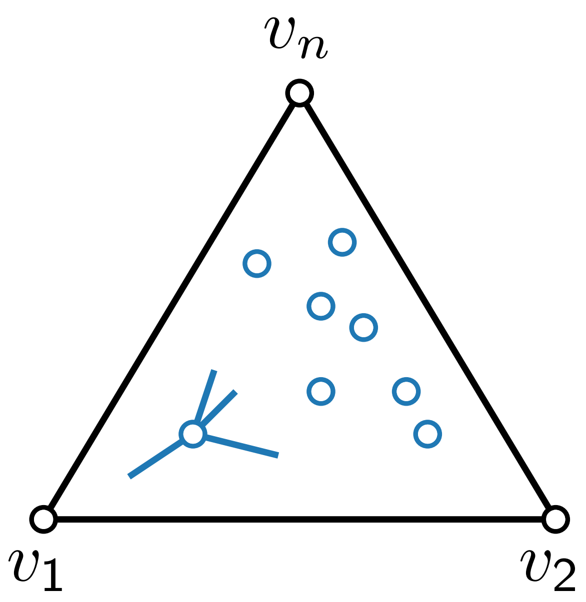

Idea.

- Fix outer triangle.
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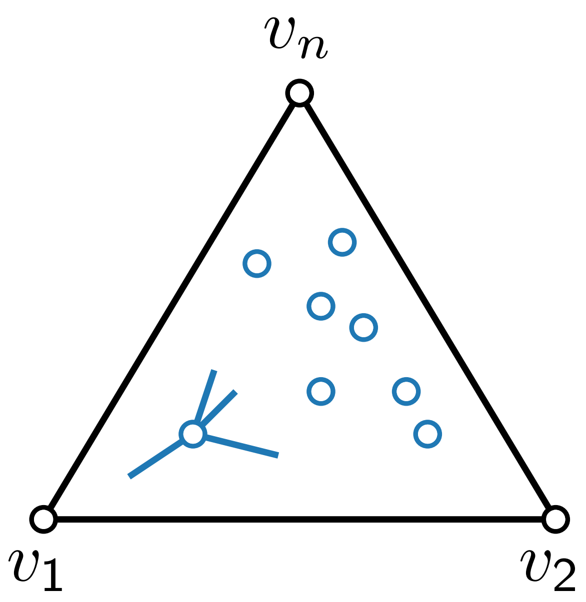
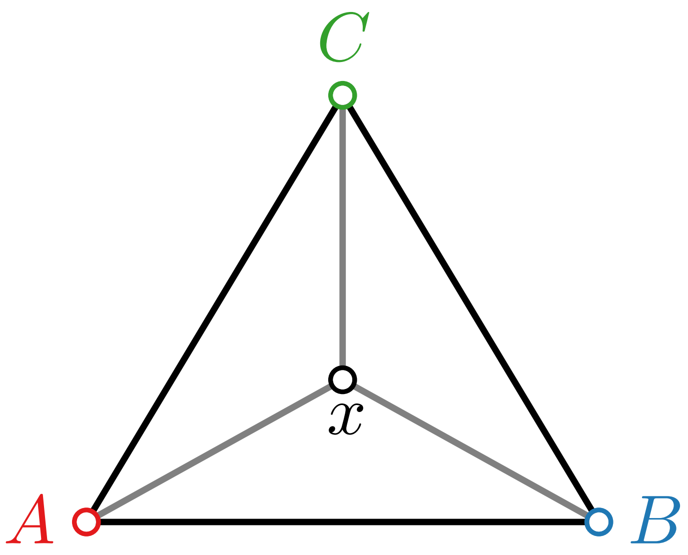


Barycentric Coordinates

Recall: barycenter$(x_1, \ldots, x_k) = \sum_{i=1}^{k} x_i / k$



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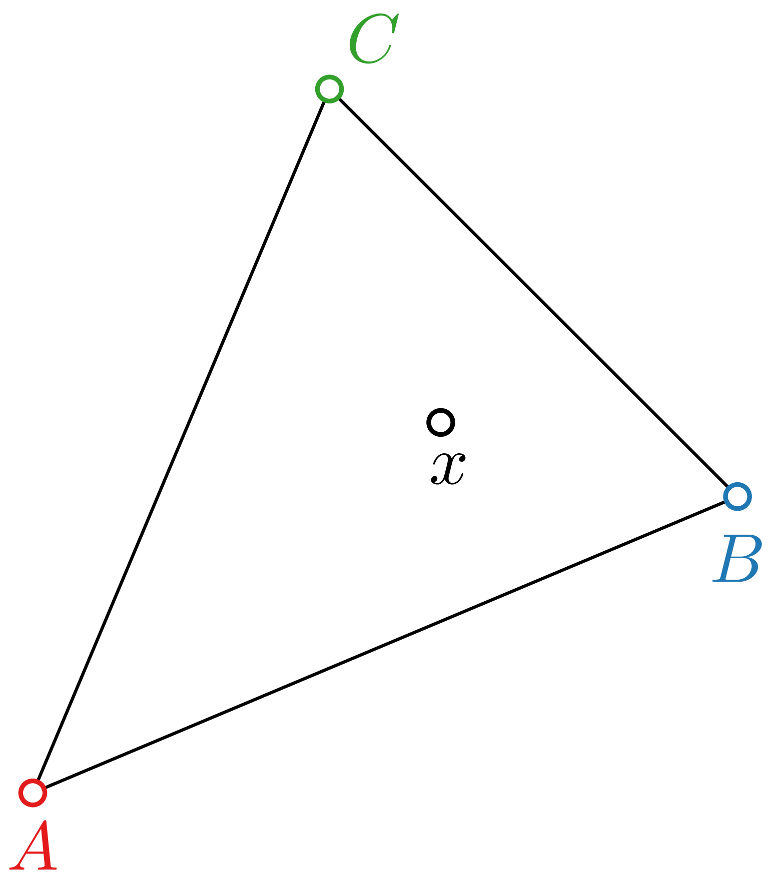

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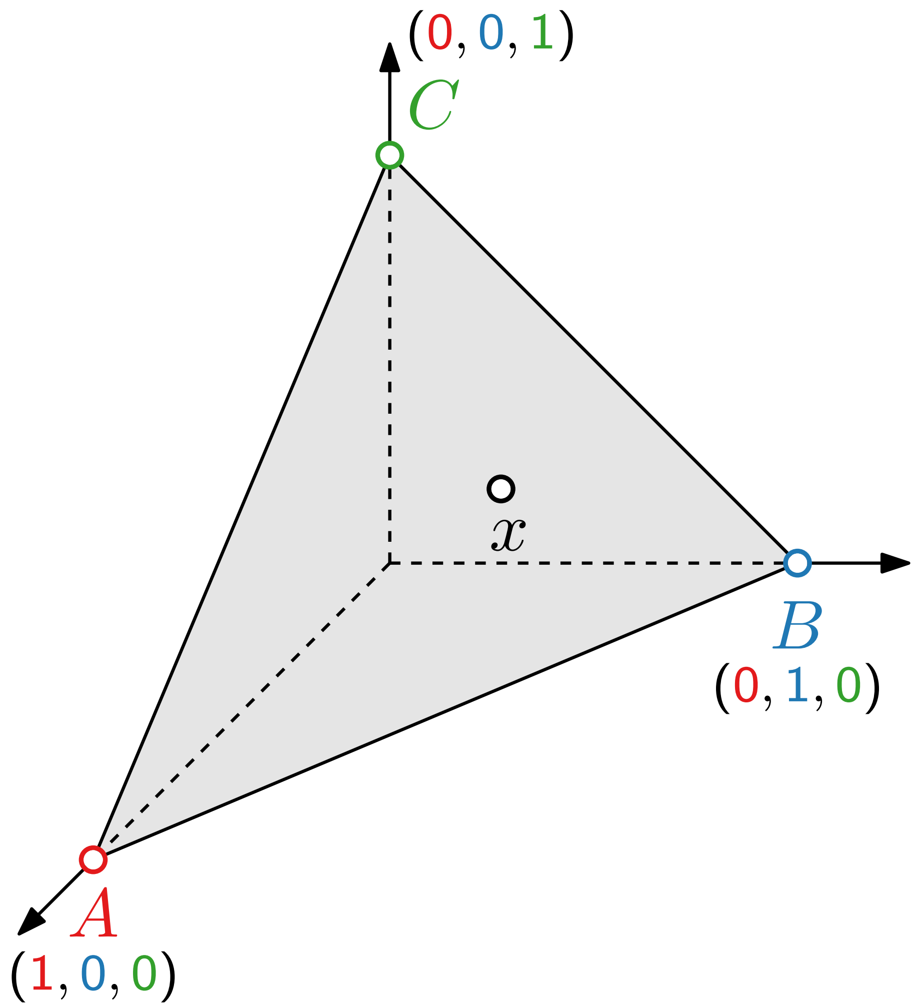

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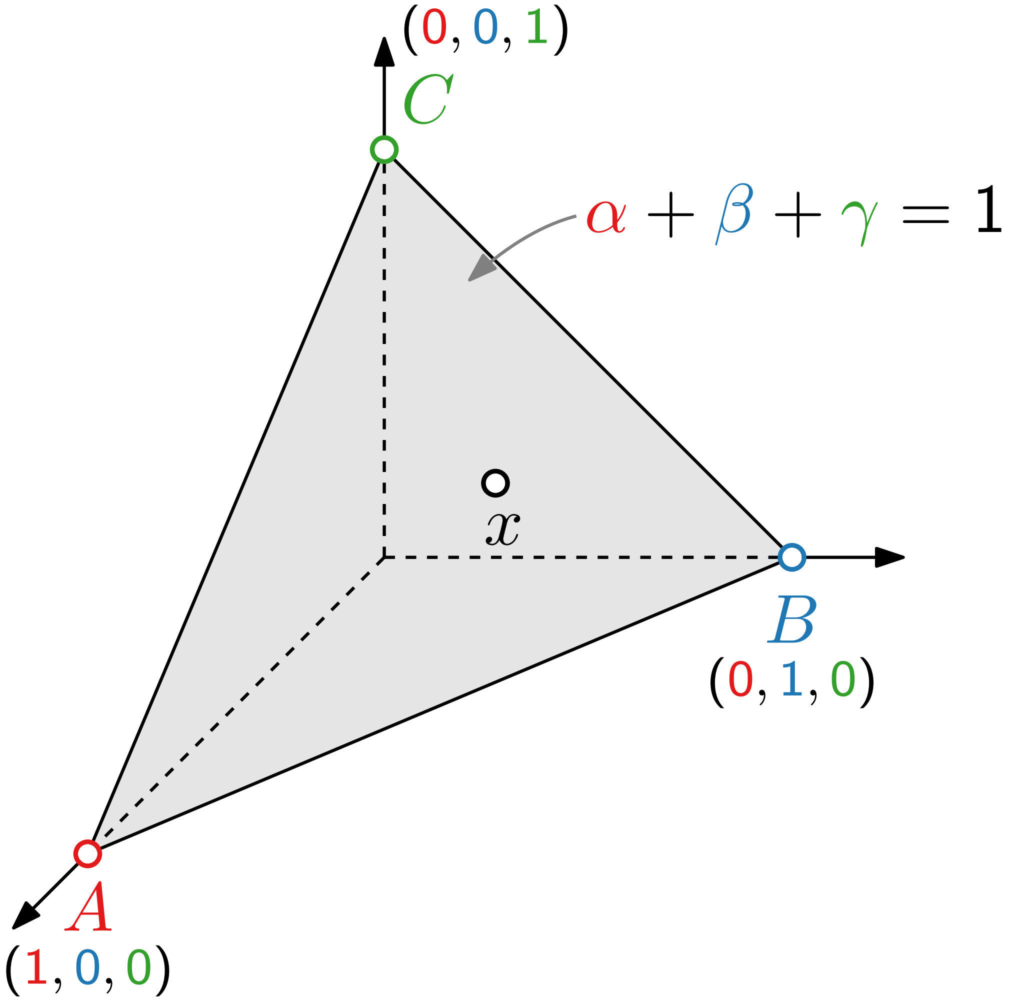

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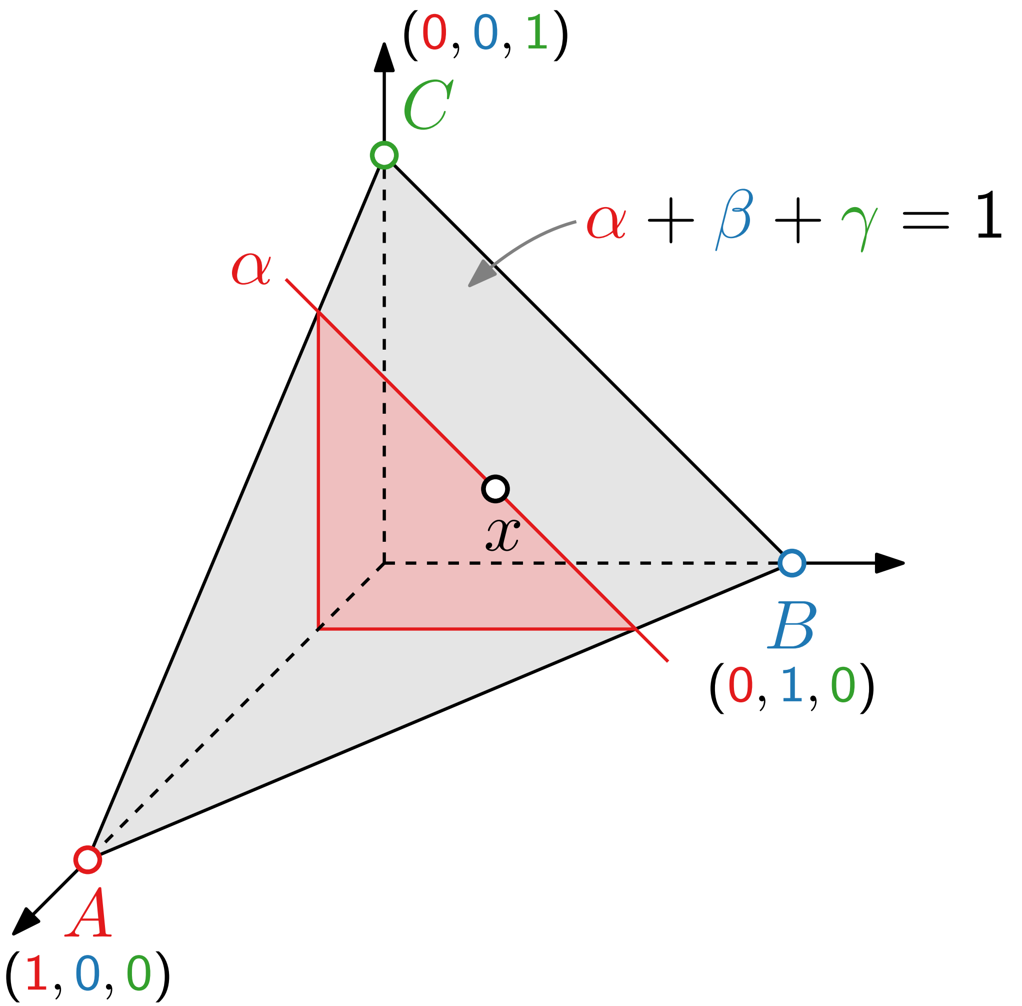

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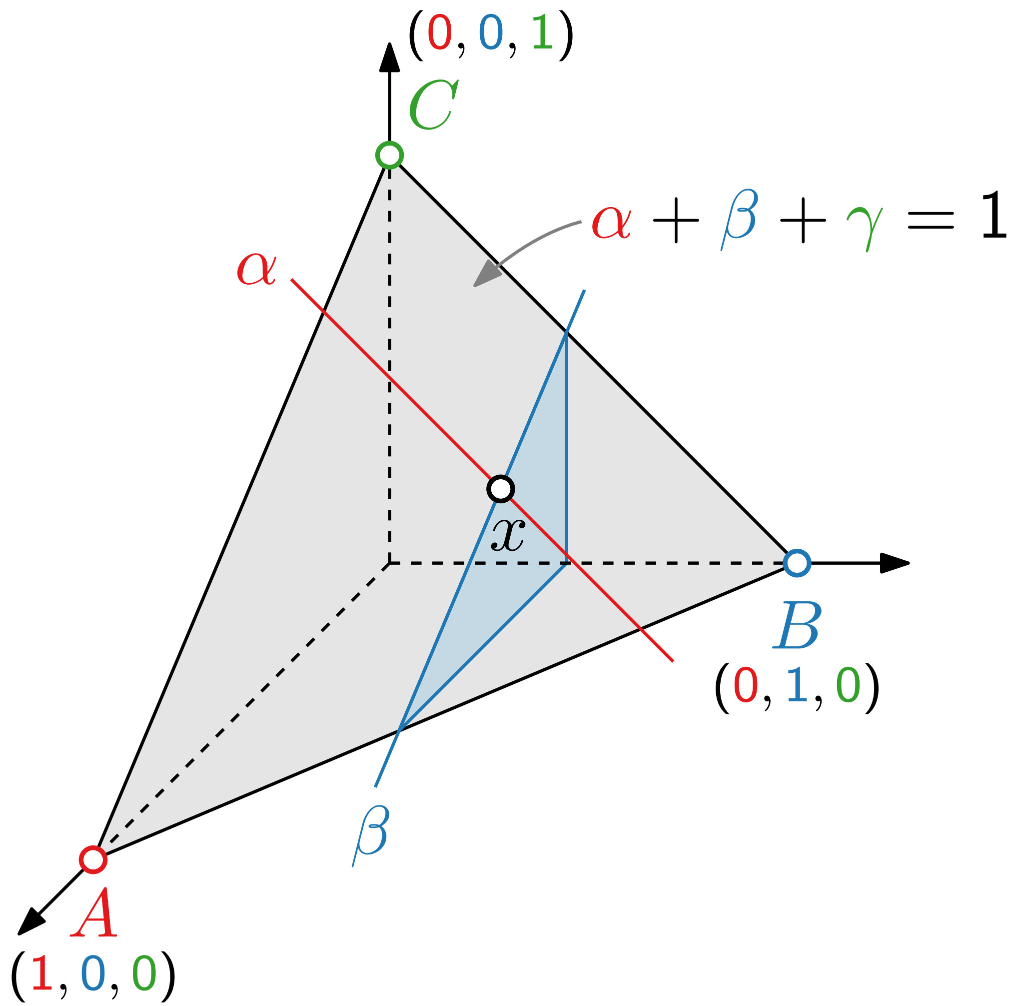

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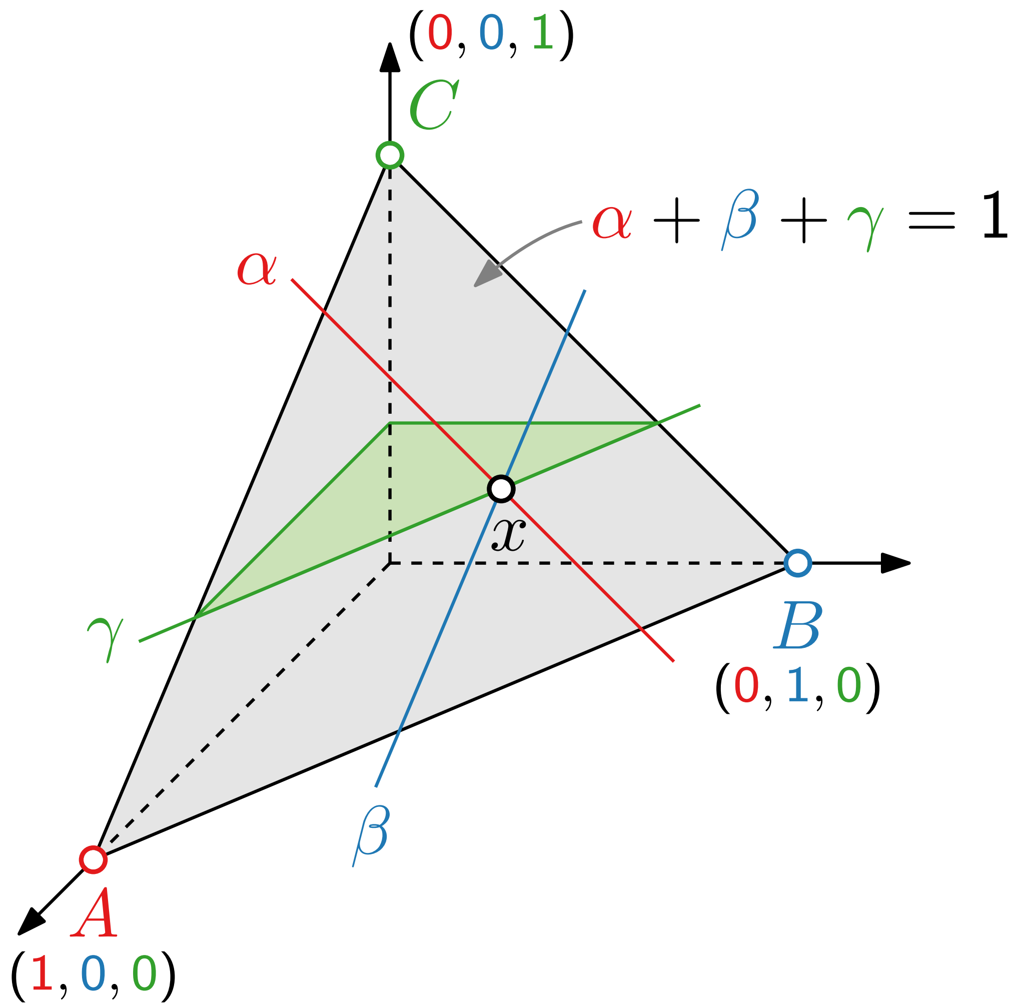

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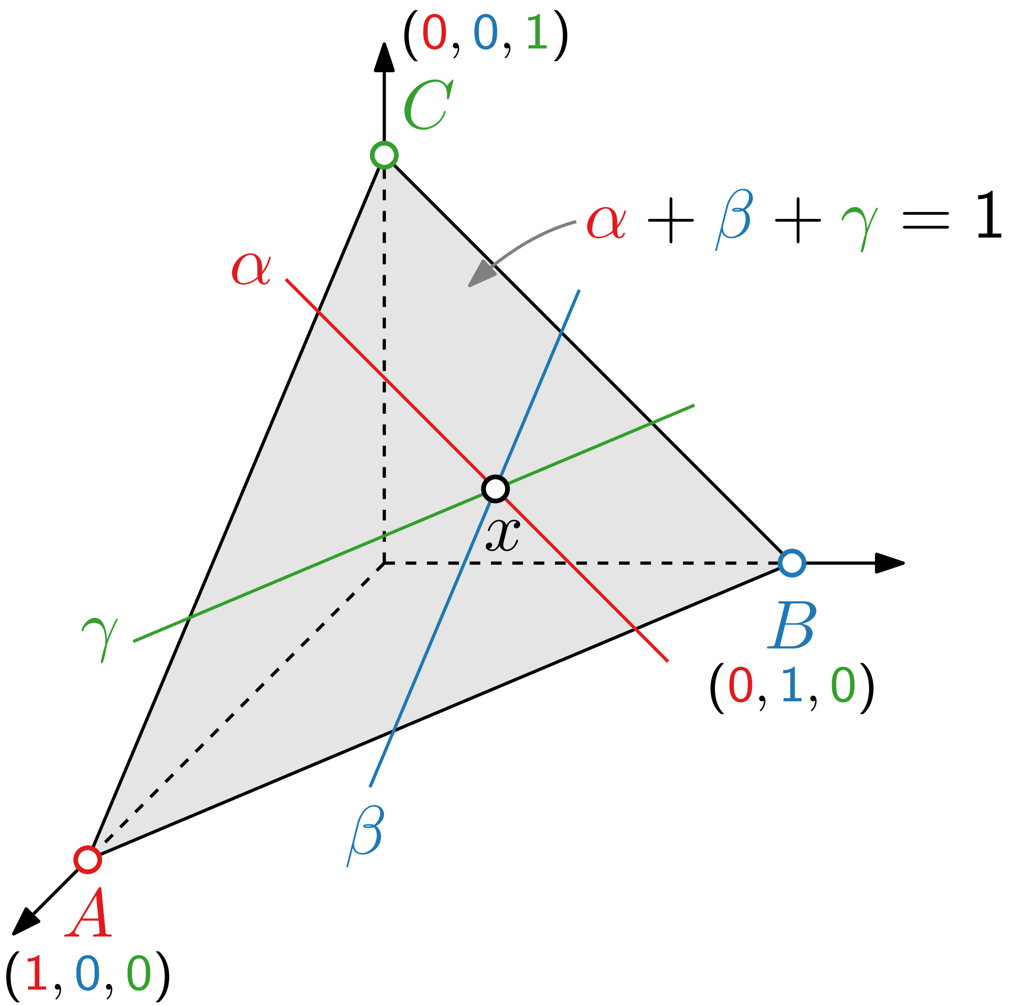

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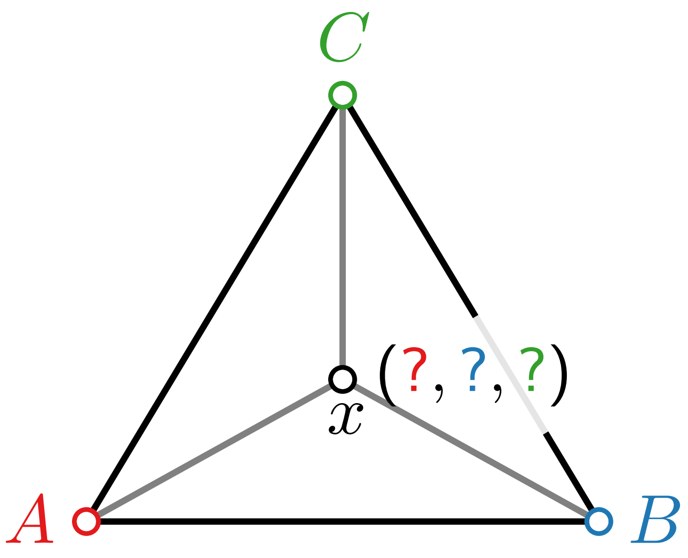

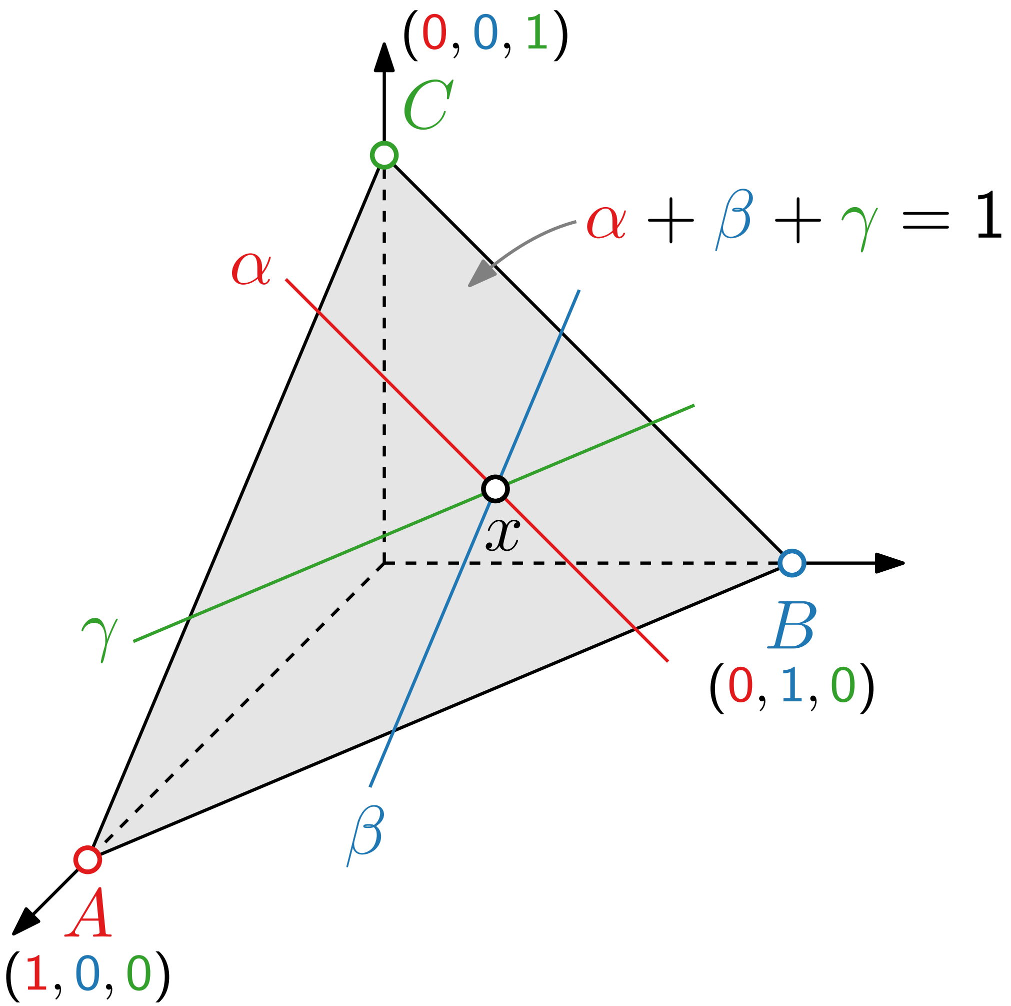

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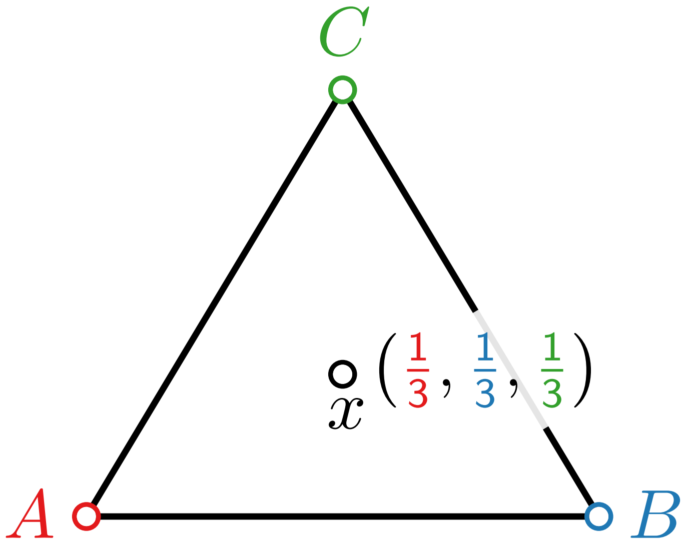

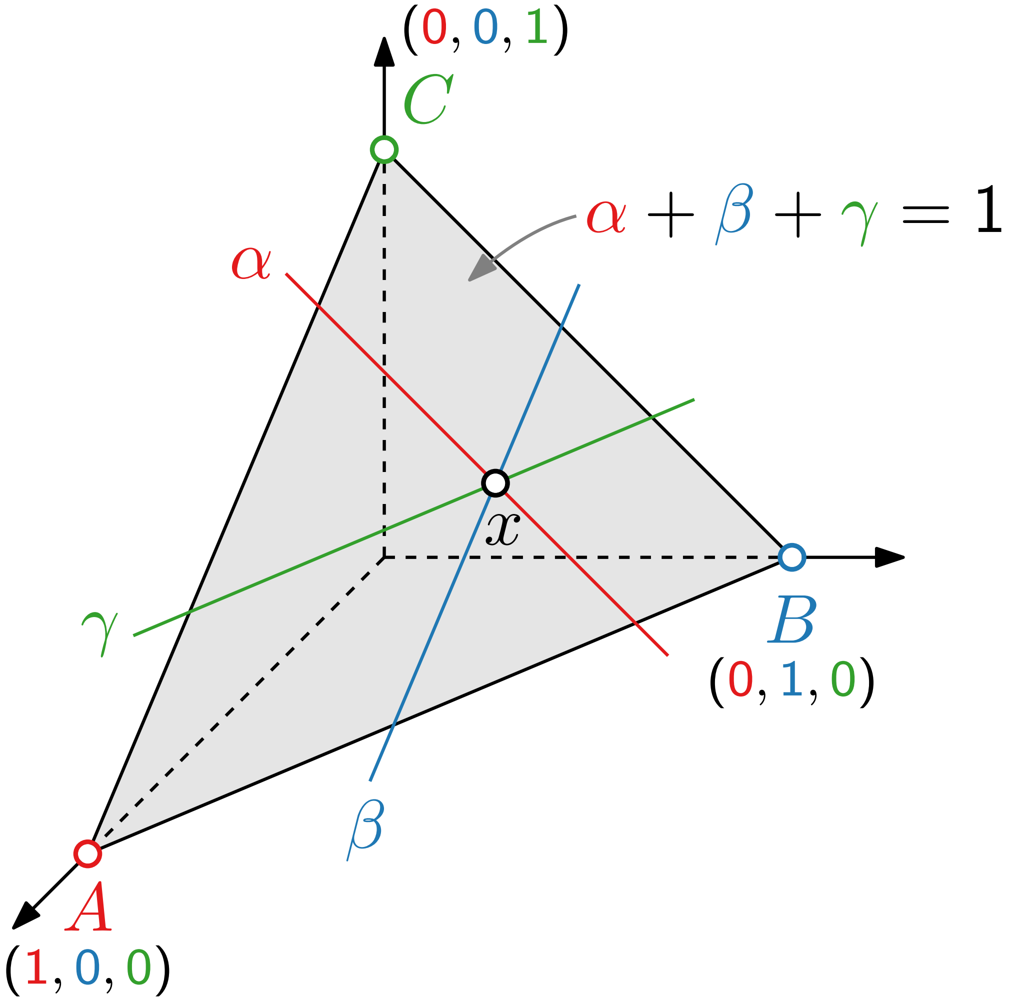

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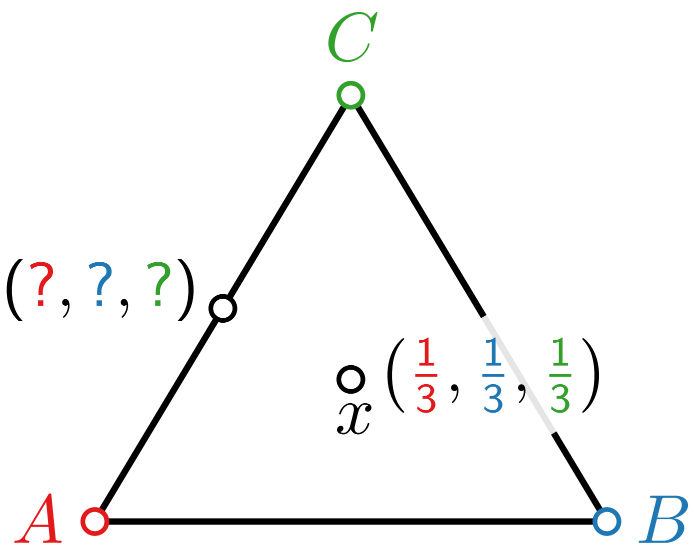

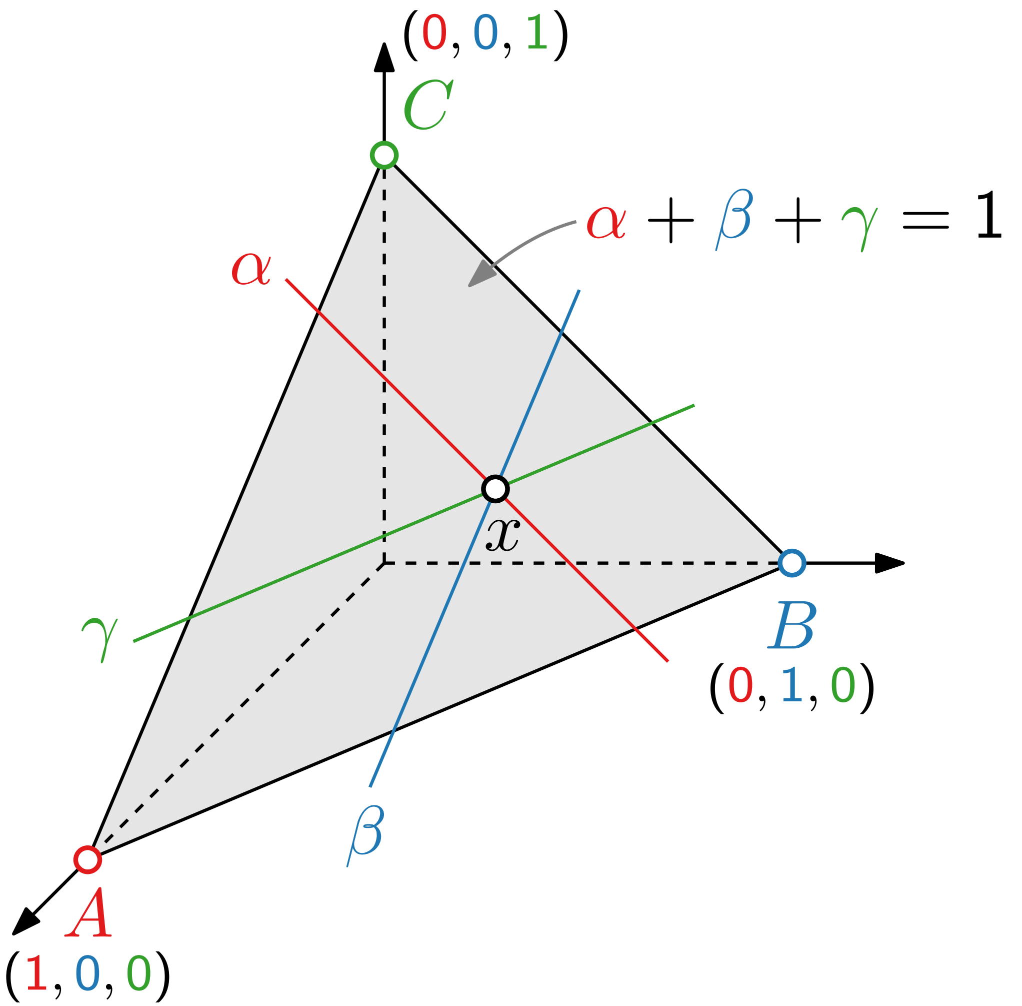

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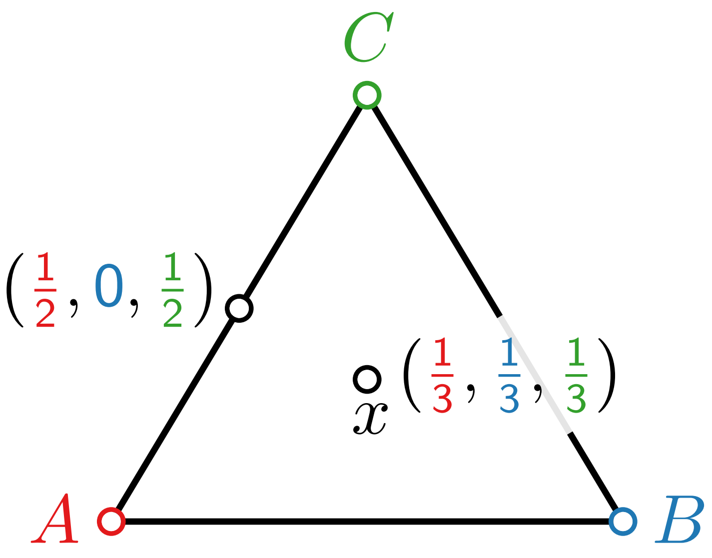

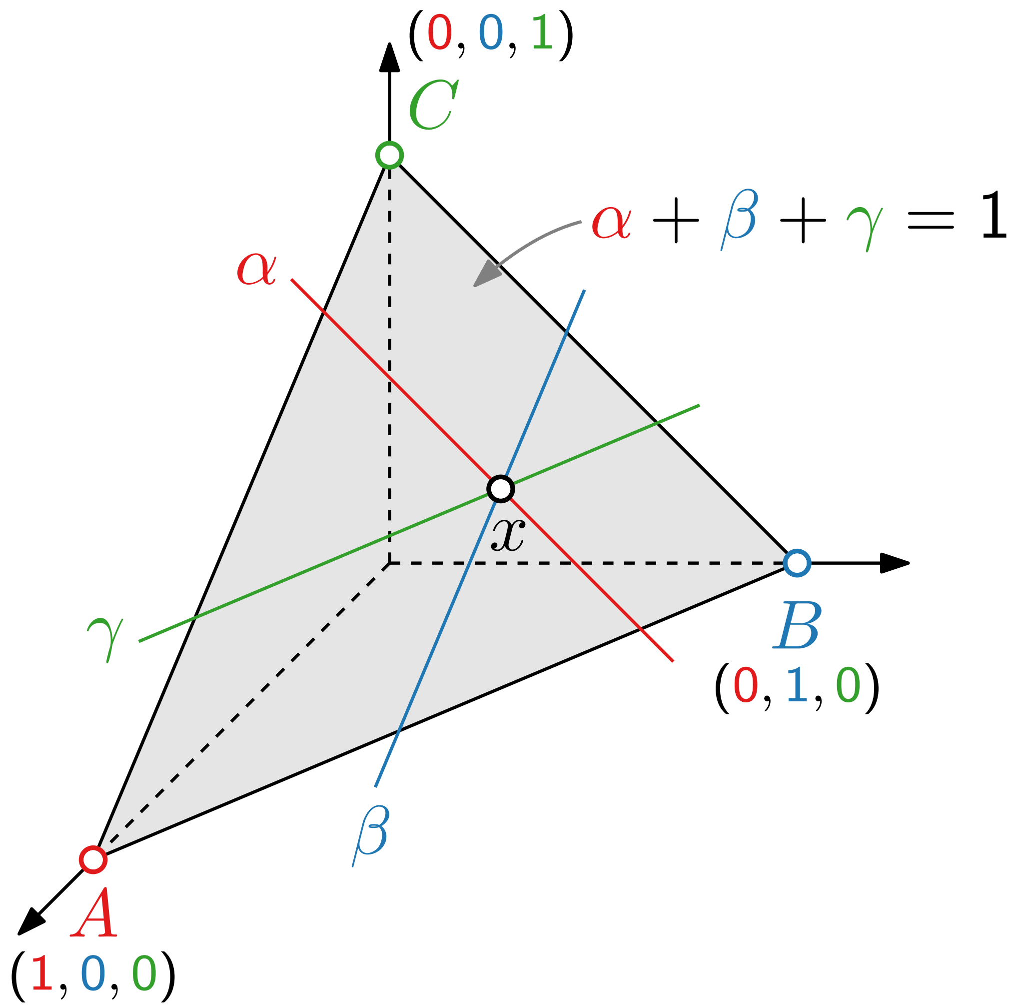

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A **barycentric representation** of a graph G is an assignment of barycentric coordinates to the vertices of G:

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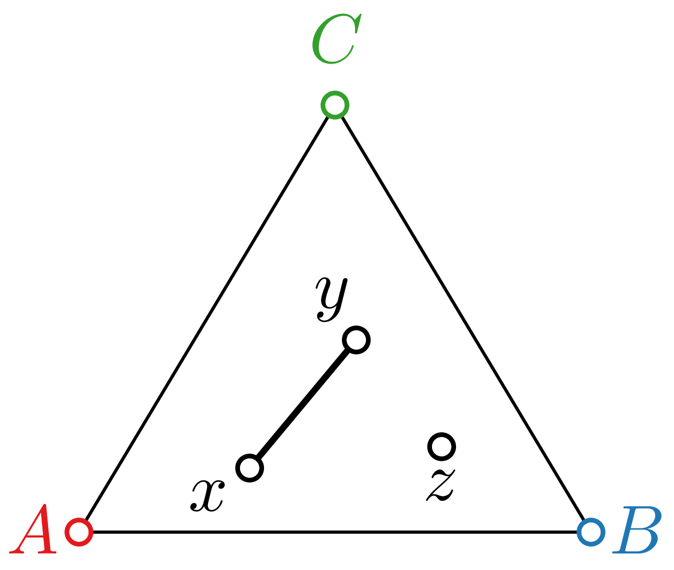

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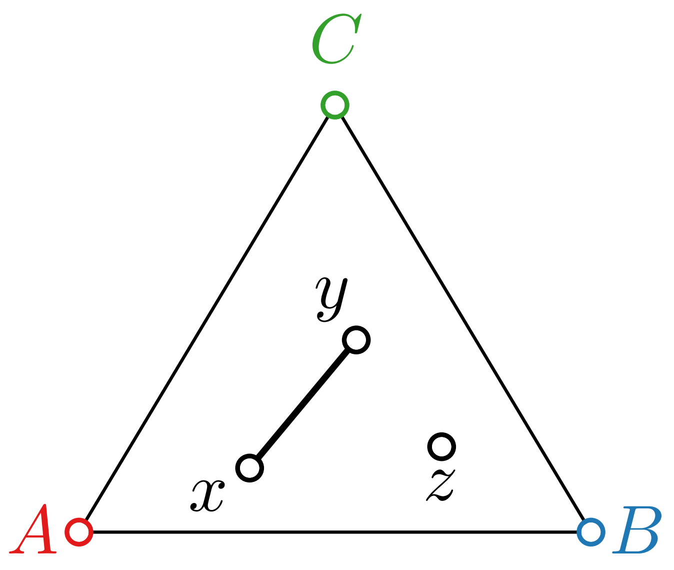

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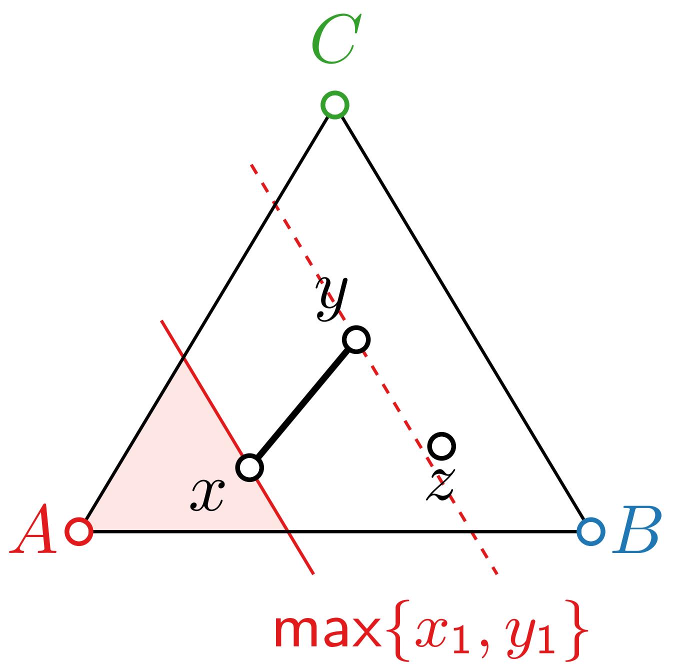

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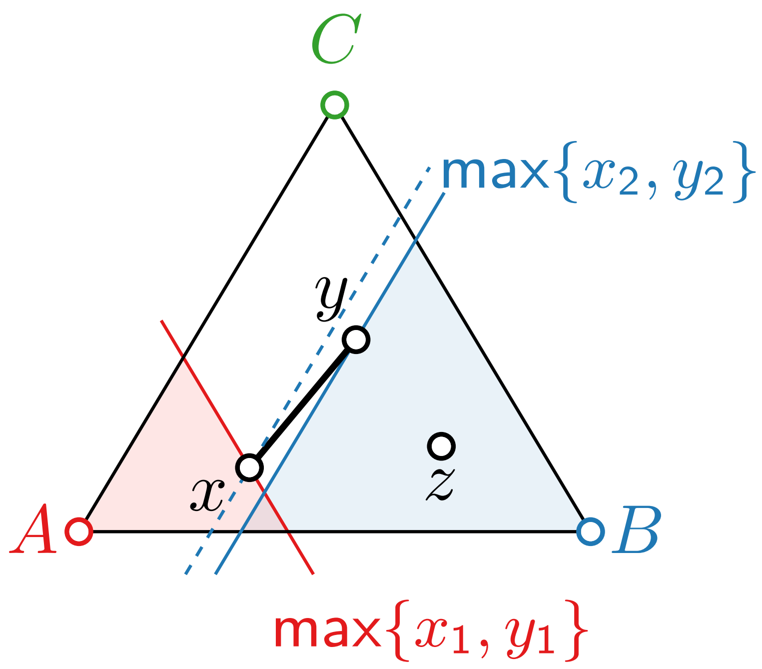

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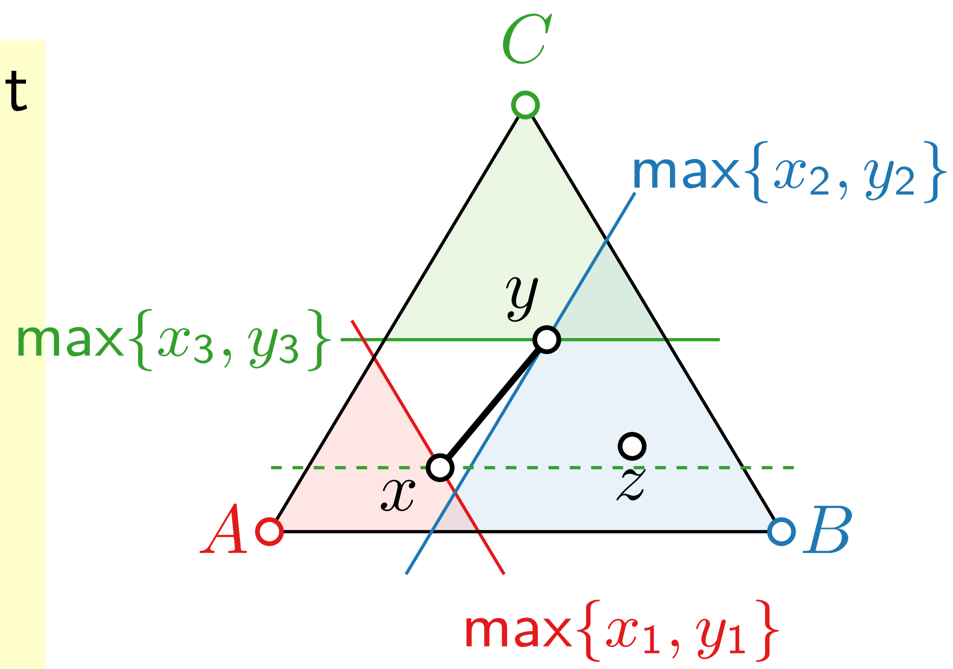

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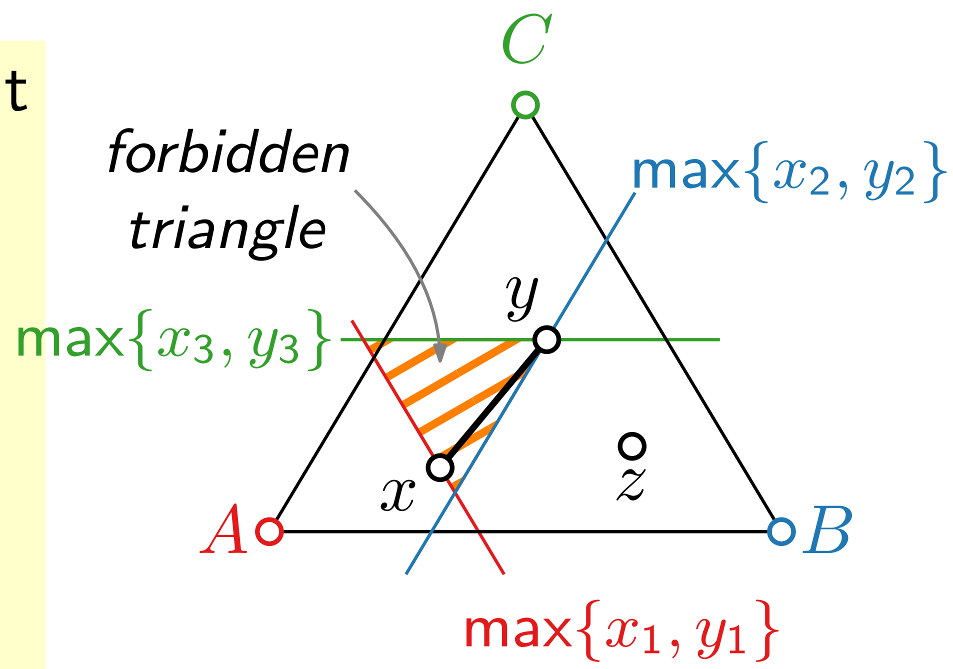

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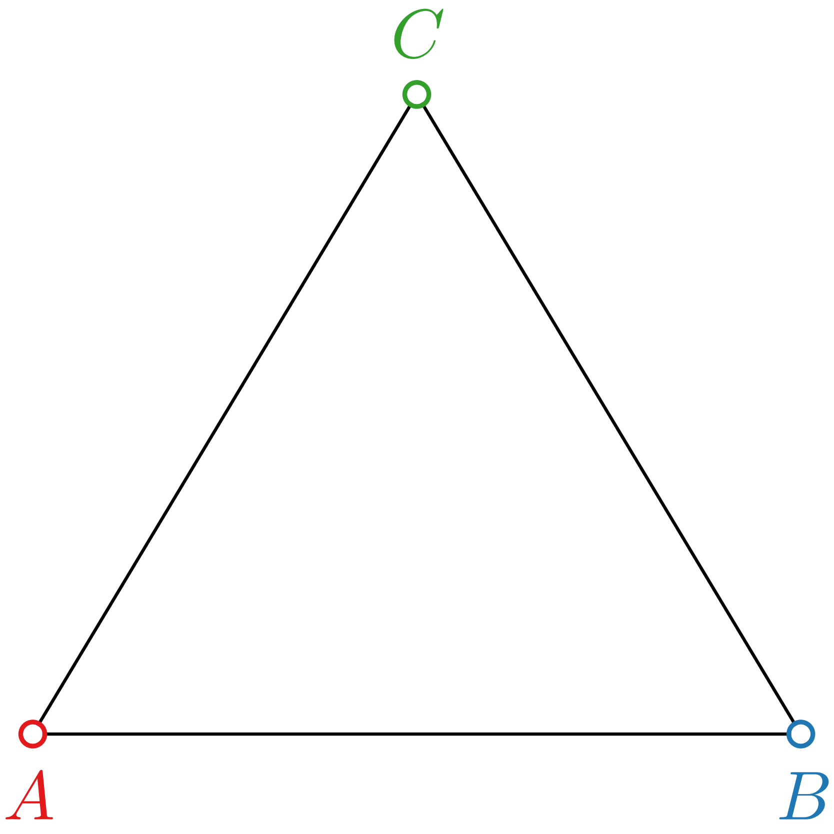

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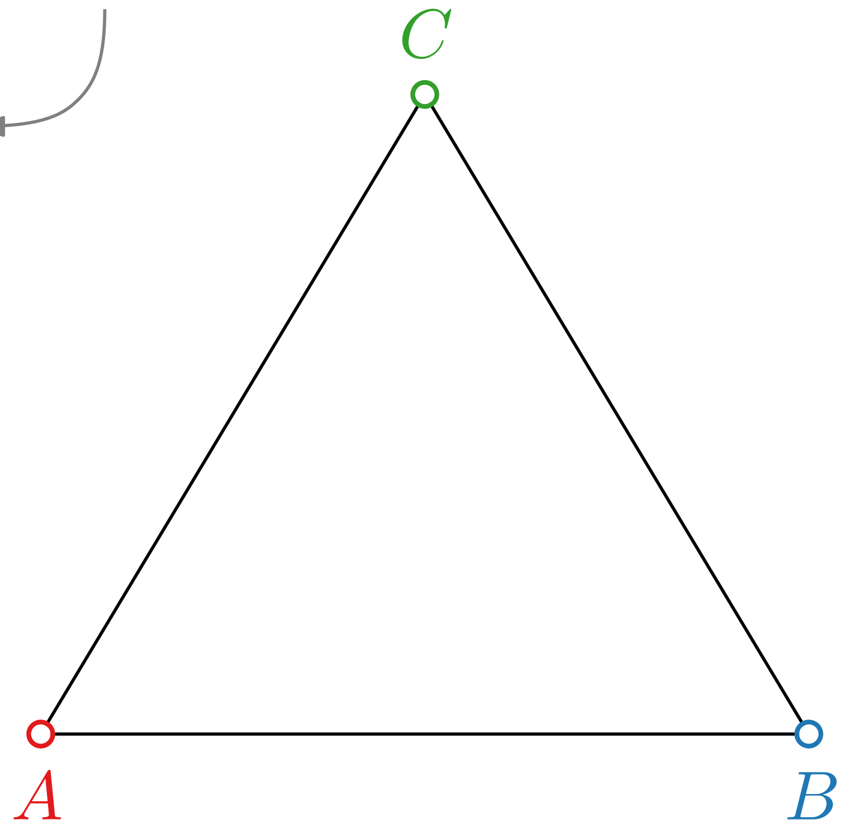
no three points on a line



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yields a **planar straight-line** drawing of G inside $\triangle ABC$.

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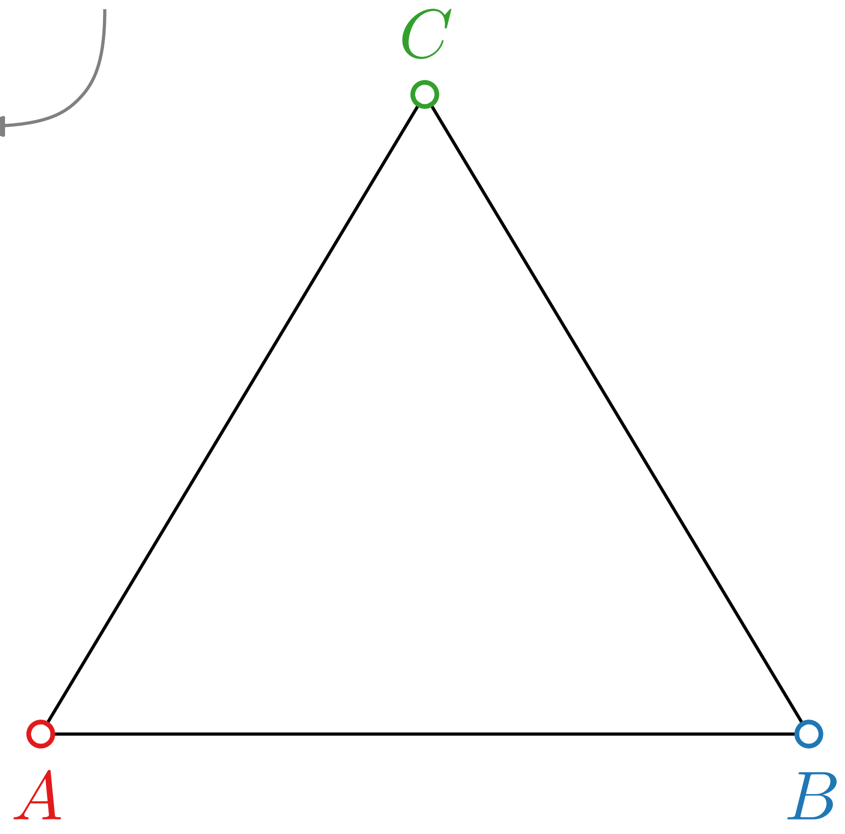

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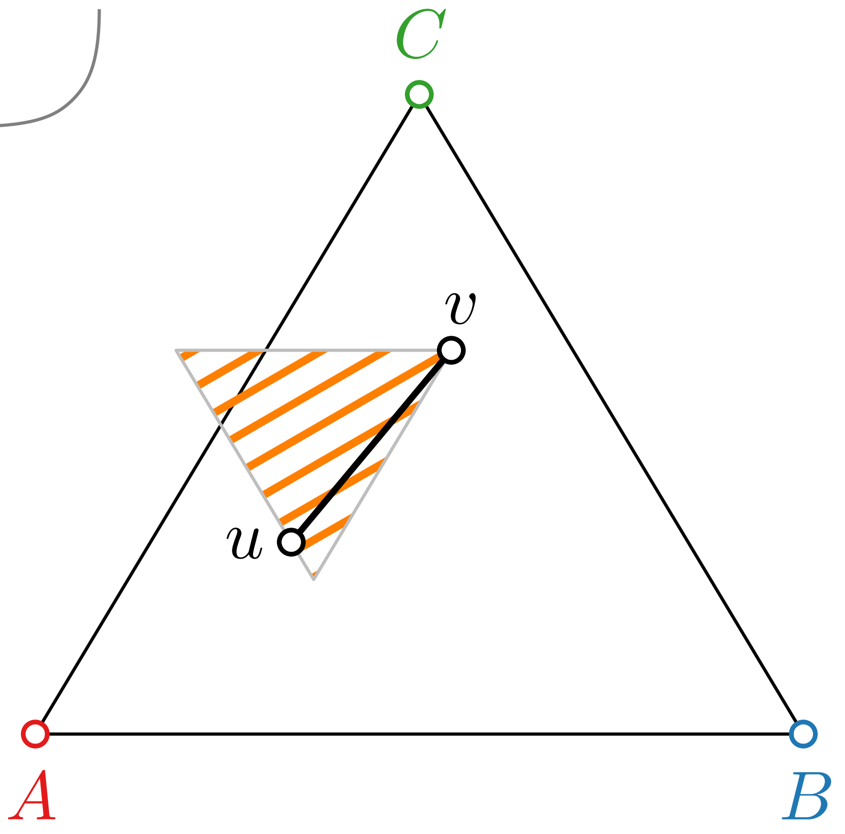

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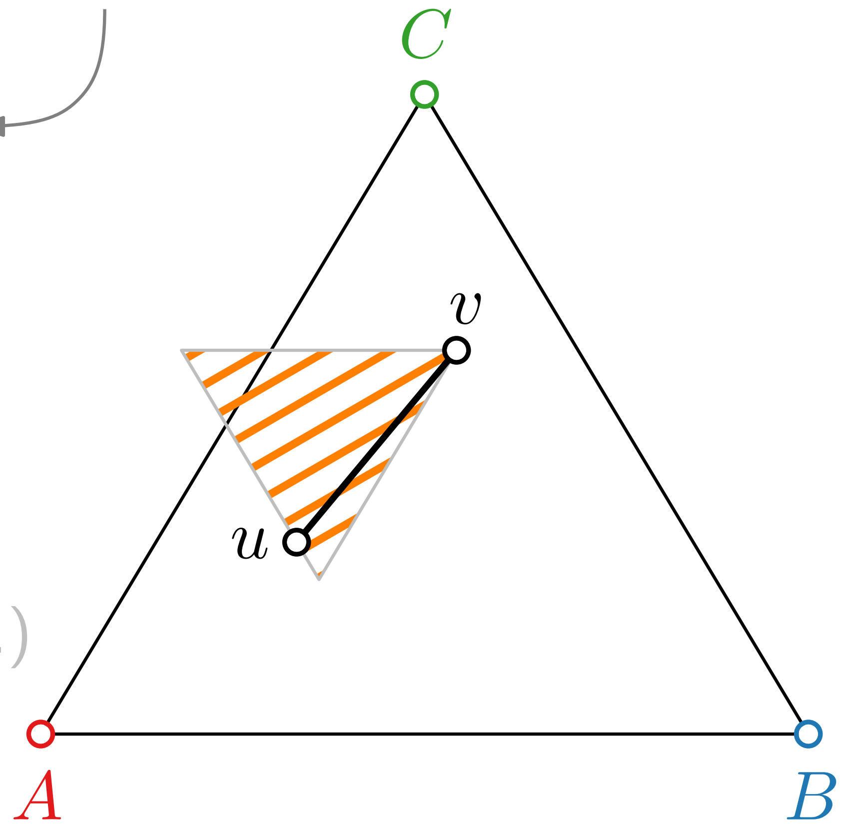

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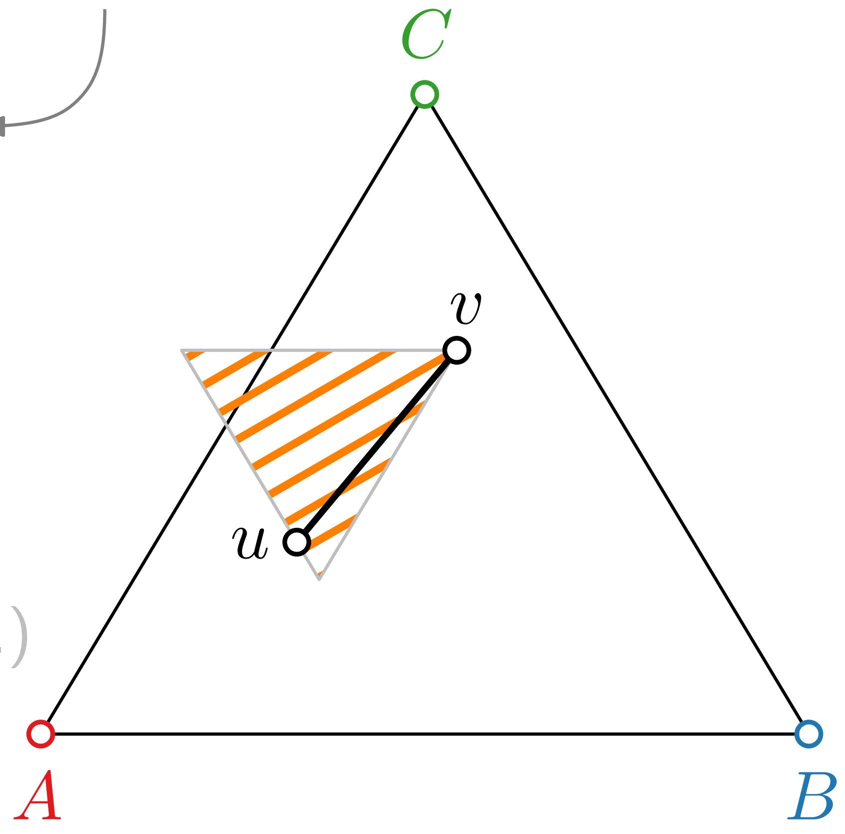

- No vertex x can lie on an edge $\{u, v\}$. (Obvious by definition.)
- No pair of edges $\{u, v\}$ and $\{u', v'\}$ crosses:

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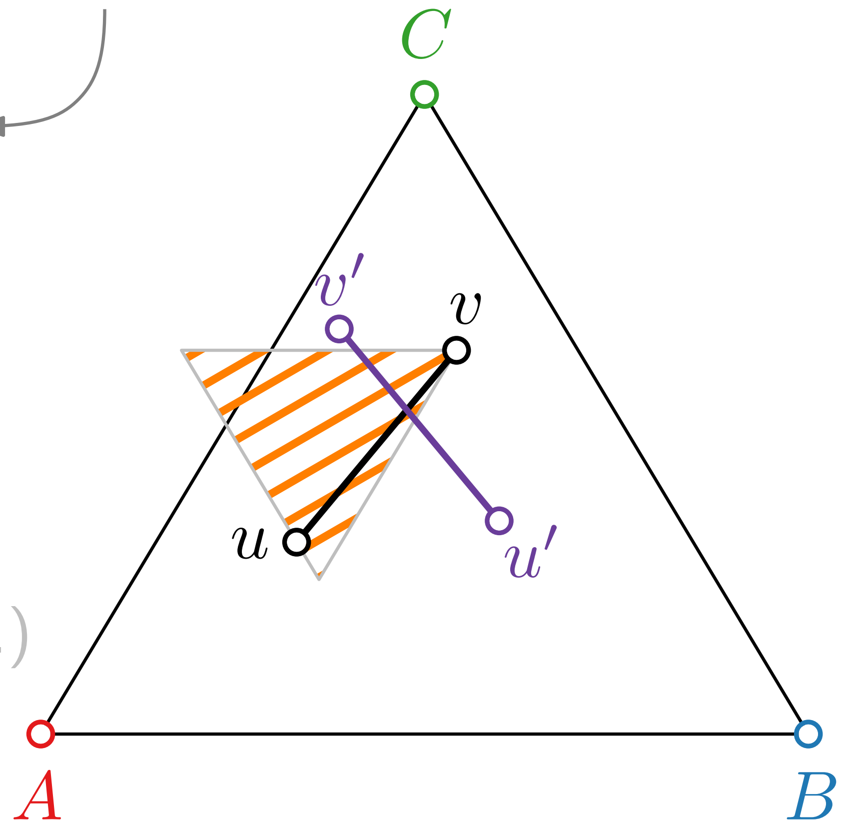

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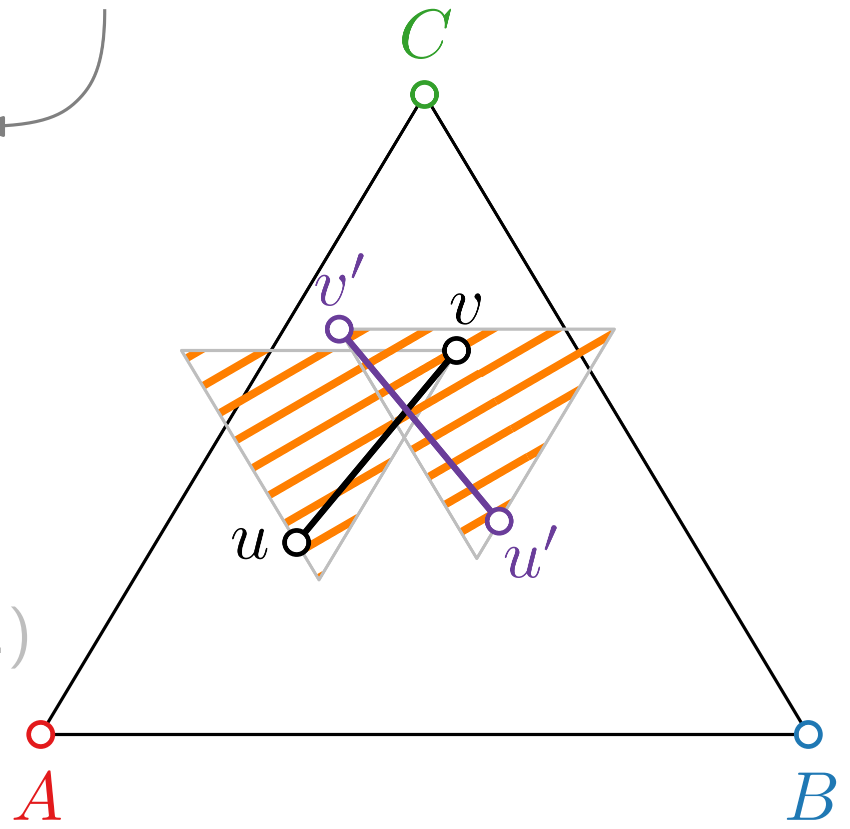

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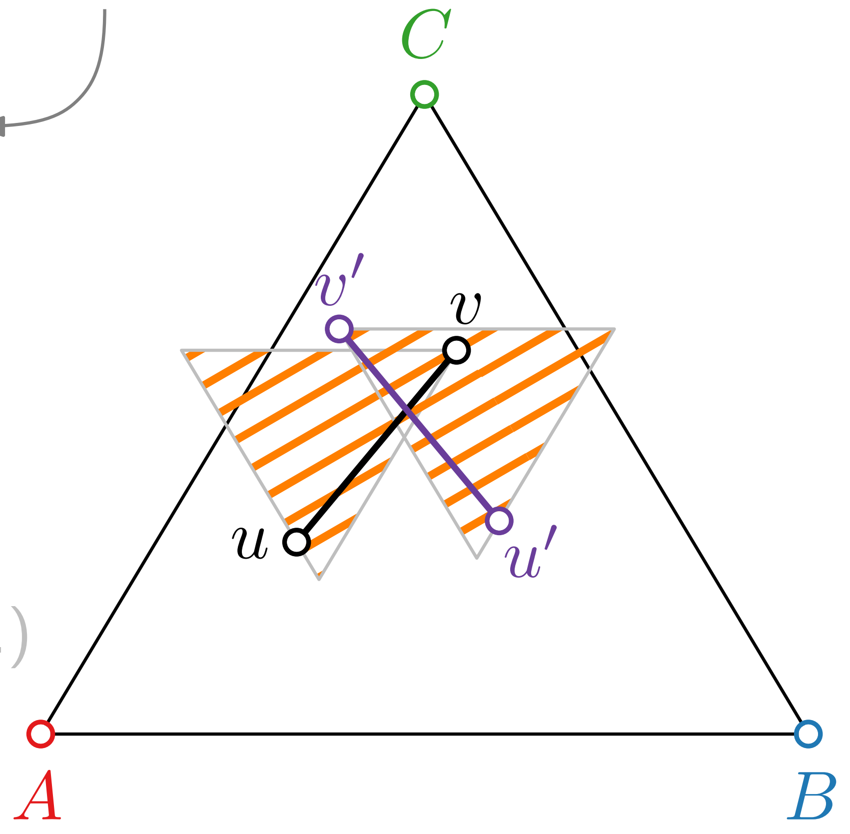

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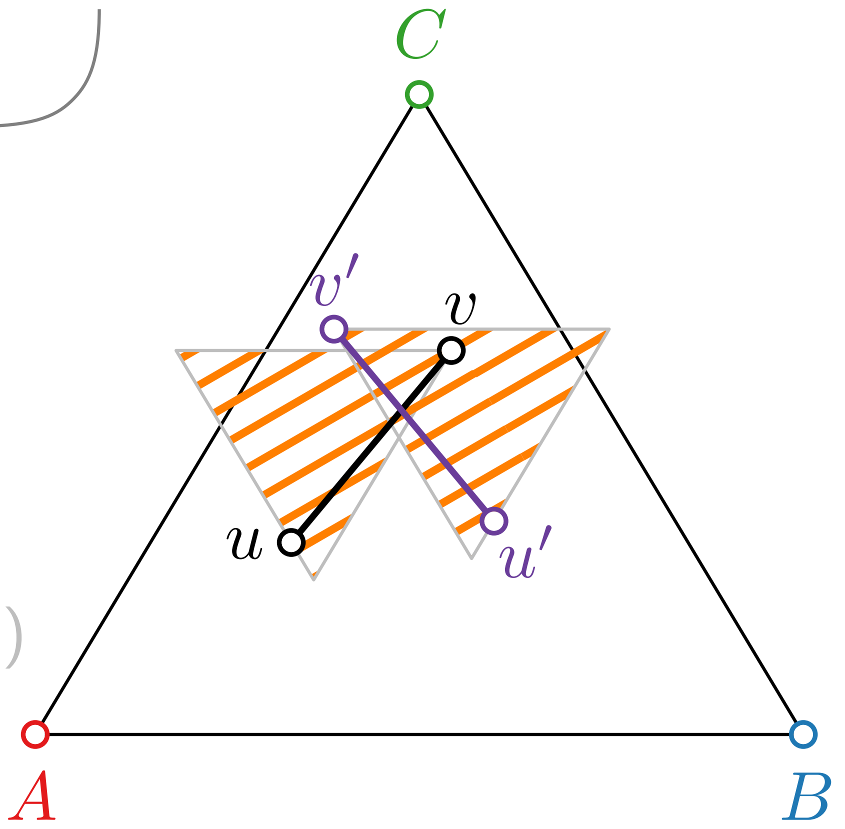

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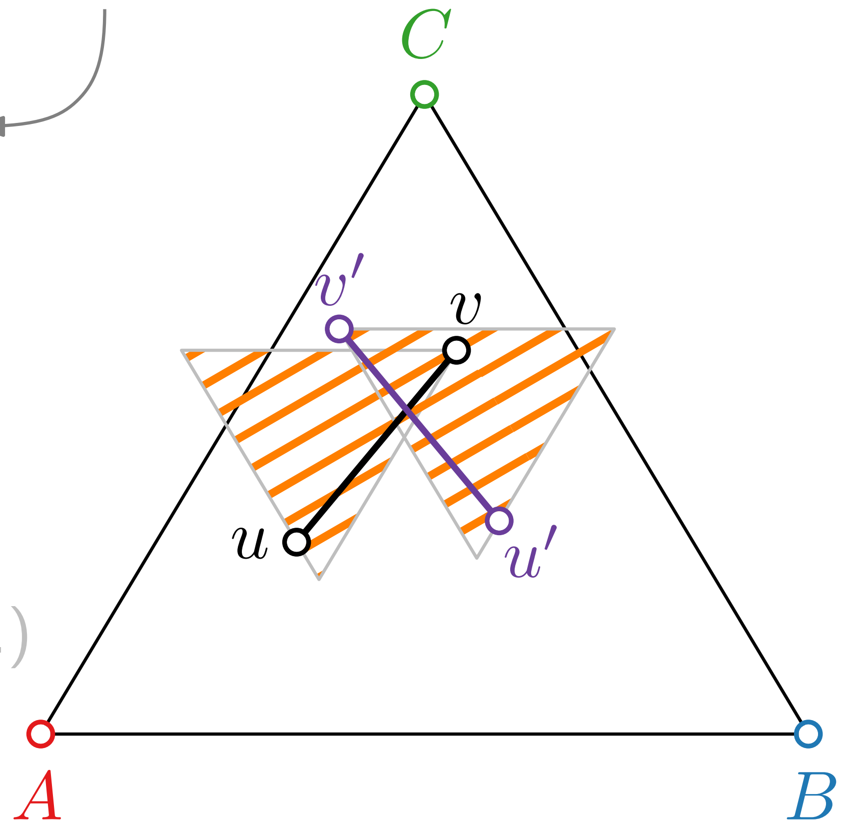

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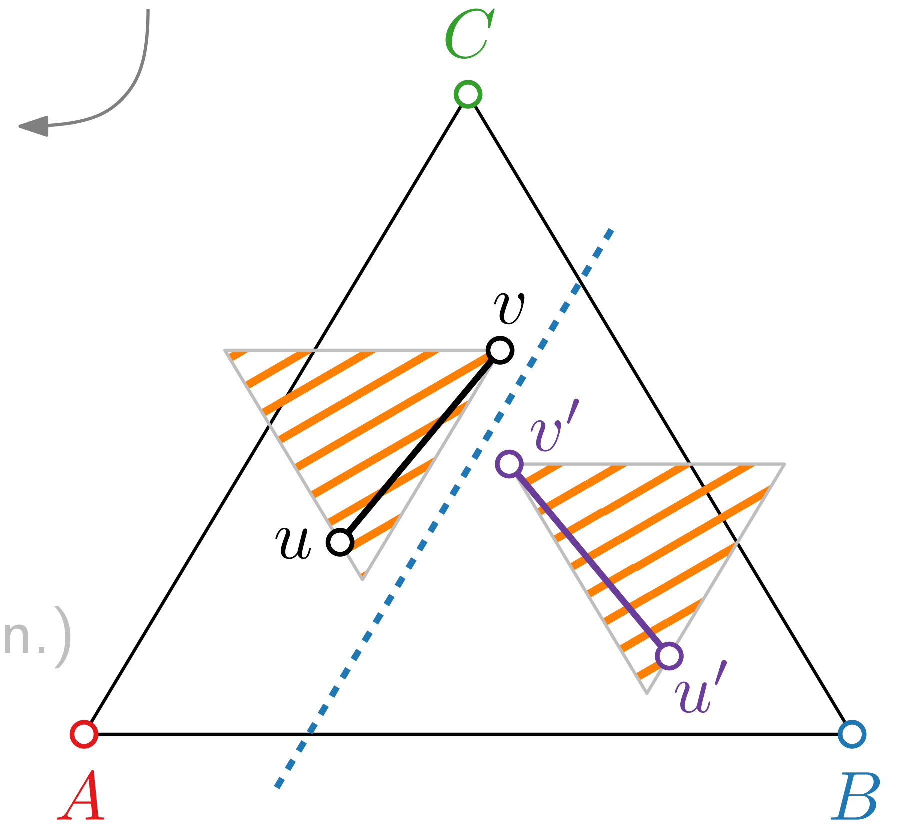

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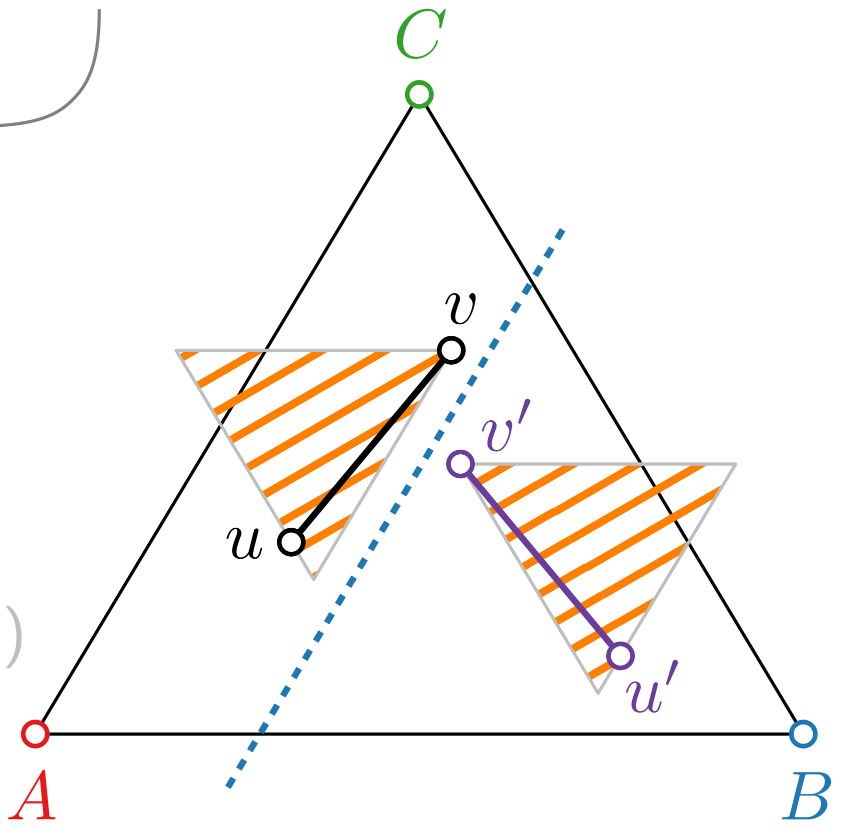

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How to find a barycentric representation?

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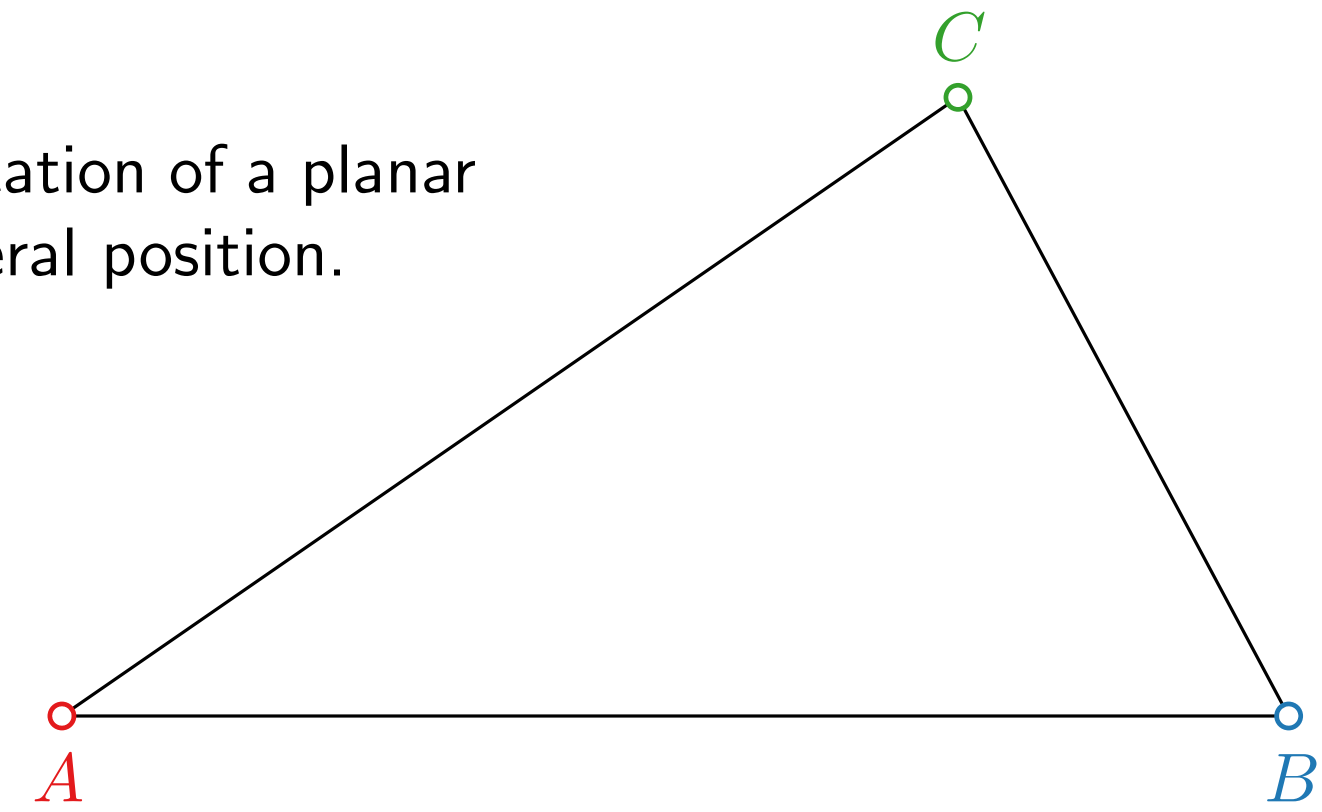

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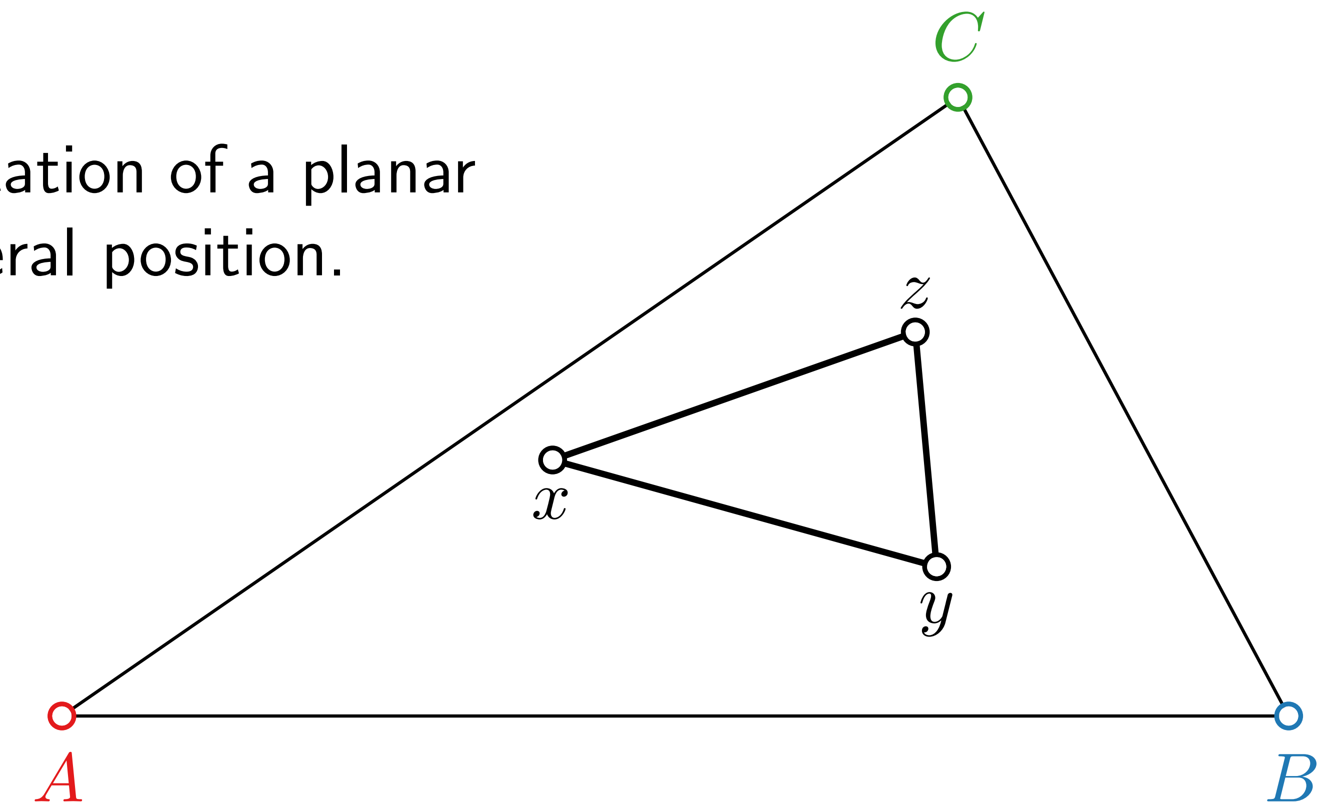

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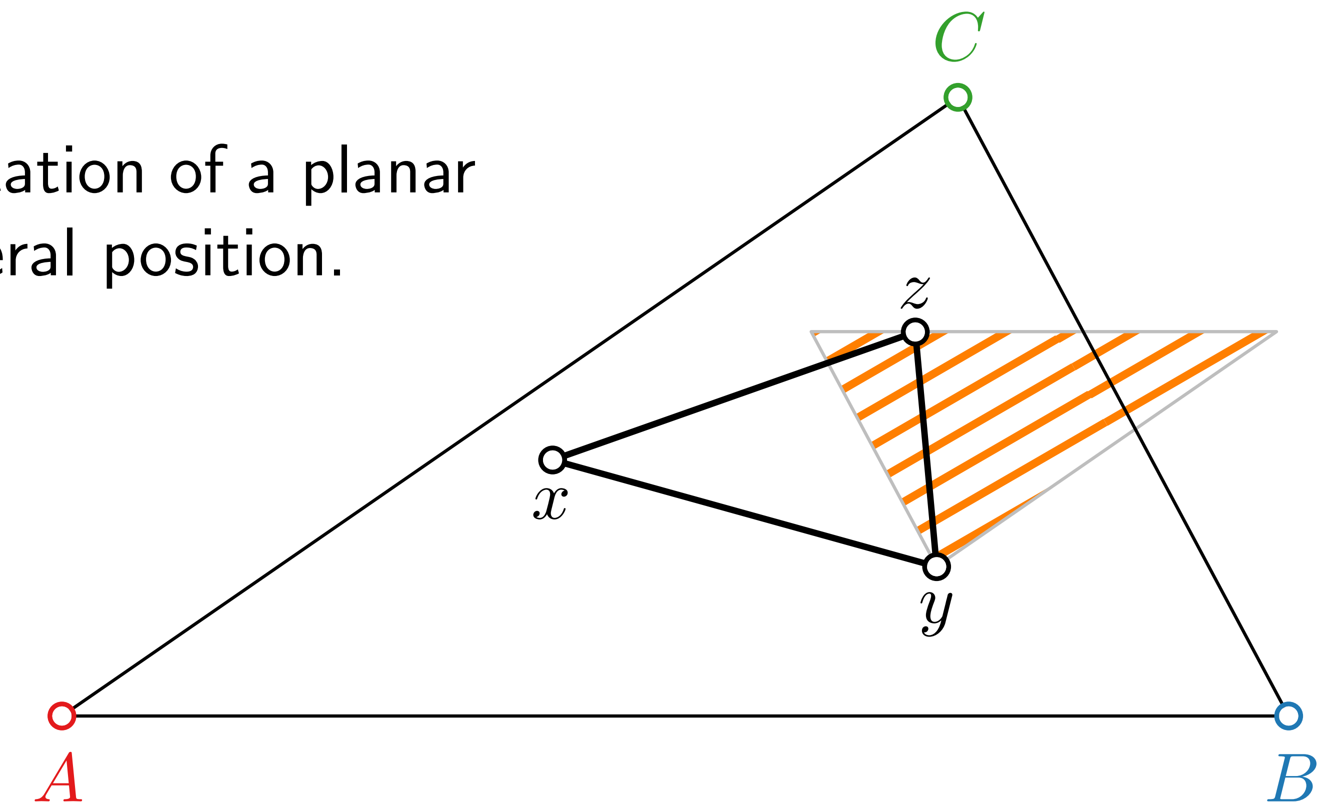

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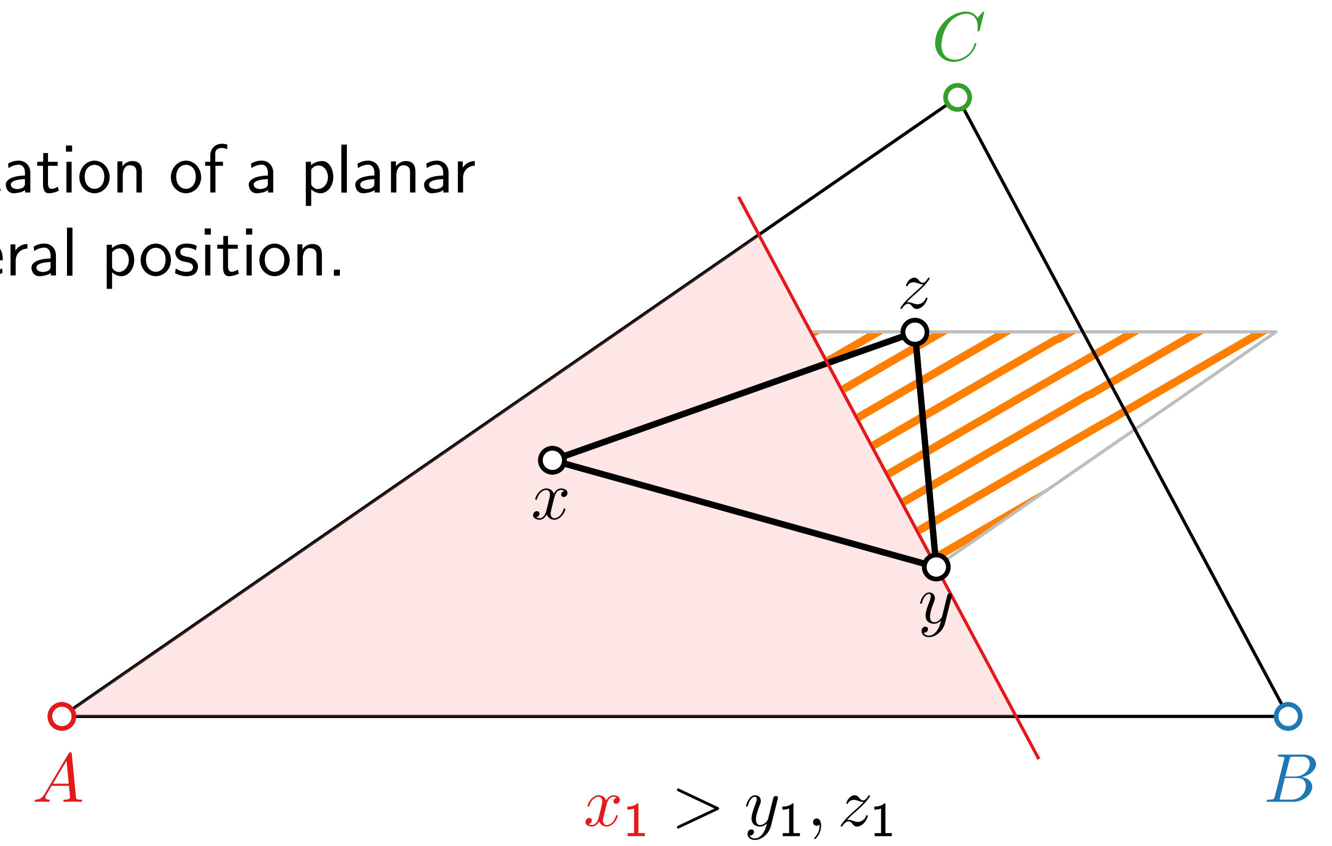

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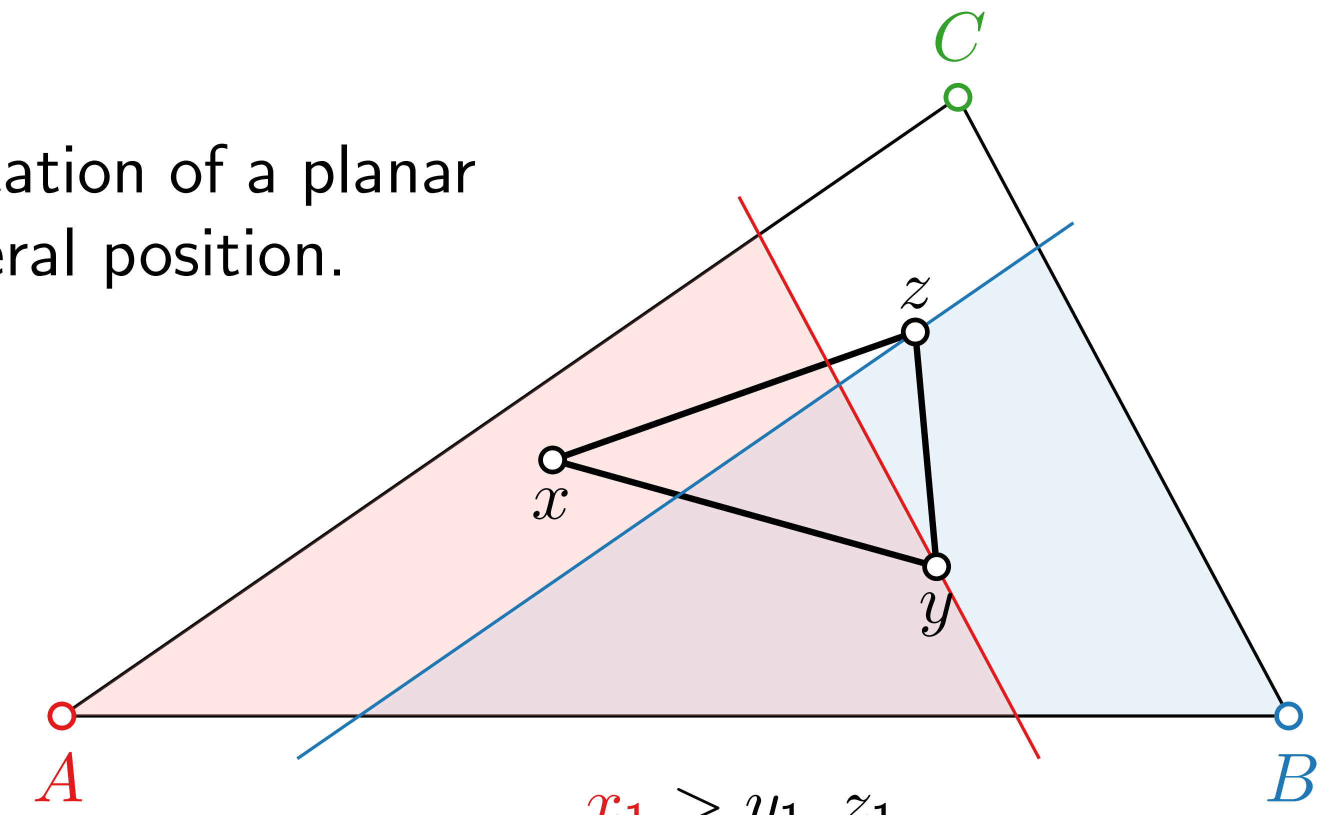

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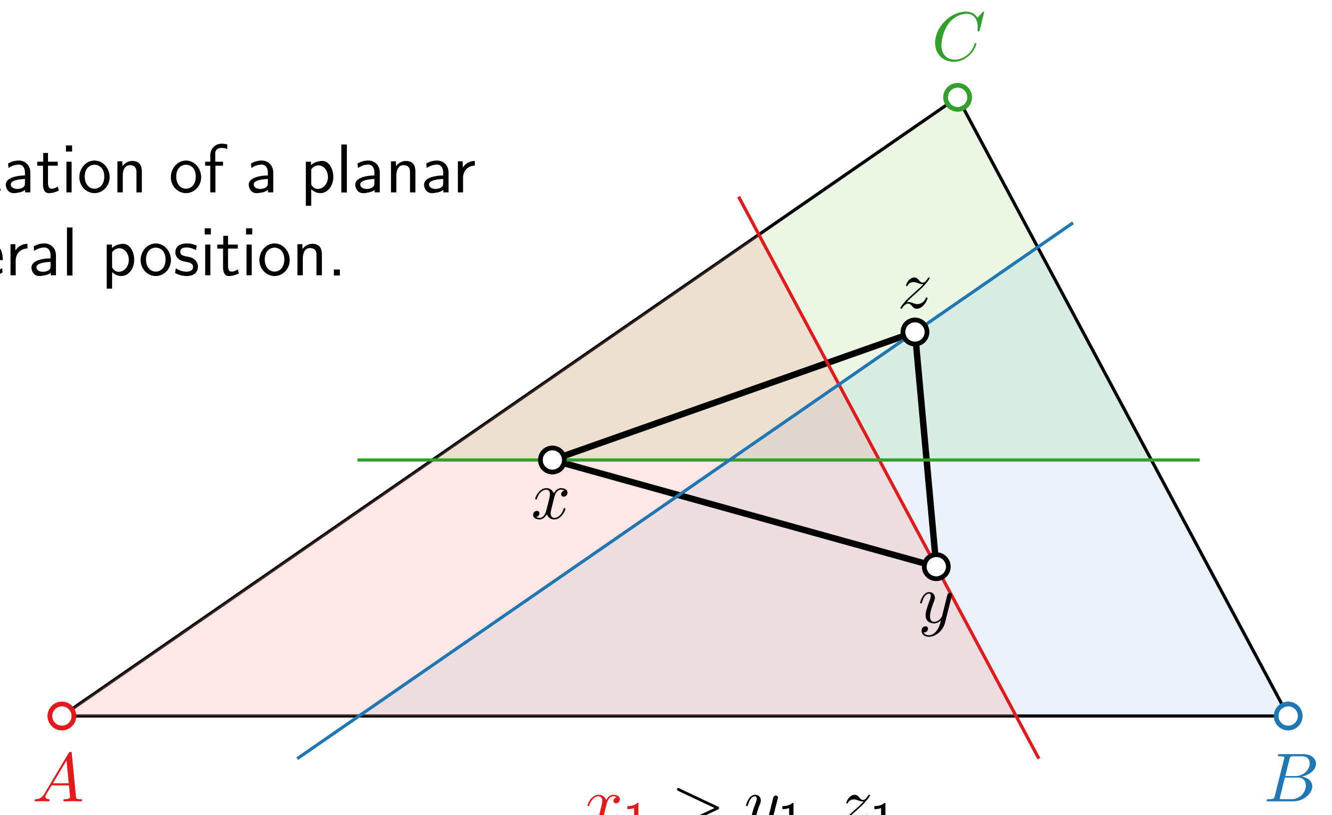

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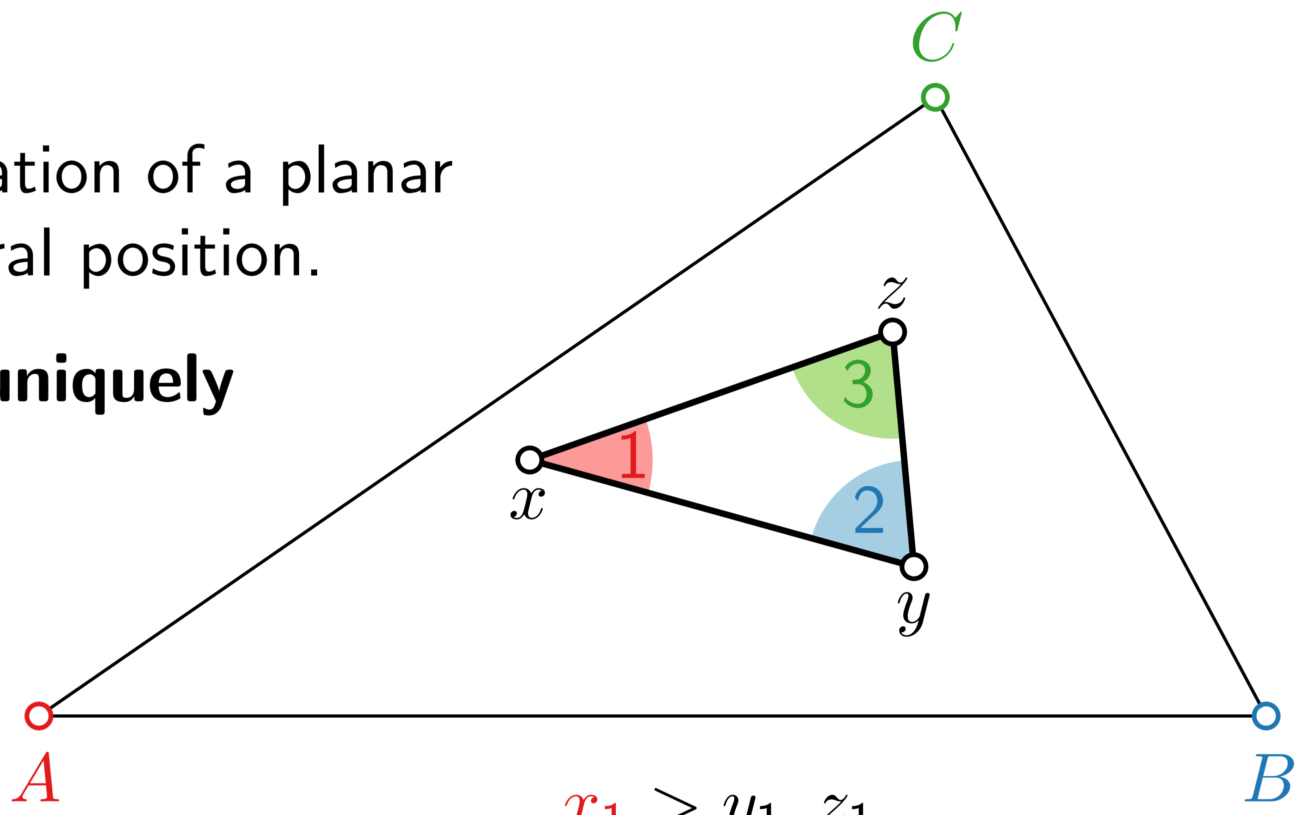

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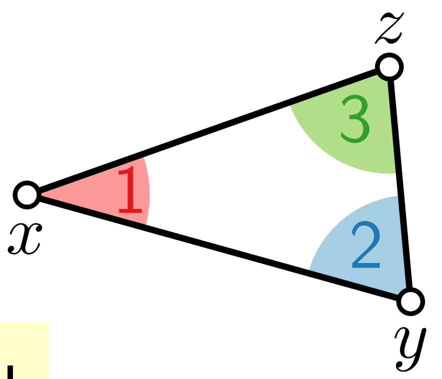

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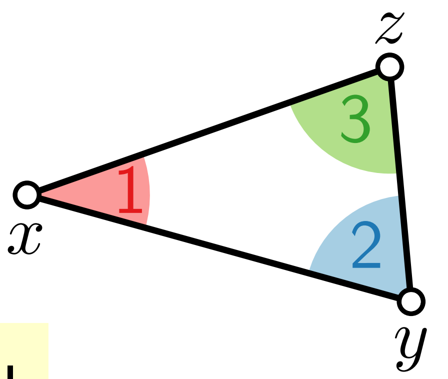

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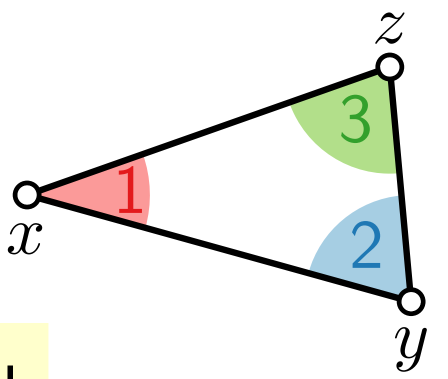

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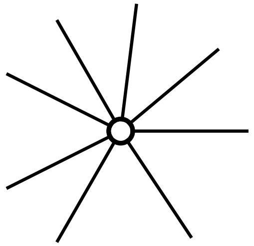

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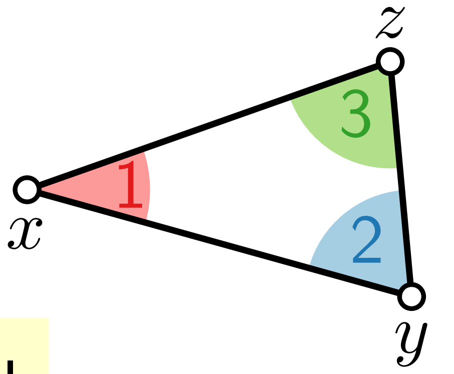

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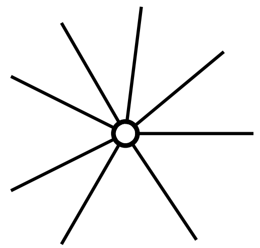

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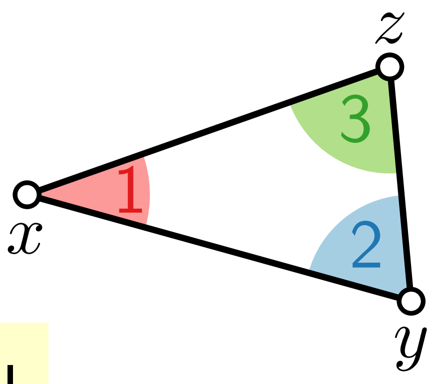

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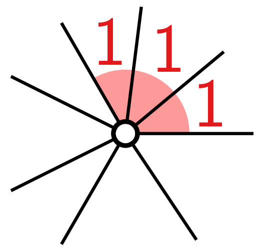

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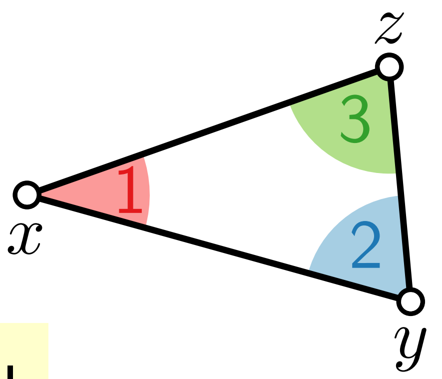

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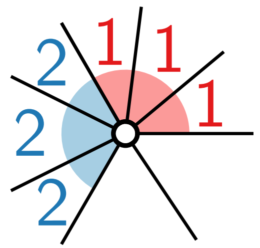

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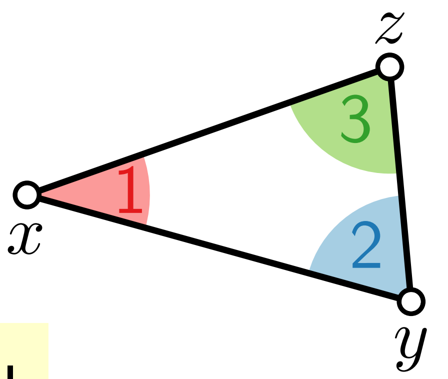

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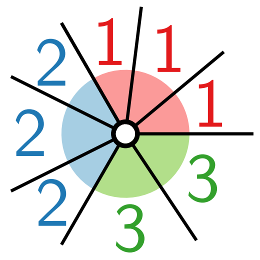

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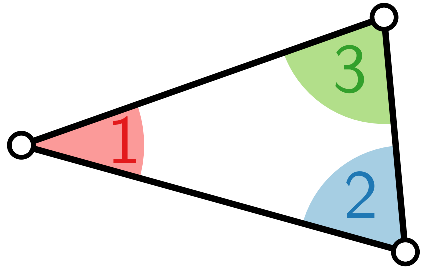
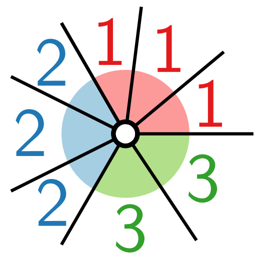
A Schnyder labeling induces an edge labeling.





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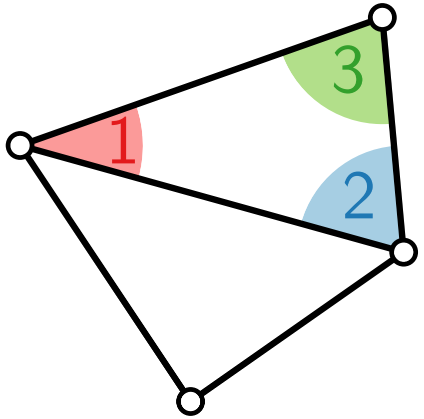

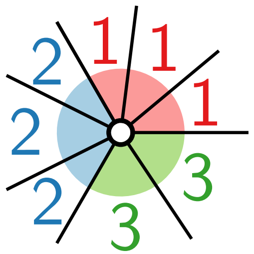

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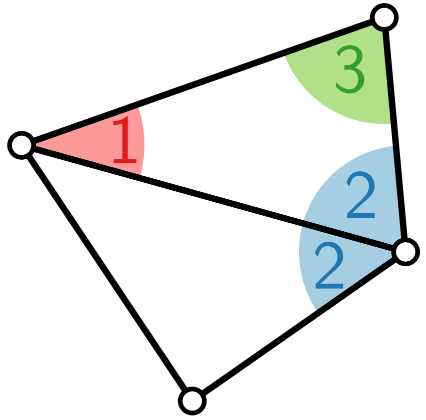

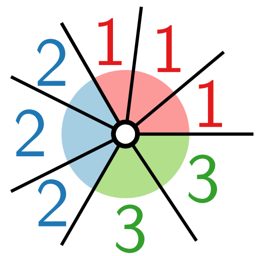

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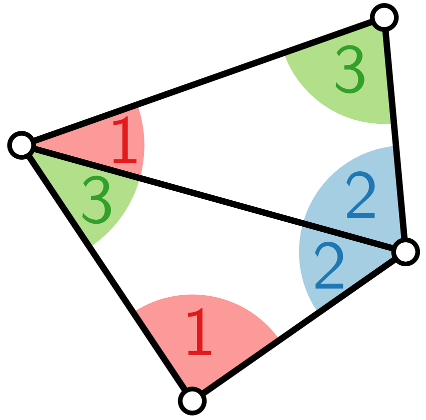

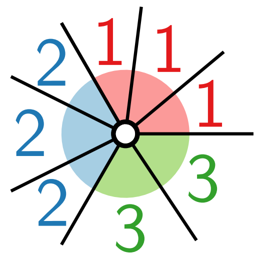

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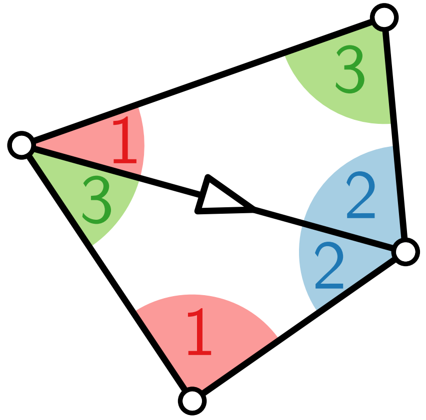

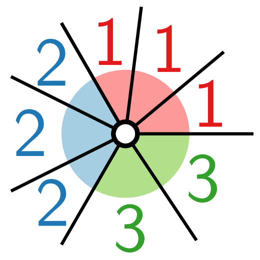

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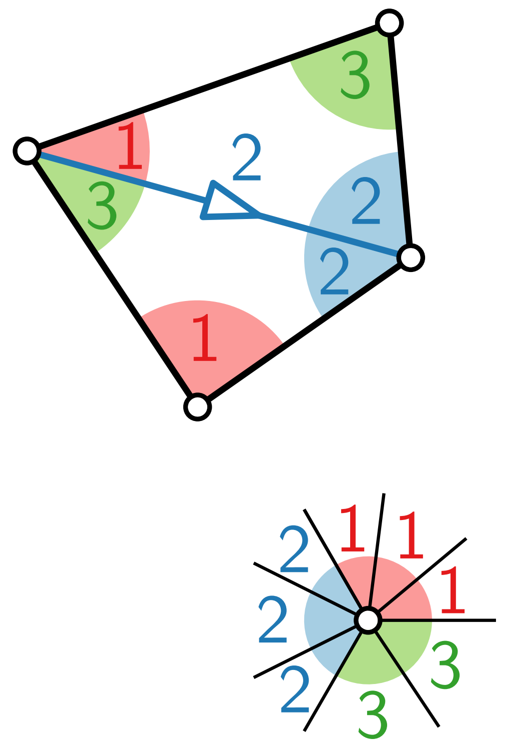

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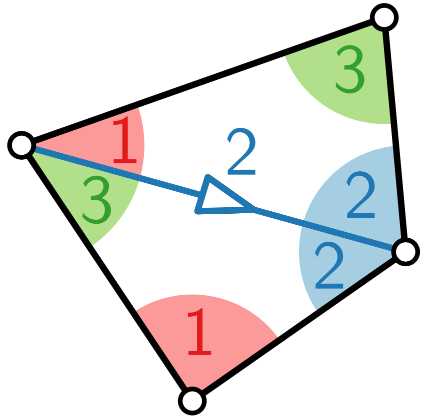

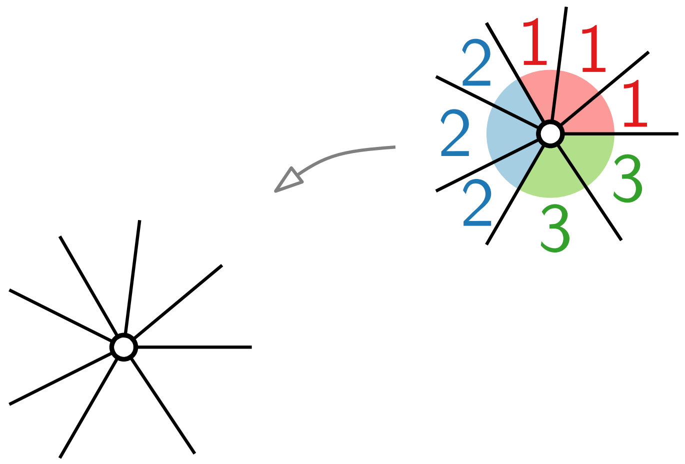

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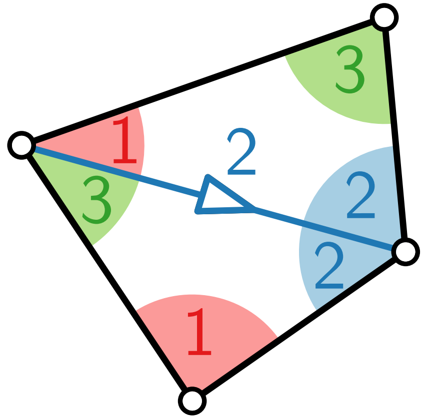

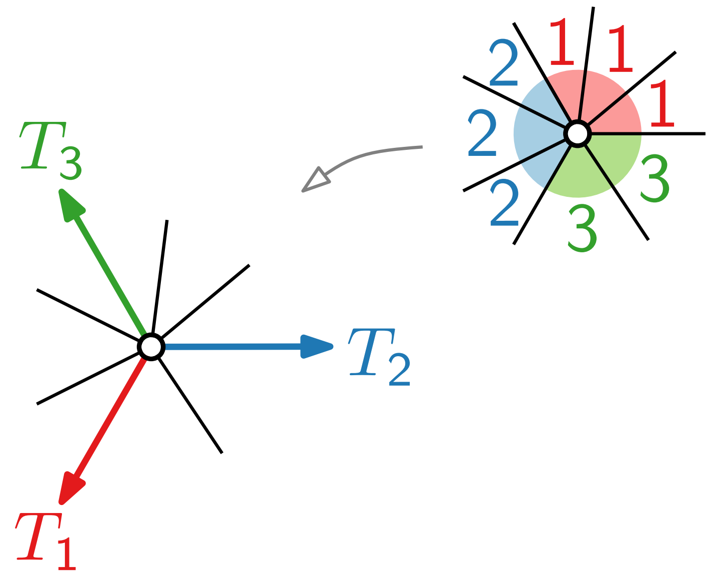

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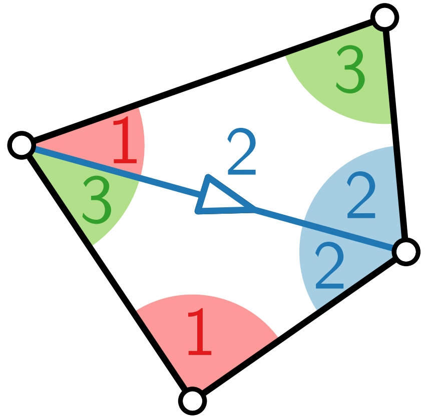

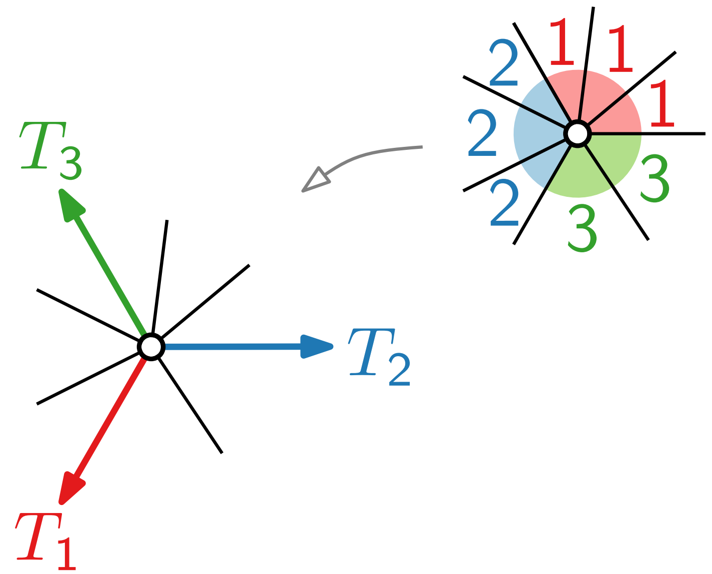

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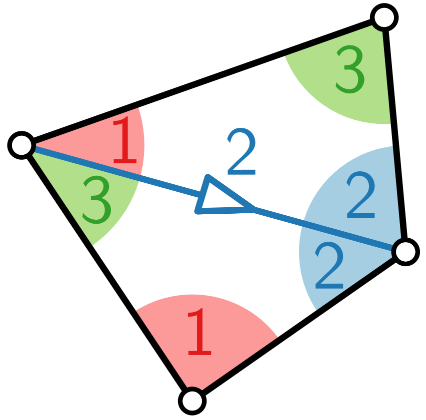

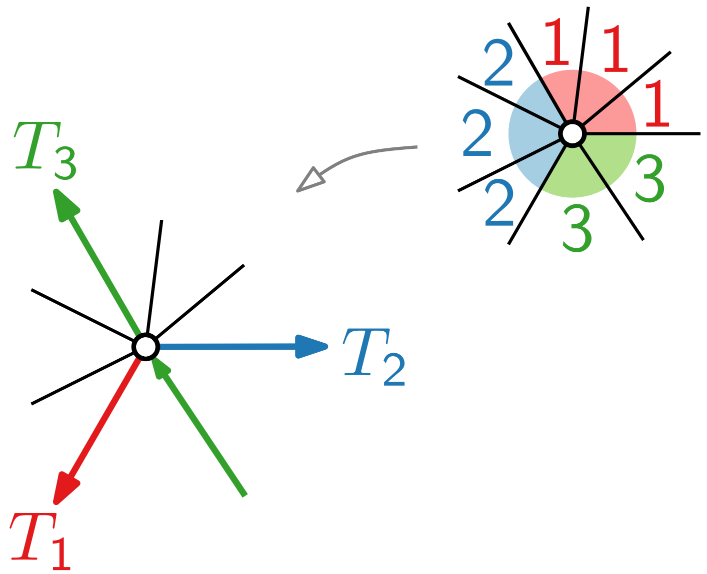

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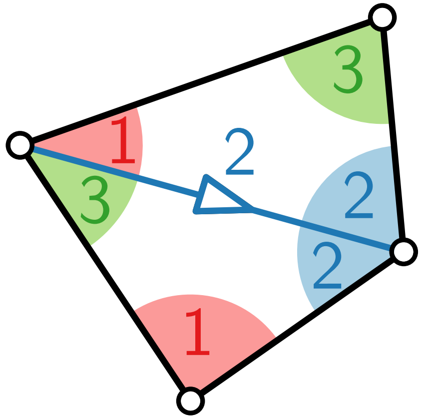

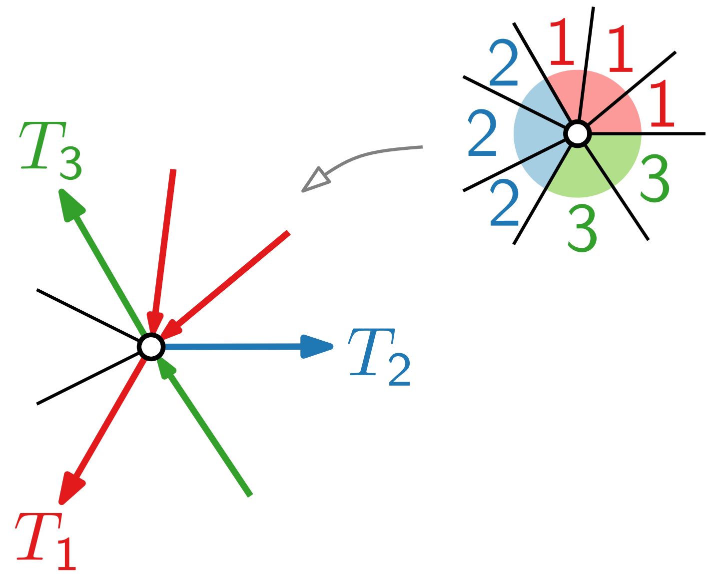

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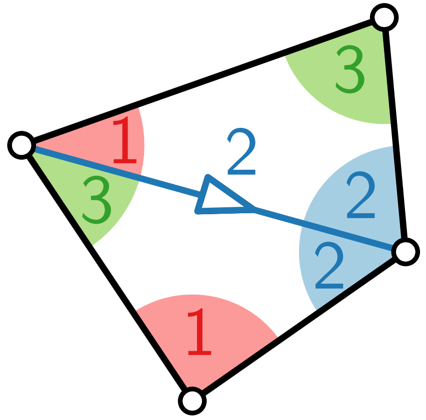
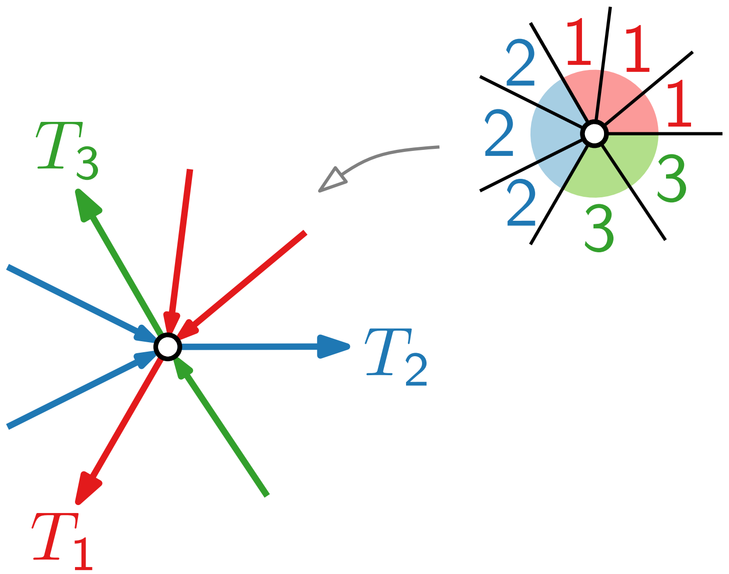
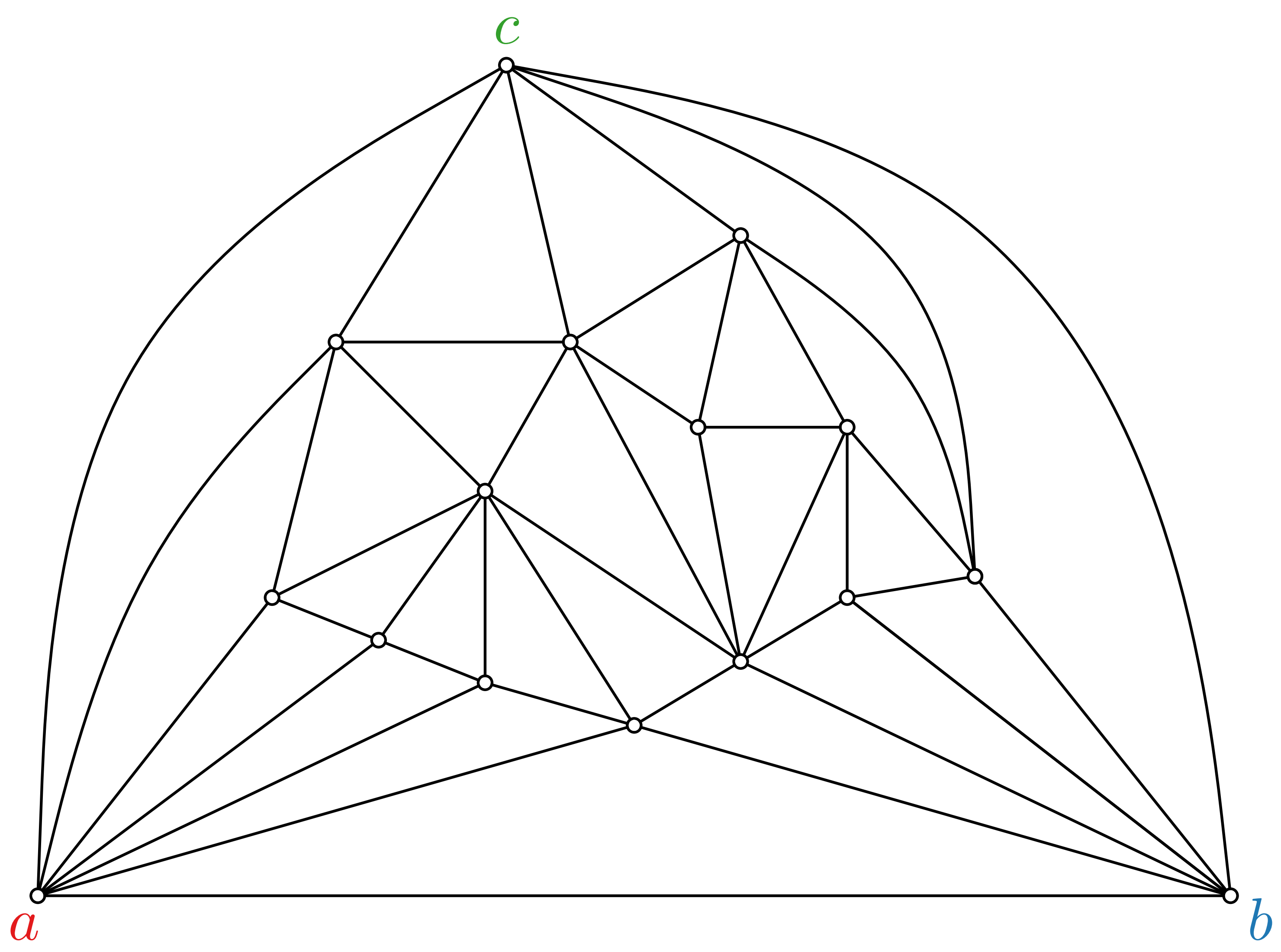
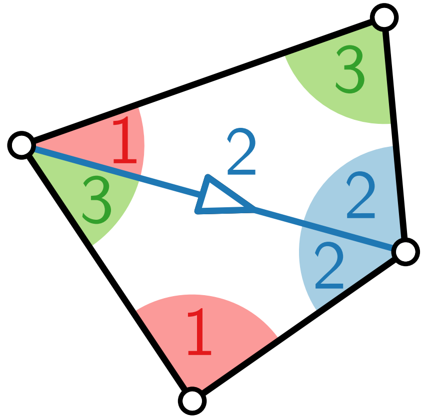
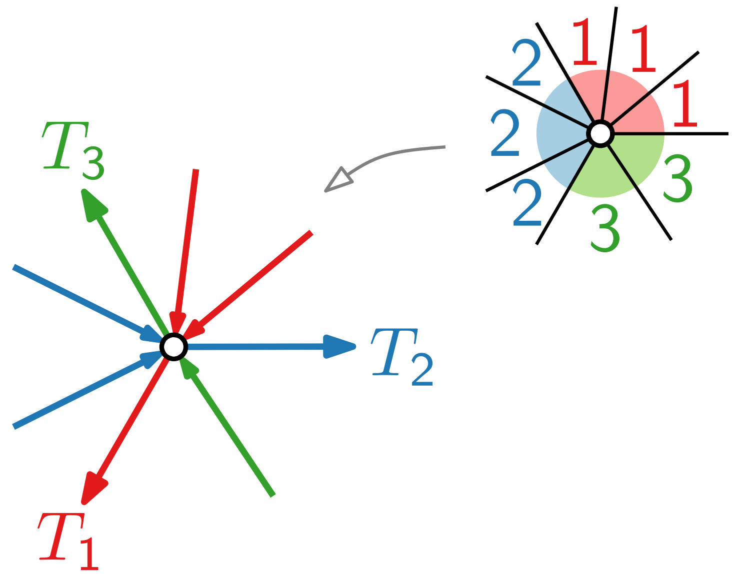




Schnyder Wood – Example and Properties







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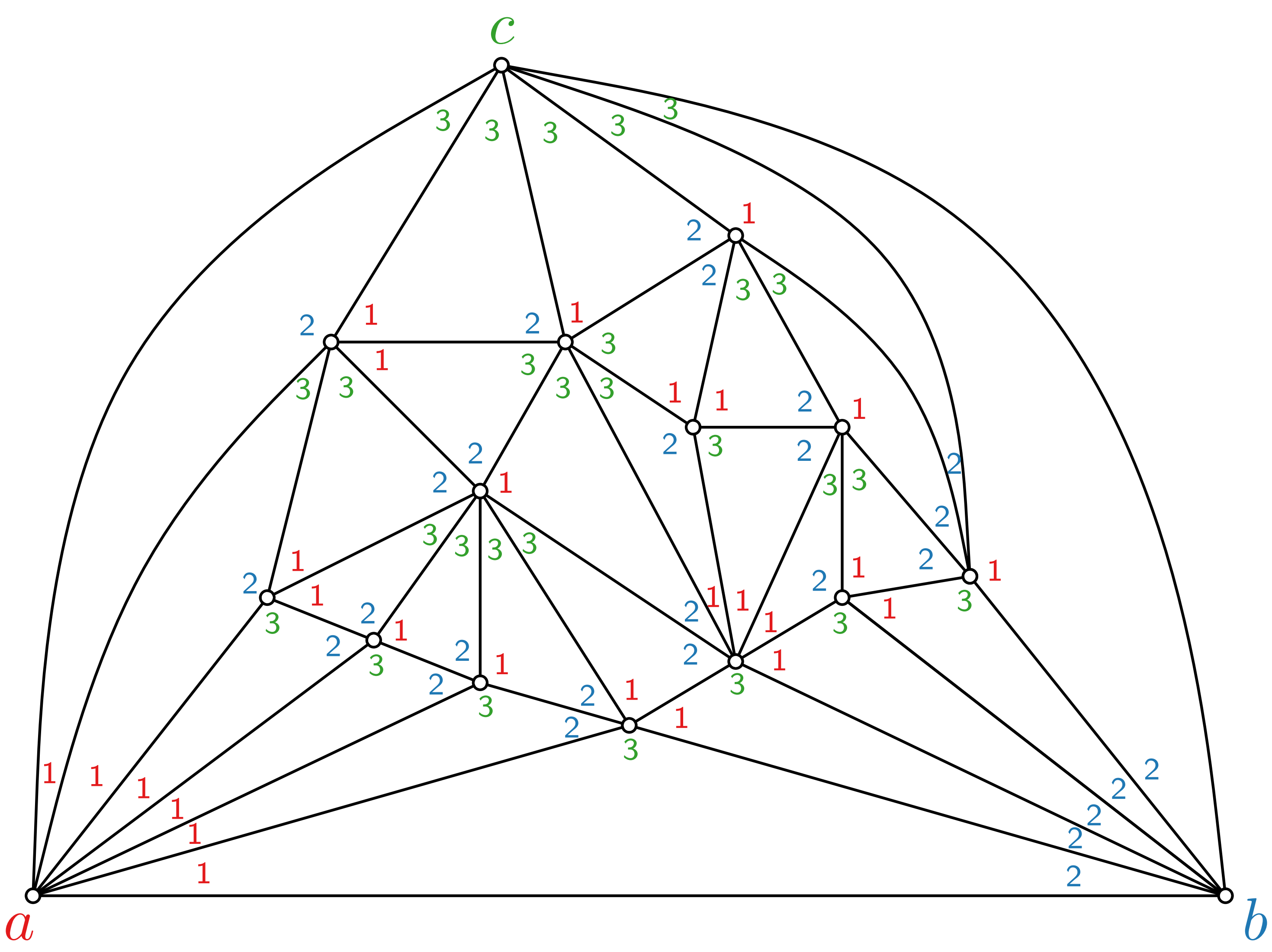

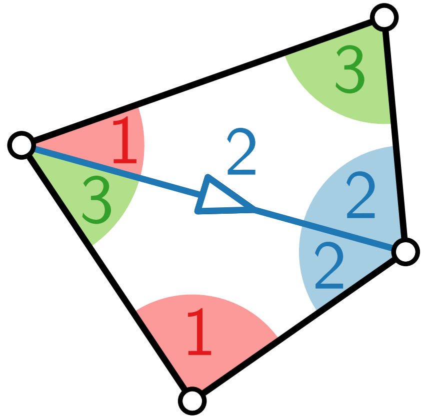

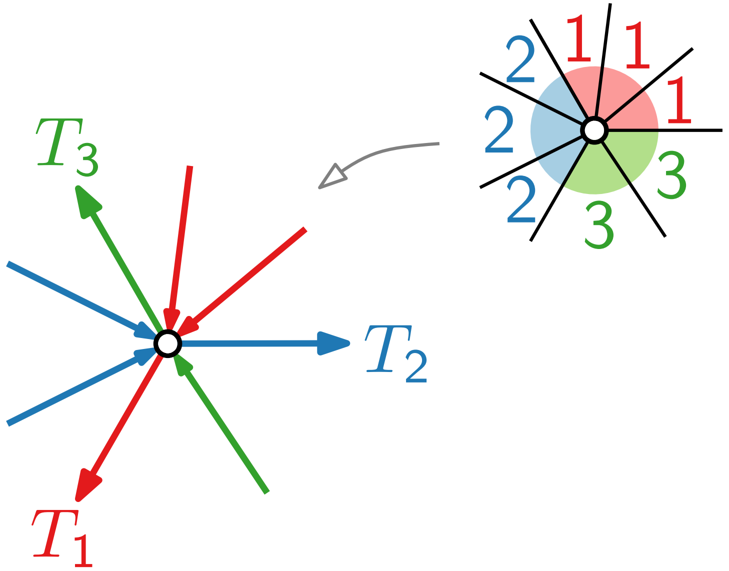

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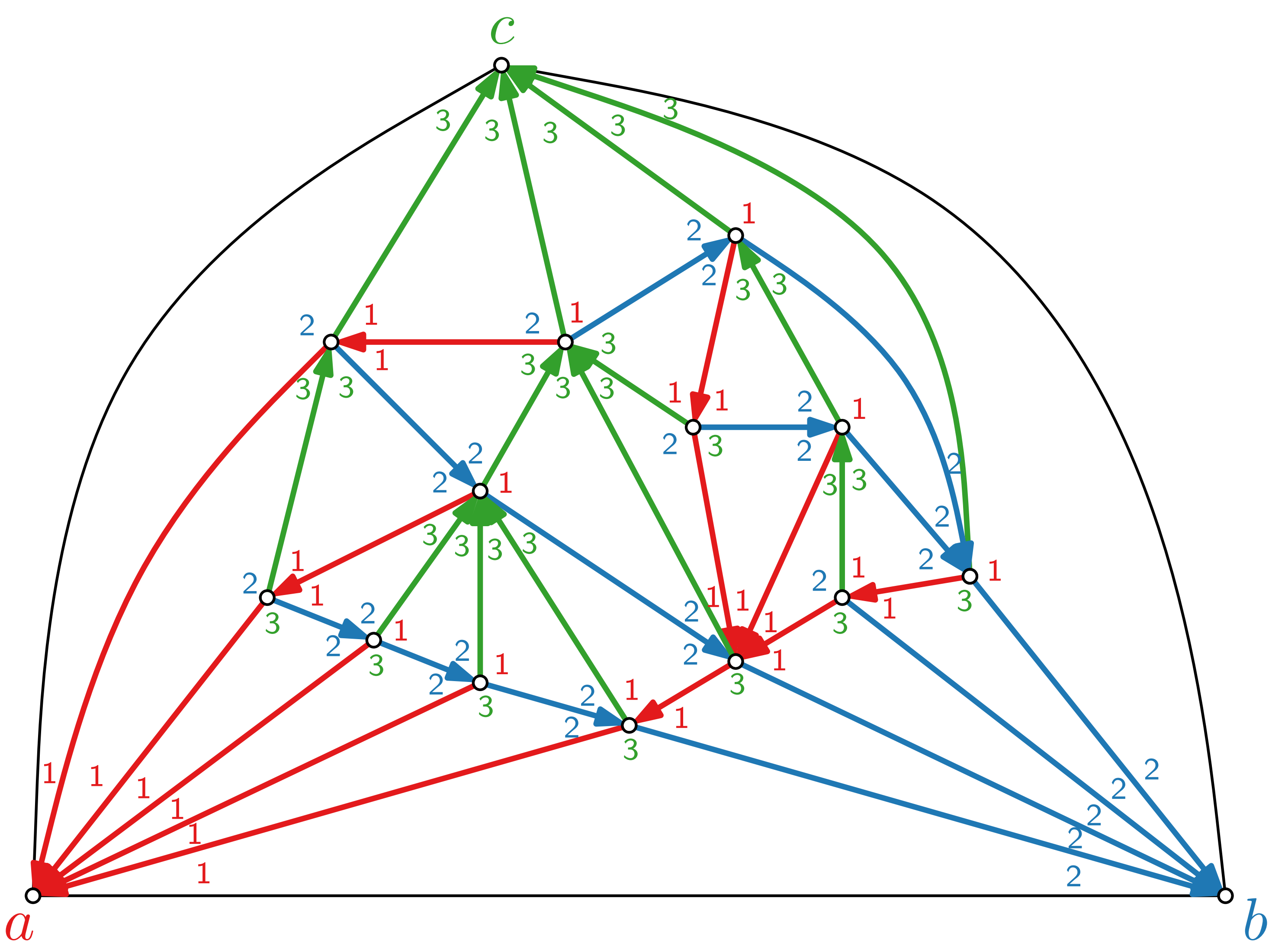

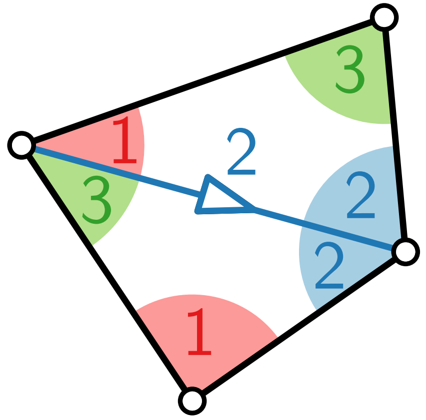

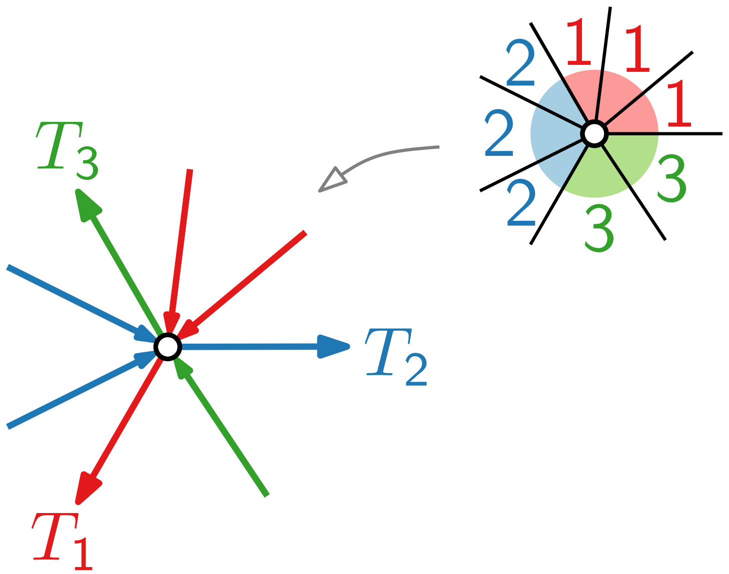

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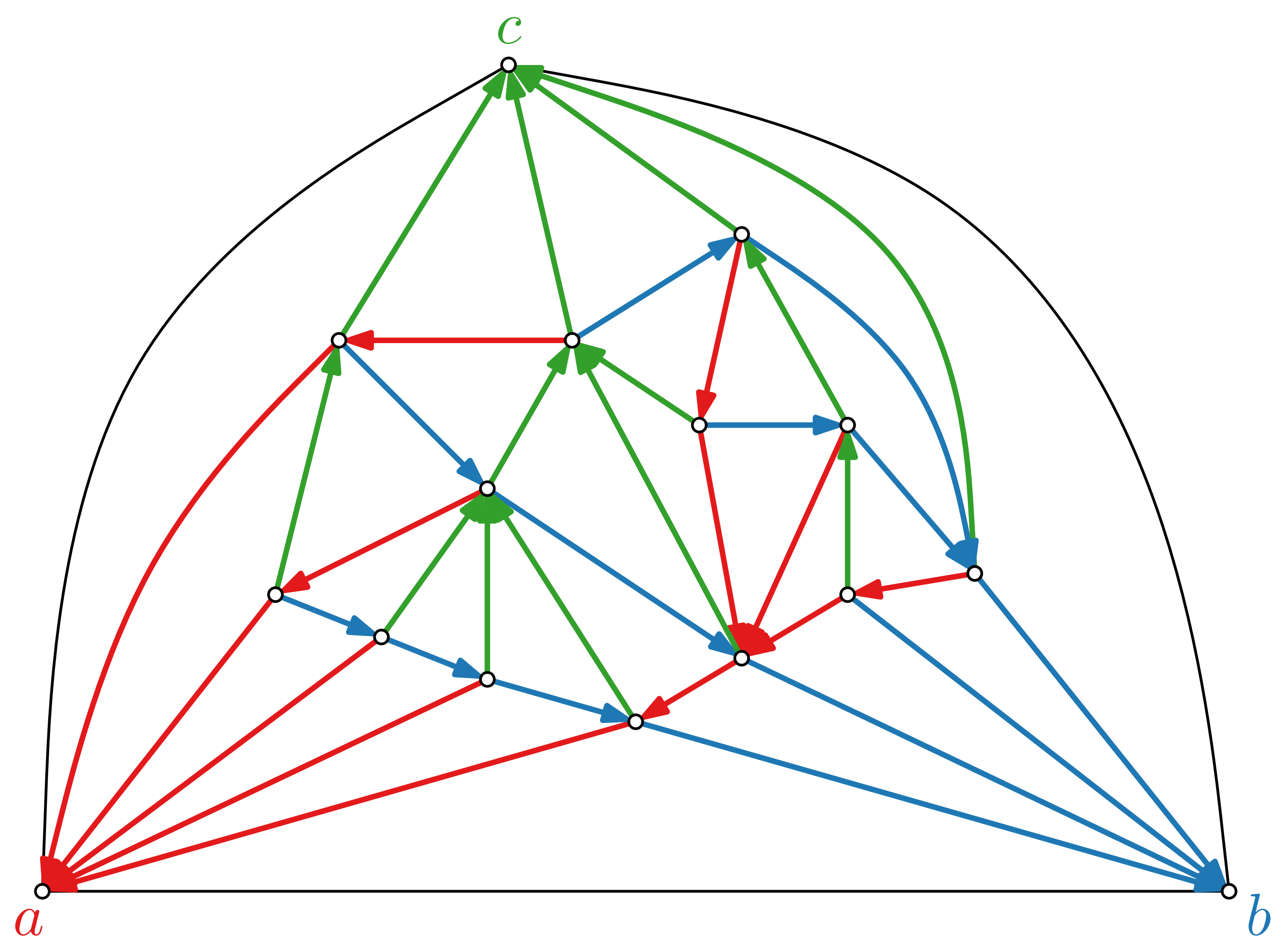

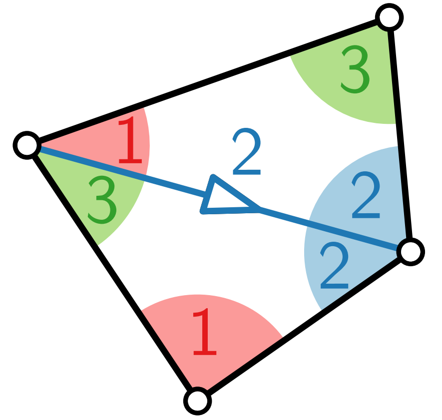

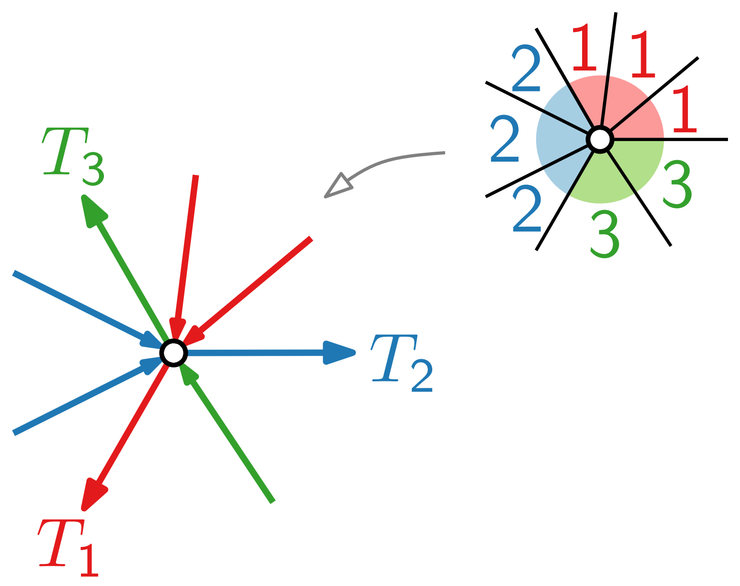

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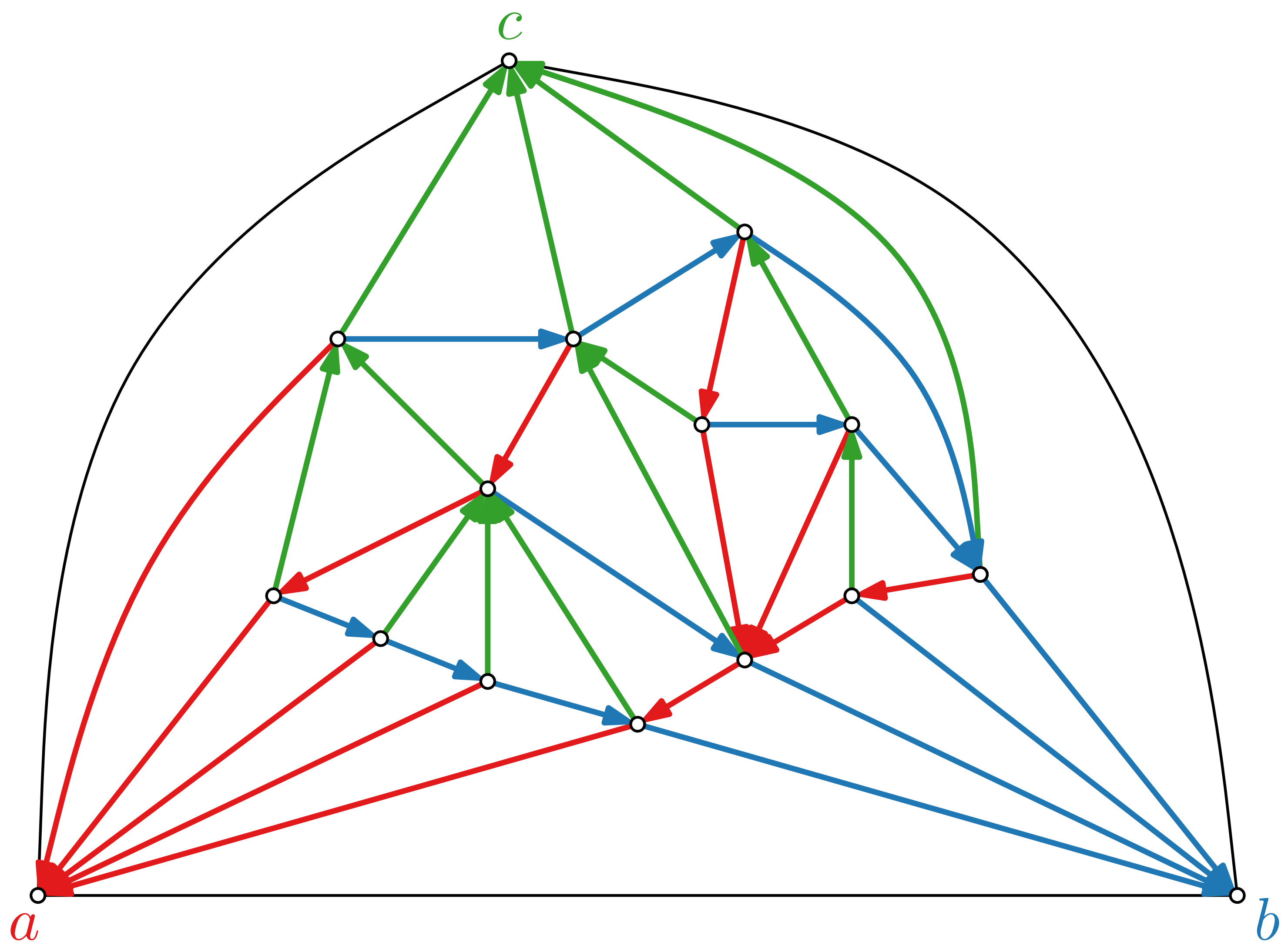
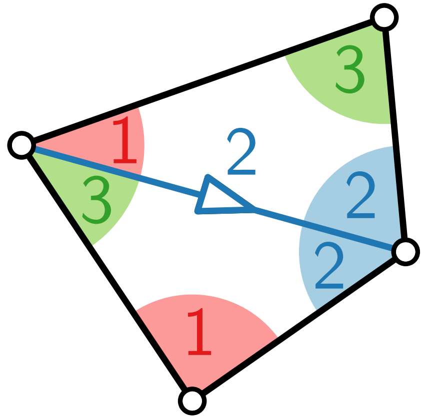
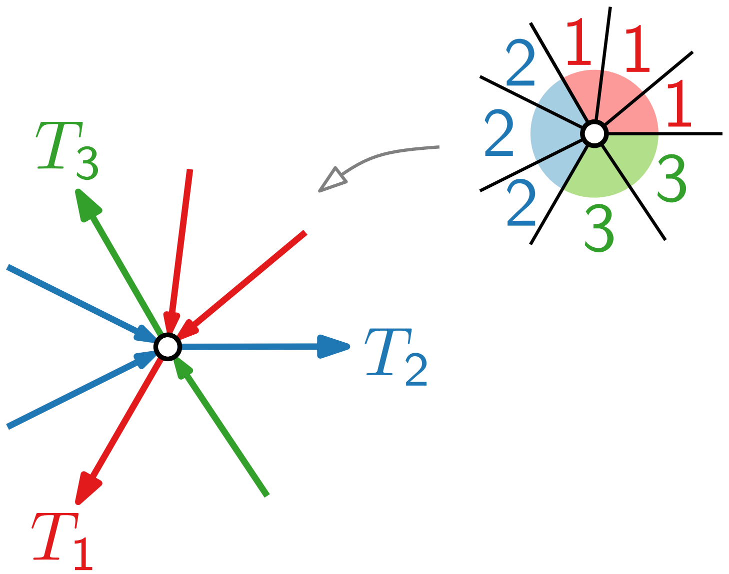


(a Schnyder labeling is not unique)





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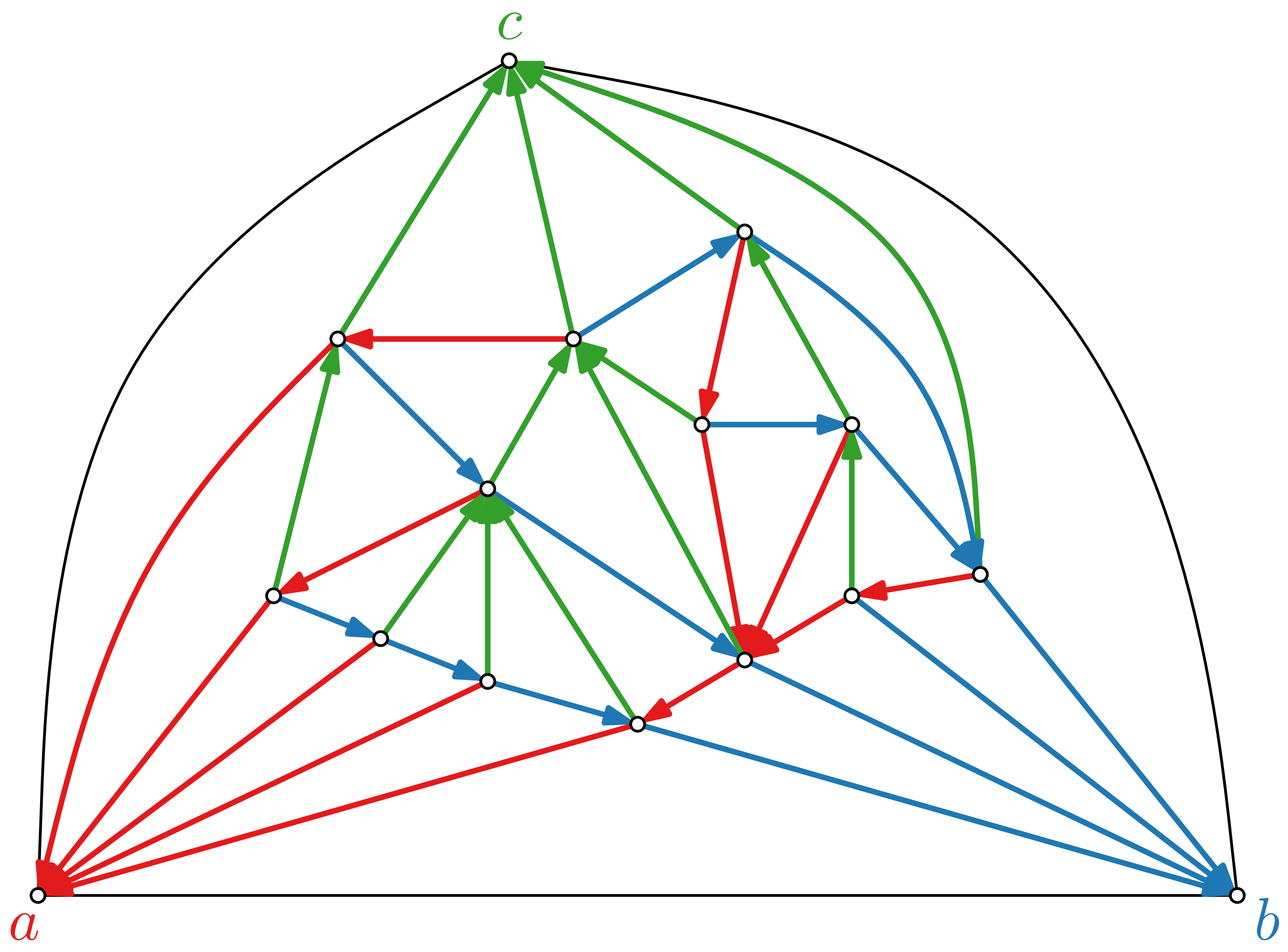

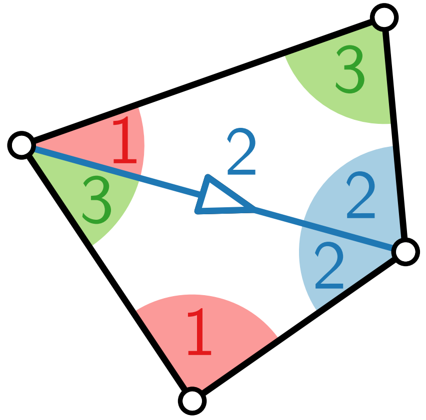

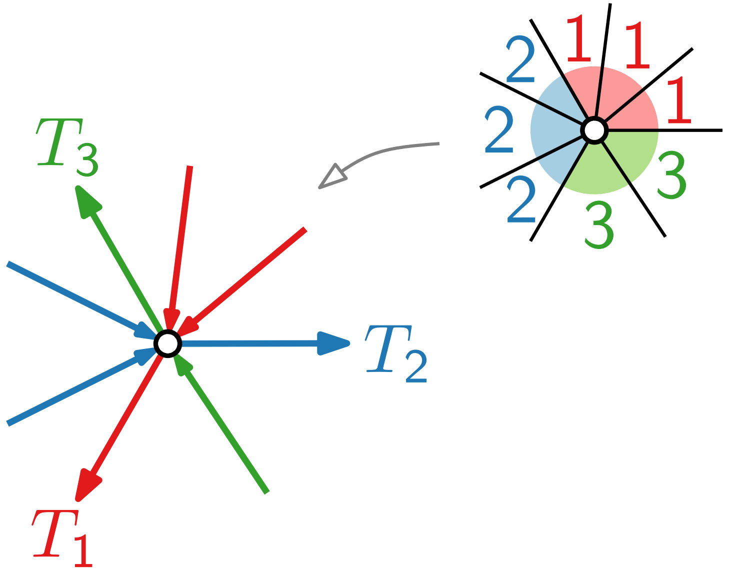

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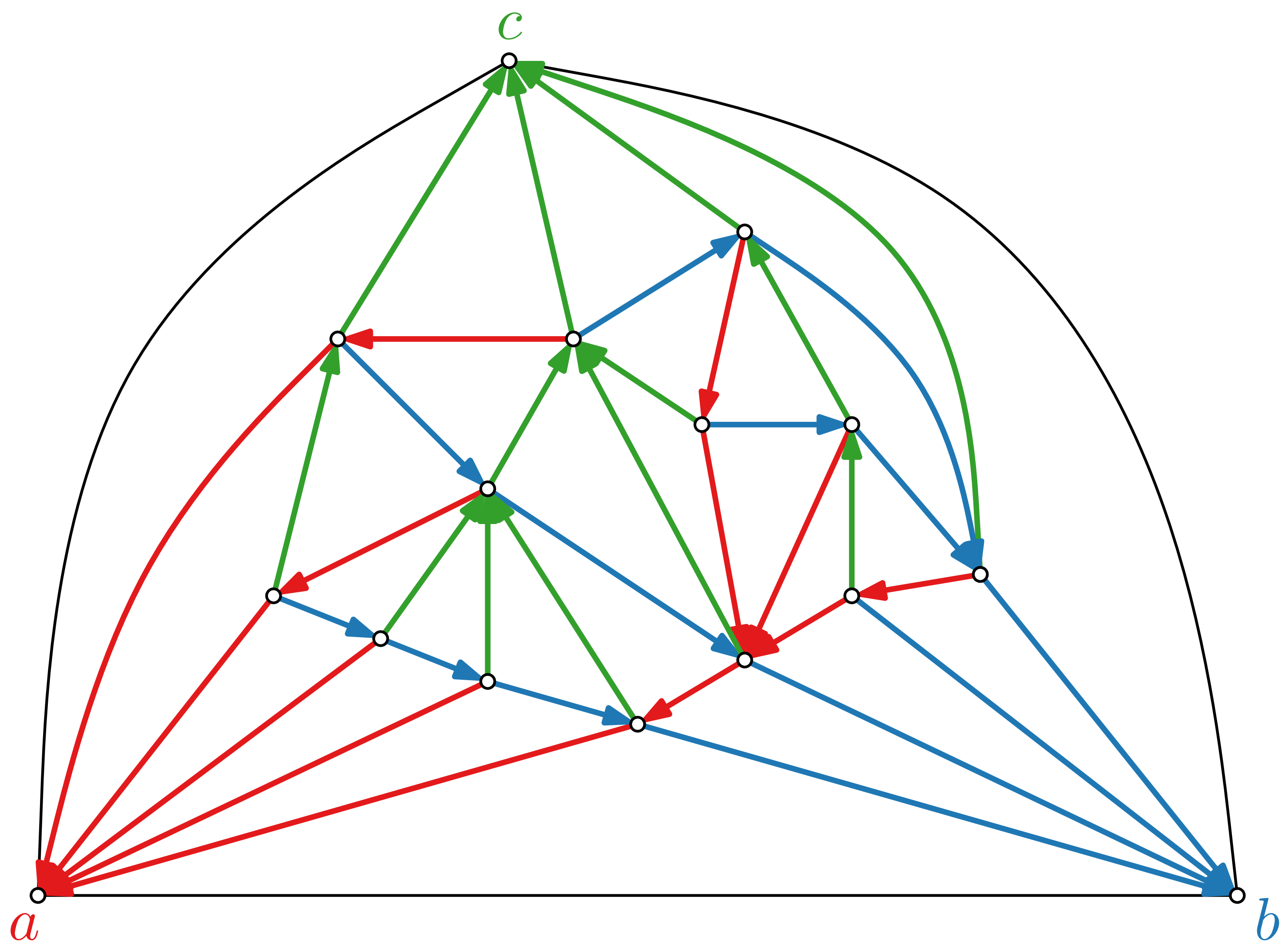

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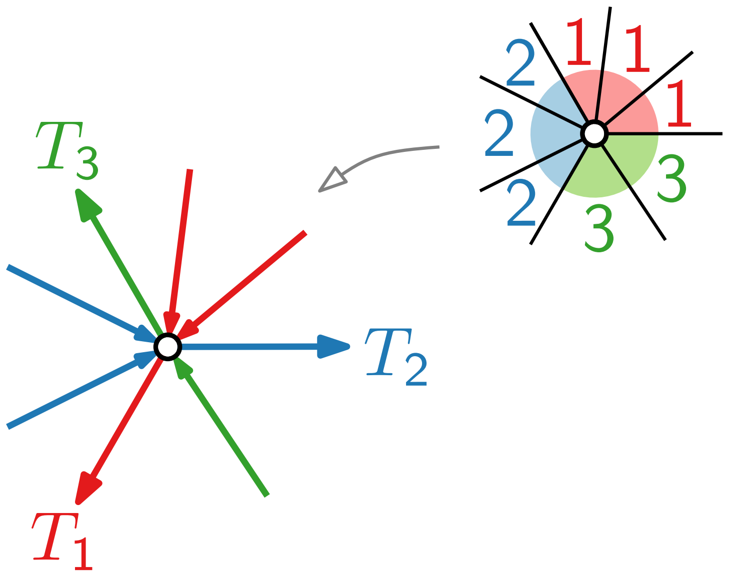

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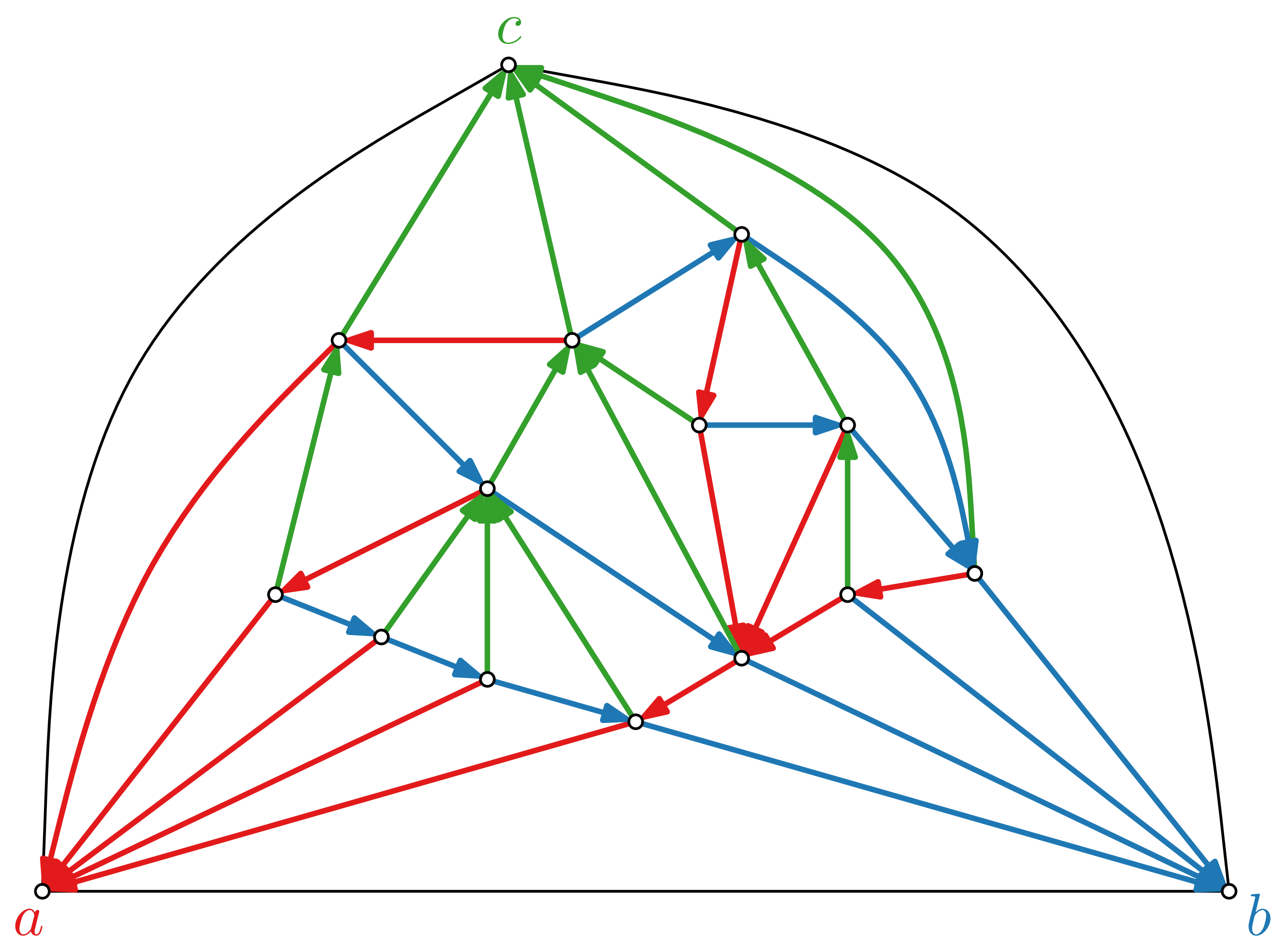
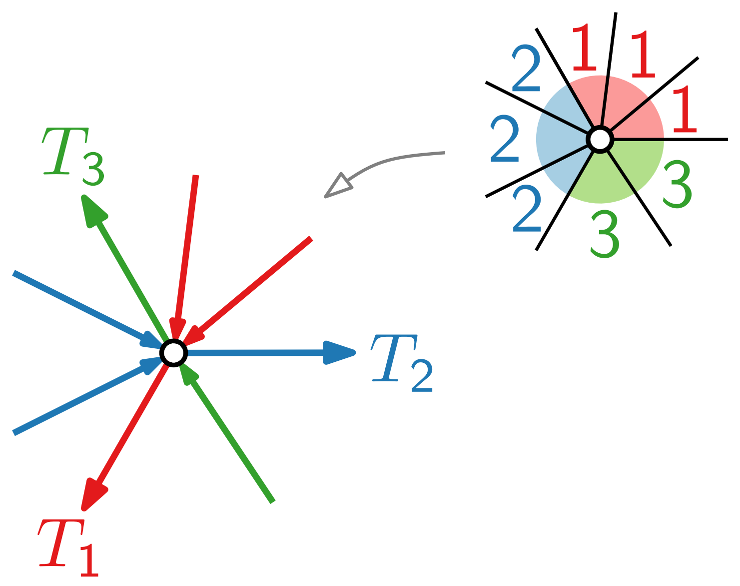


- All inner edges incident to a, b, and c are incoming in the same set (color).
- T_1, T_2, and T_3 are trees. Each spans all inner vertices and one outer vertex (its root).



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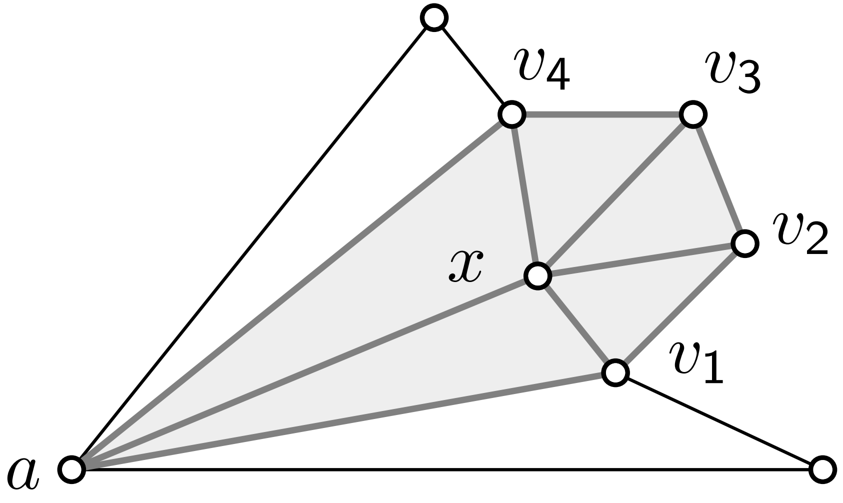

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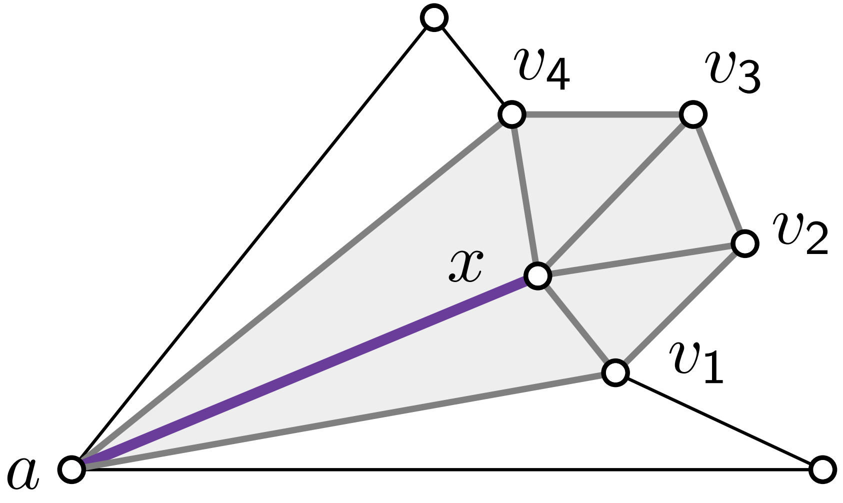

$\{a, x\}$

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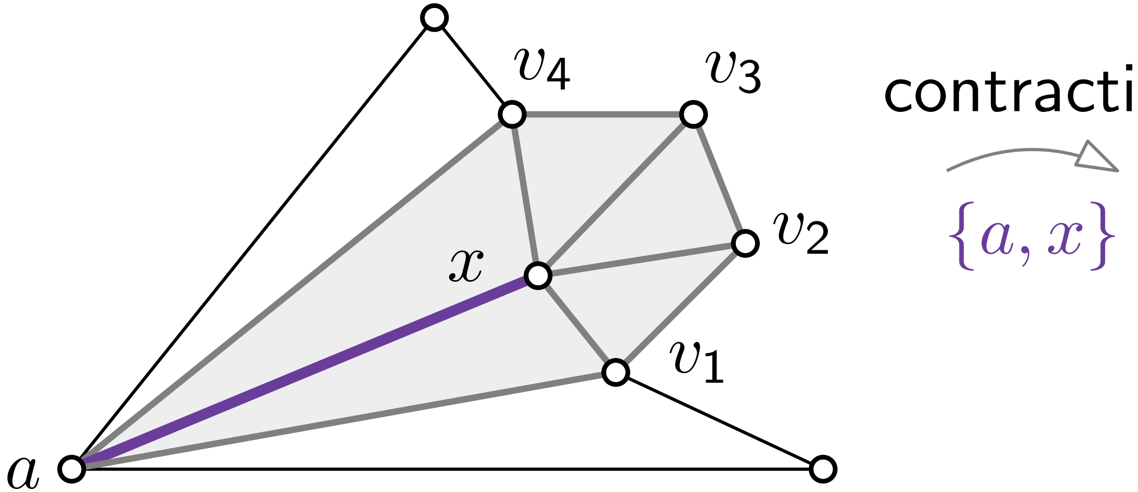

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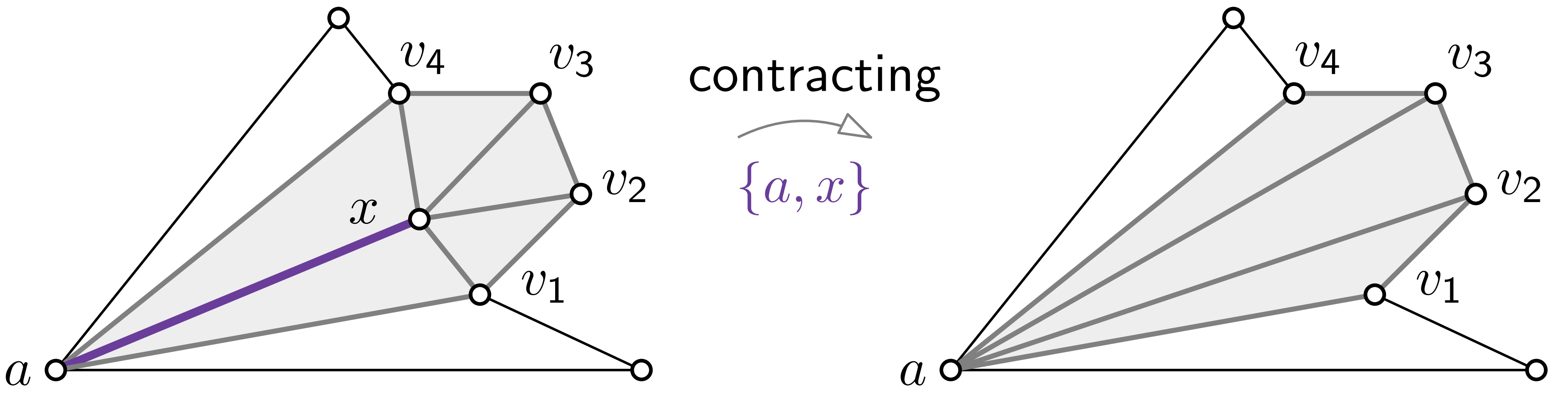

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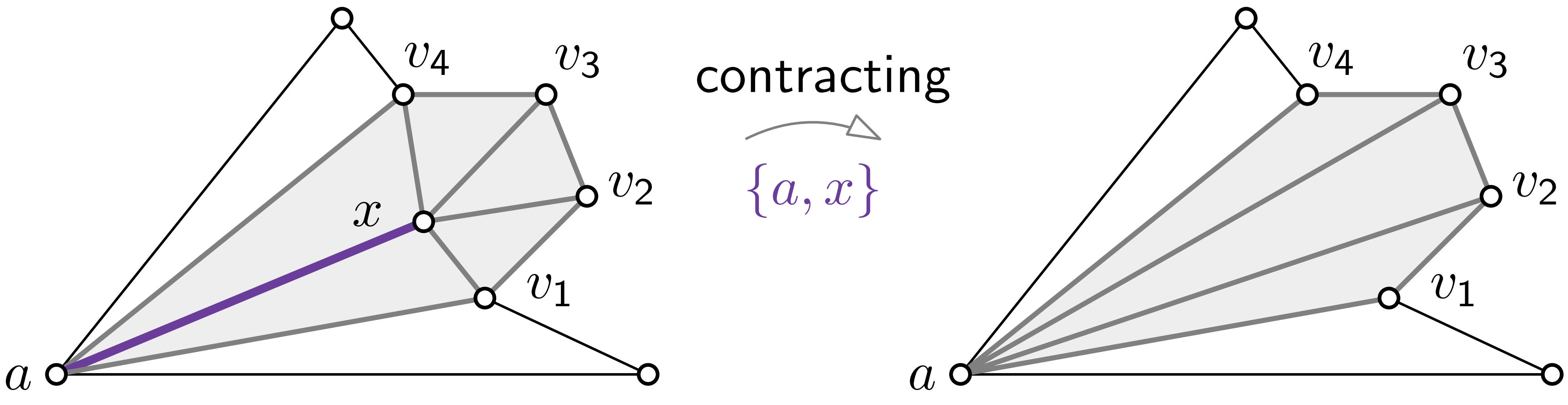

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Let G be a plane triangulation with vertices a, b, c on the outer face.
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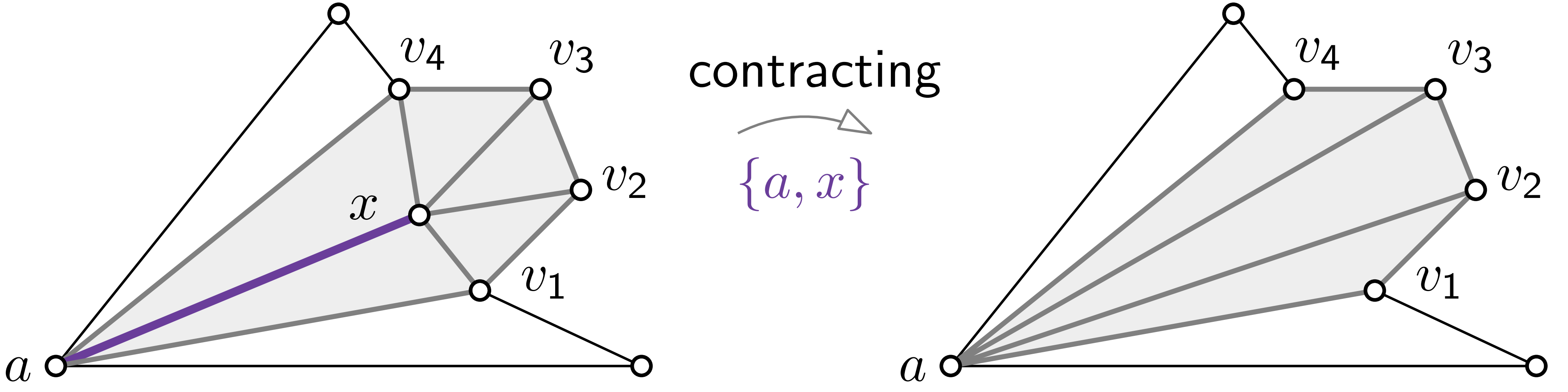

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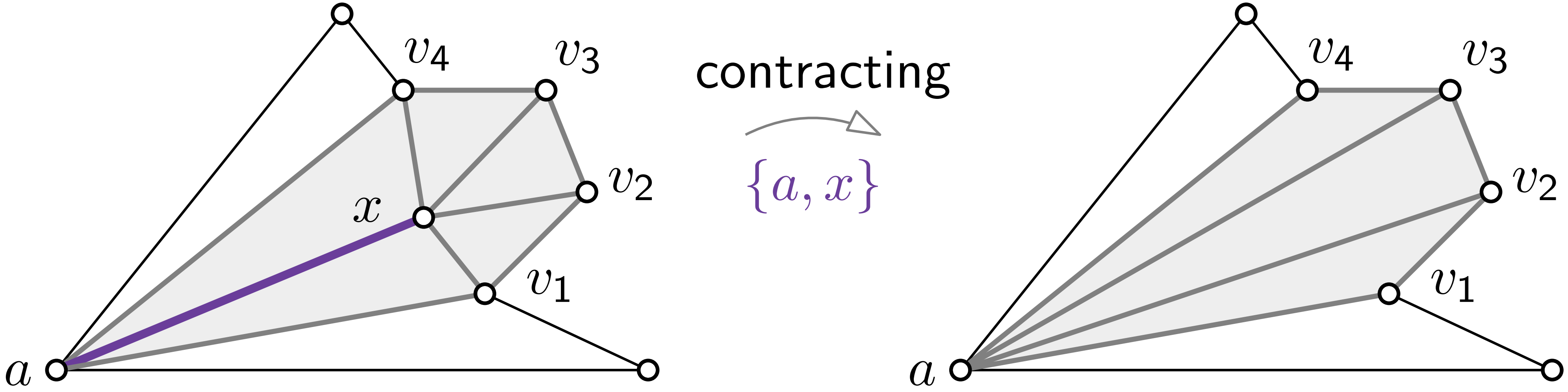

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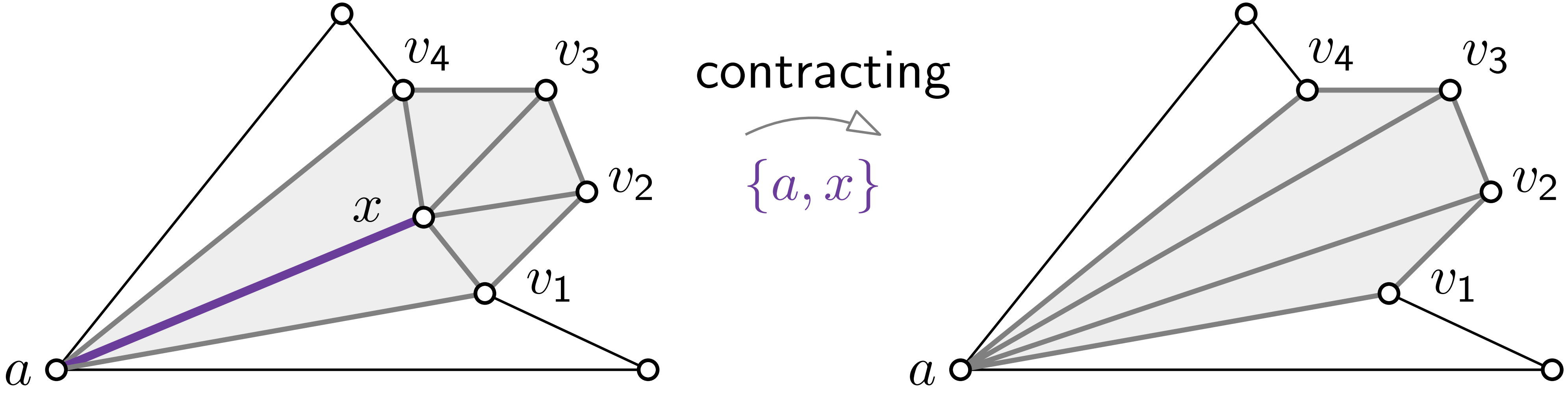

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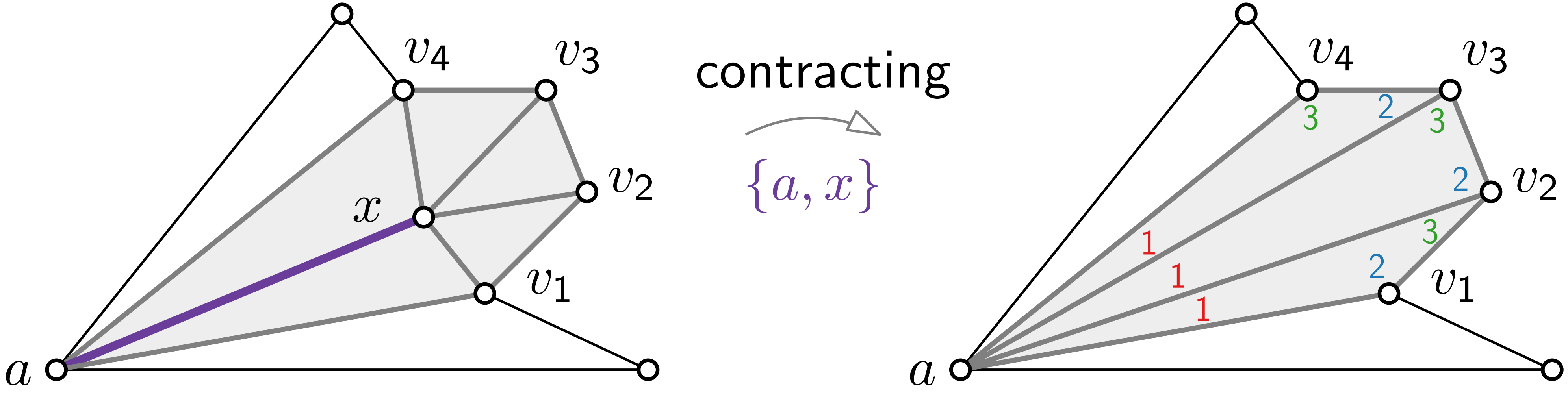

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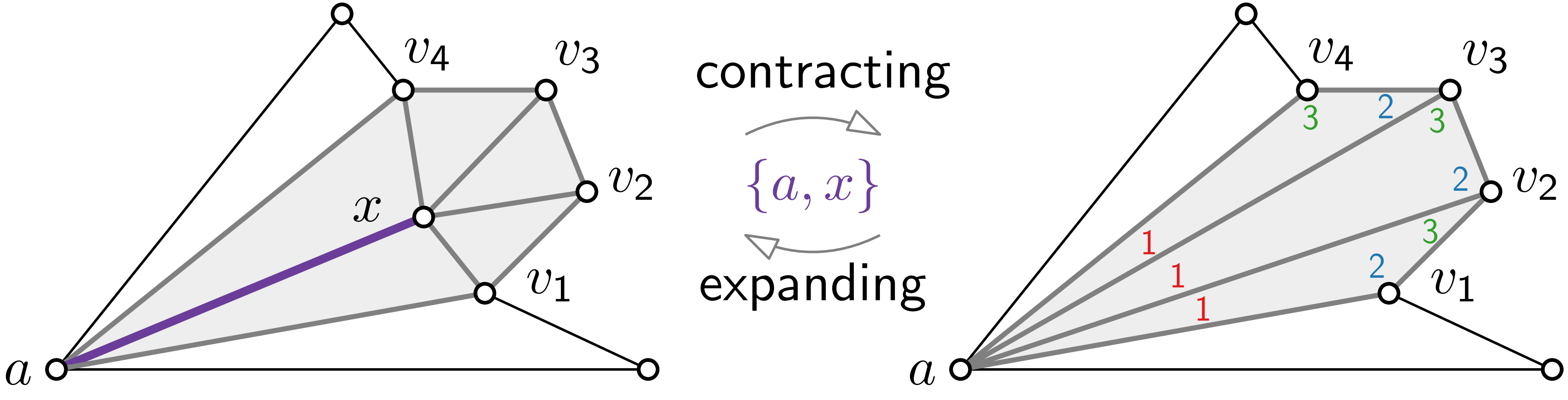

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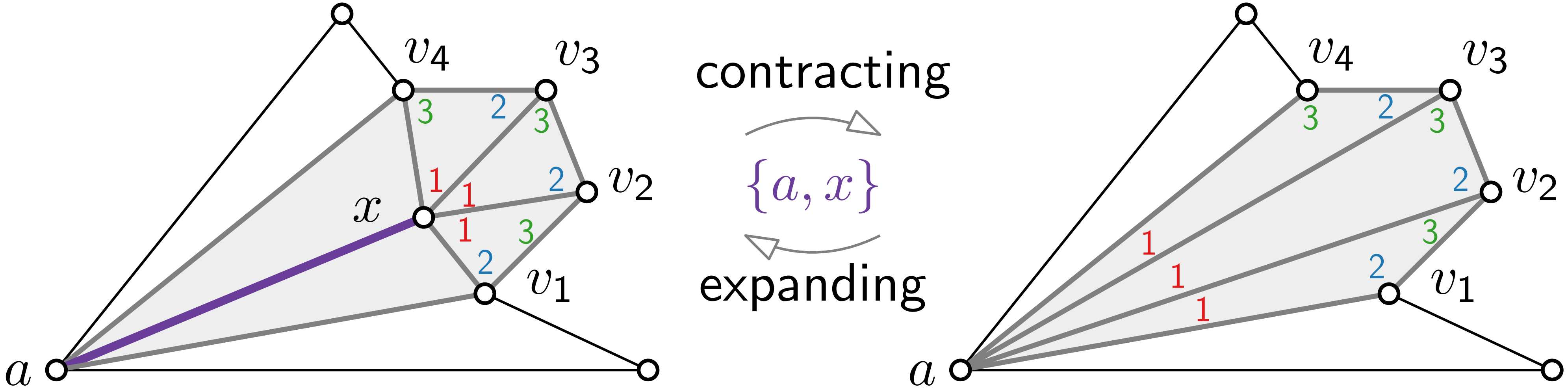

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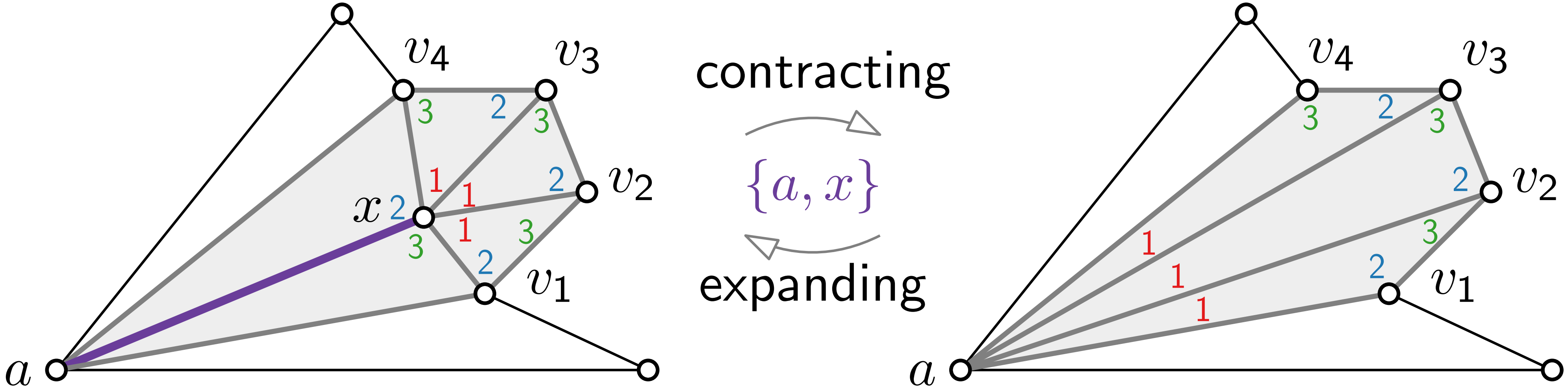

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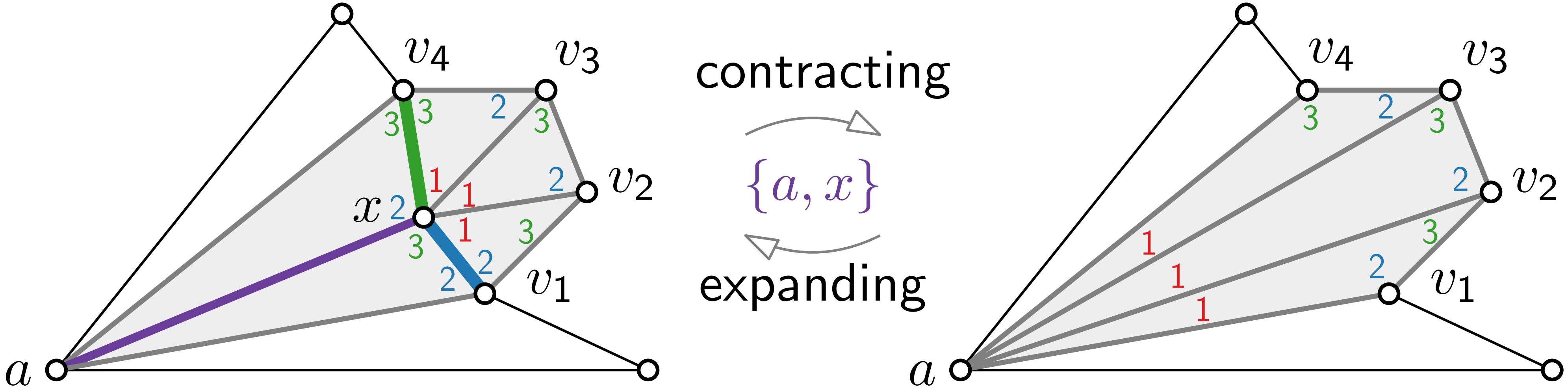

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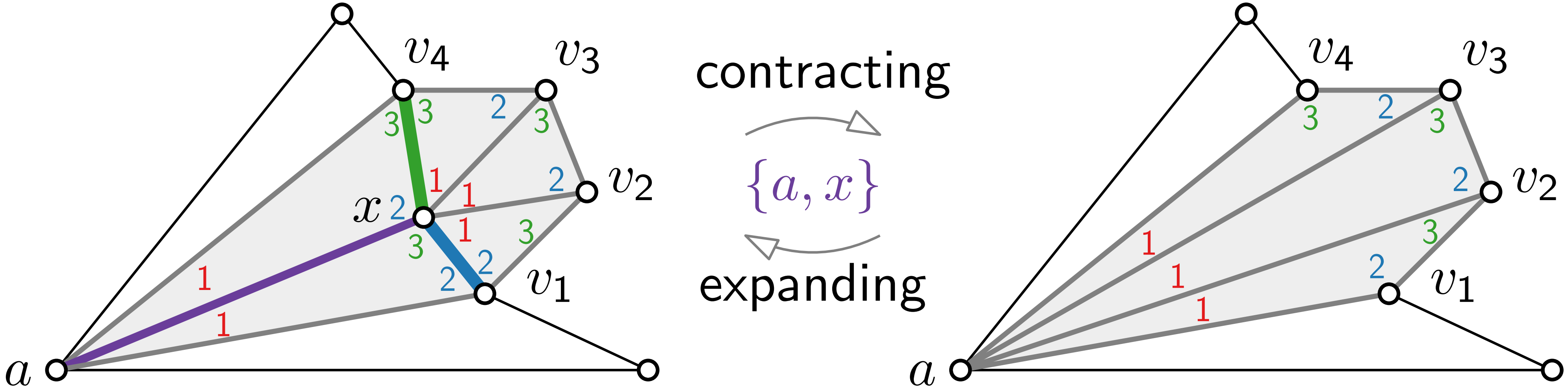

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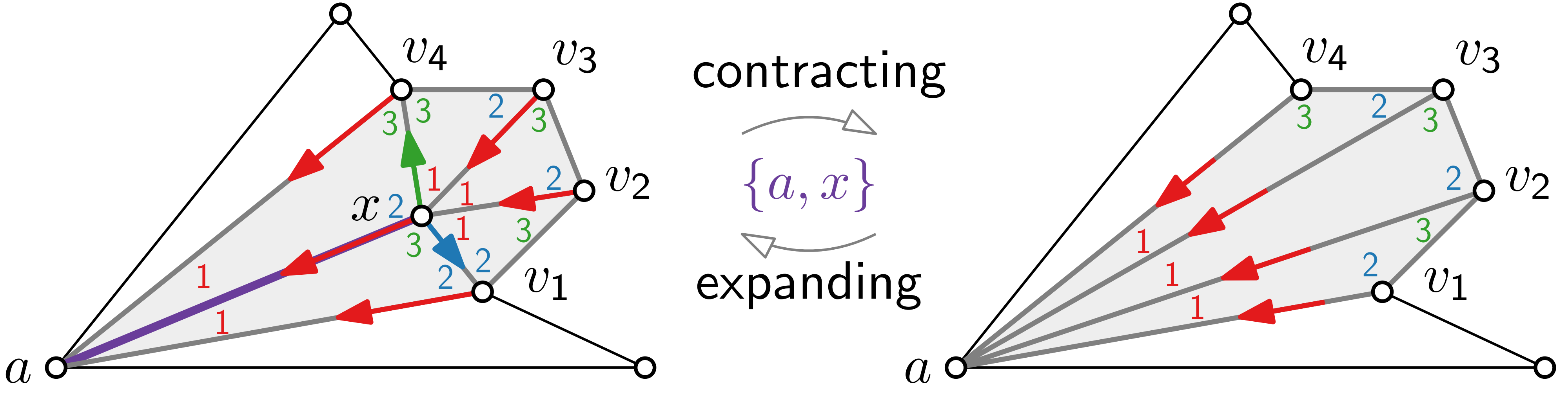

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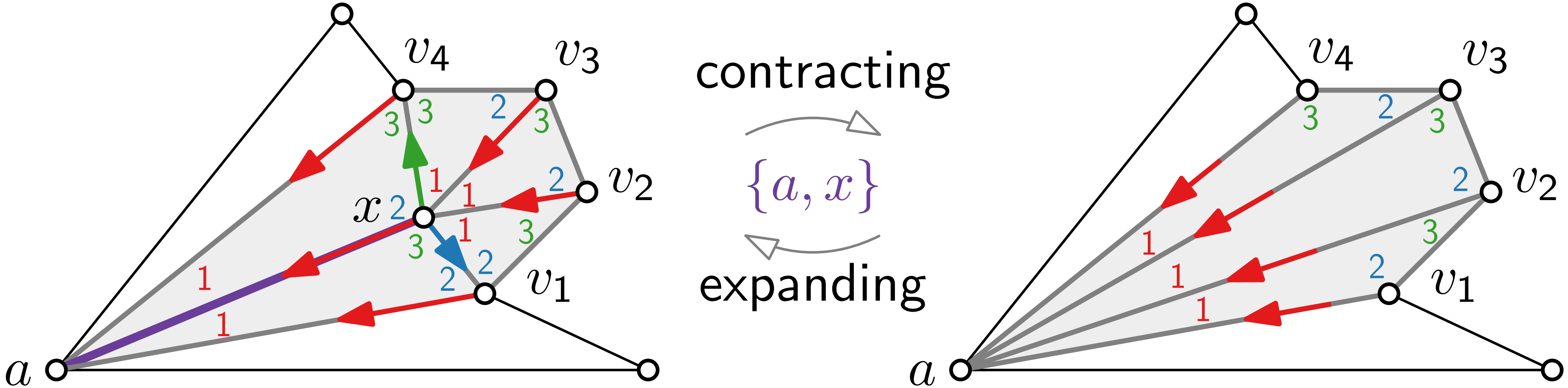

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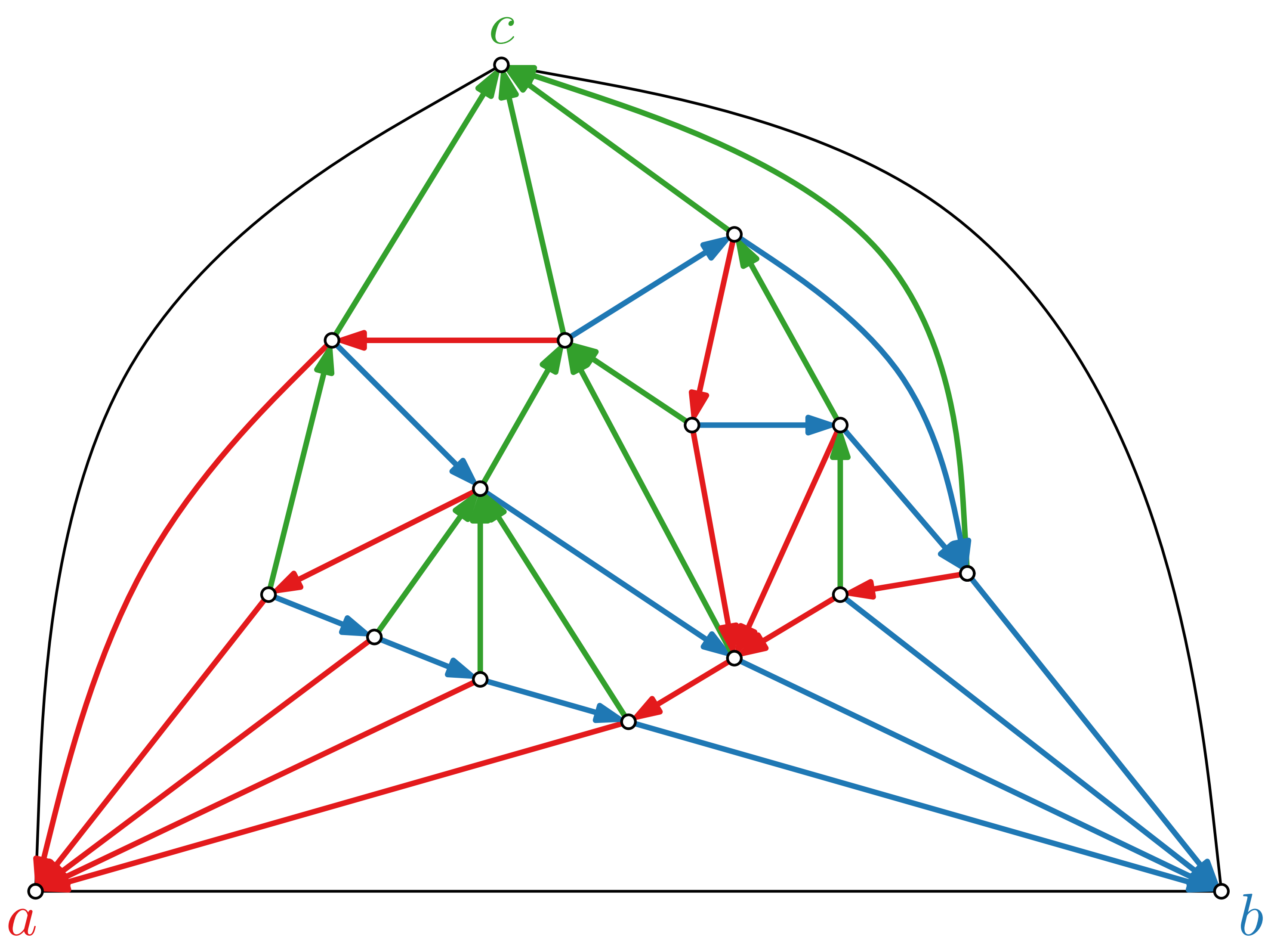
This constructive proof yields an algorithm for computing a Schnyder labeling. It can be implemented to run in $\mathcal{O}(n)$ time.

→ *Exercise*

Schnyder Wood – More Properties



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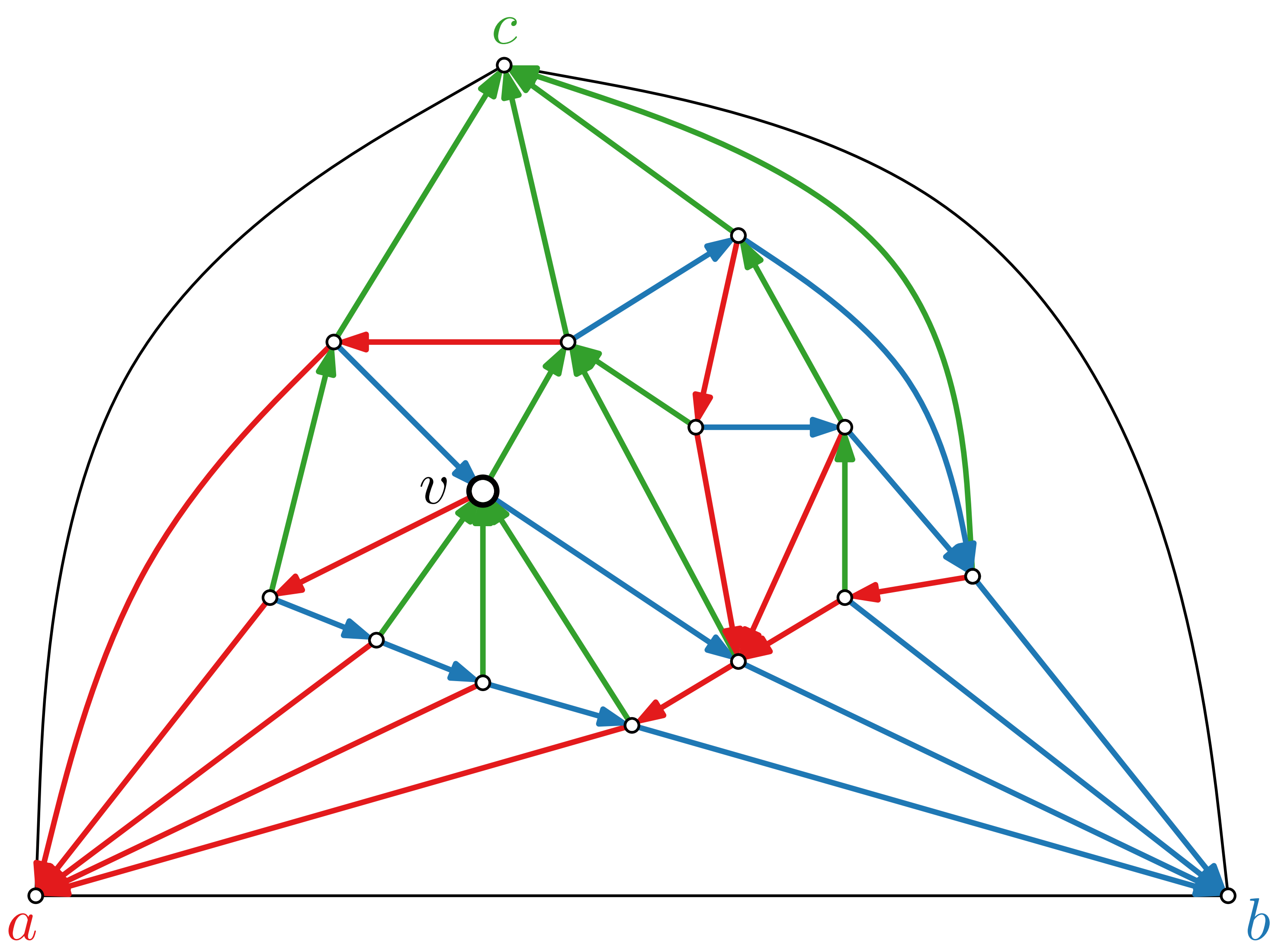

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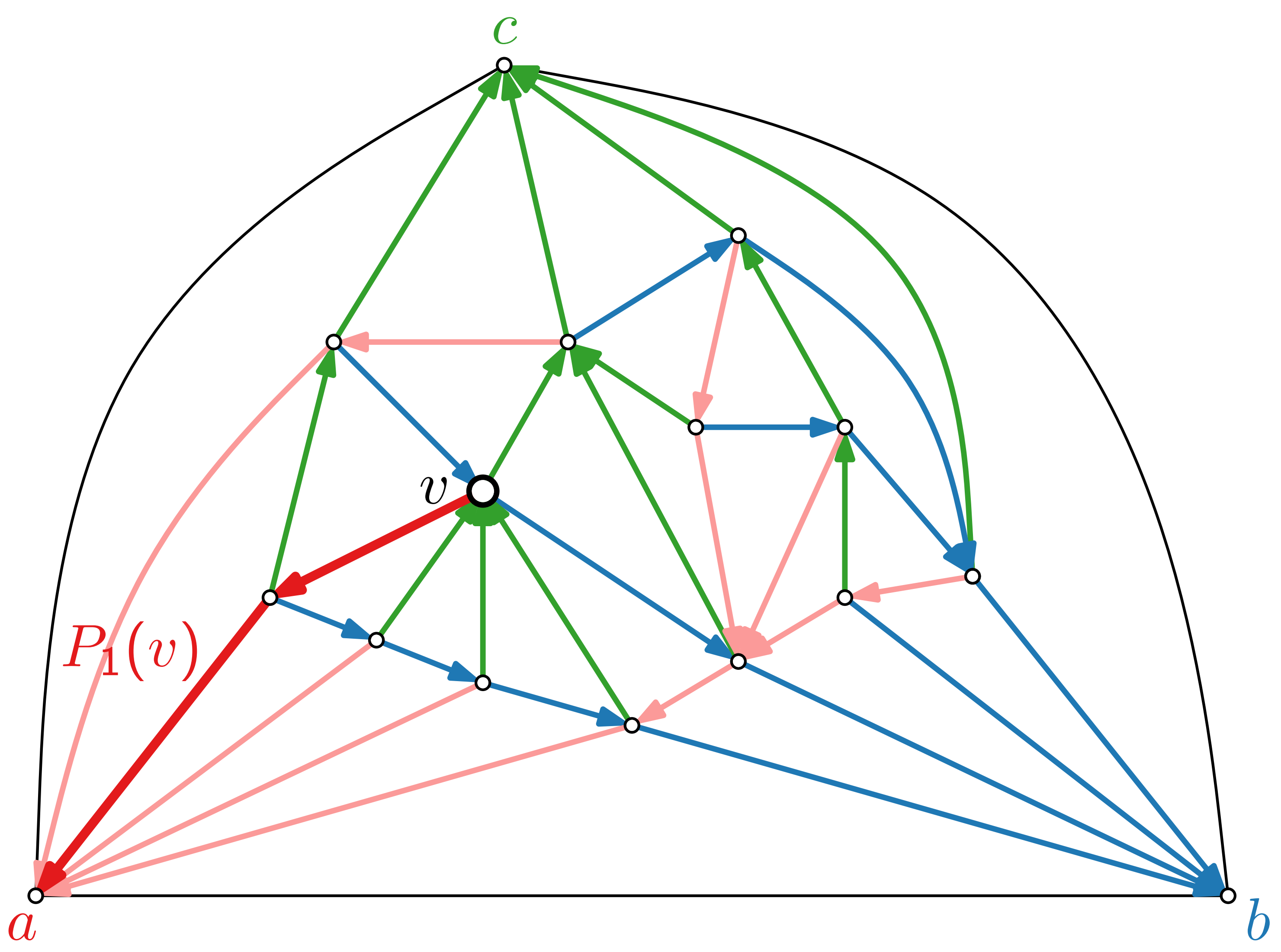

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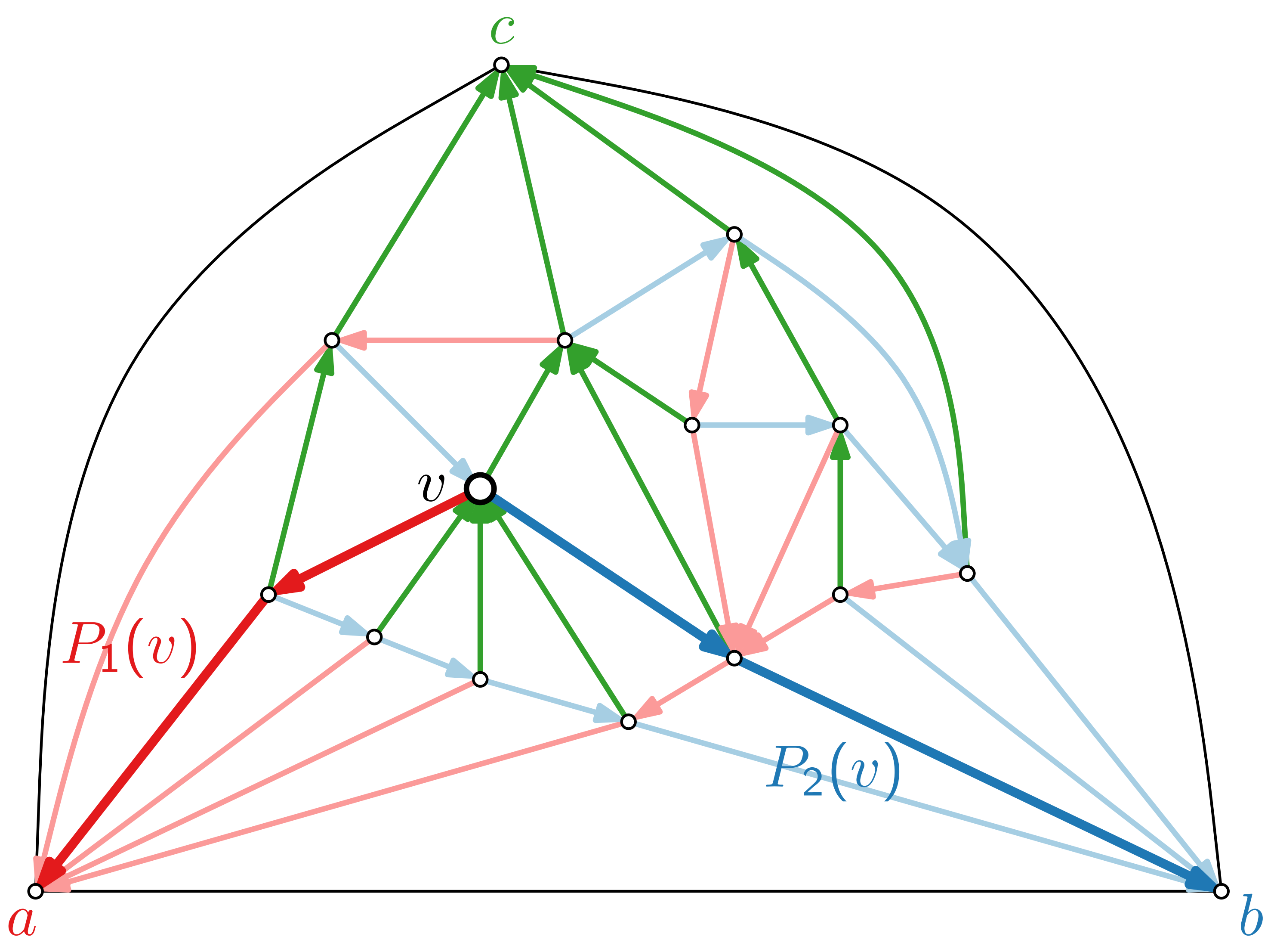

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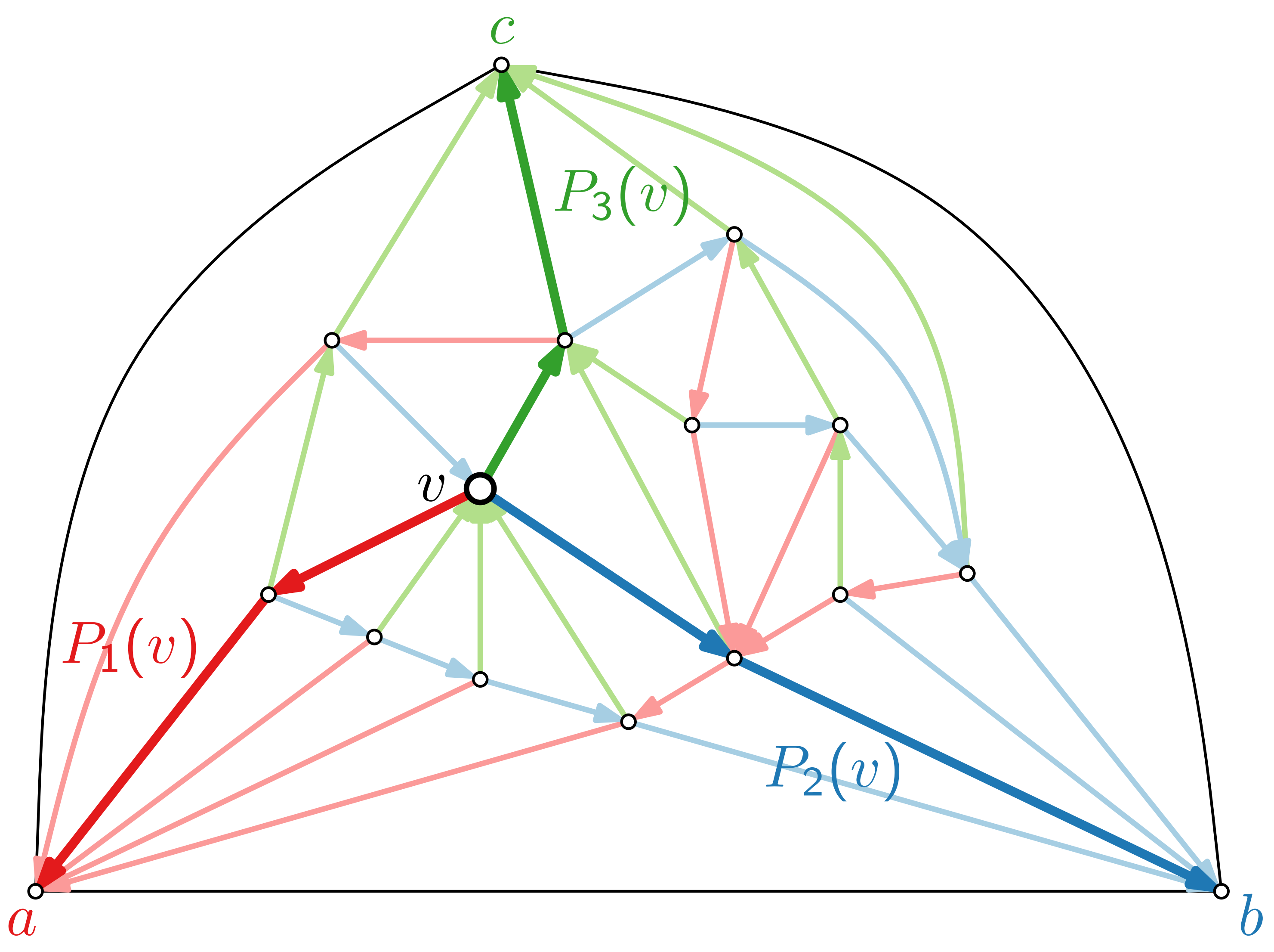

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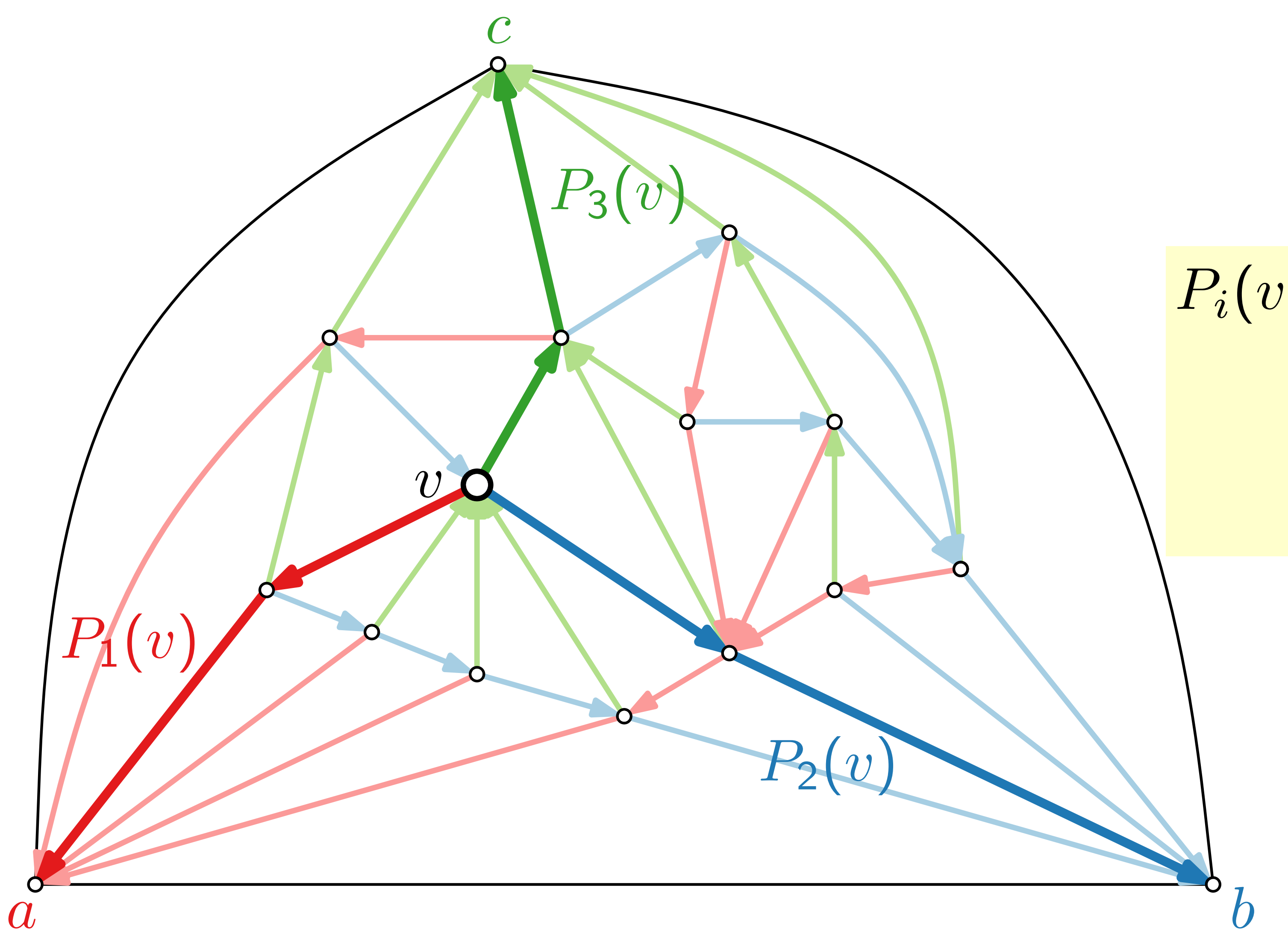

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$P_i(v)$: unique path from v to root of T_i

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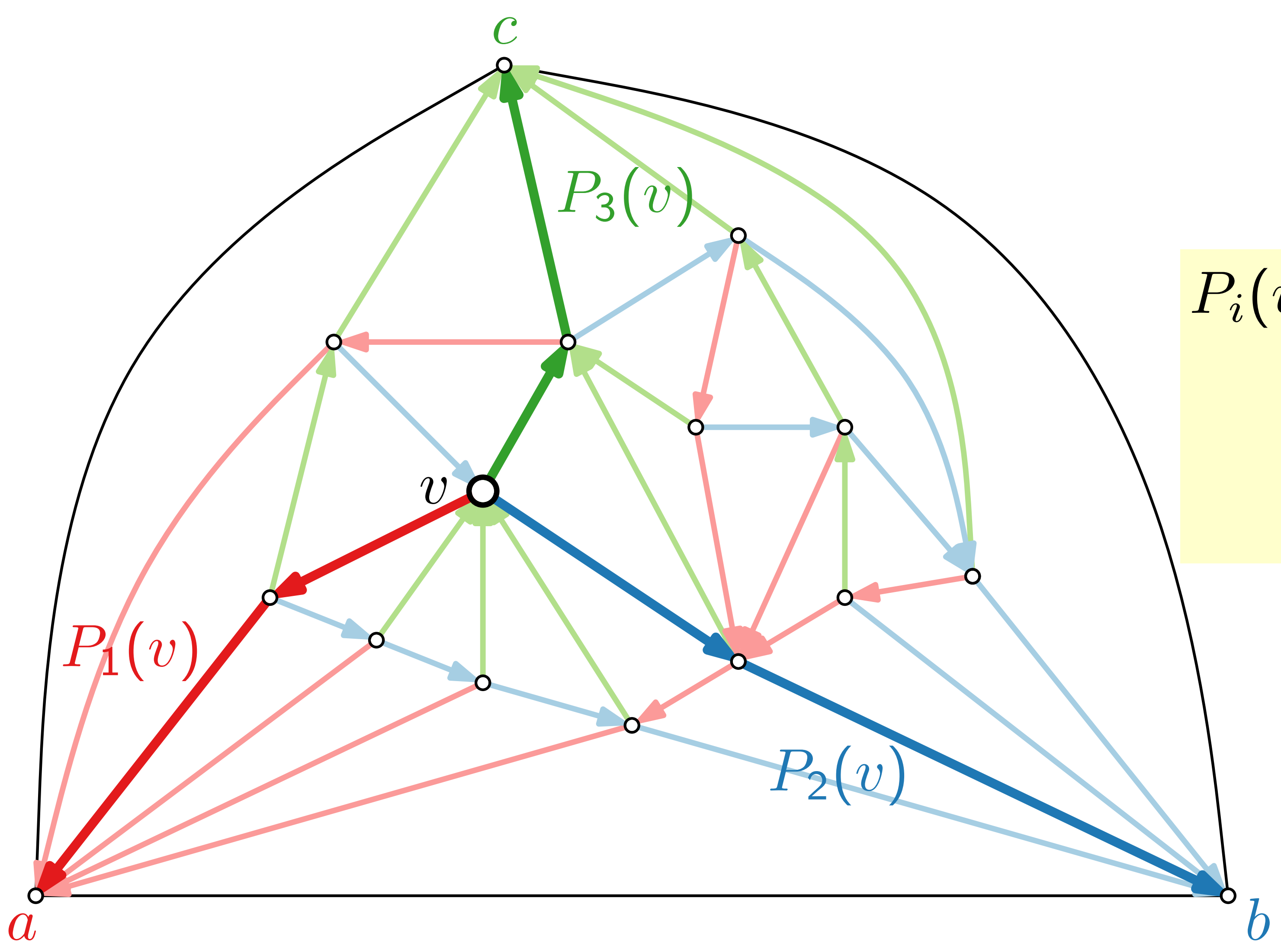

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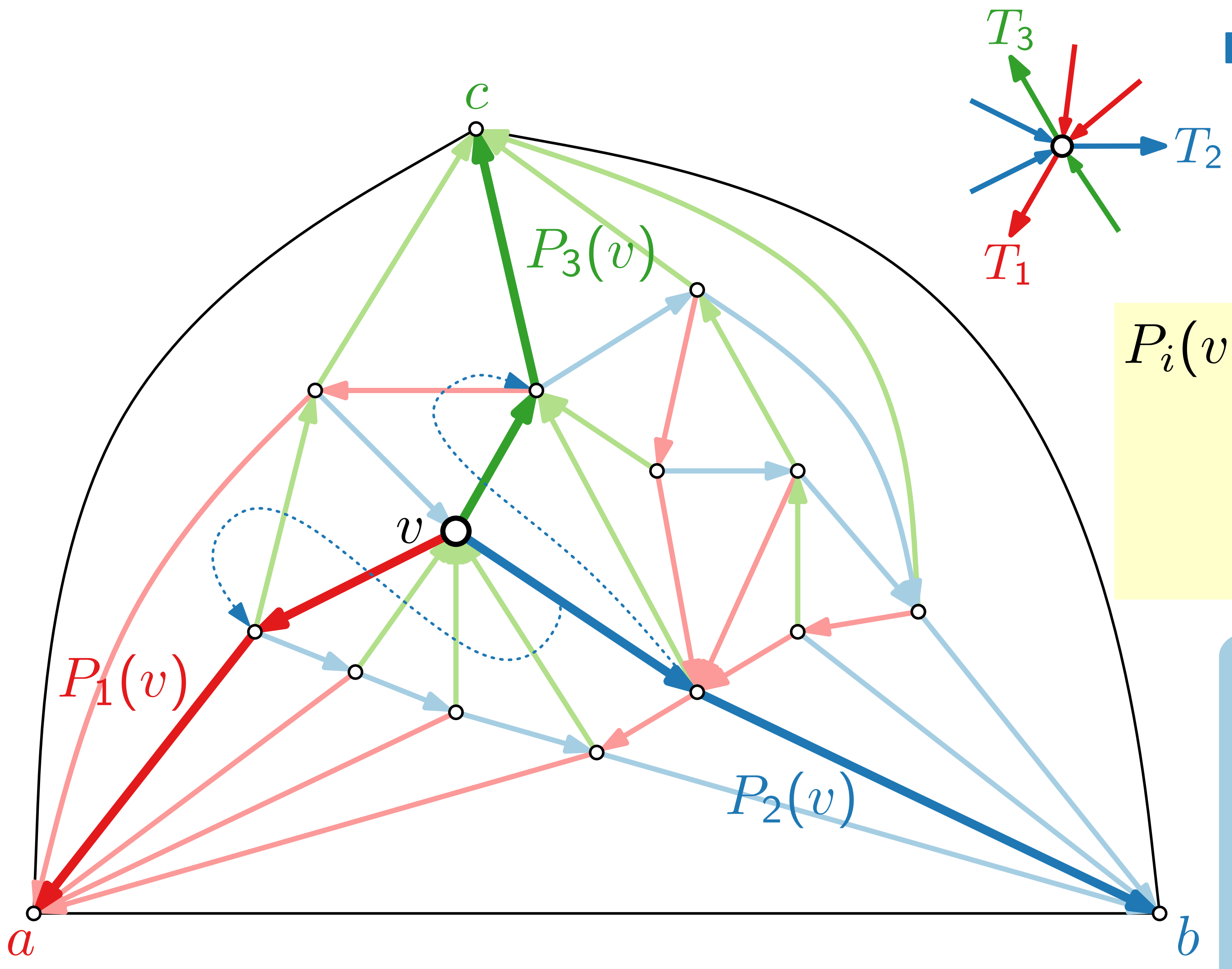
Lemma.

- $P_1(v)$, $P_2(v)$, $P_3(v)$ cross only at v.

Schnyder Wood – More Properties



- From each vertex v there exists a unique
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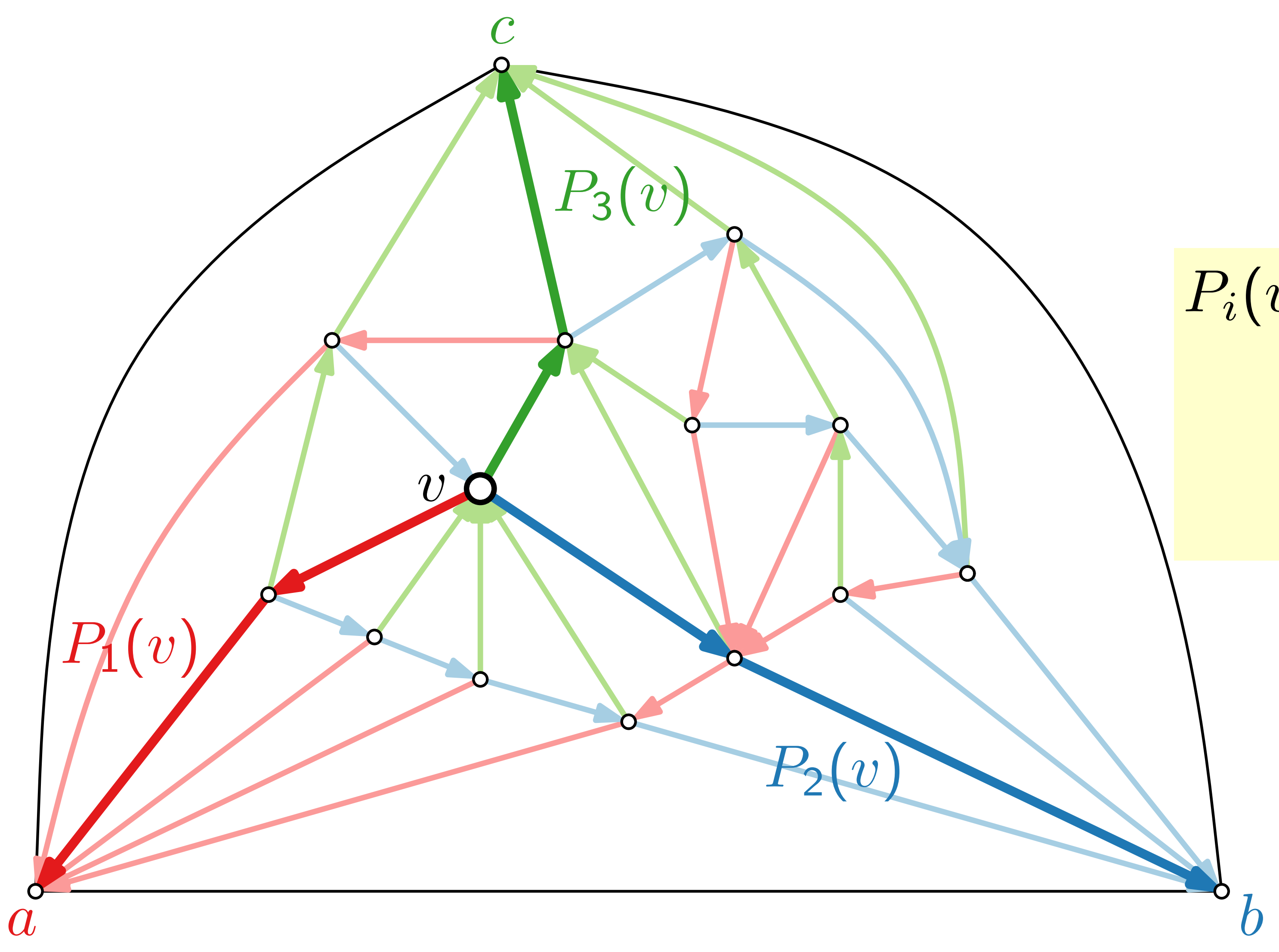

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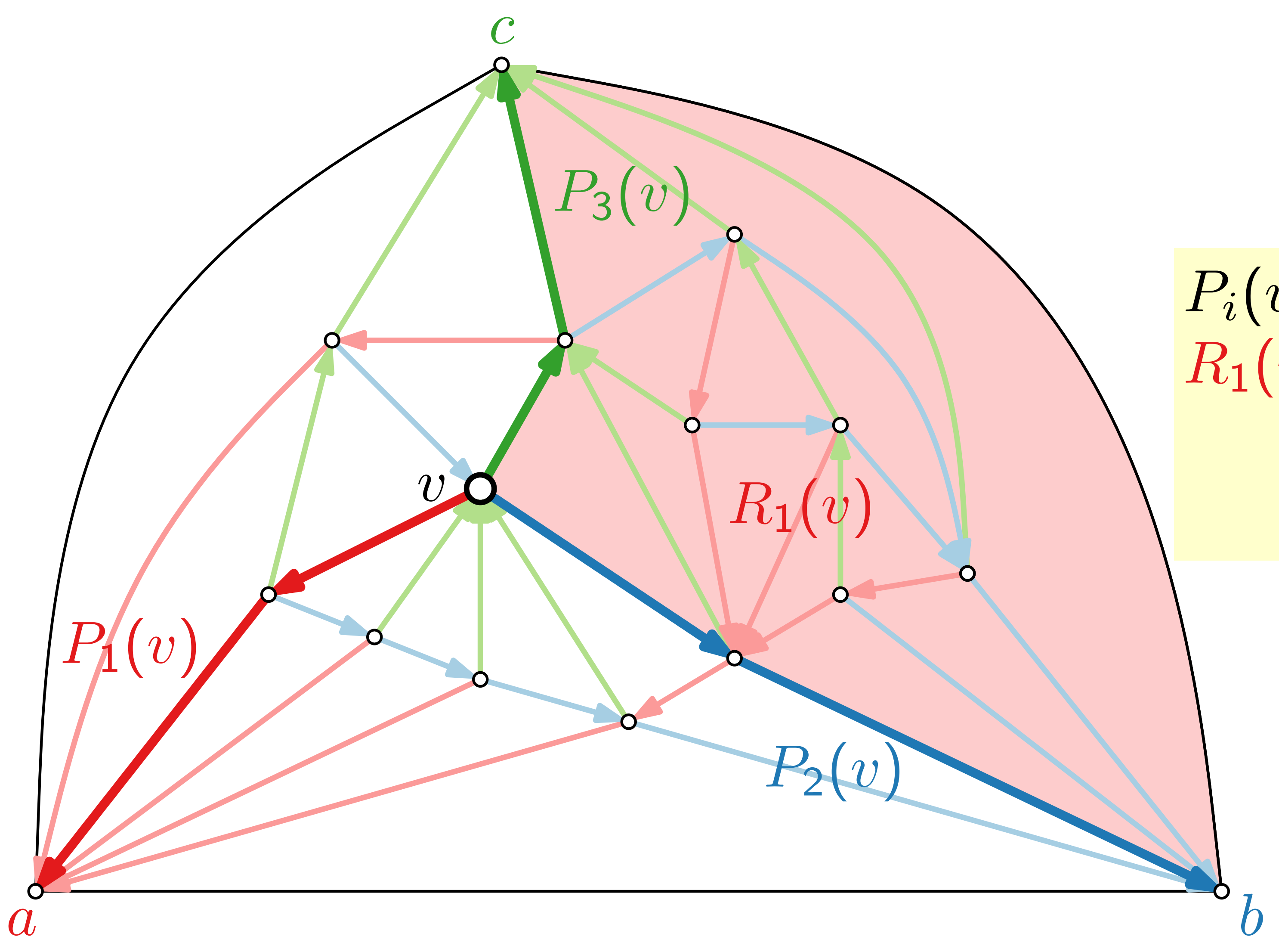

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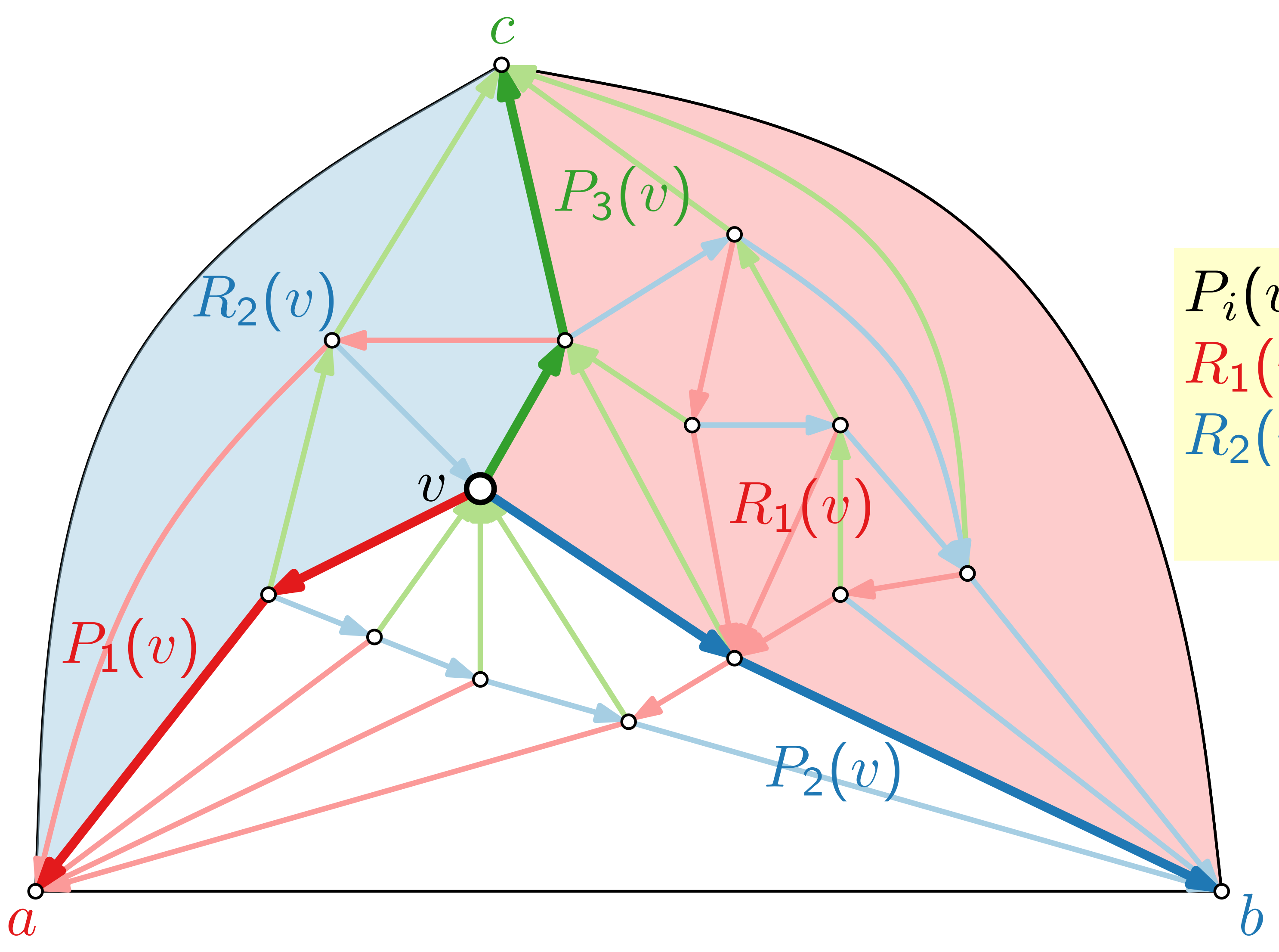

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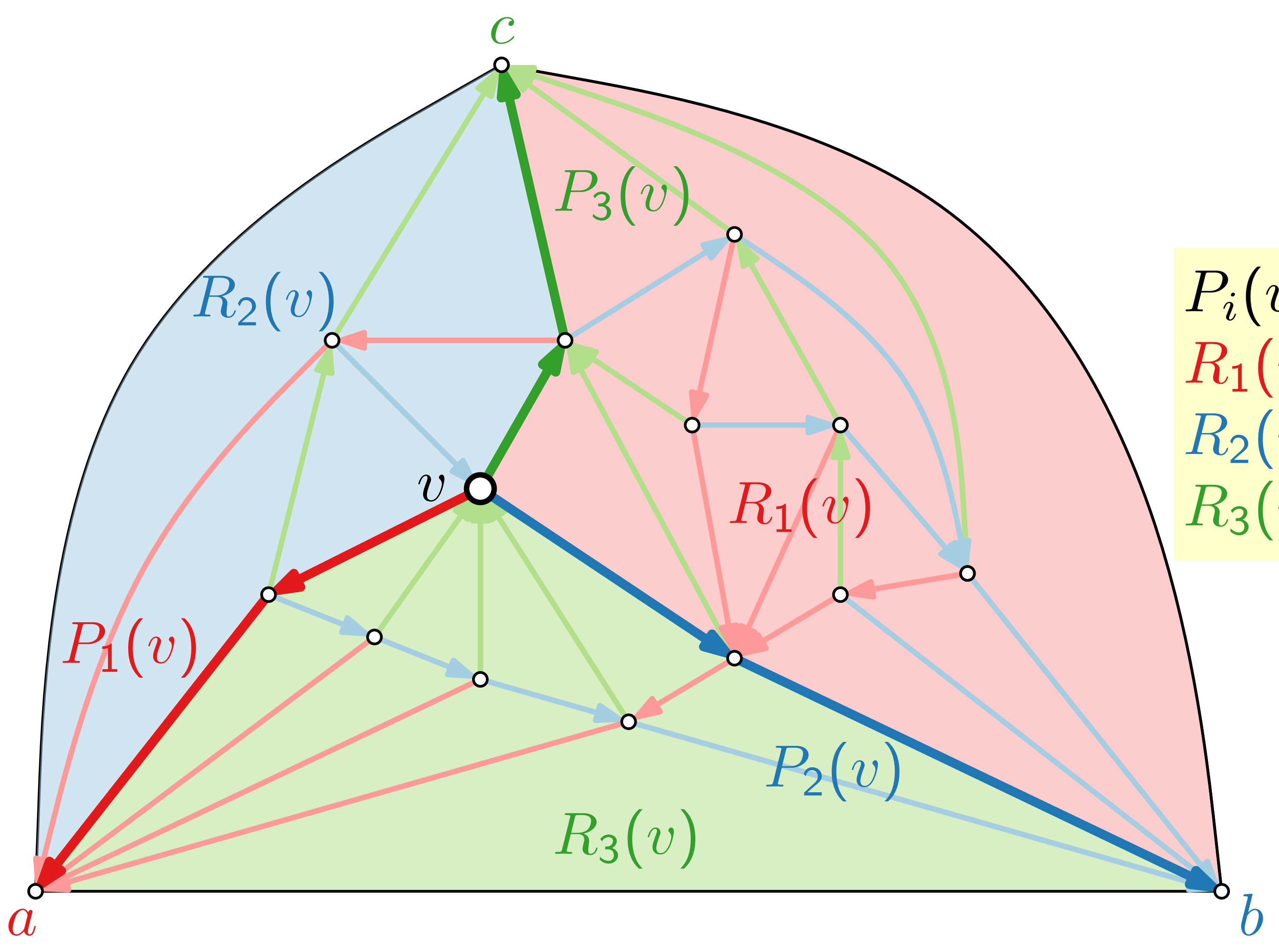

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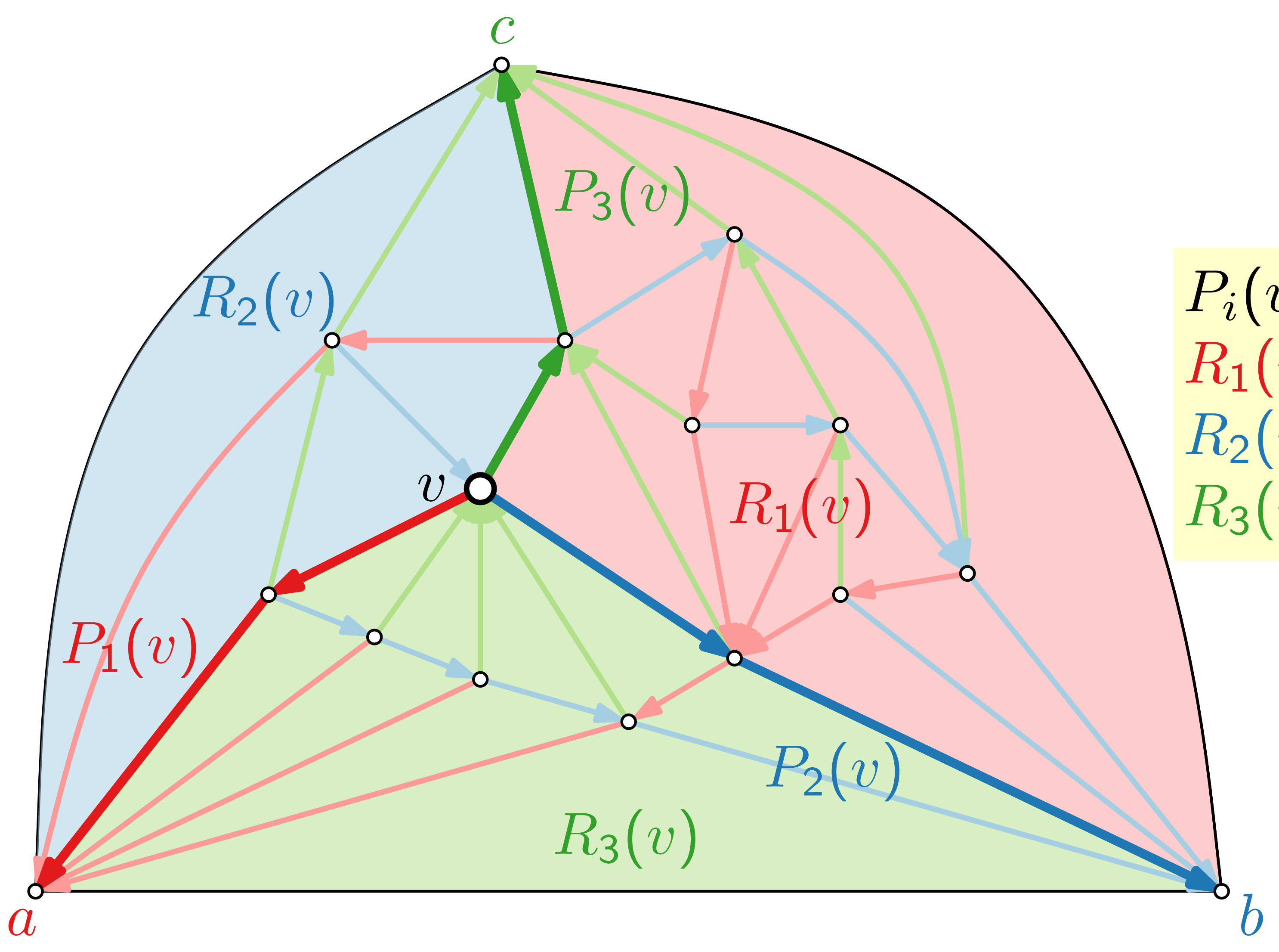

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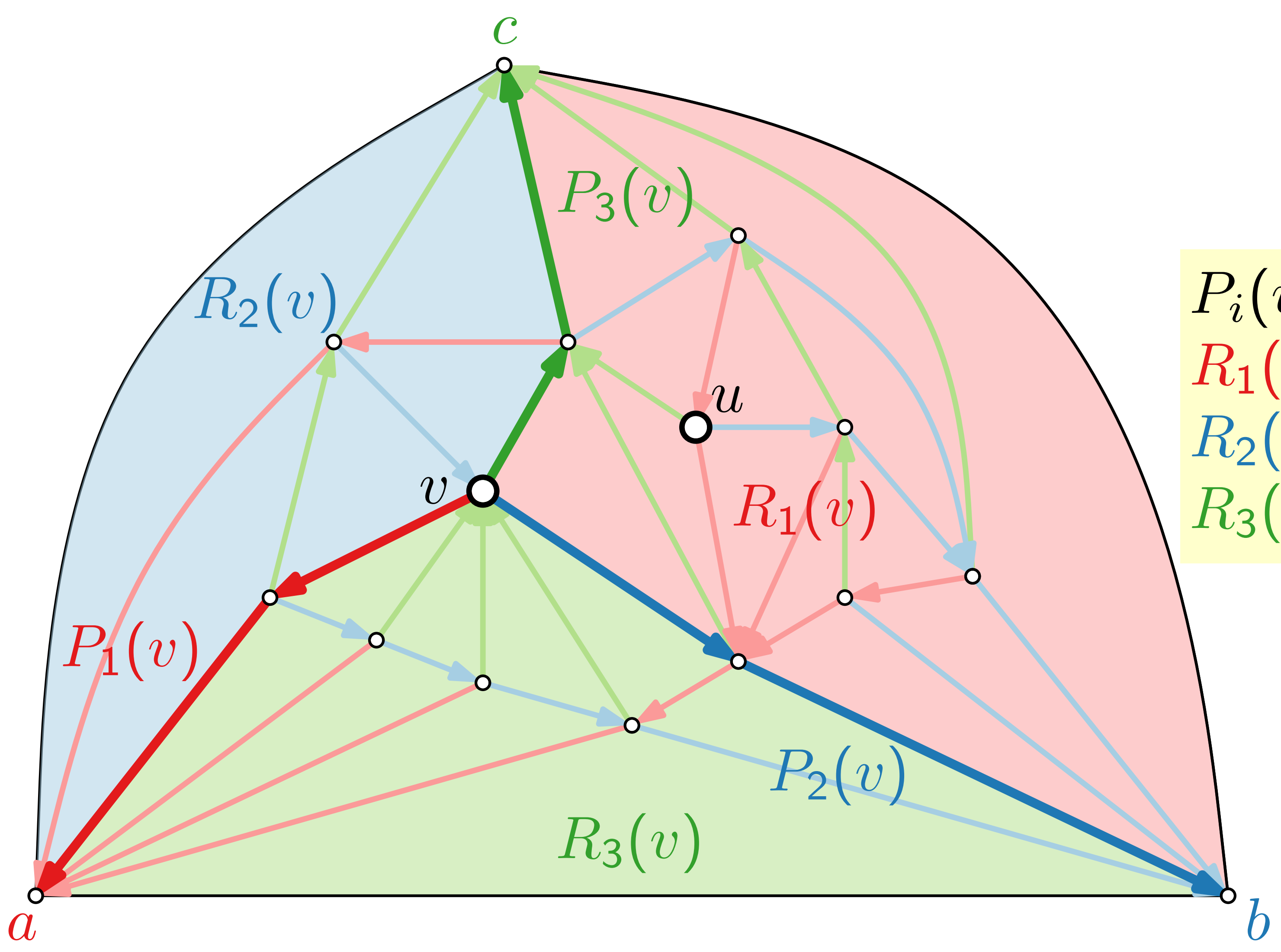

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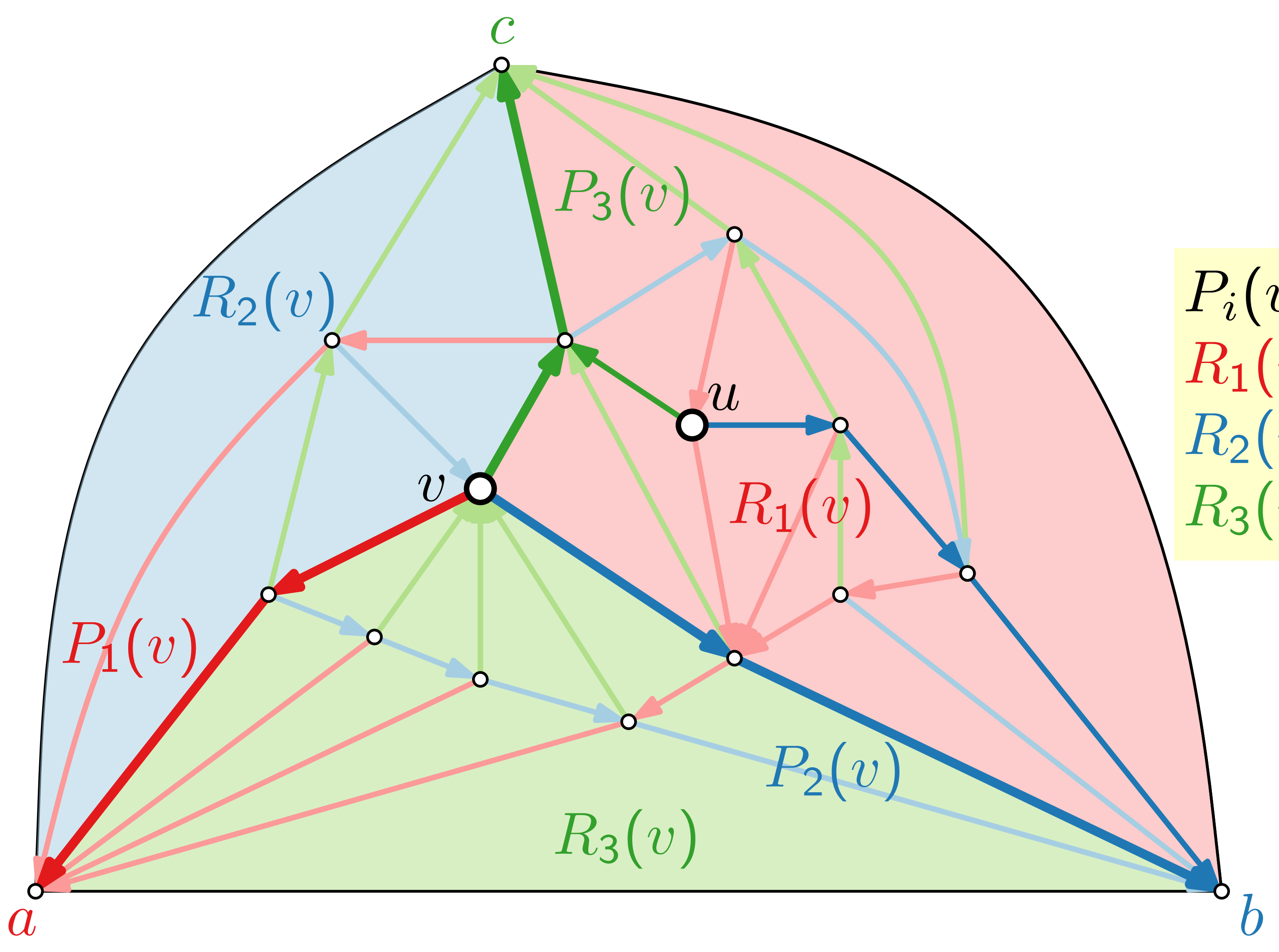

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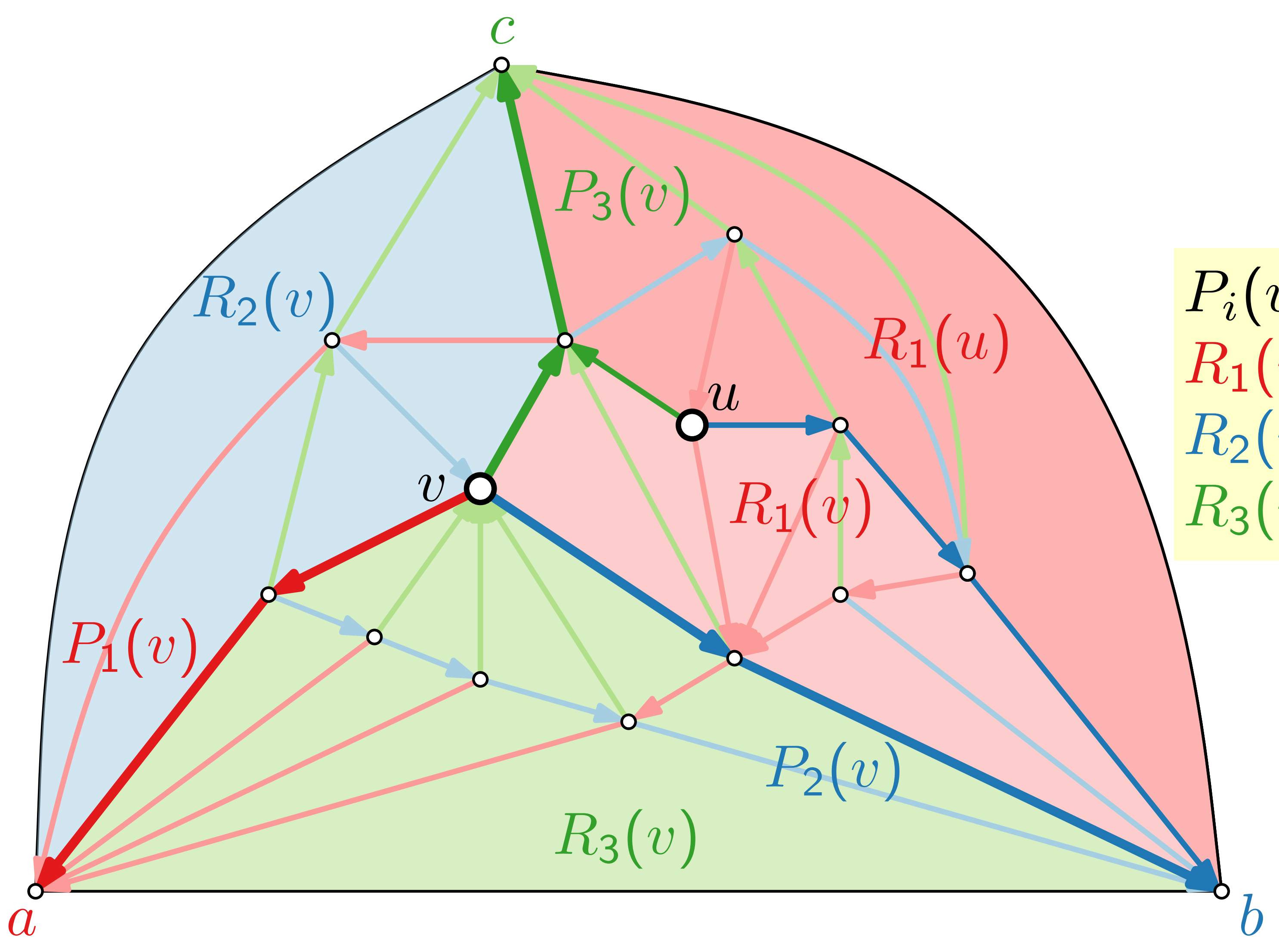

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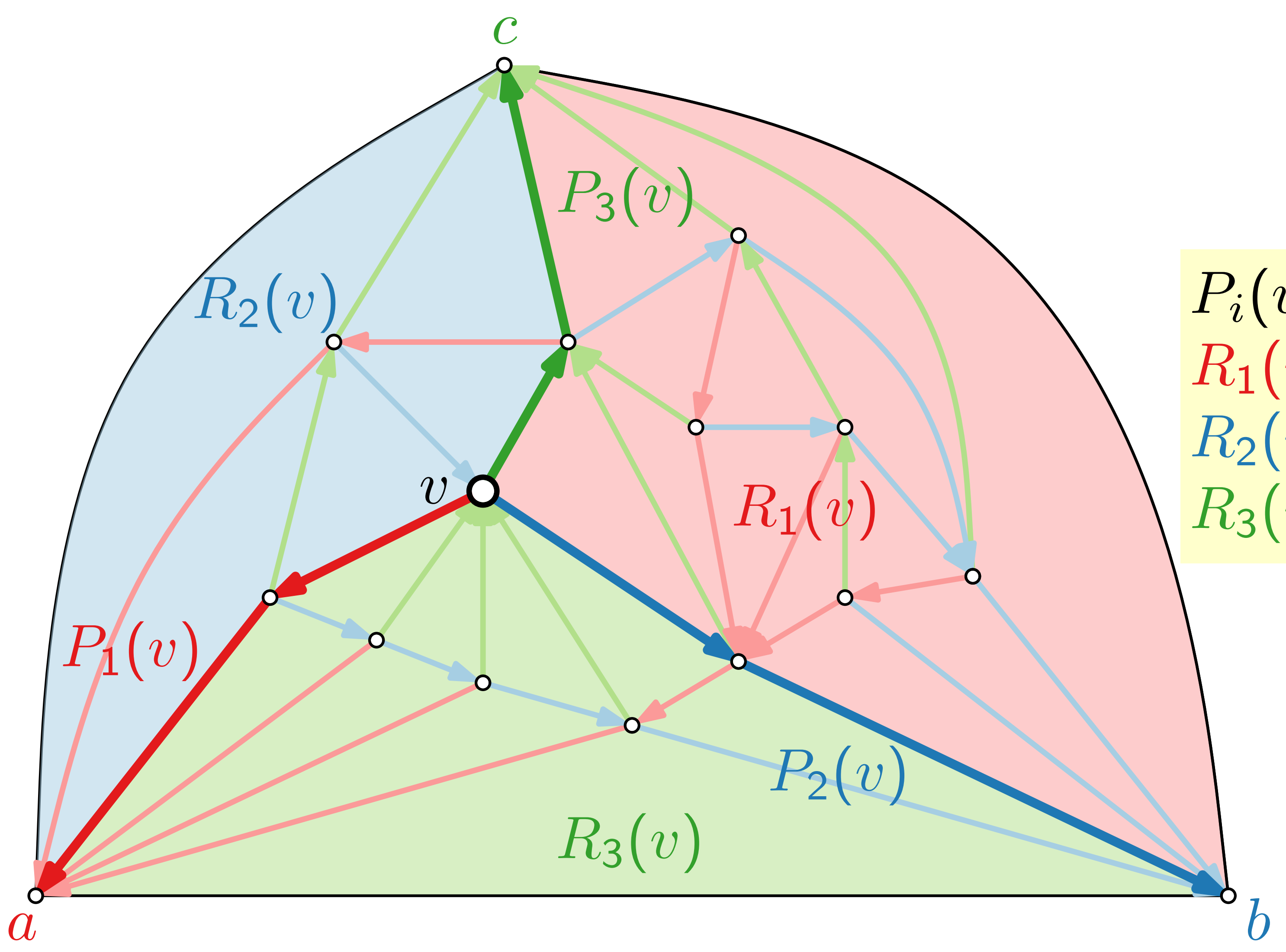

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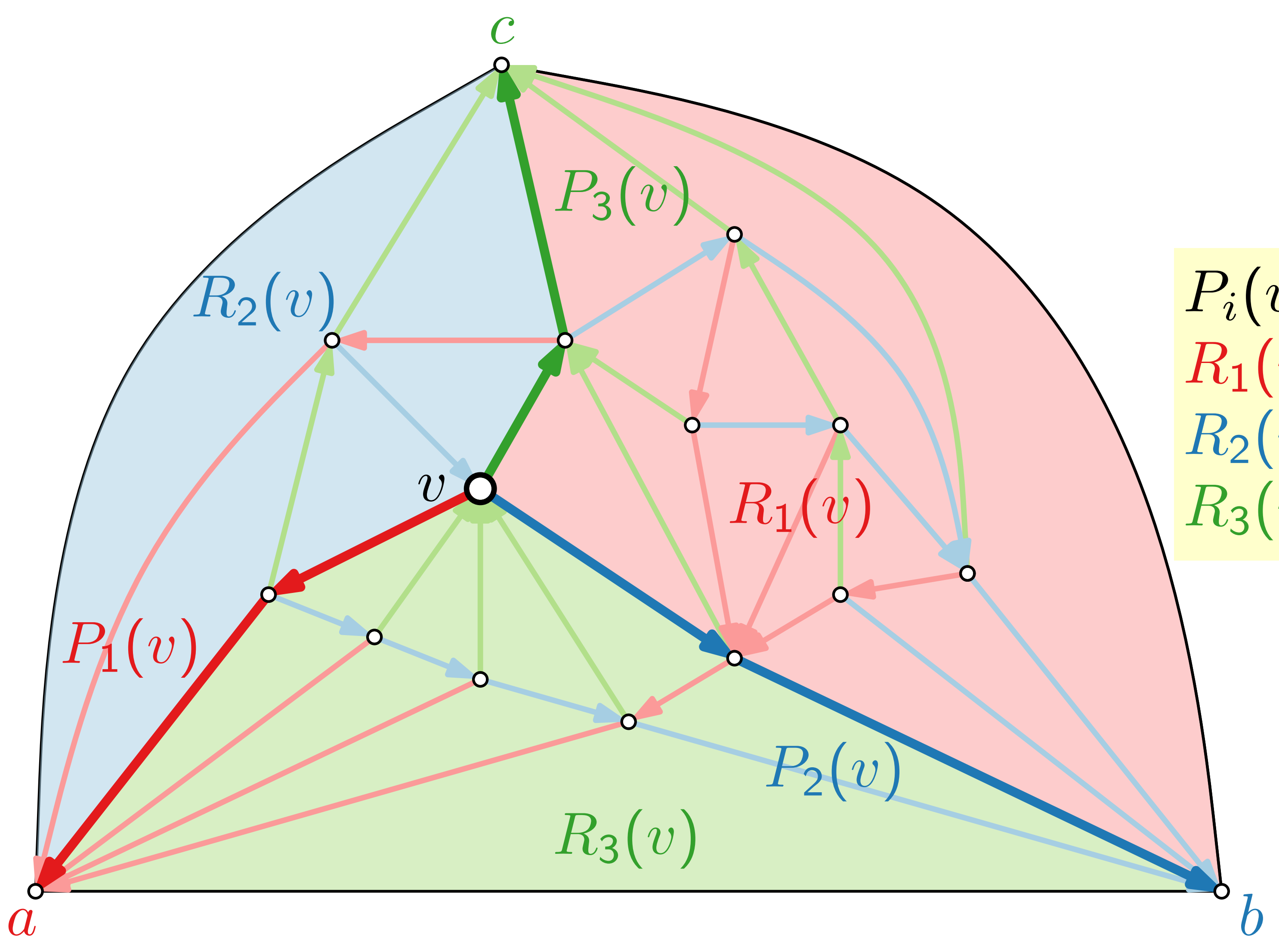

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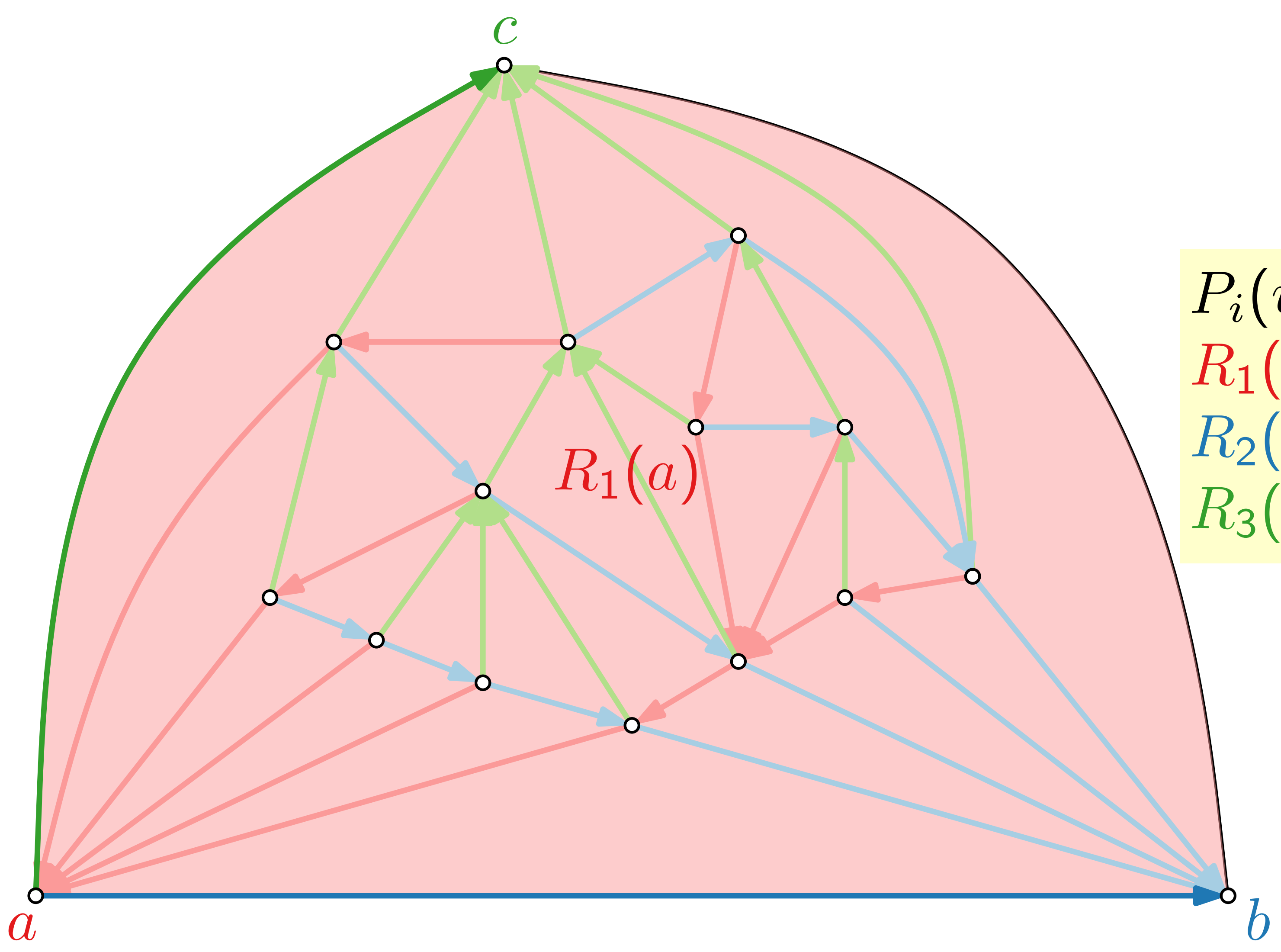

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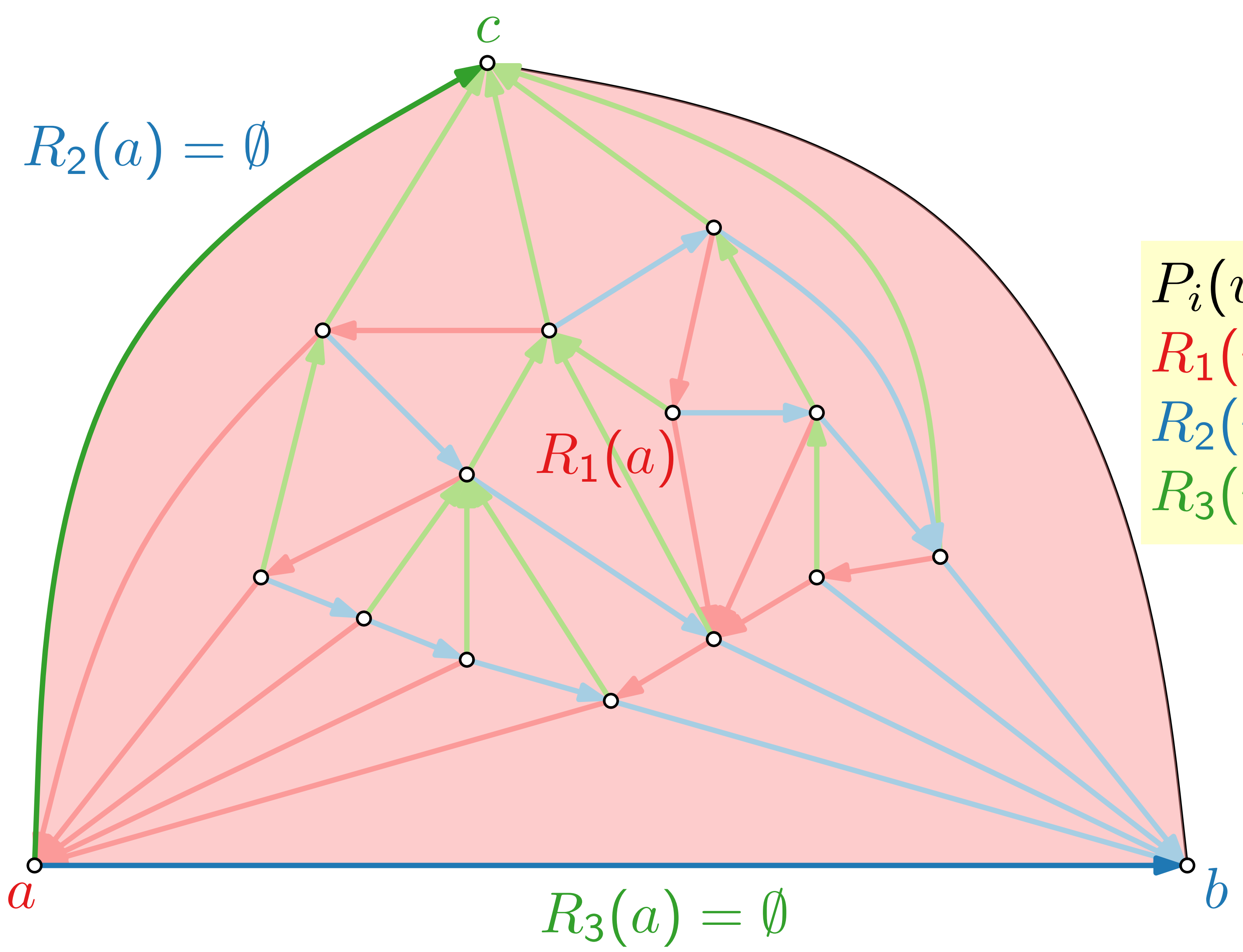

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Theorem. [Schnyder '90]

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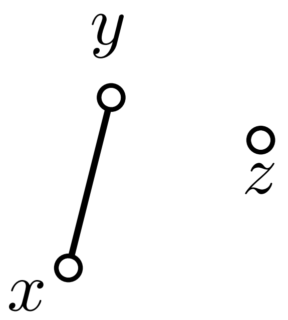

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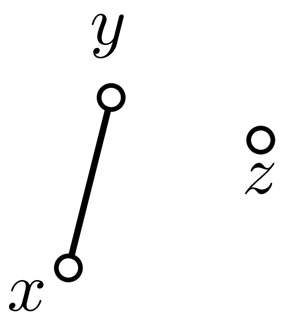

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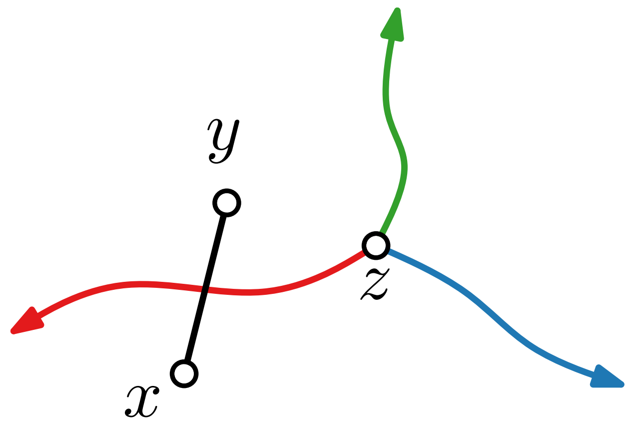

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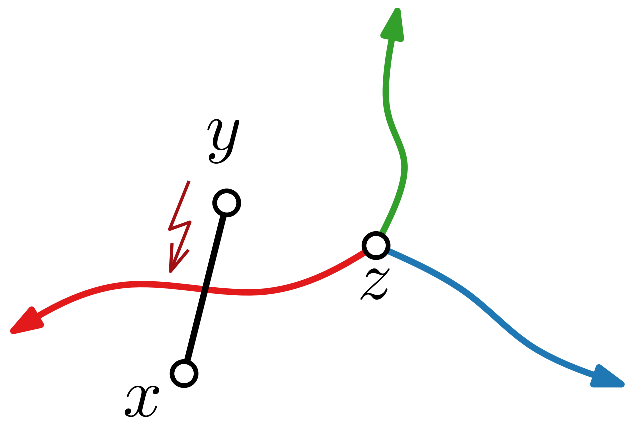

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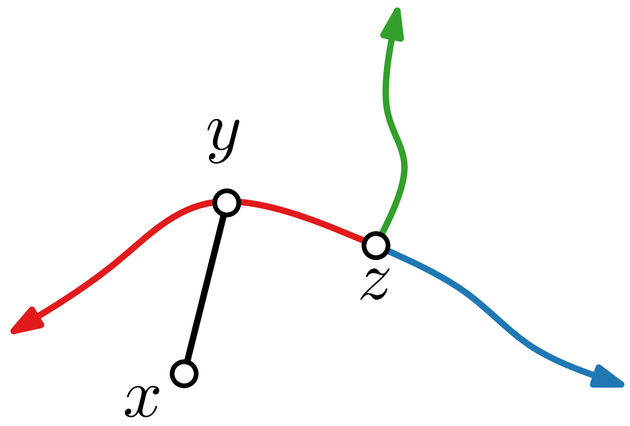

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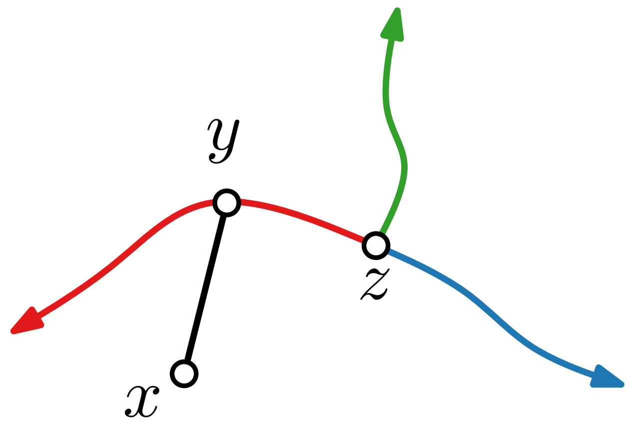

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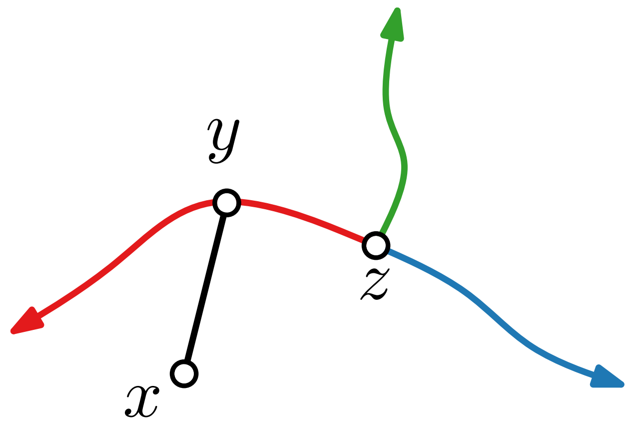

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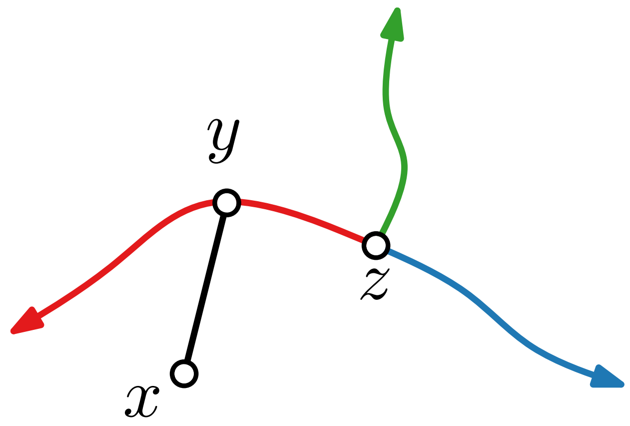

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Set $A = (0, 0)$, $B = (2n-5, 0)$, and $C = (0, 2n-5)$.

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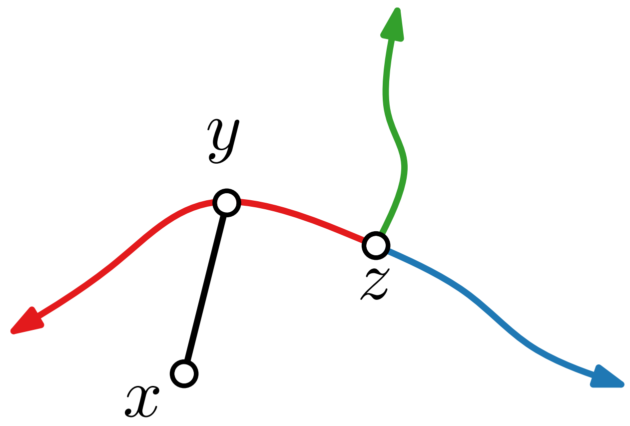


Schnyder Drawing

Set $A = (0,0)$, $B = (2n-5, 0)$, and $C = (0, 2n-5)$.

Theorem. [Schnyder '90]
For a plane triangulation G, the mapping

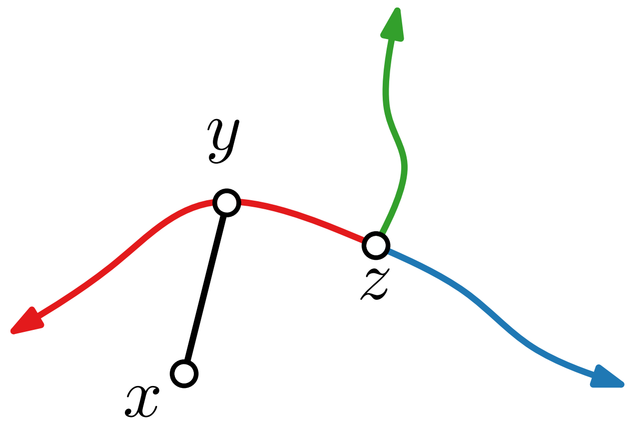
$$f\colon v \mapsto (v_1, v_2, v_3) = \frac{1}{2n-5}(|R_1(v)|, |R_2(v)|, |R_3(v)|)$$

is a barycentric representation of G and, thus, yields a planar straight-line drawing of G on the $(2n-5) \times (2n-5)$ grid.

(B1) $v_1 + v_2 + v_3 = 1$ for all $v \in V(G)$ ✓

(B2) for each $\{x, y\} \in E(G)$ and each $z \in V(G) \setminus \{x, y\}$ there exists $k \in \{1, 2, 3\}$ with $x_k < z_k$ and $y_k < z_k$ ✓

- $\{x, y\}$ must lie in $R_i(z)$ for some $i \in \{1, 2, 3\}$
- $x, y \in R_i(z) \Rightarrow R_i(x), R_i(y) \subsetneq R_i(z)$
  $\Rightarrow |R_i(x)|, |R_i(y)| < |R_i(z)|$



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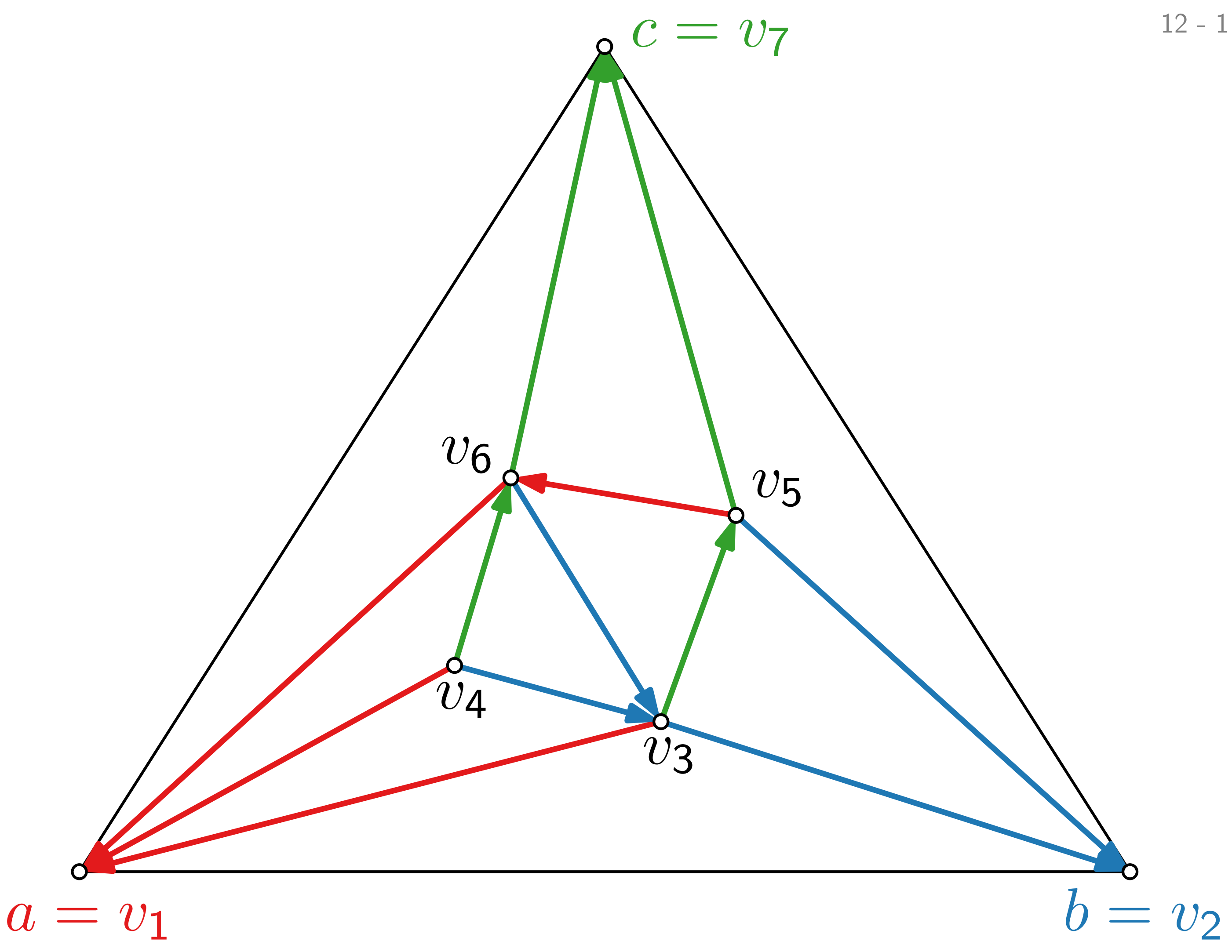

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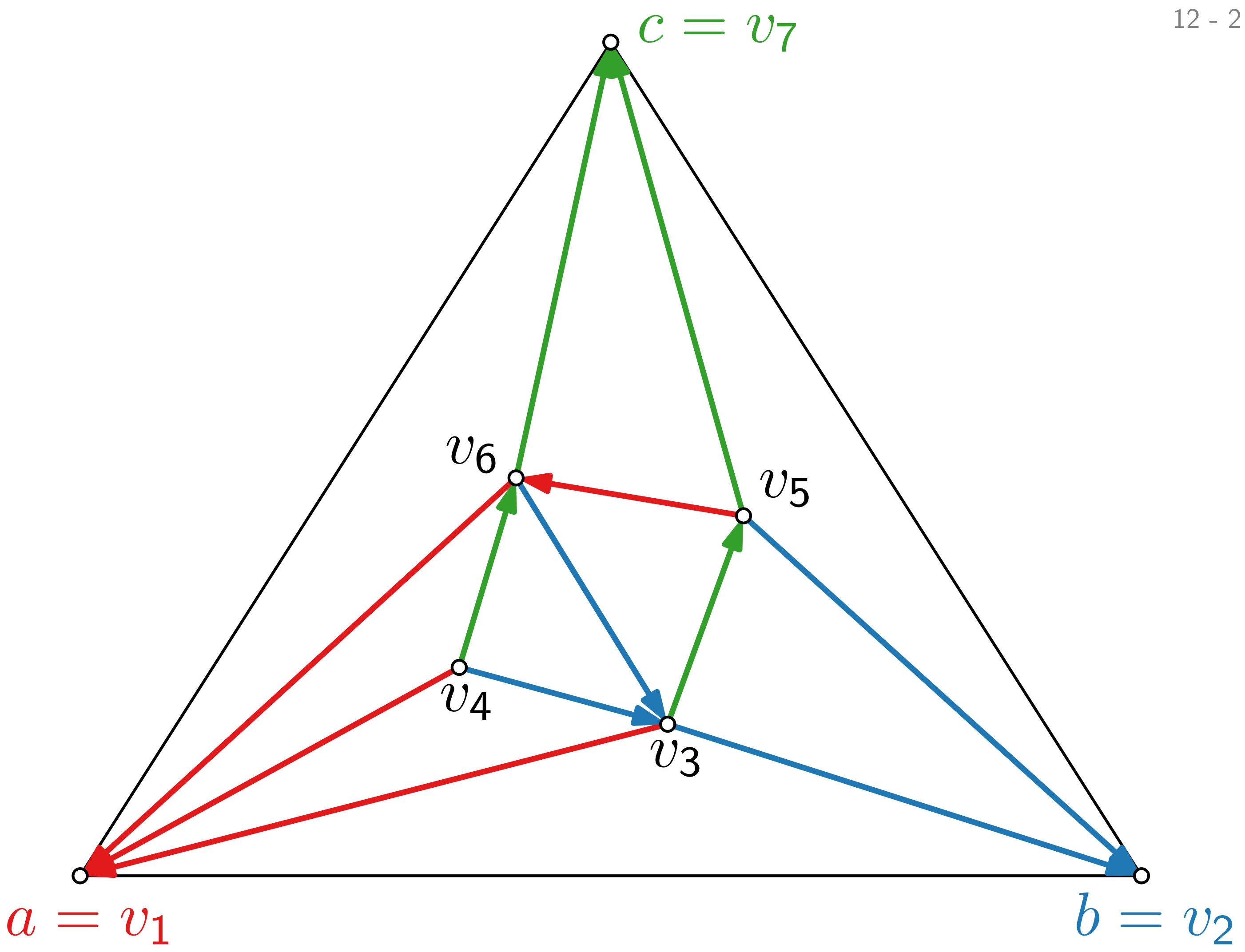

$n = 7; \quad 2n - 5 = 9$

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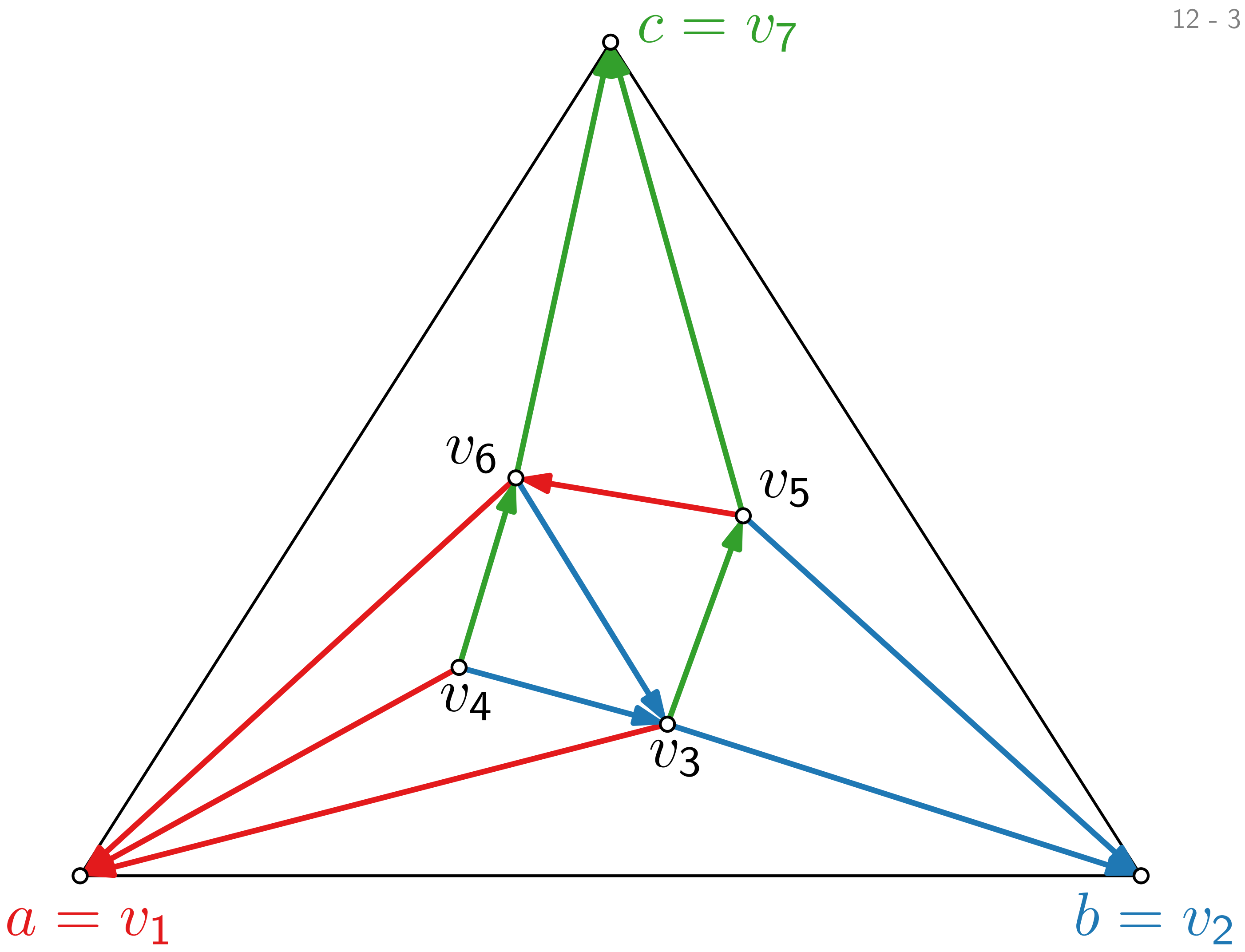

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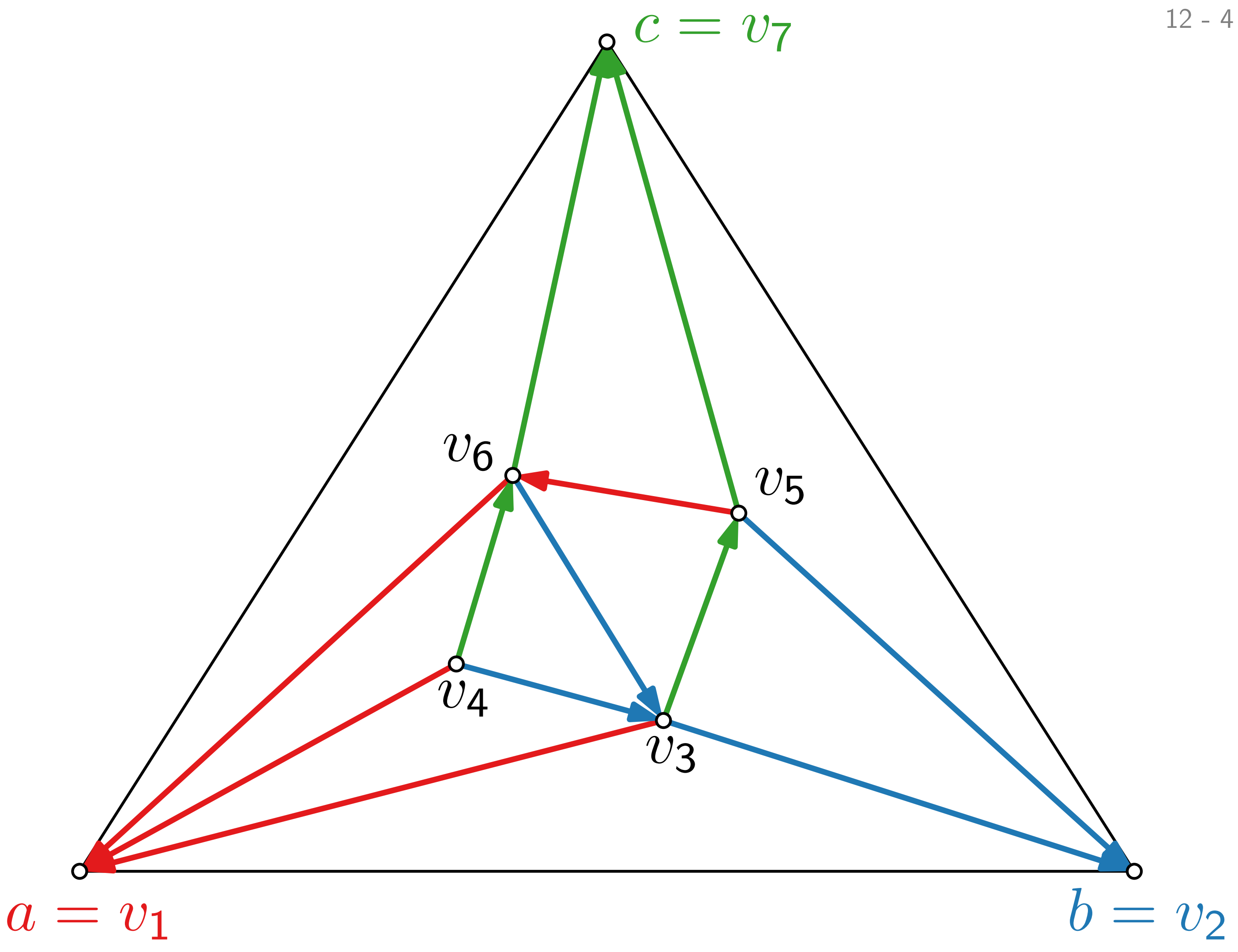

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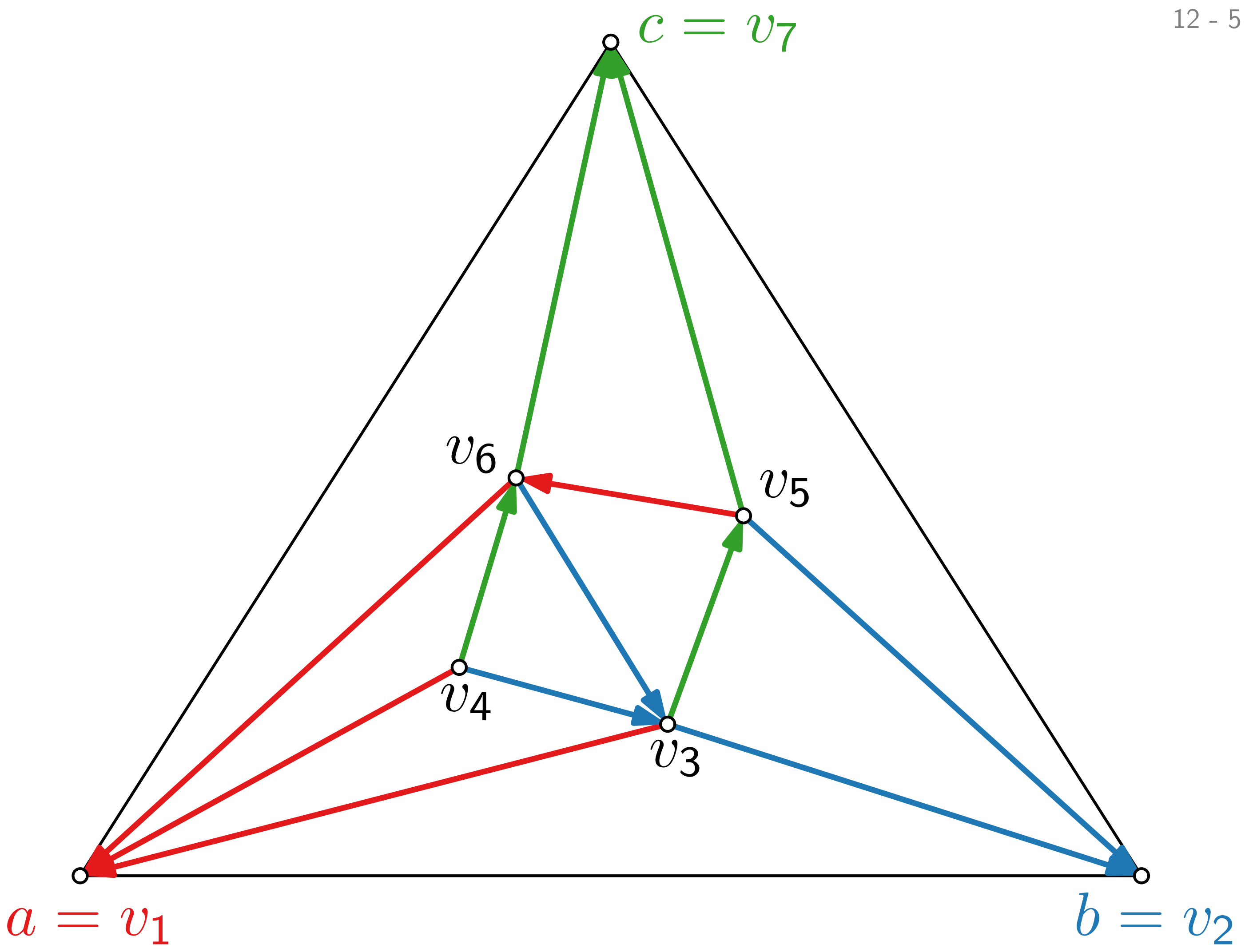

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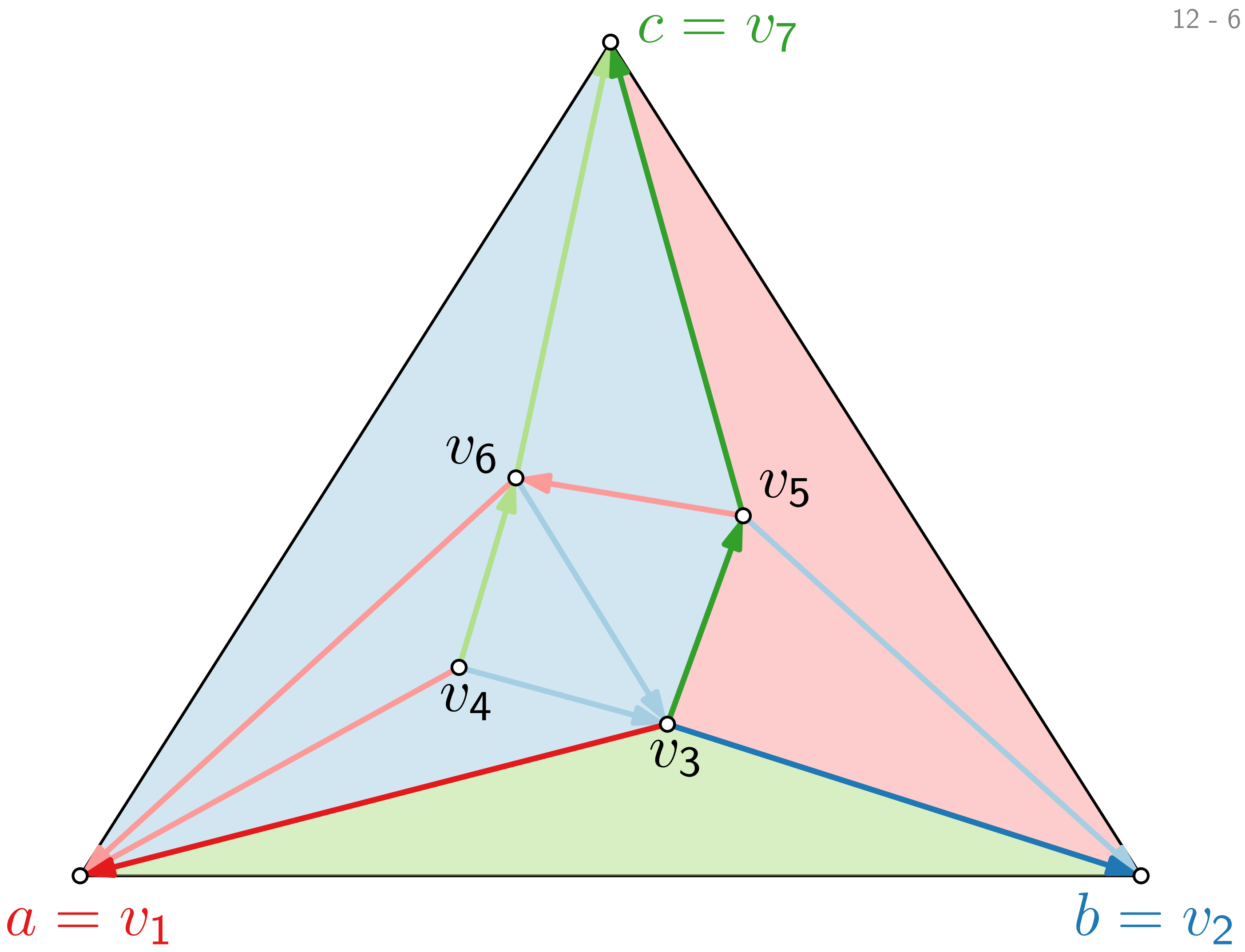

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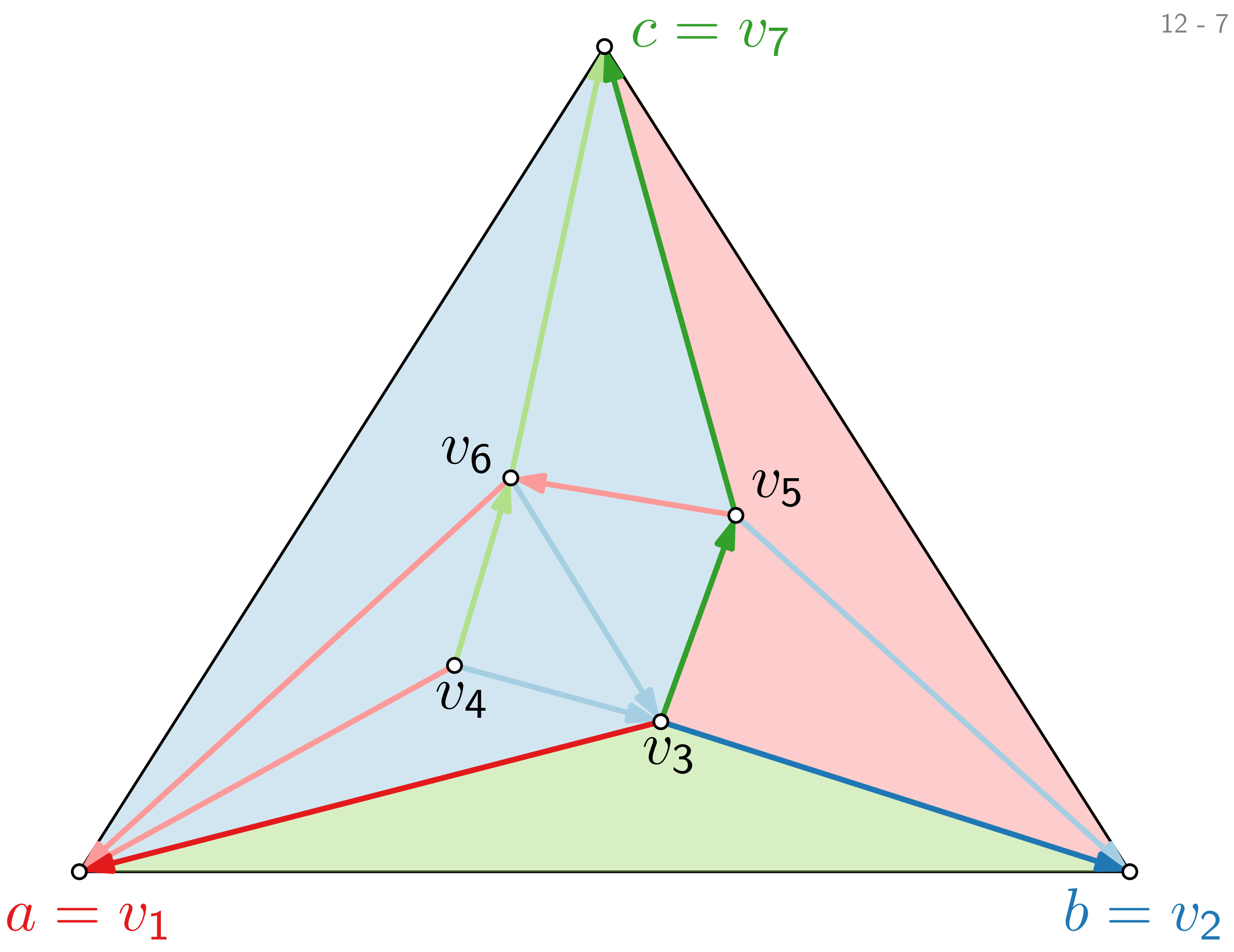

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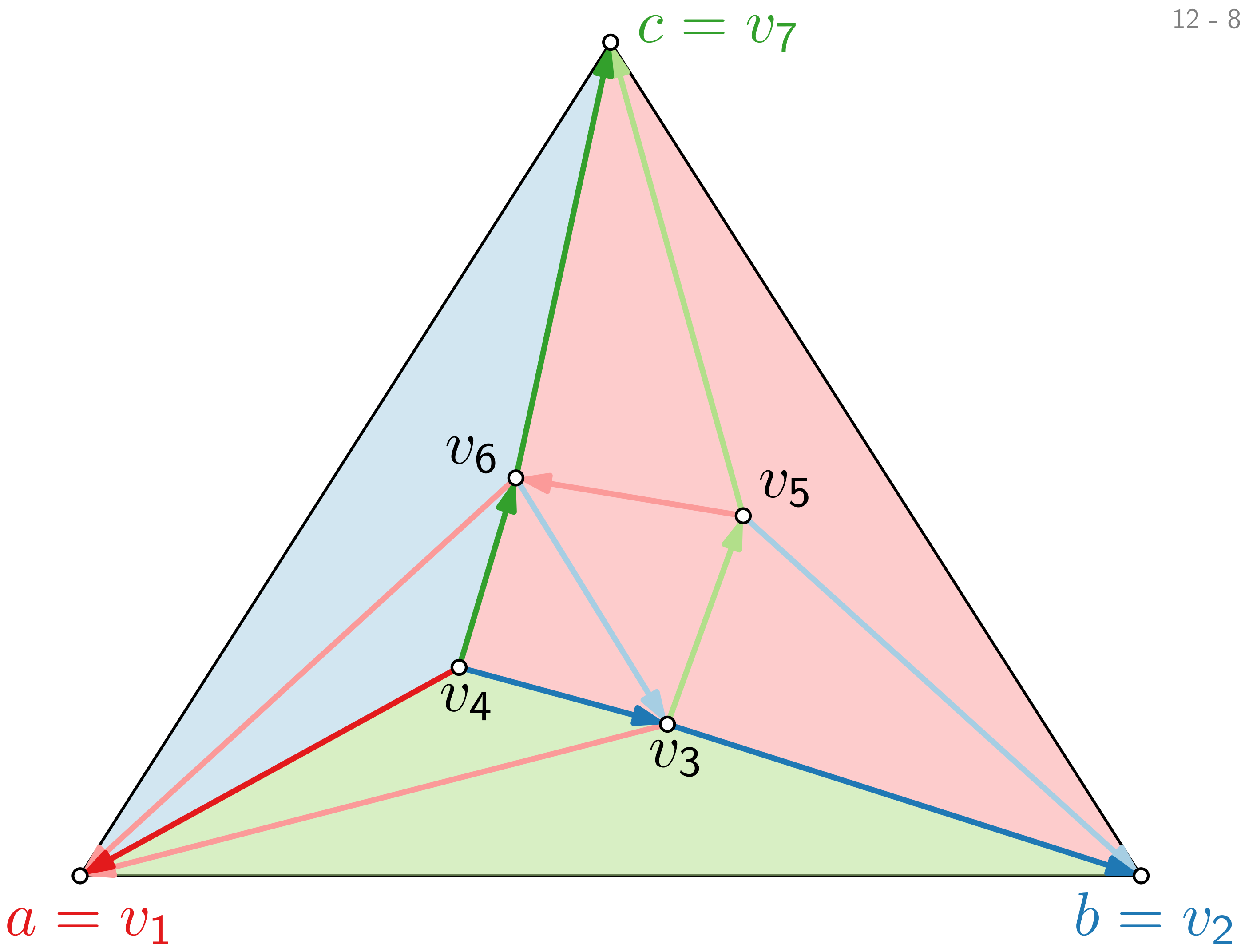

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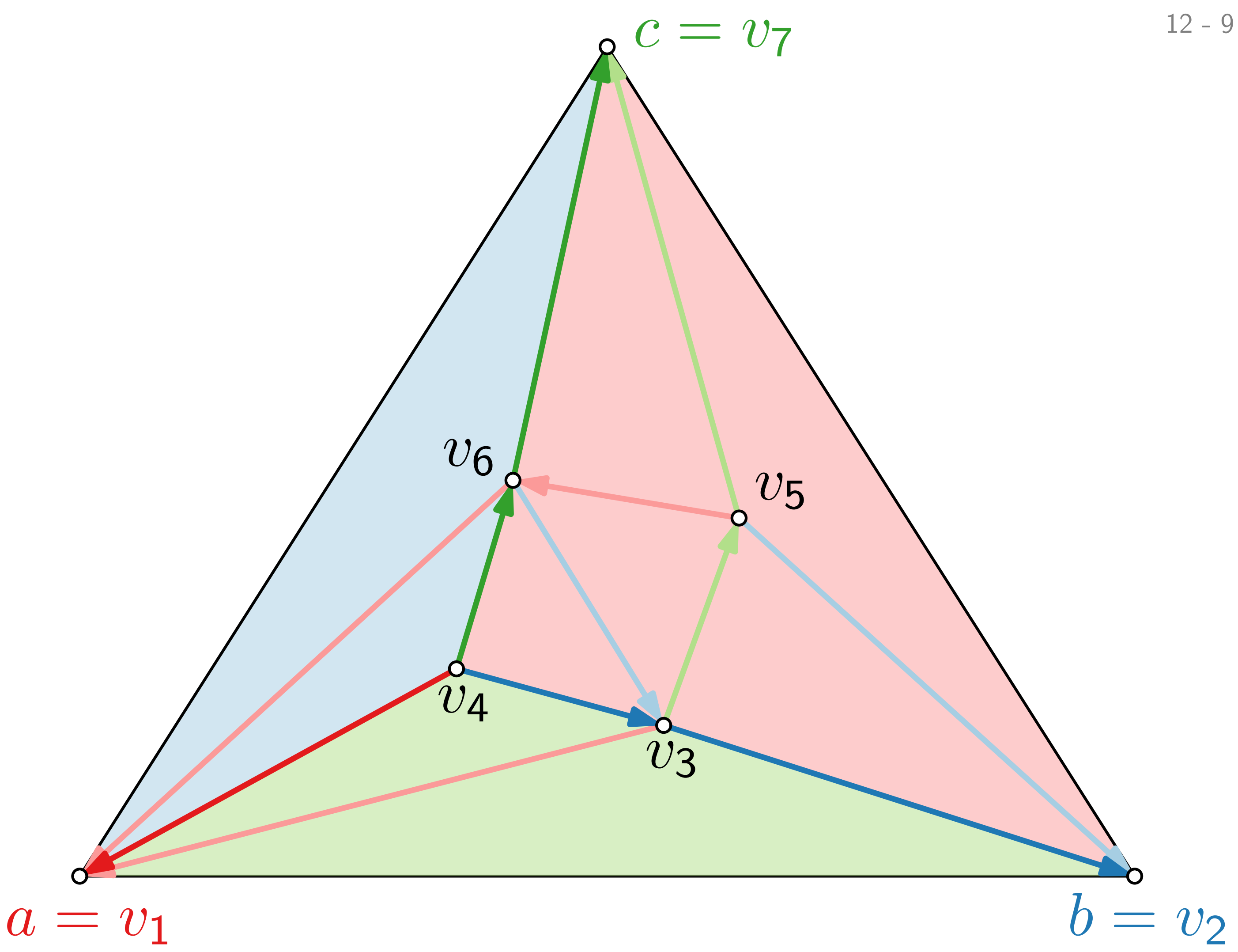

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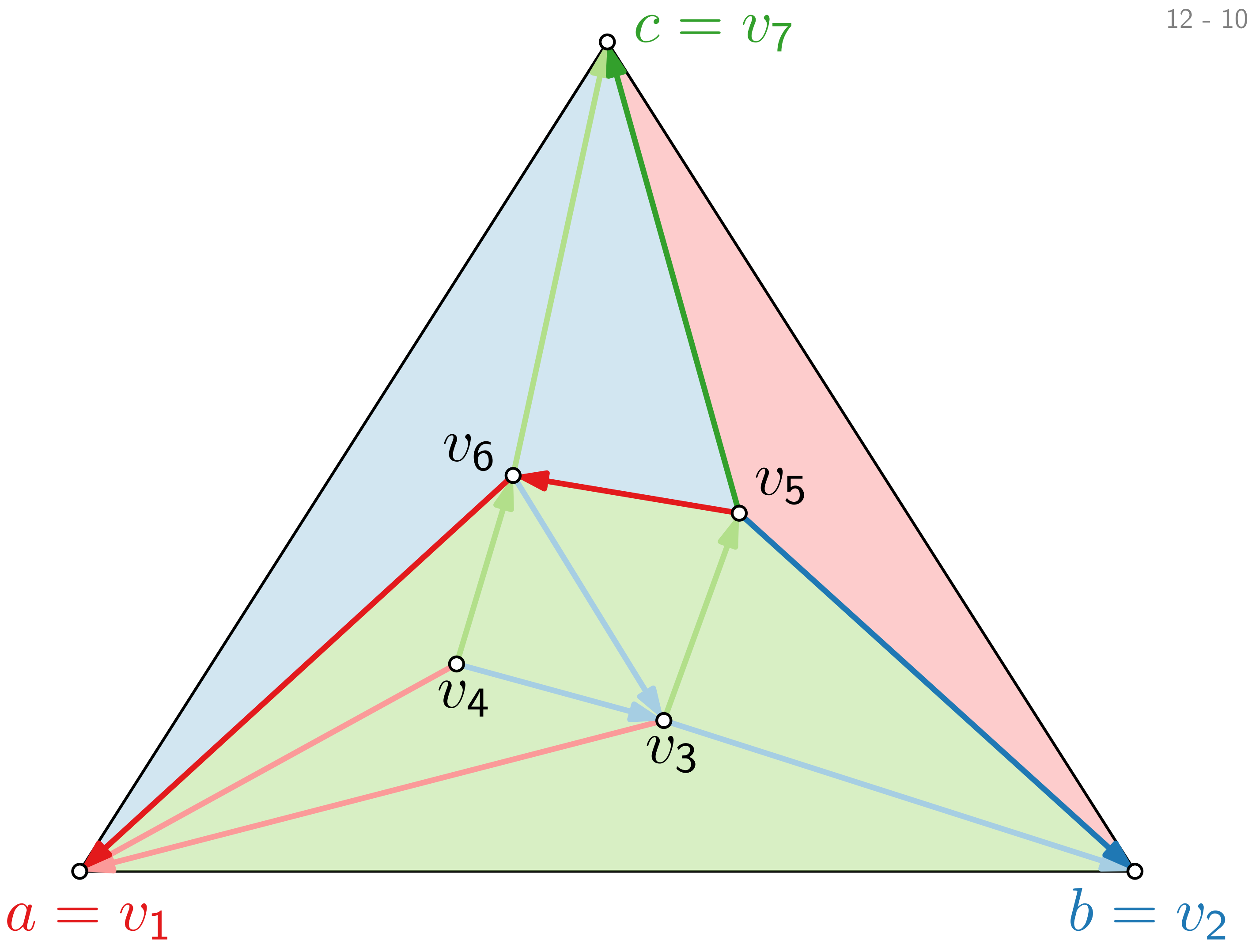

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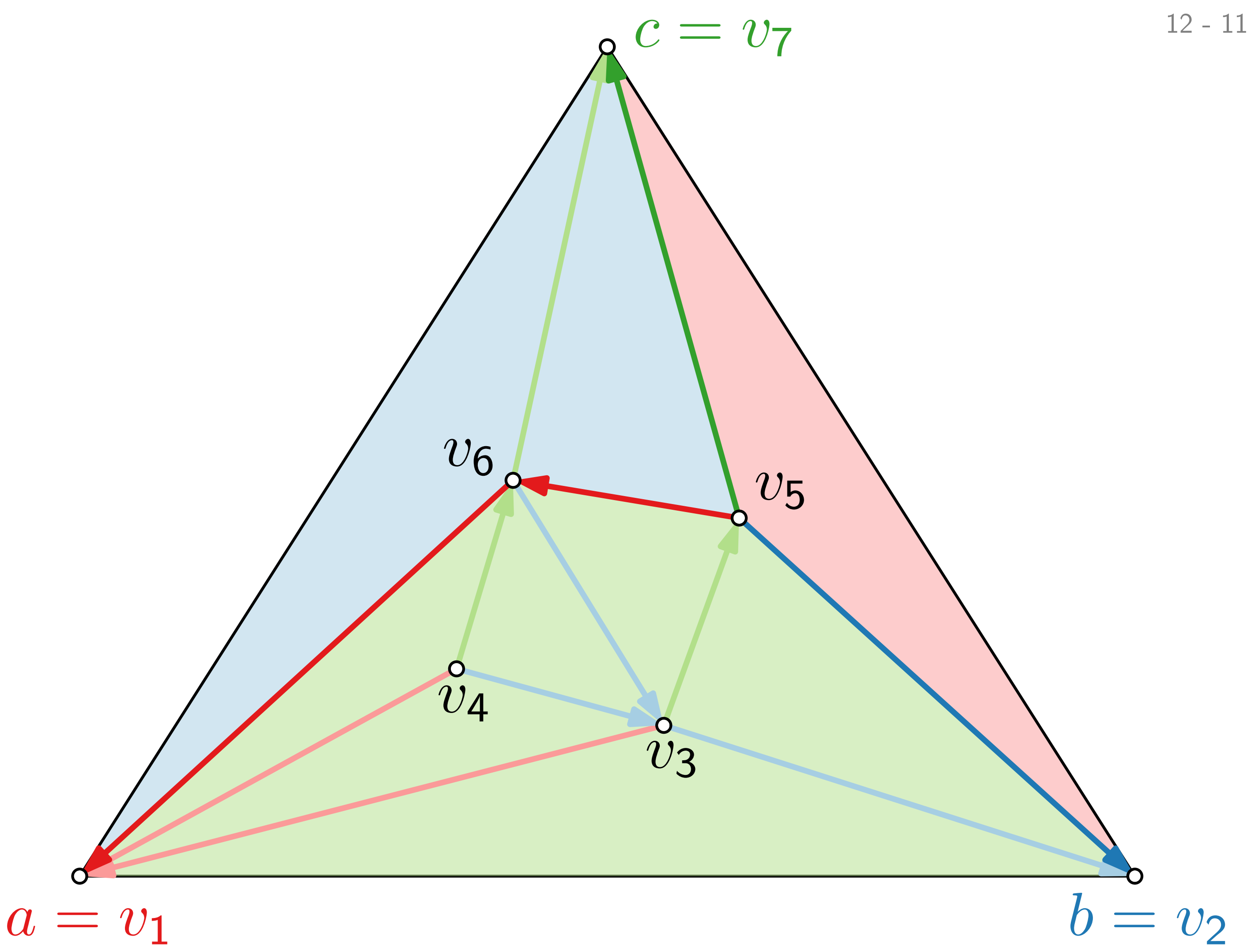

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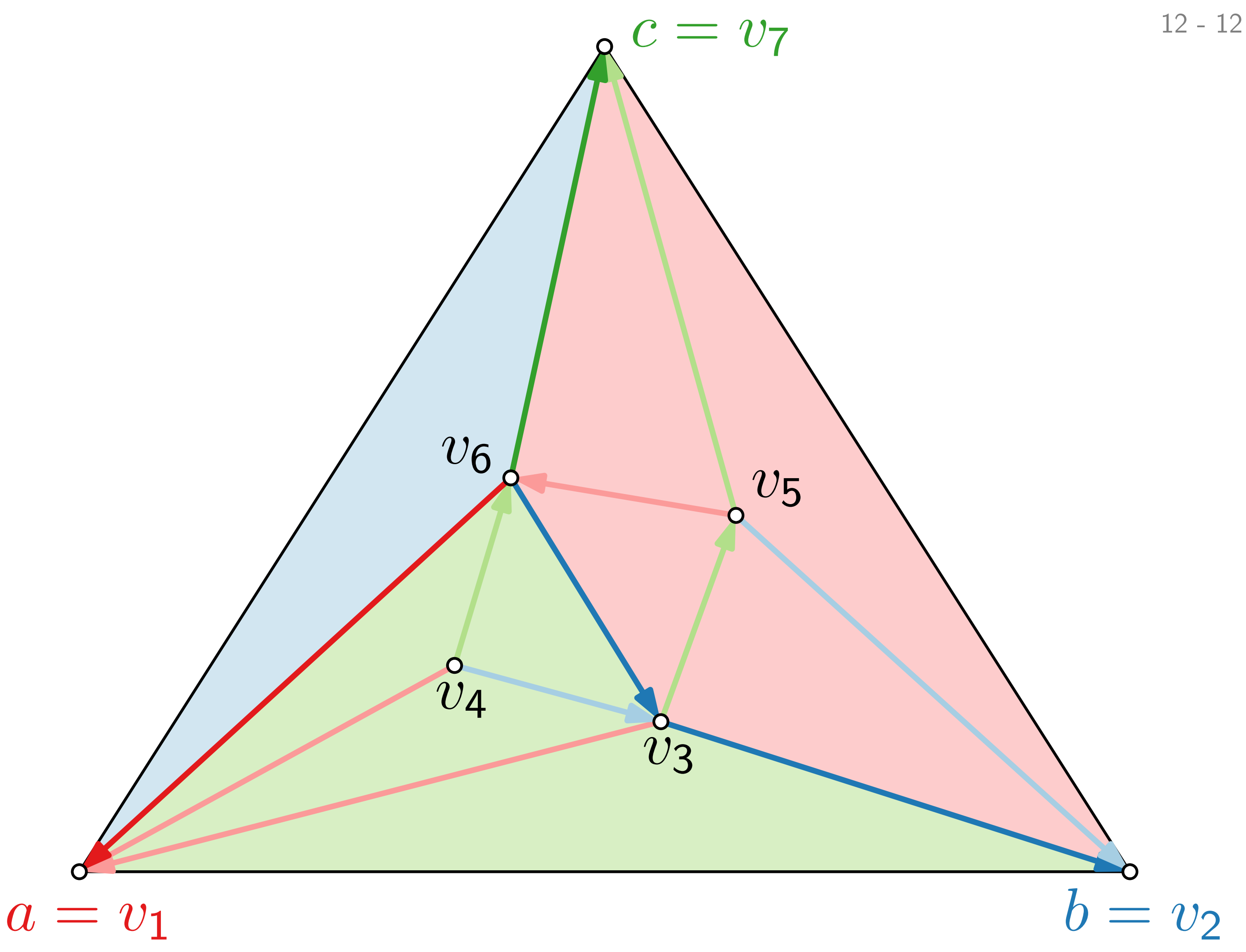

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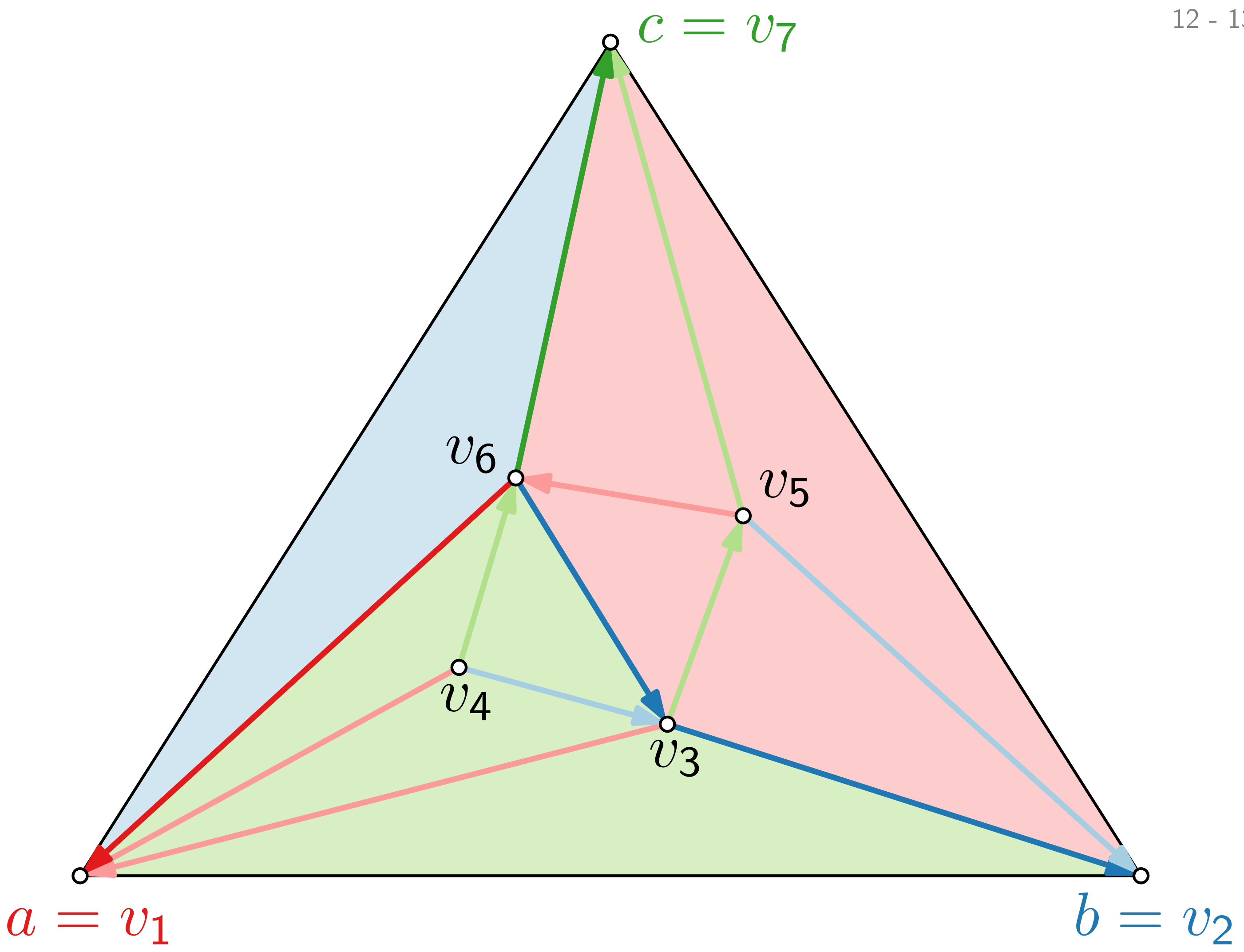

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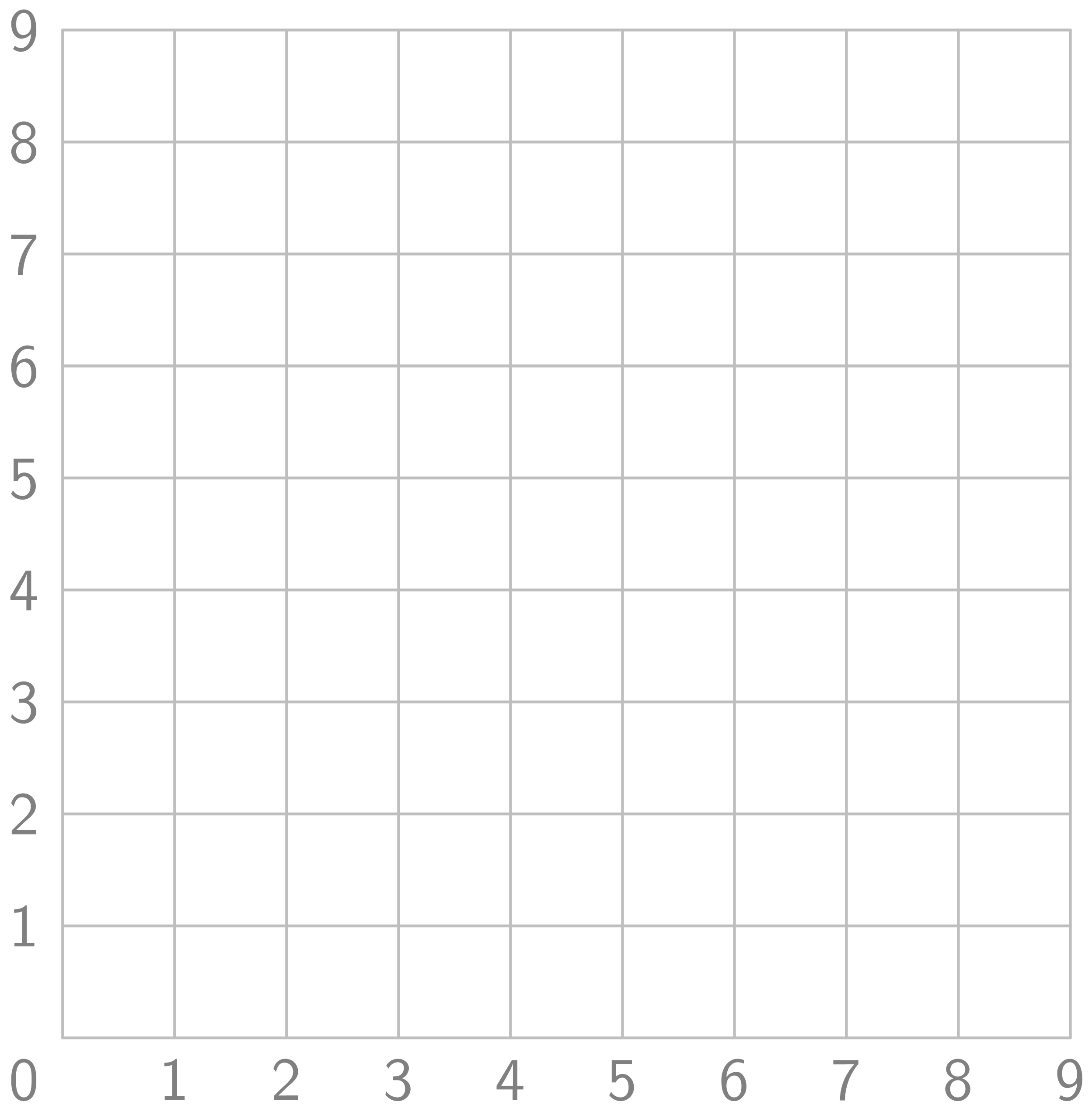

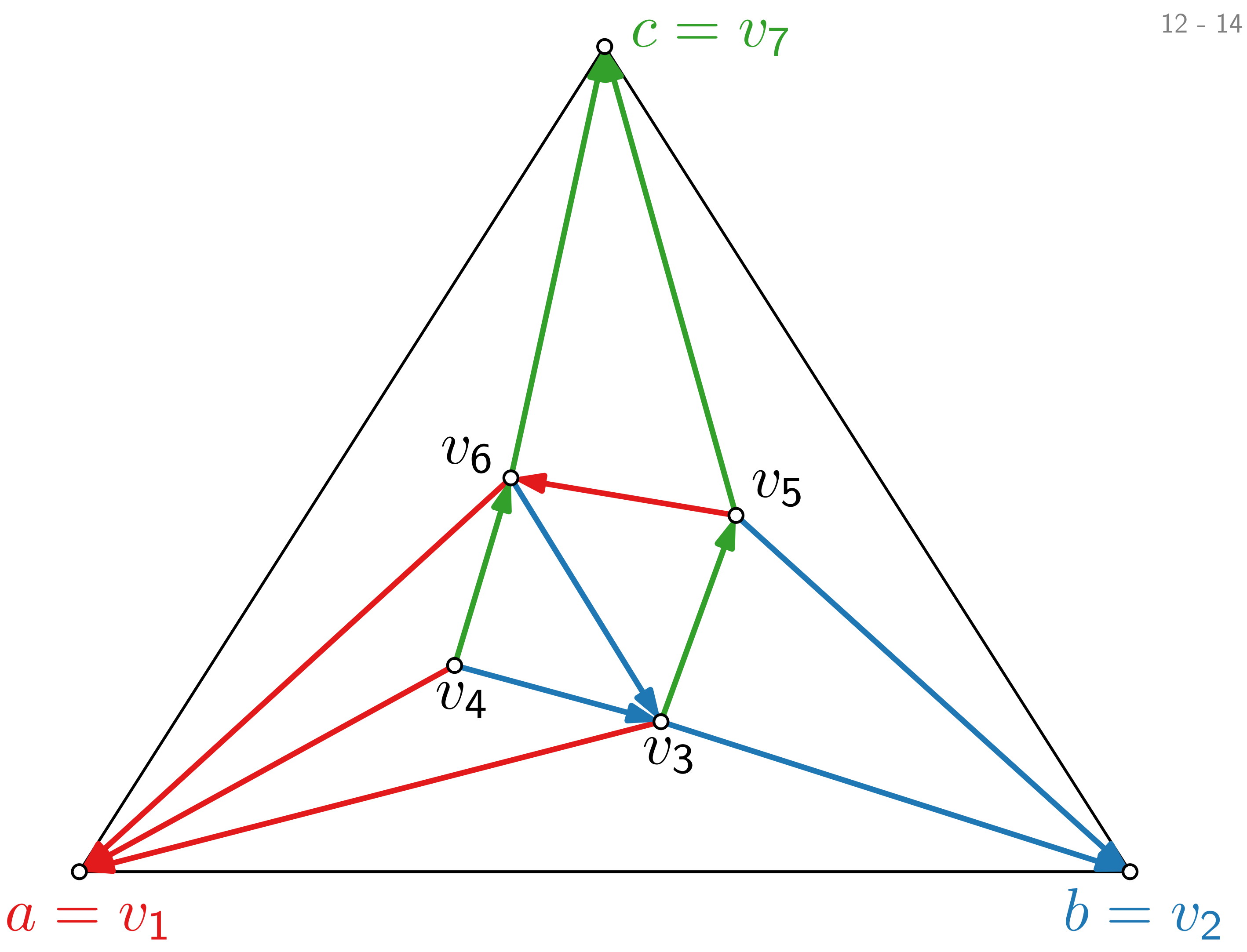

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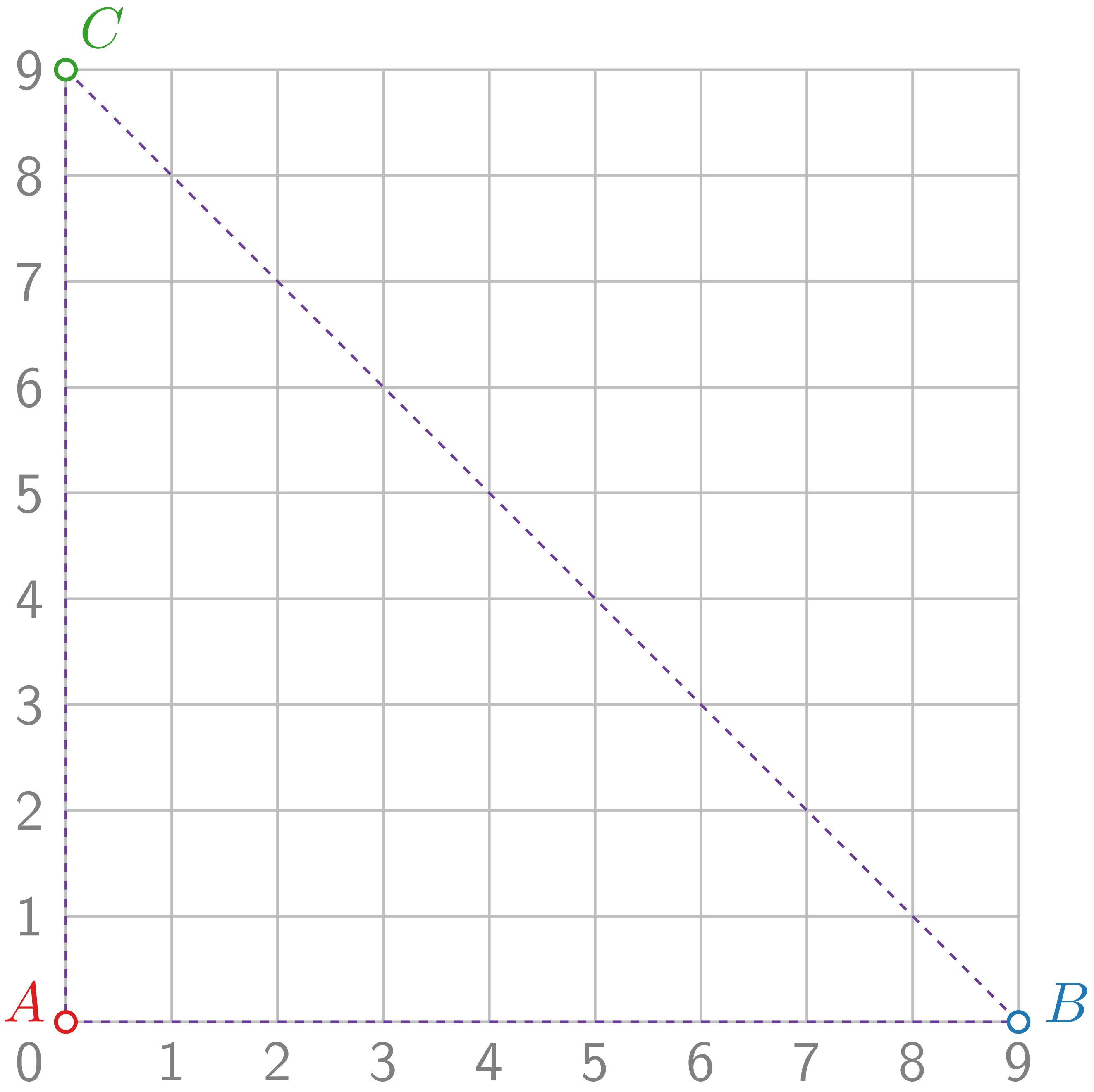

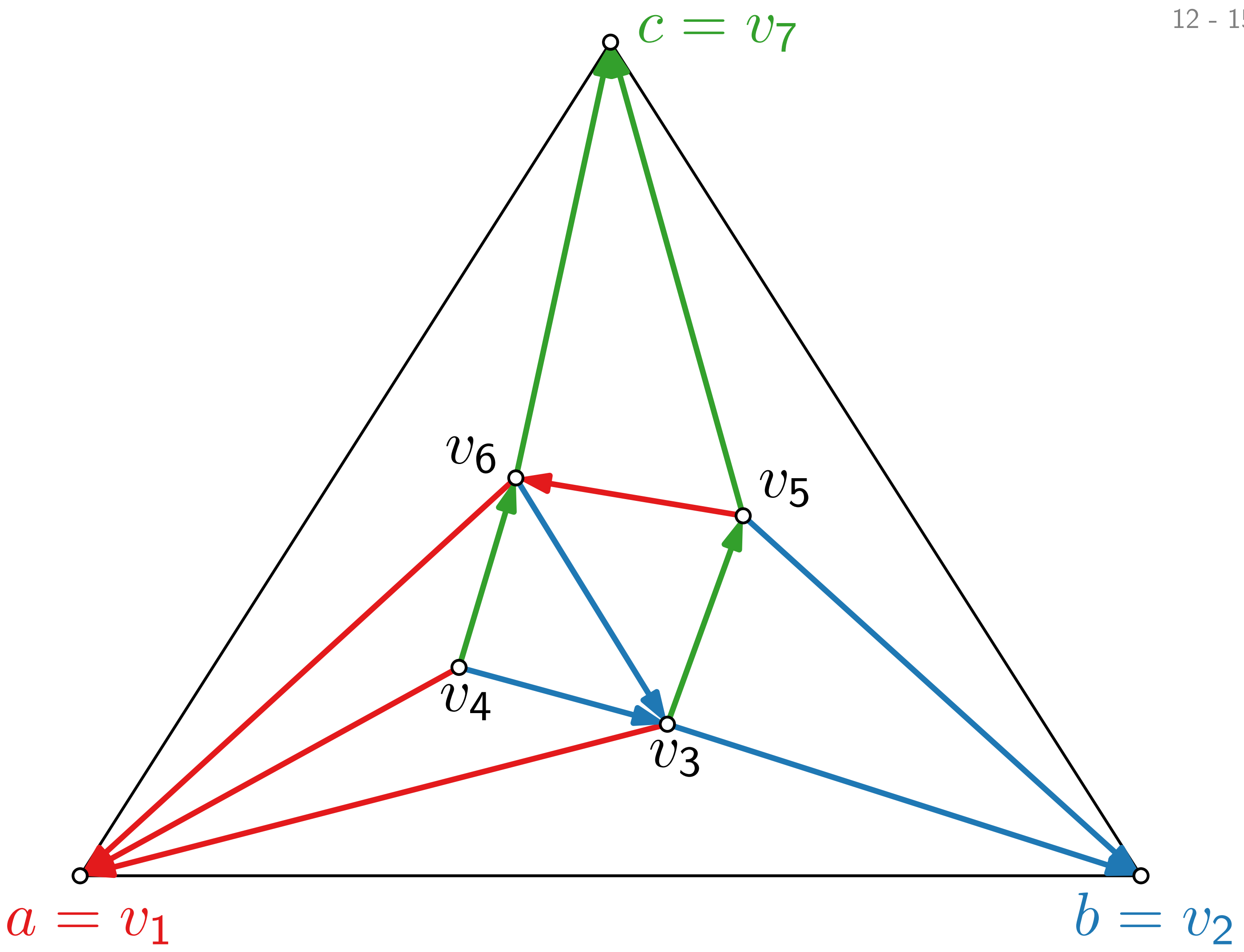

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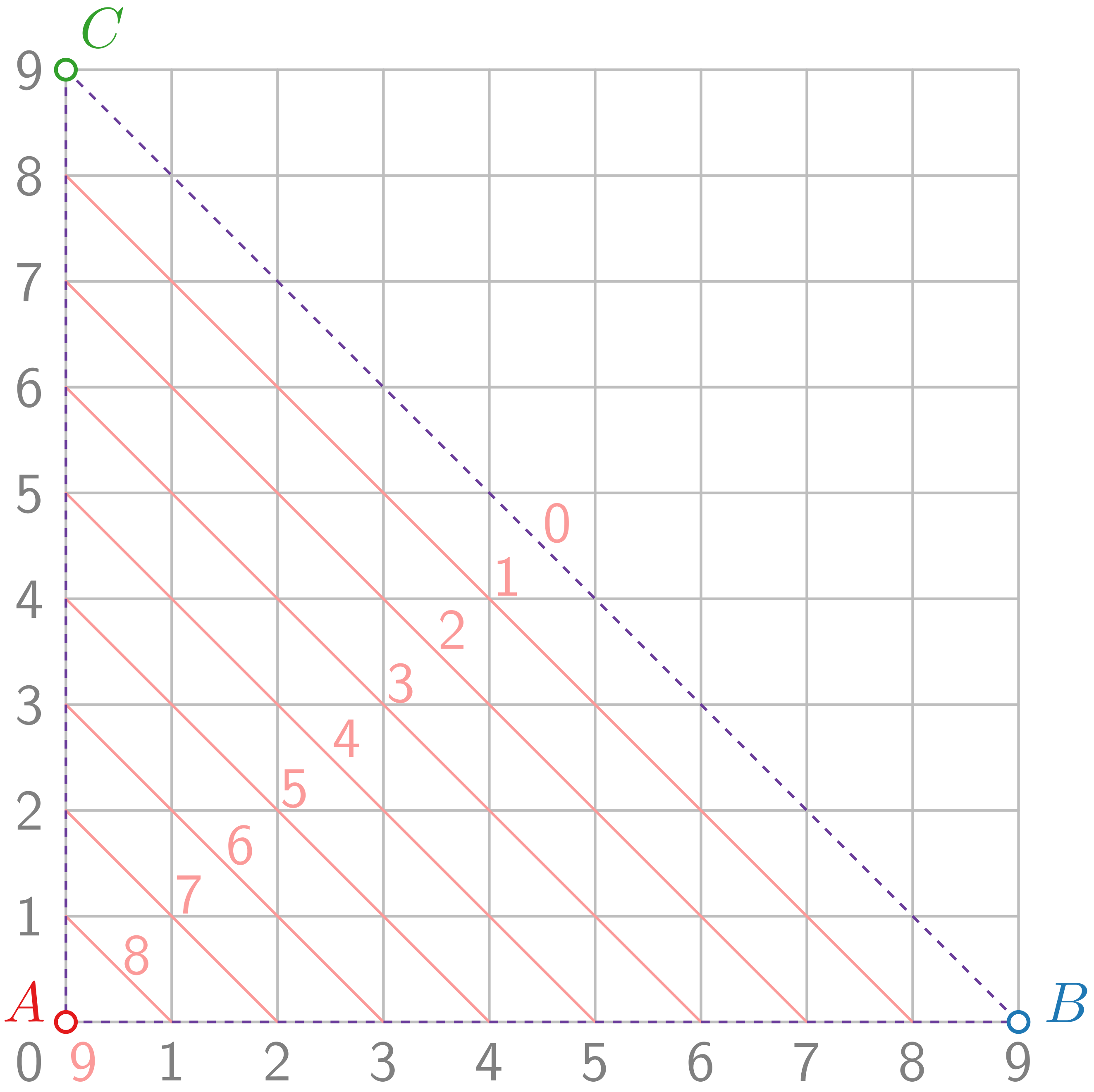

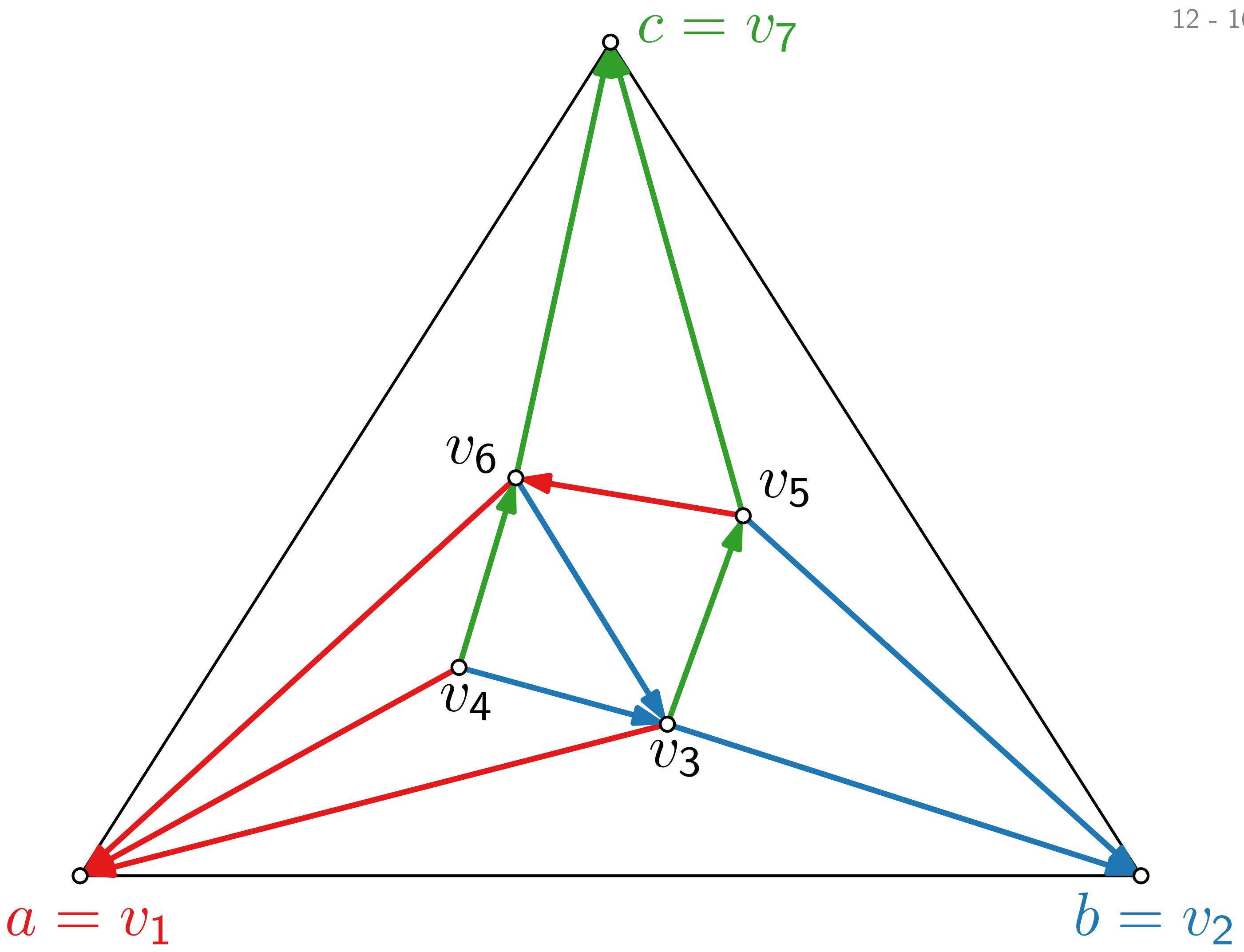

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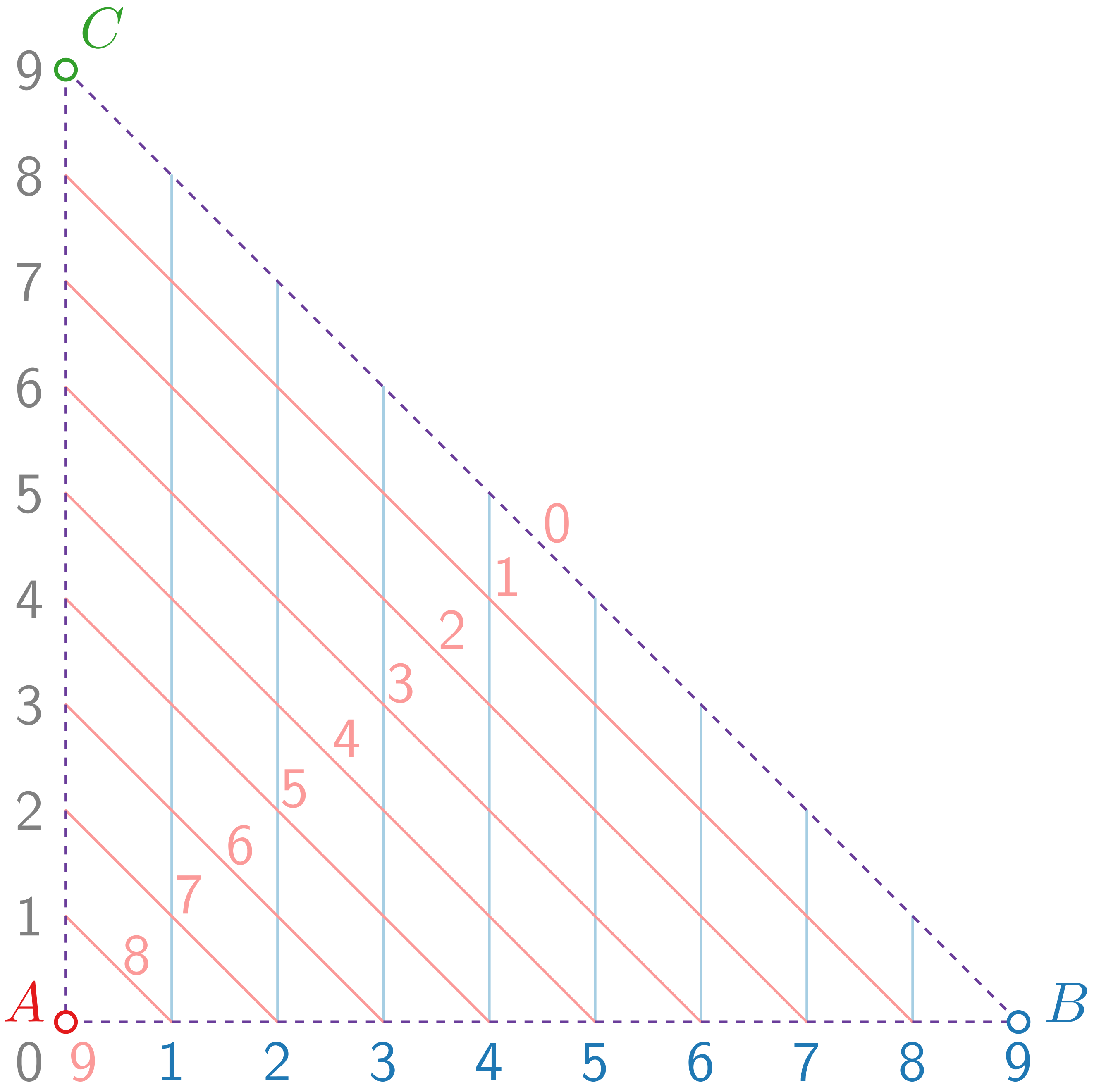

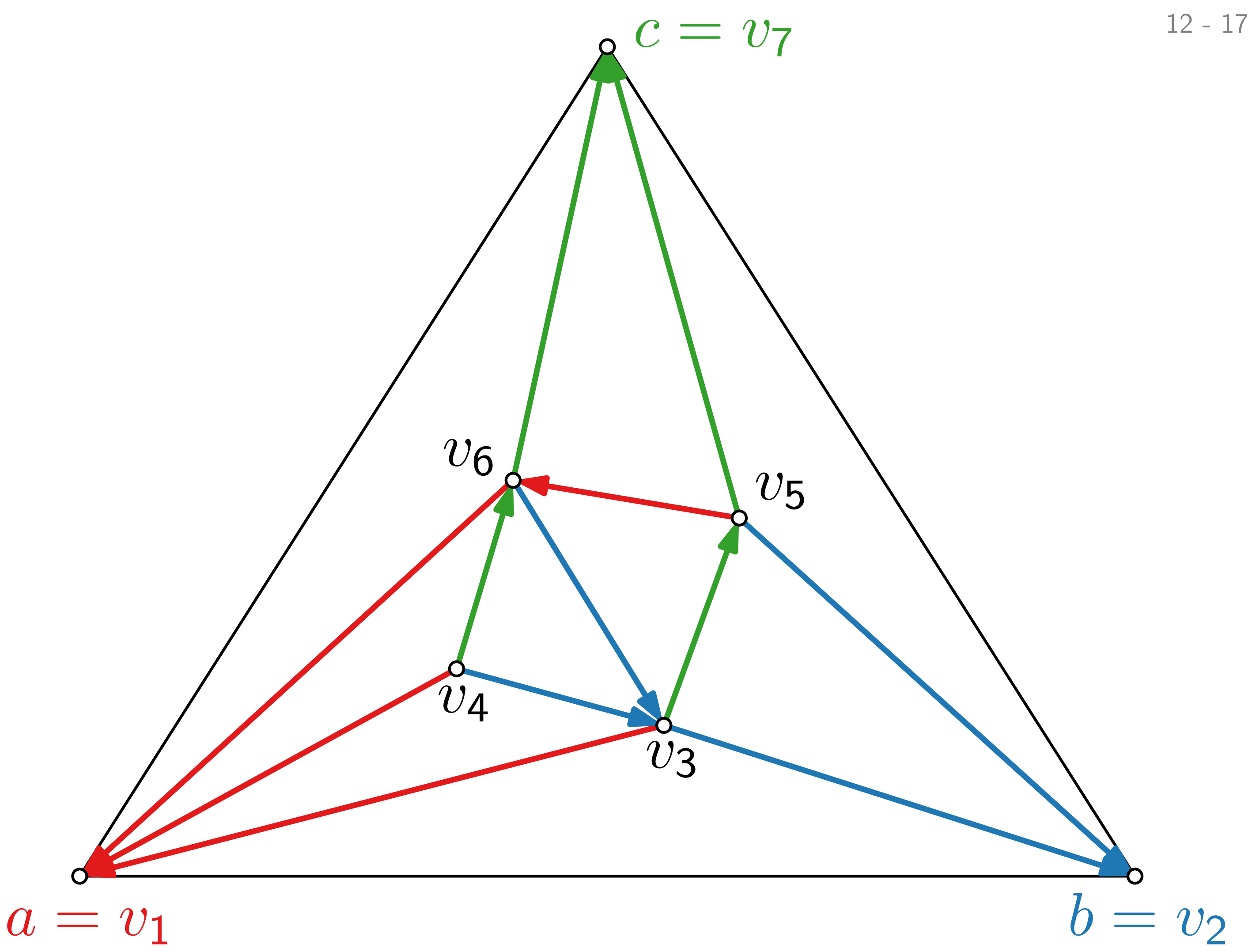

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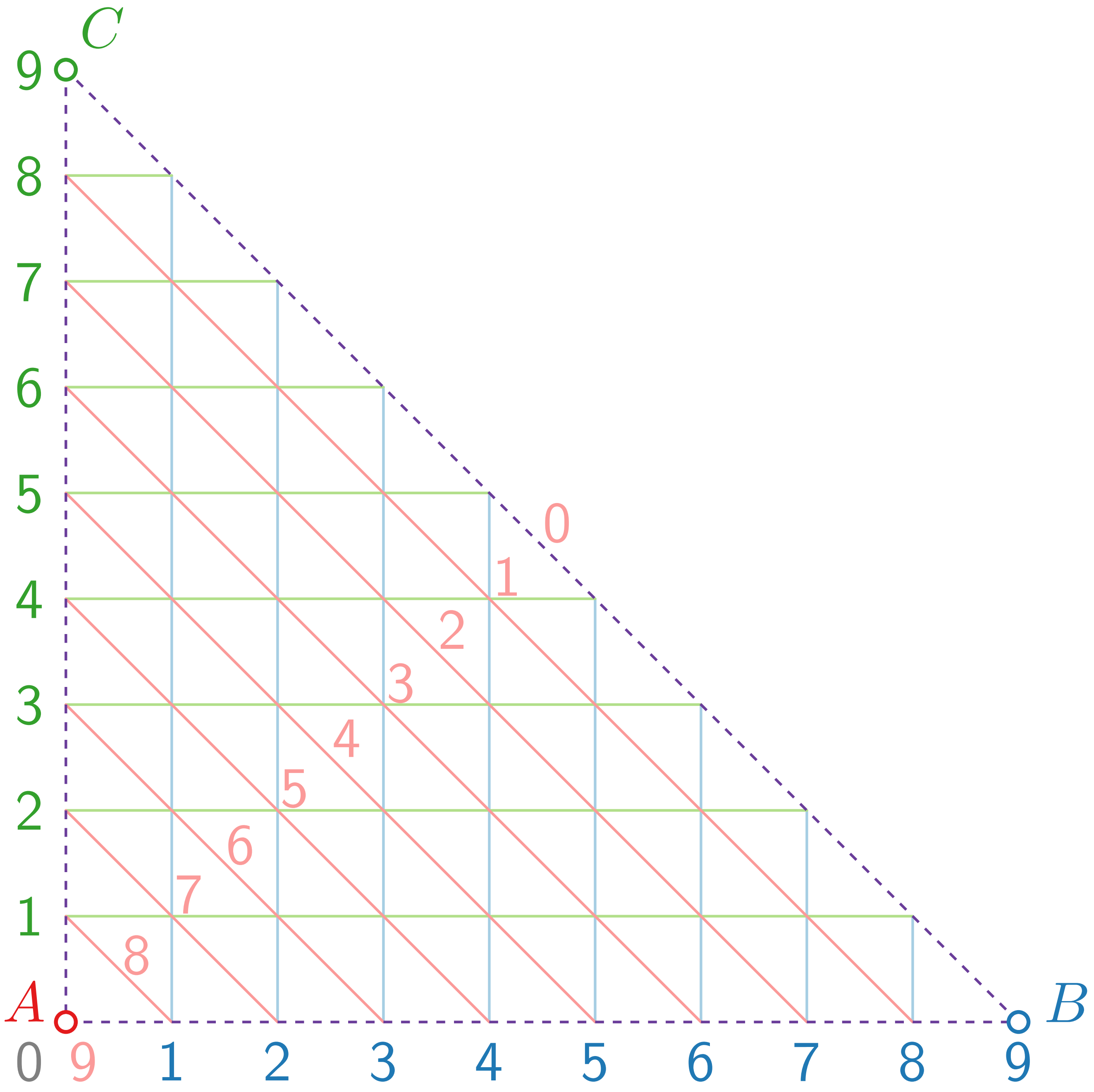

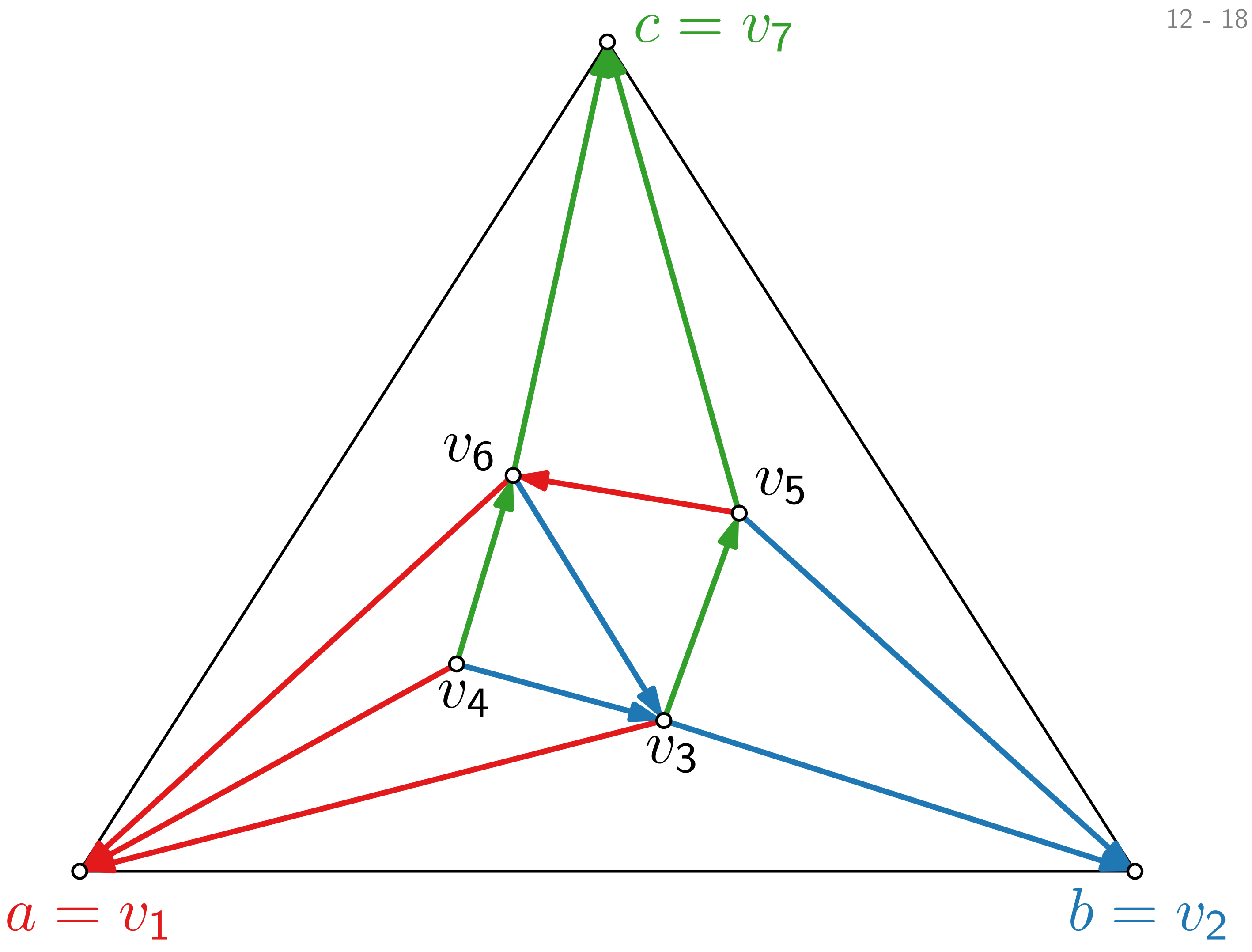

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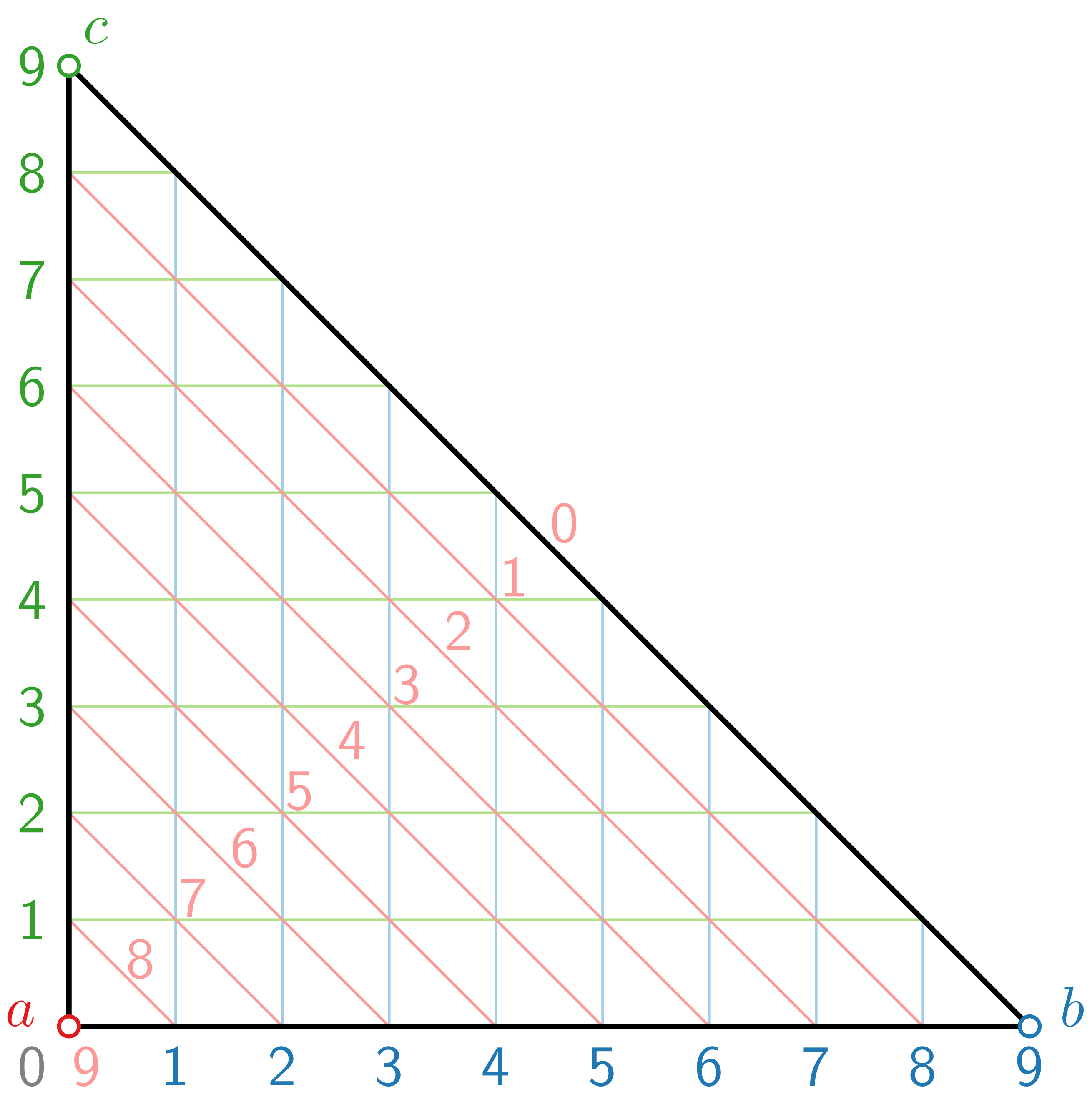

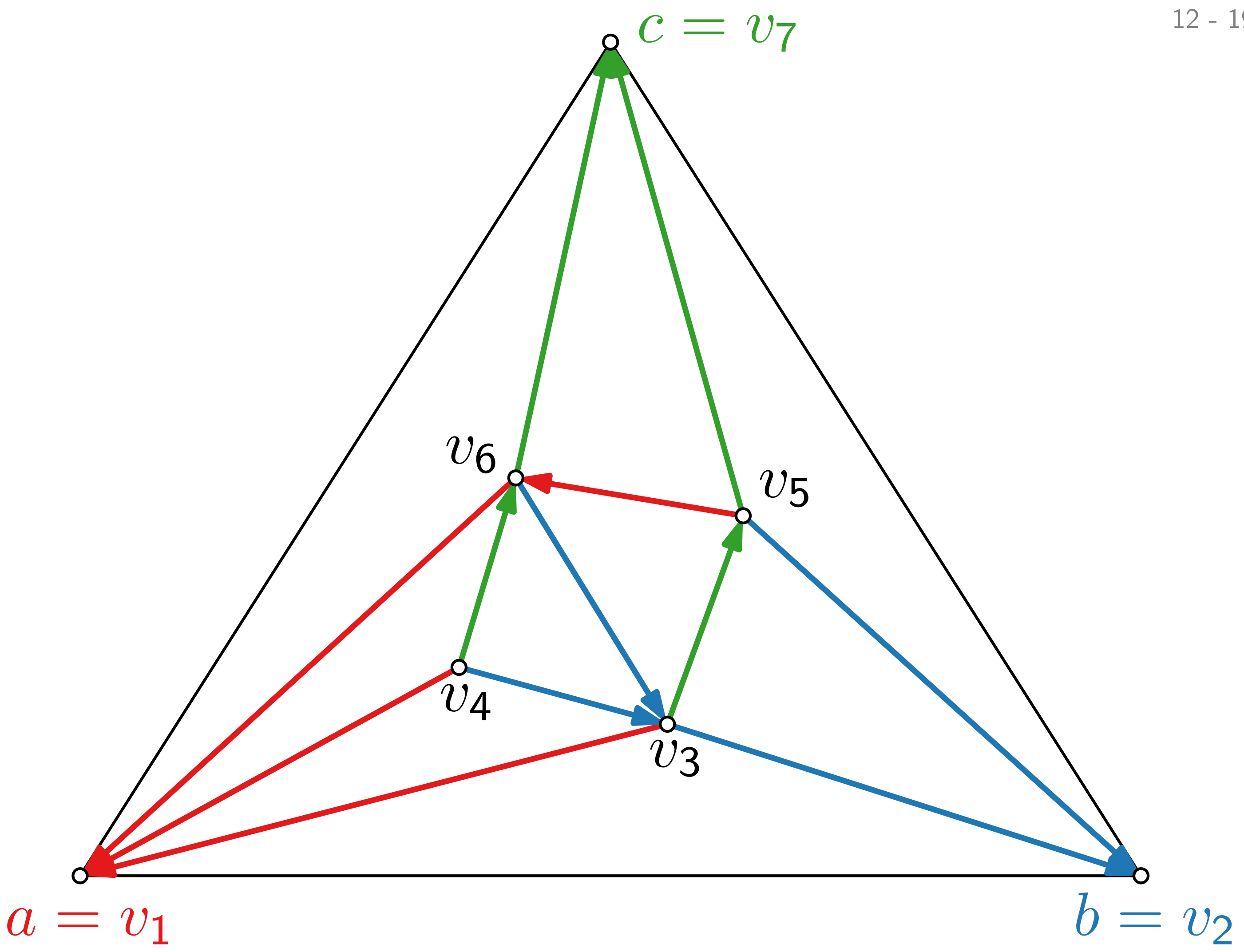

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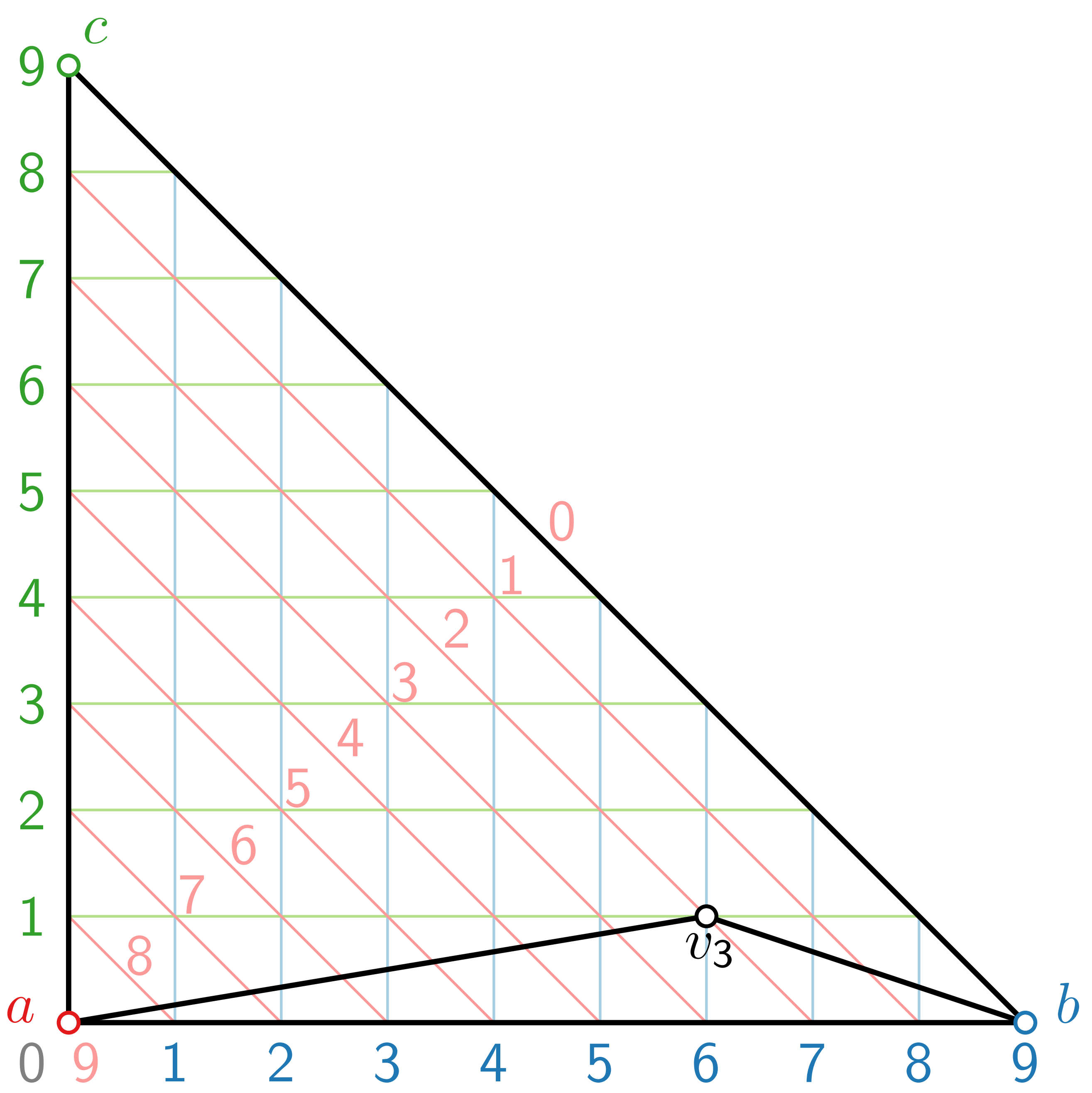

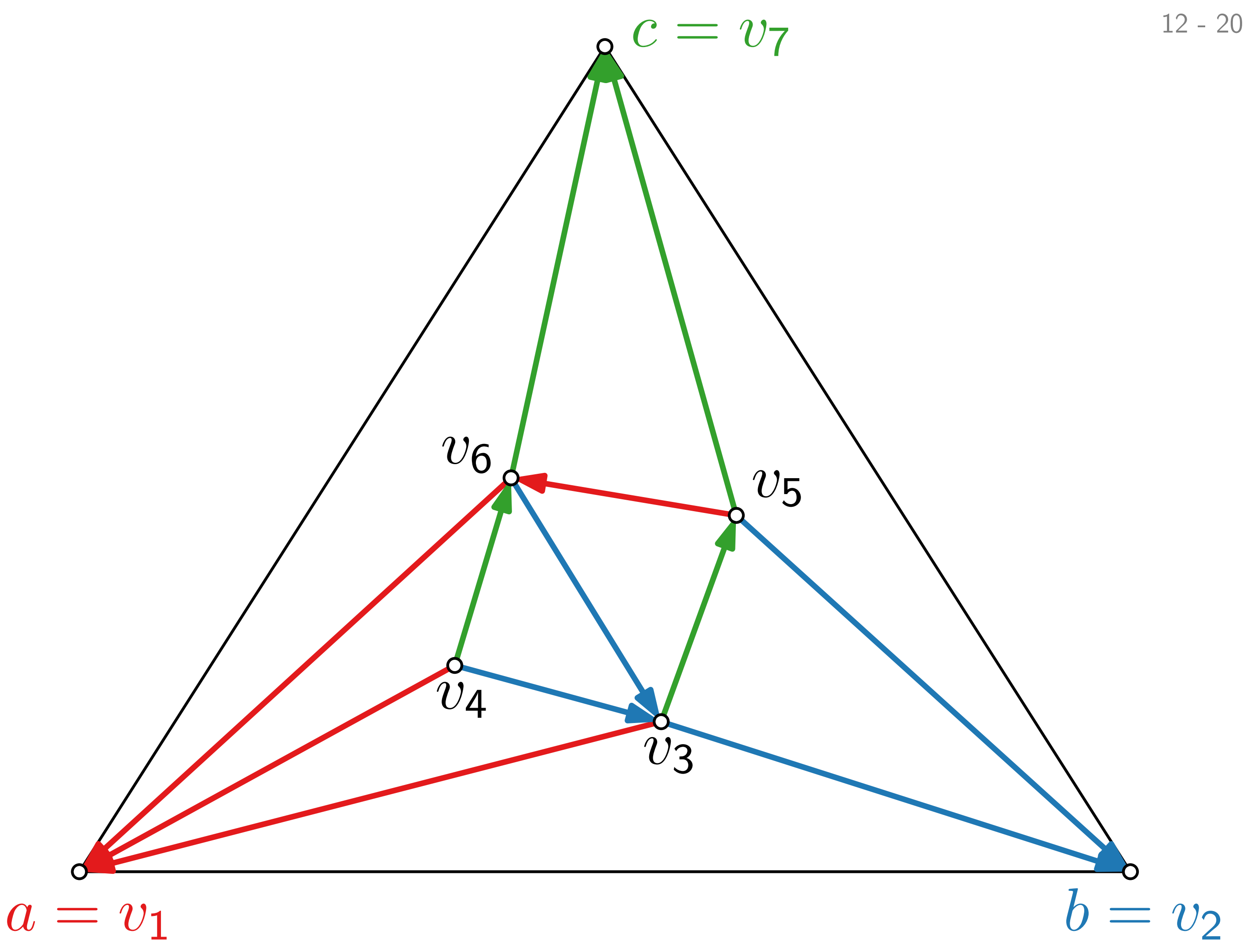

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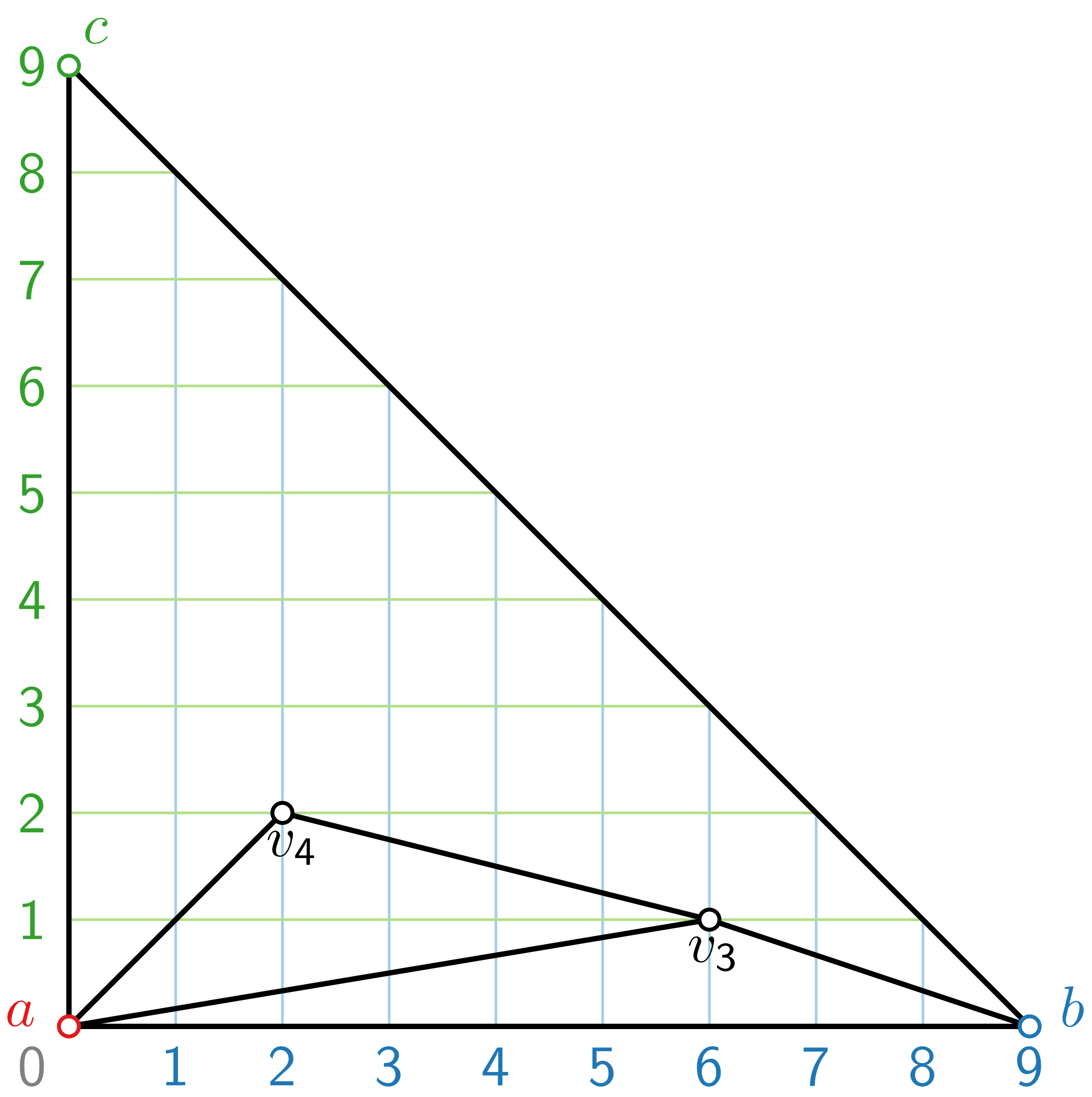

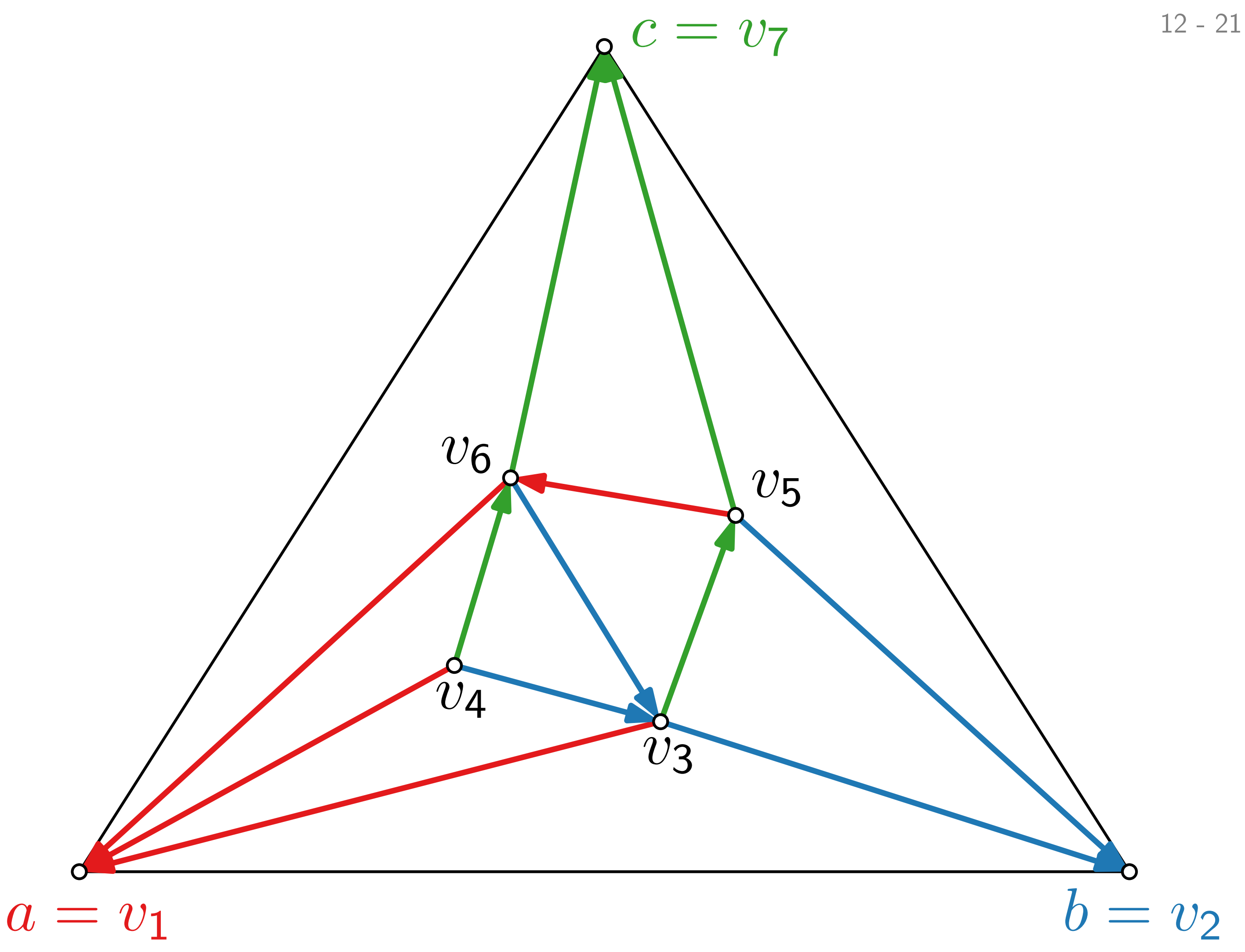

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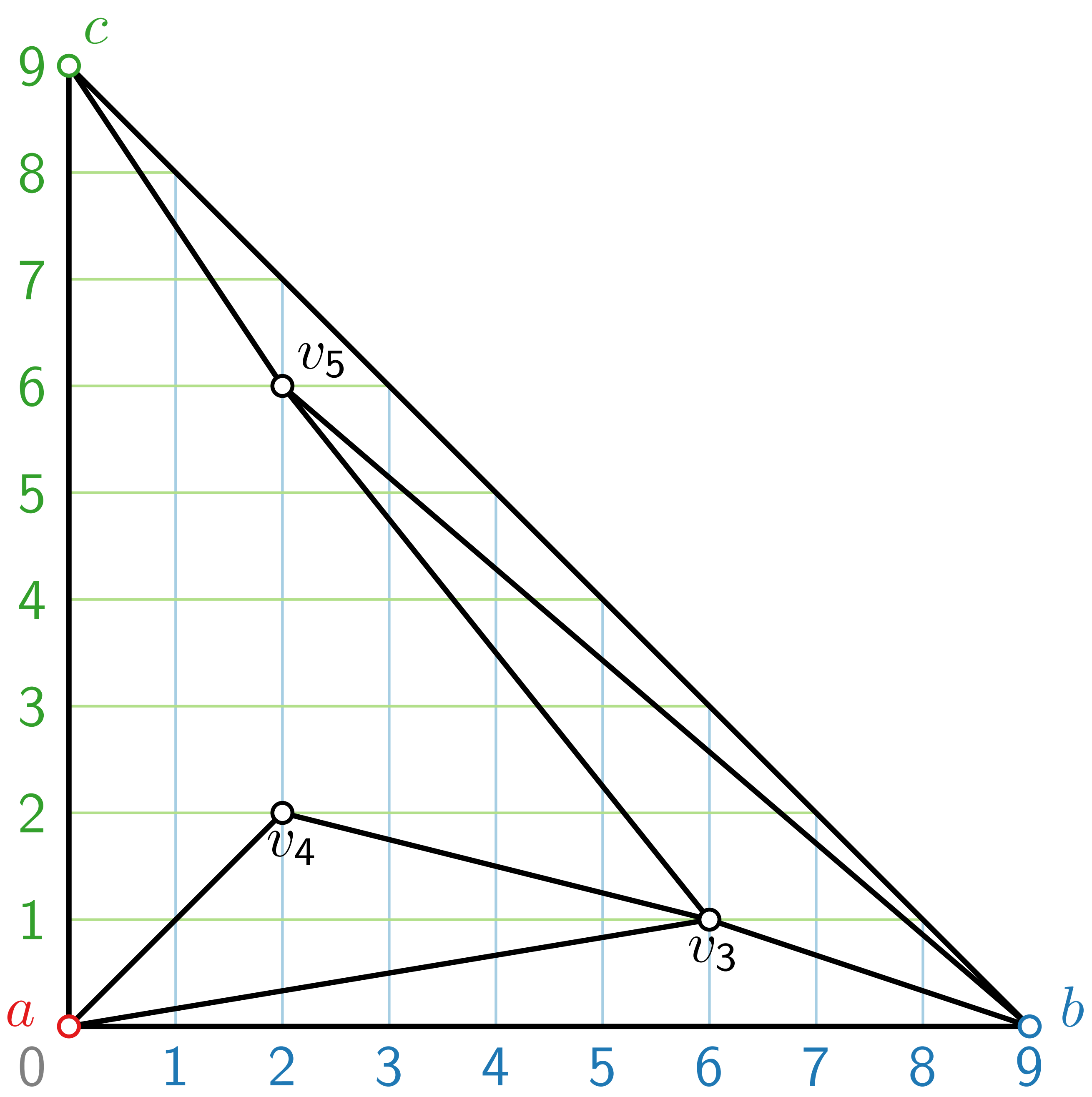

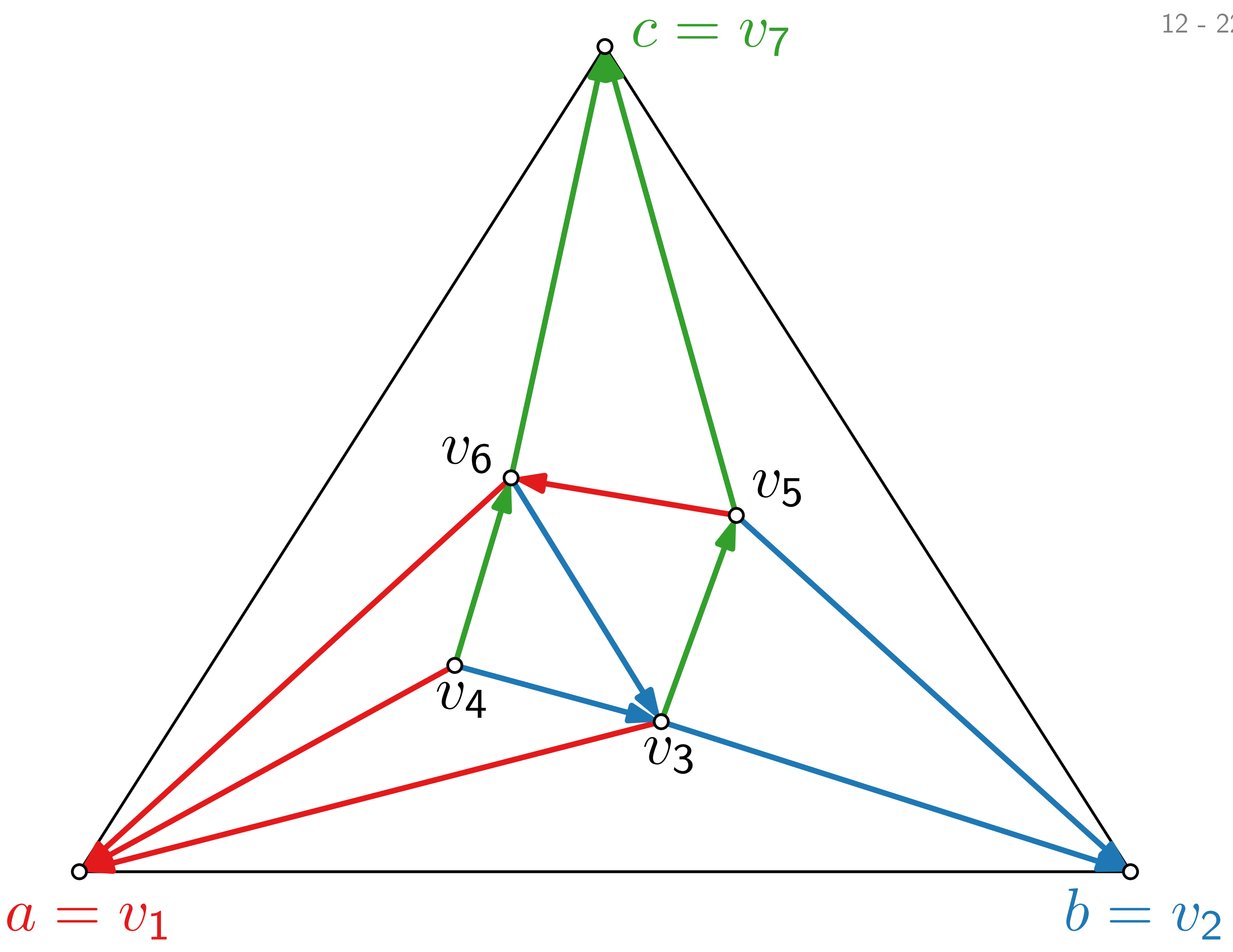

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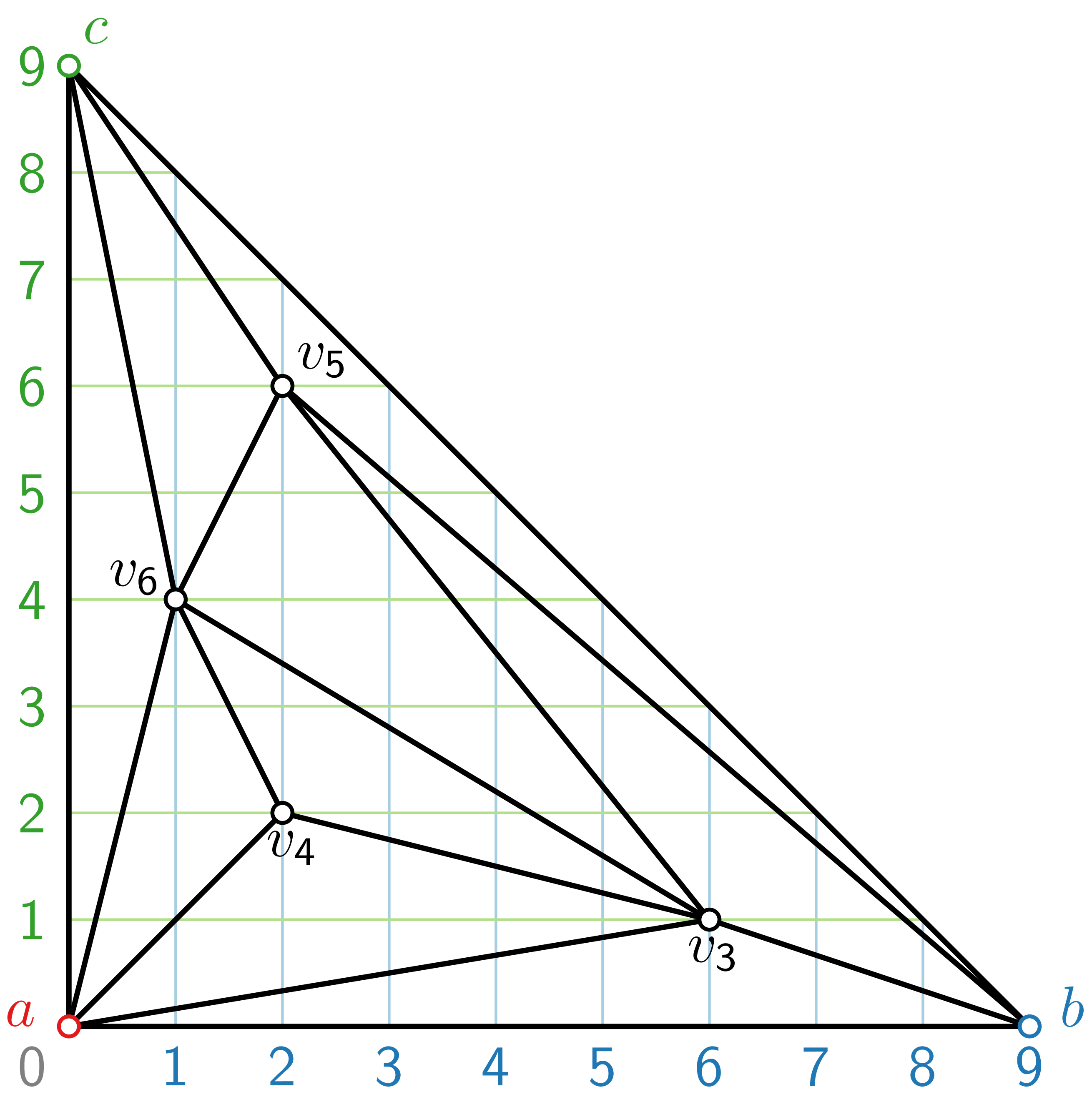

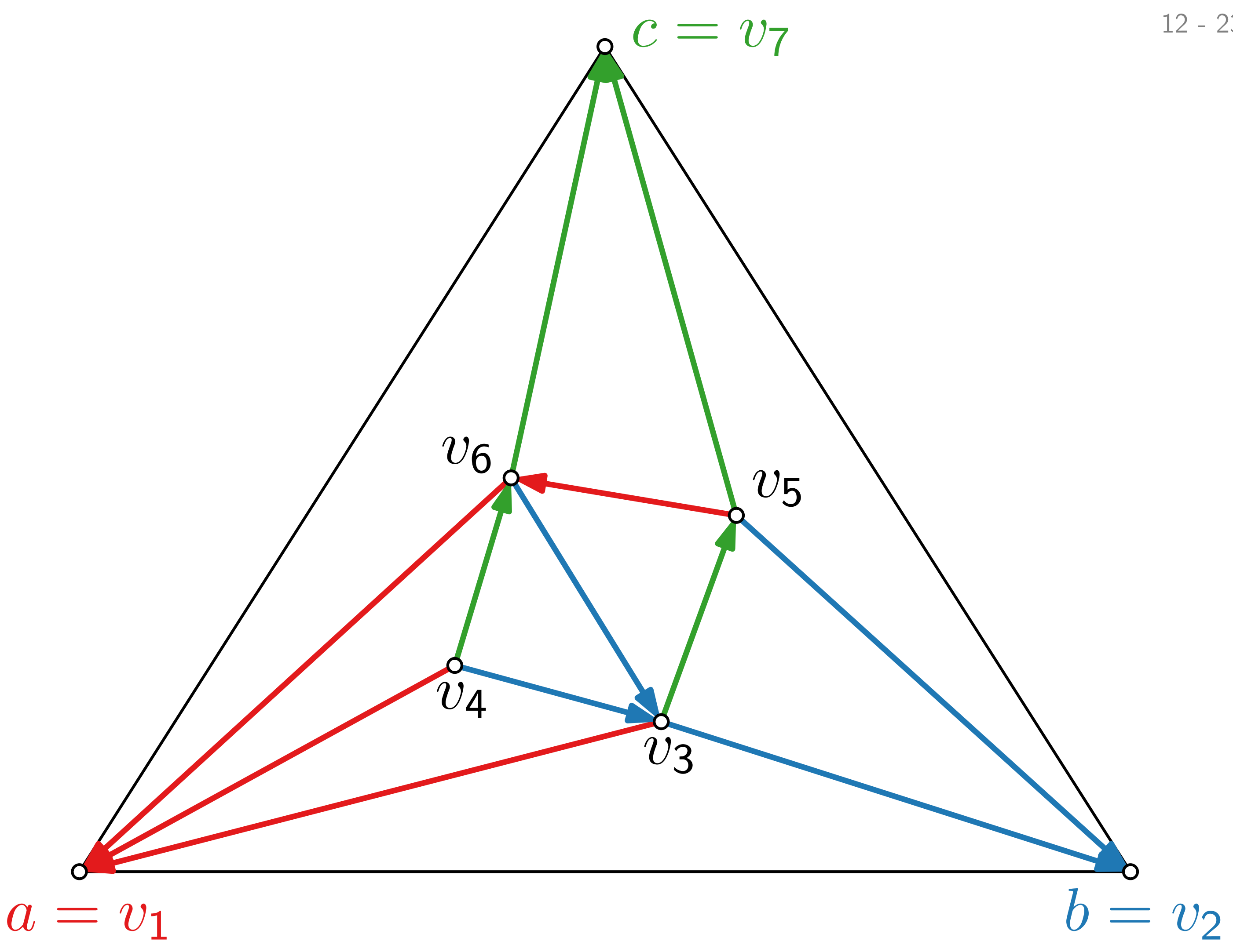

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A **weak barycentric representation** of a graph G is an assignment of barycentric coordinates to $V(G)$:

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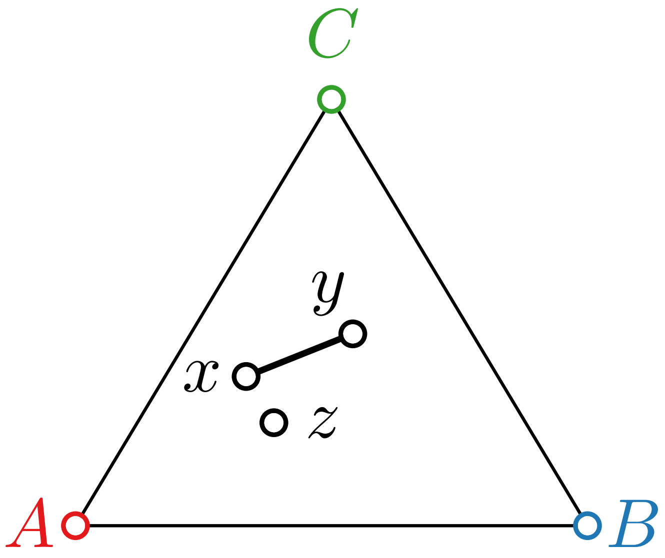

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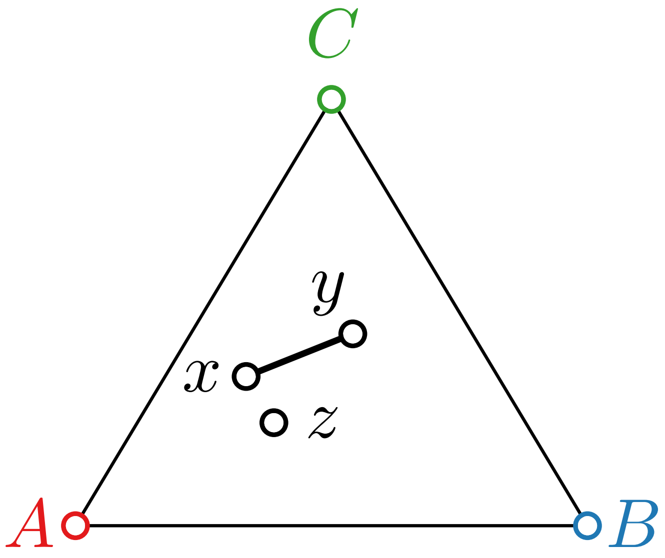

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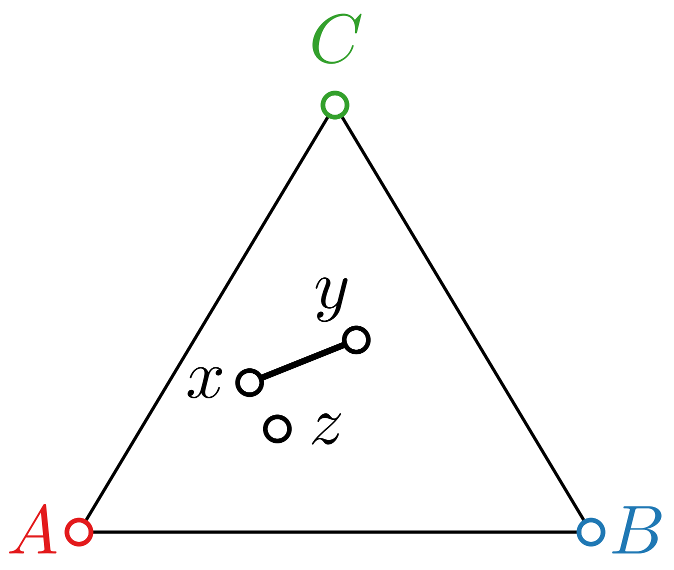

i.e., either $y_k < z_k$ or $y_k = z_k$ and $y_{k+1} < z_{k+1}$

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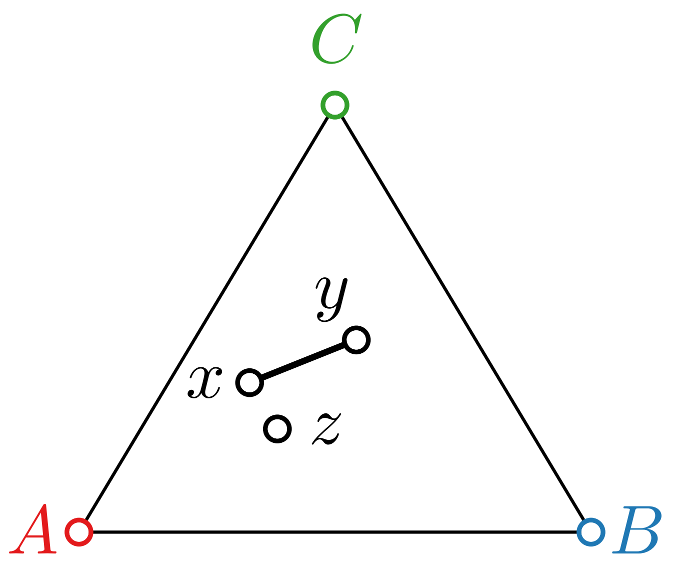

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A **weak barycentric representation** of a graph G is an assignment of barycentric coordinates to $V(G)$:

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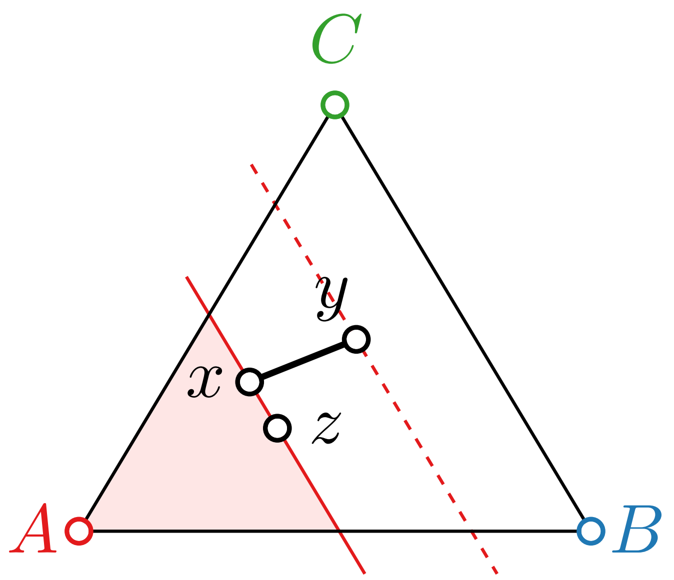

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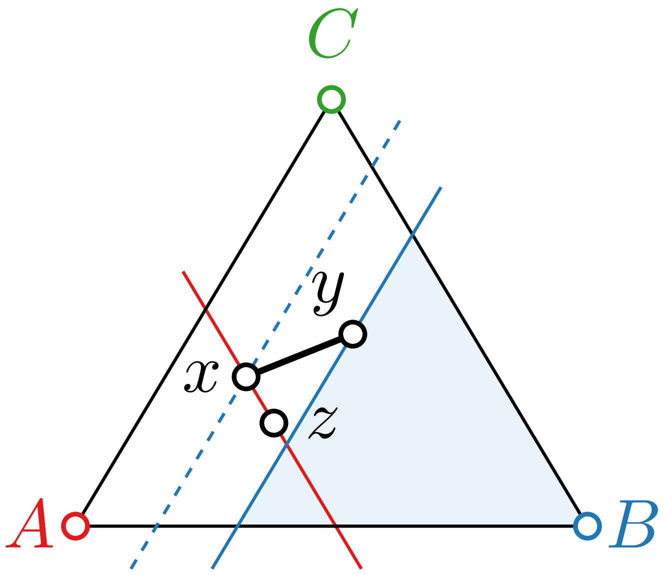

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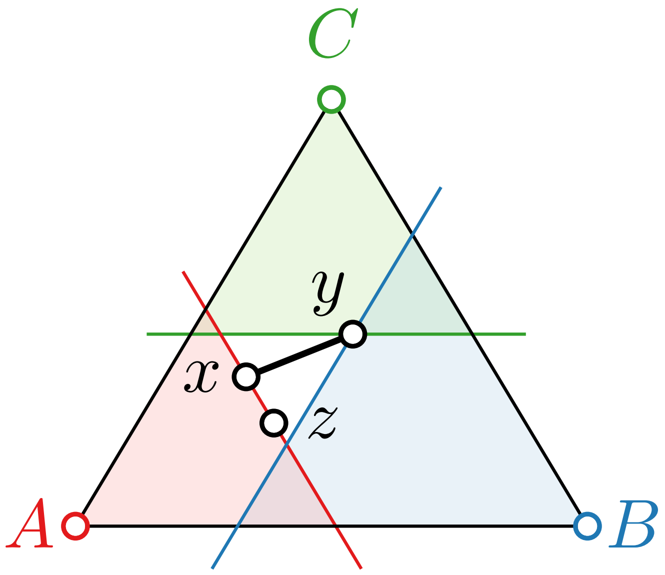

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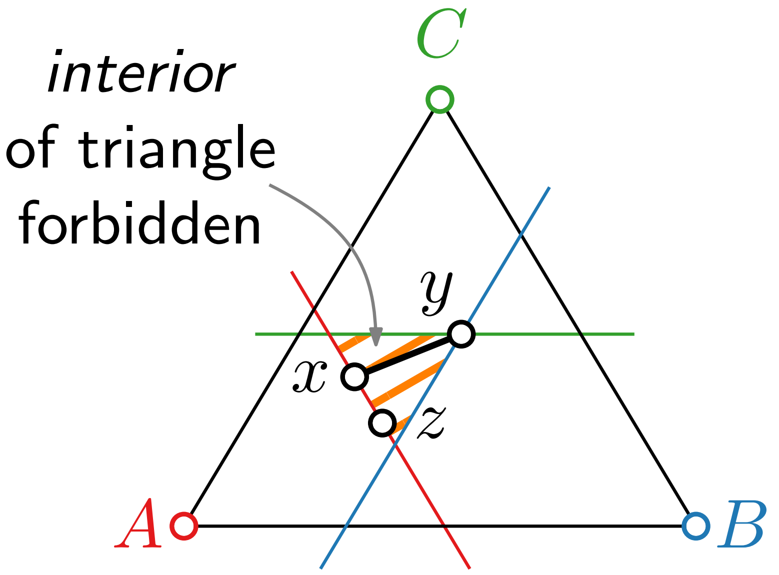

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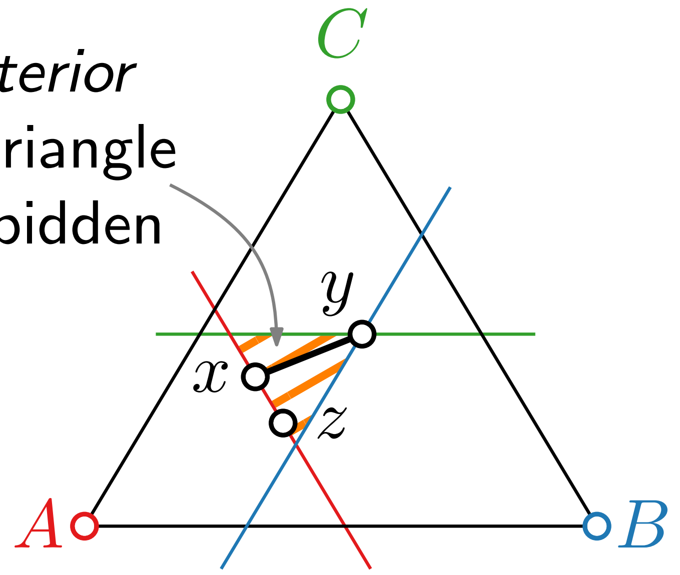

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$$v \mapsto v_1 A + v_2 B + v_3 C$$

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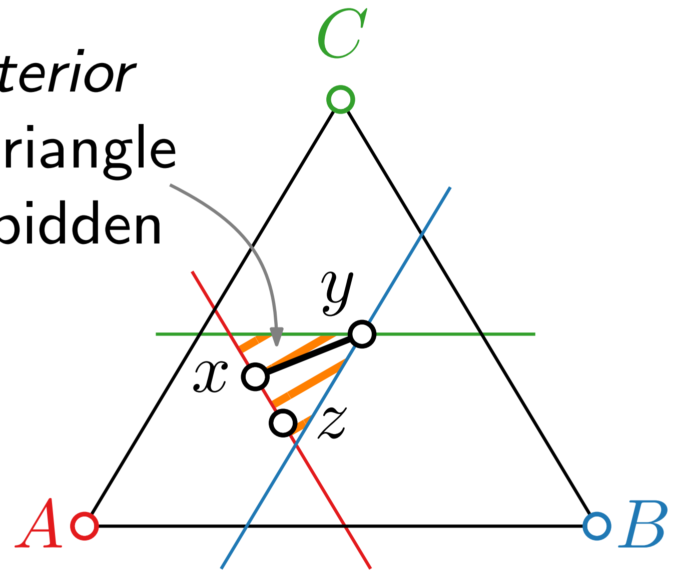

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Proof. $\to$ *Exercise!*

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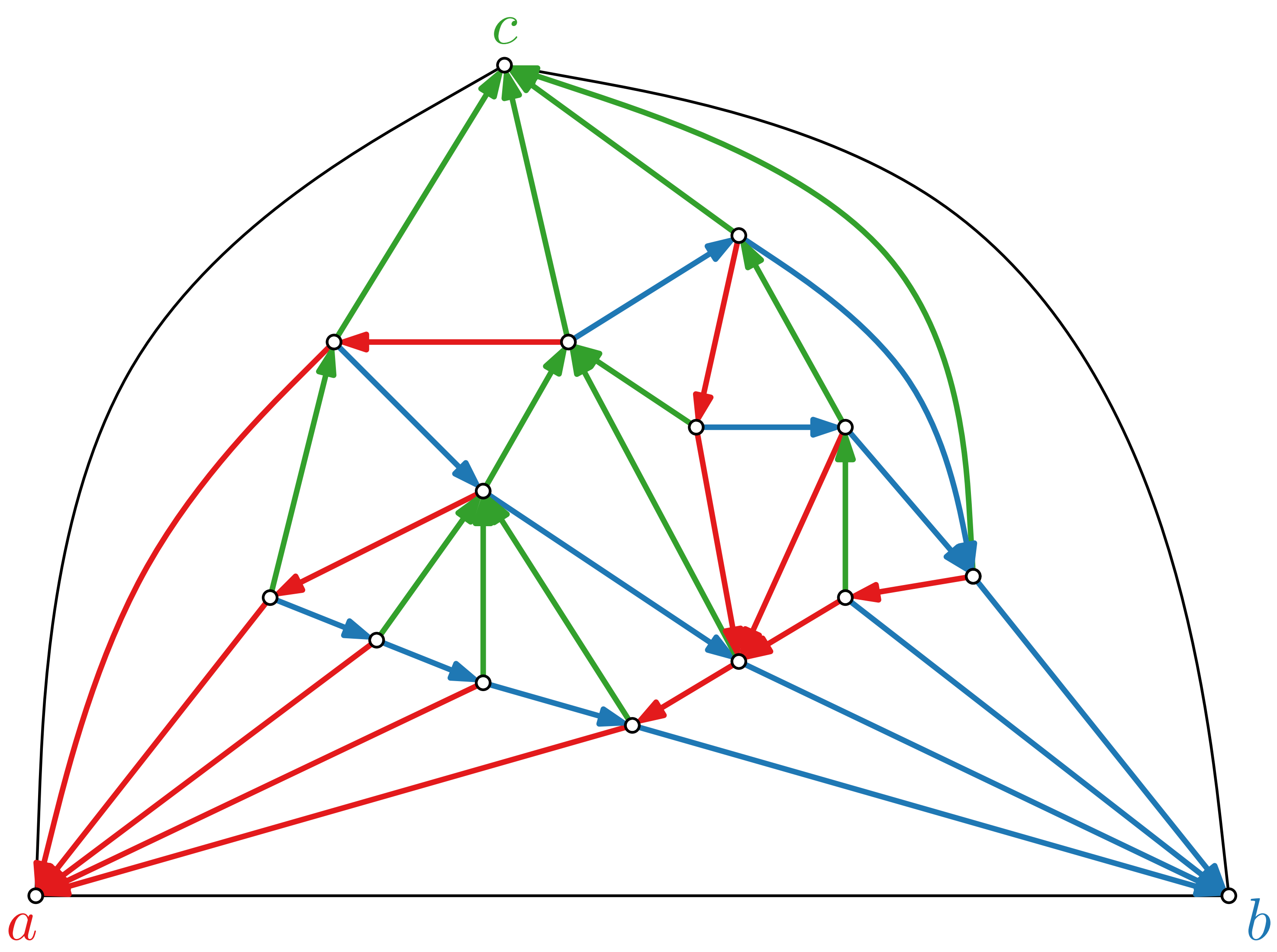

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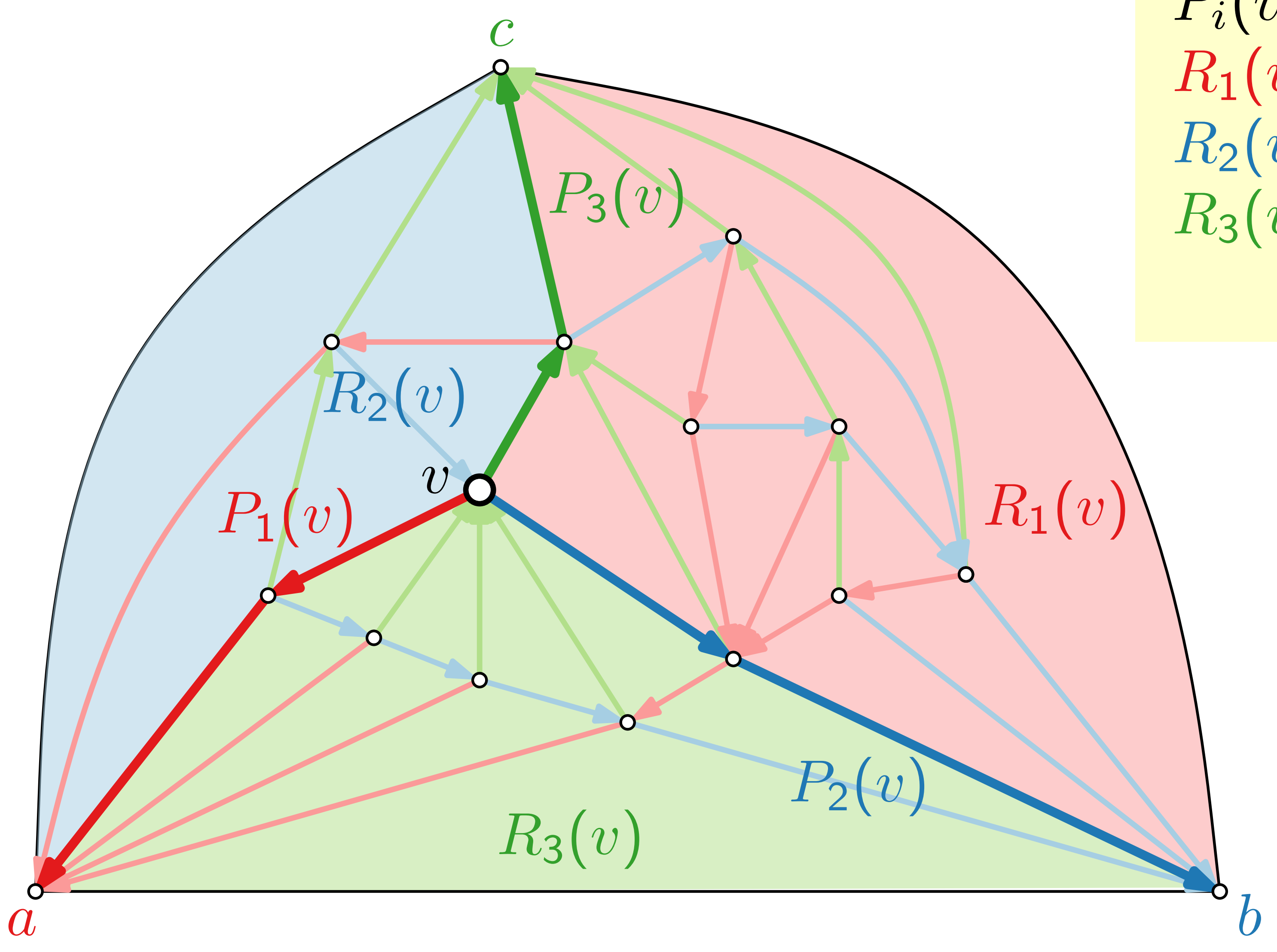

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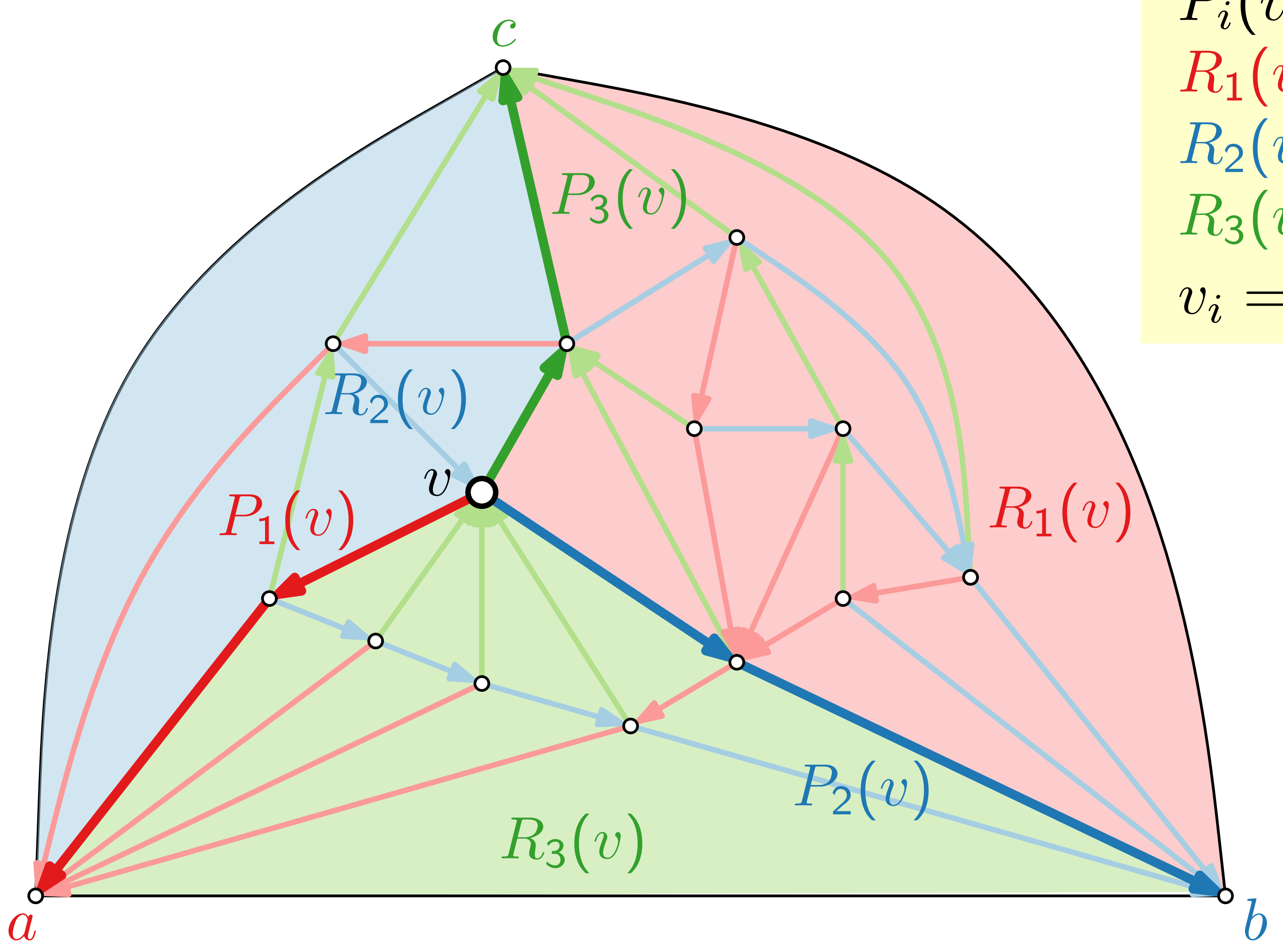

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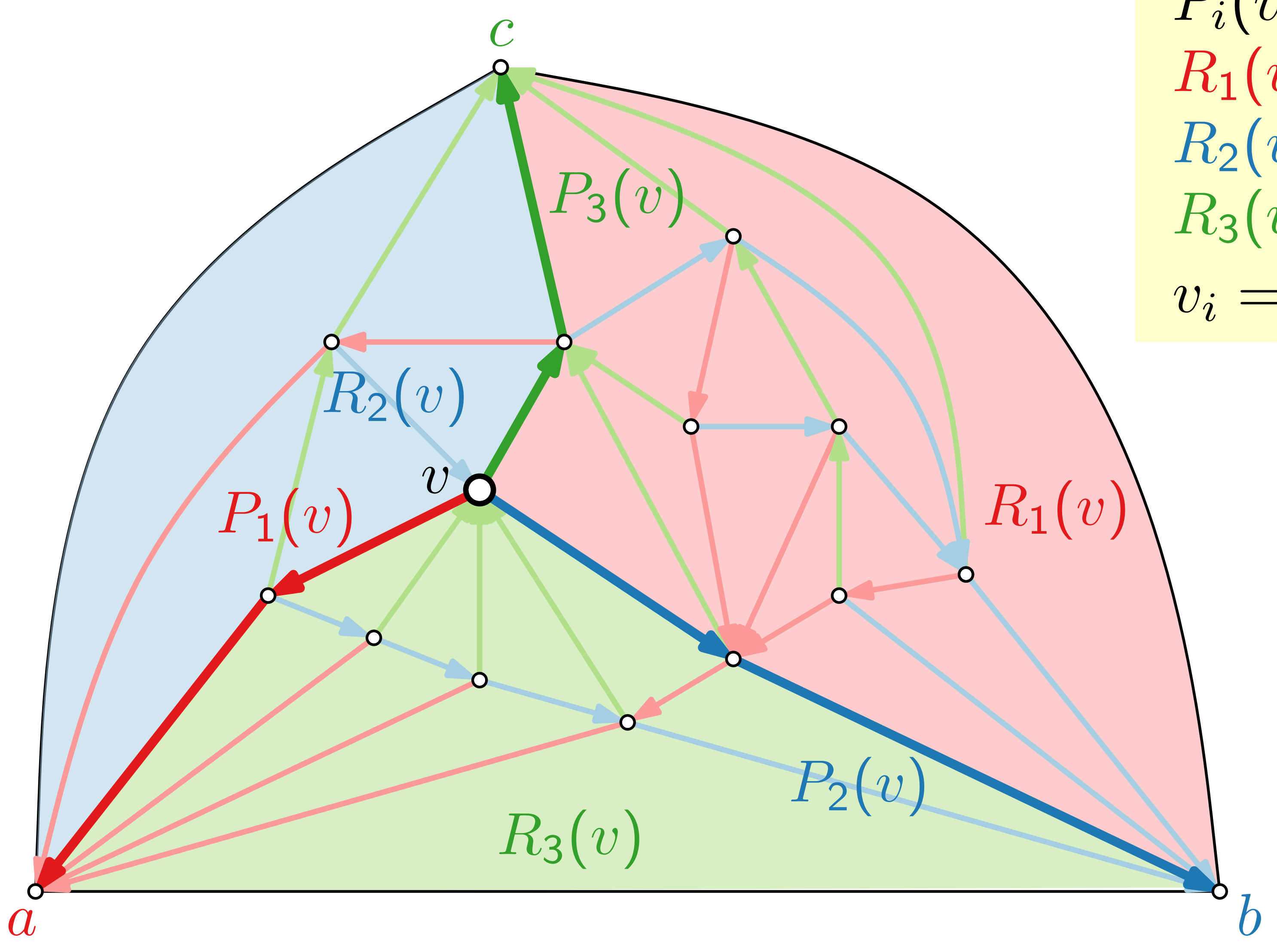

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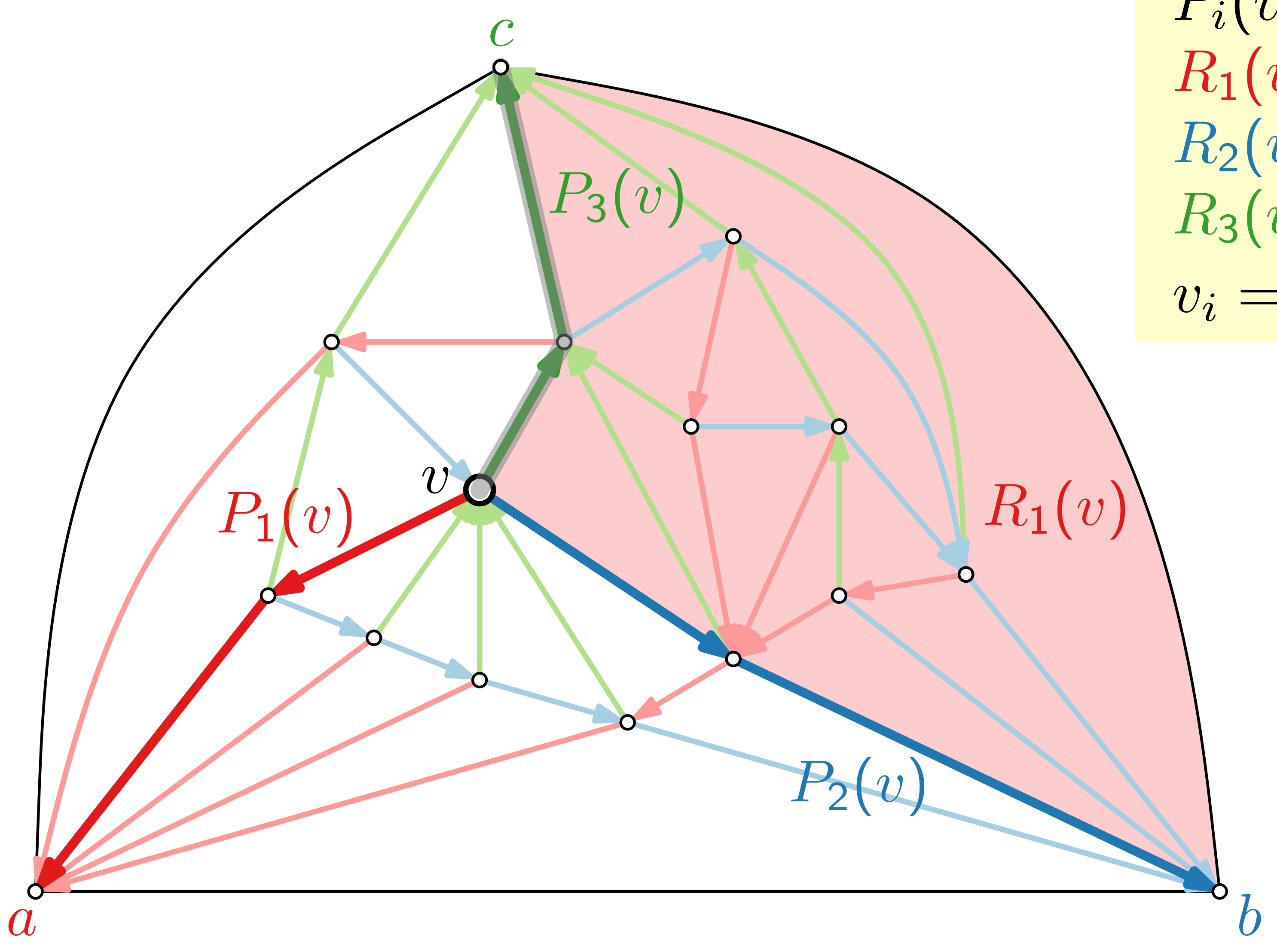

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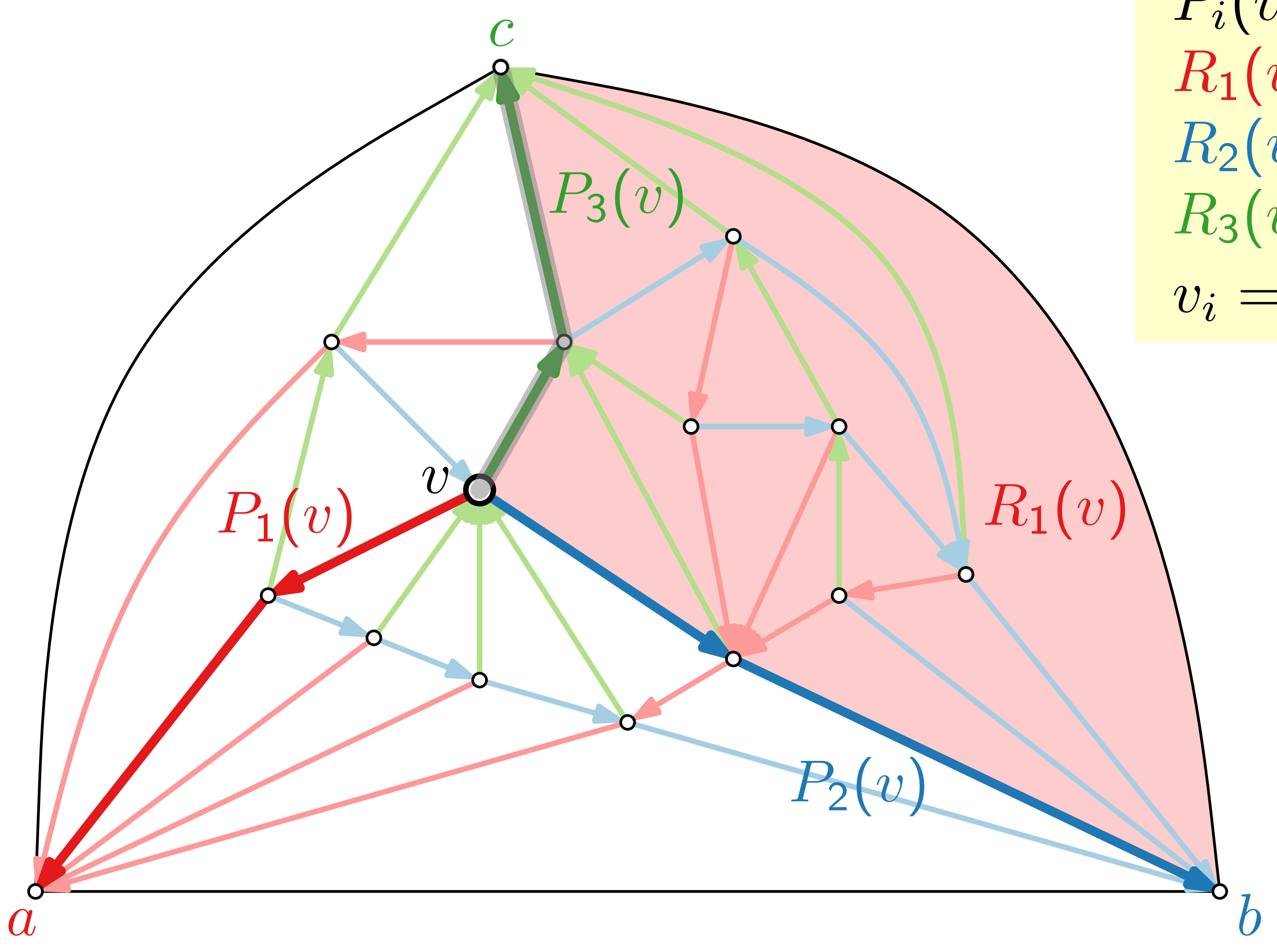

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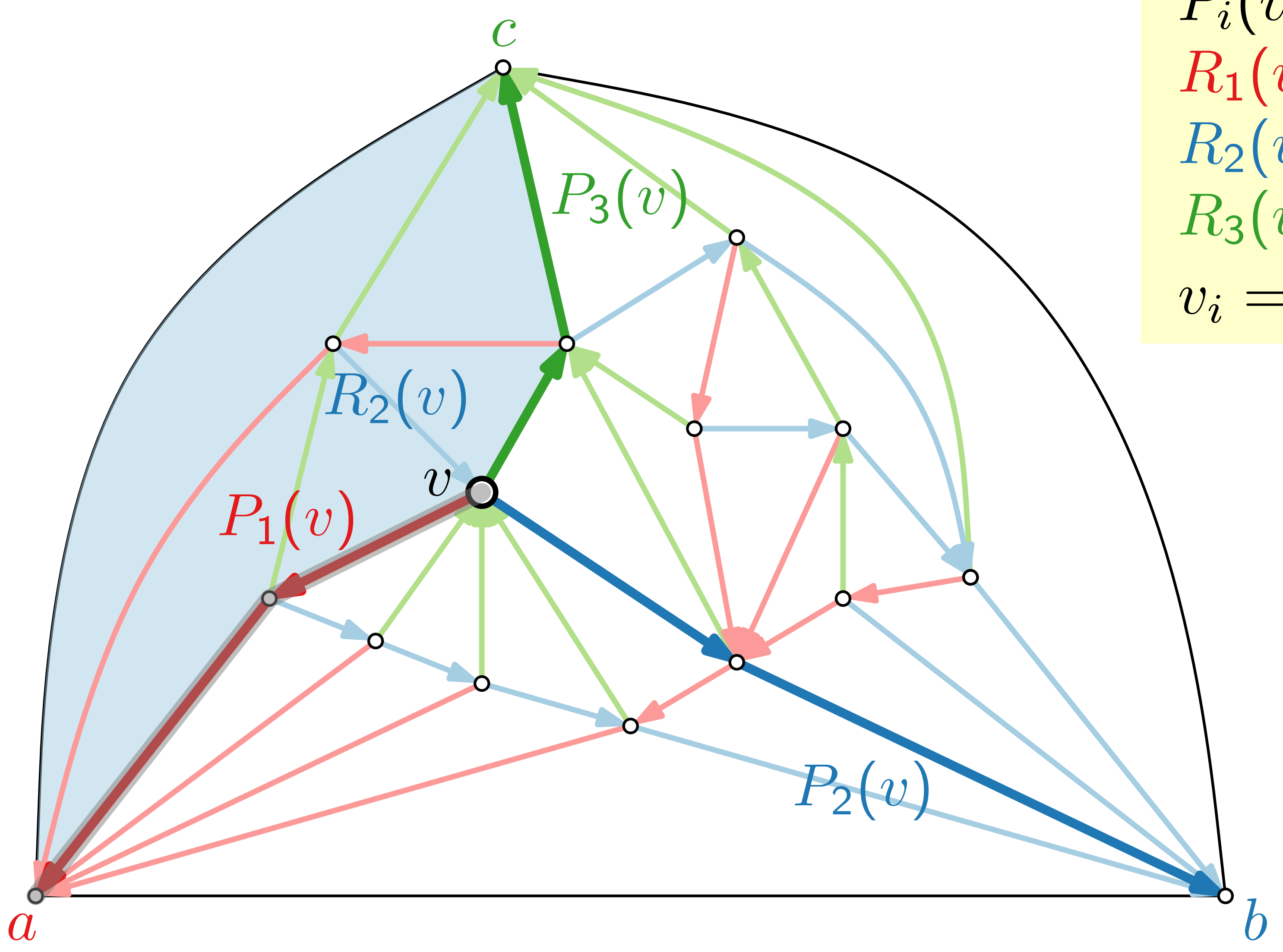

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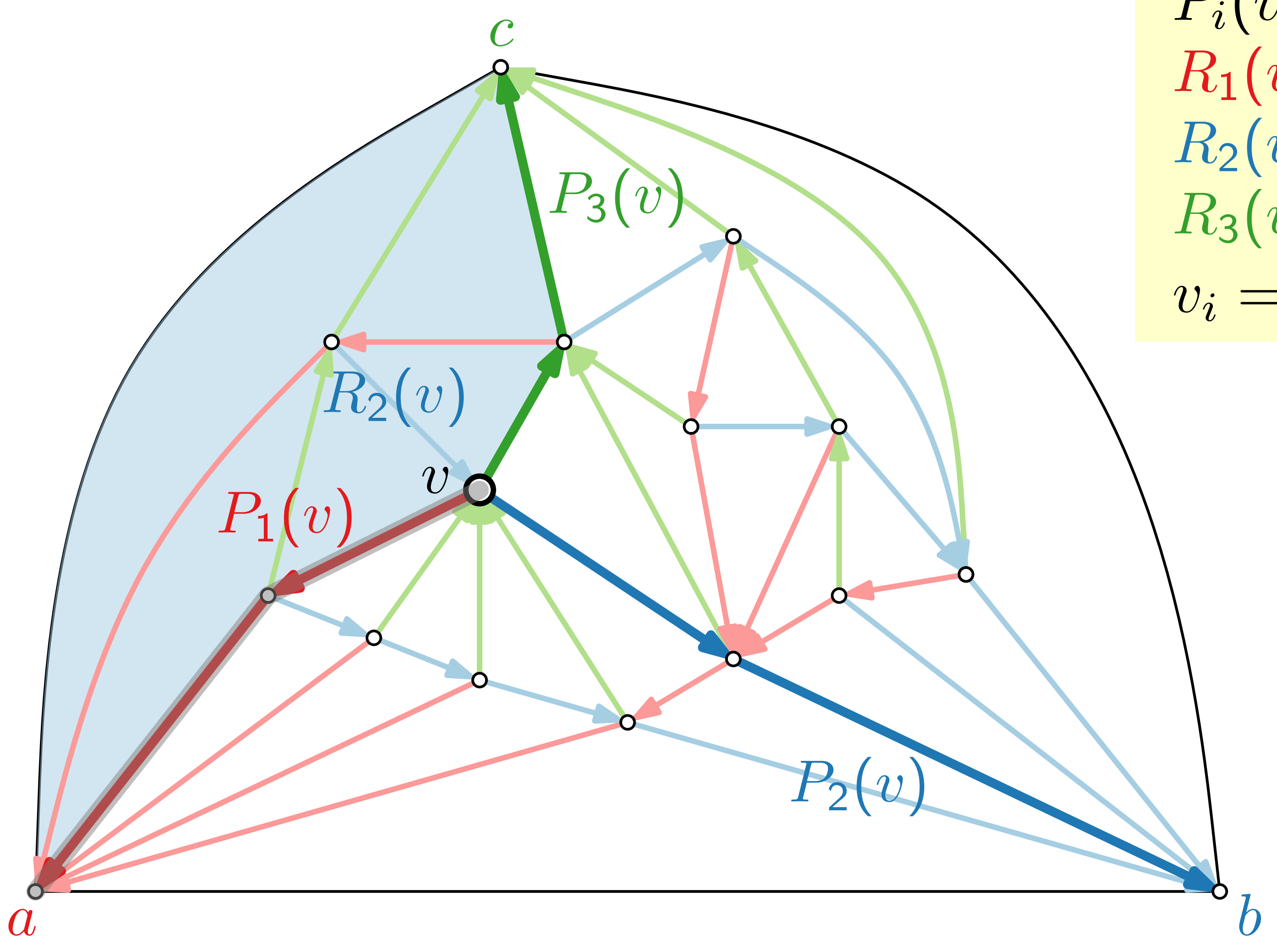

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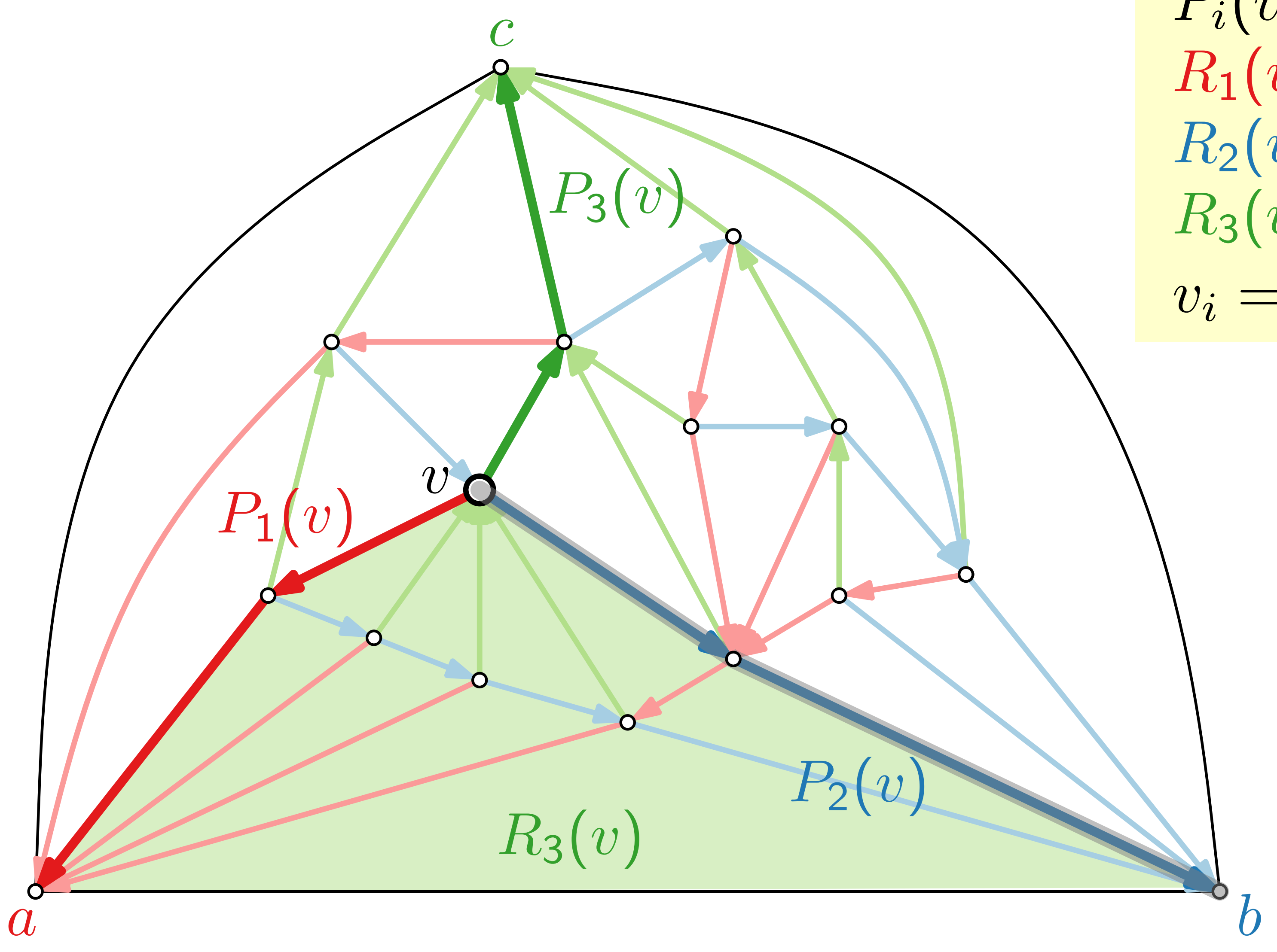

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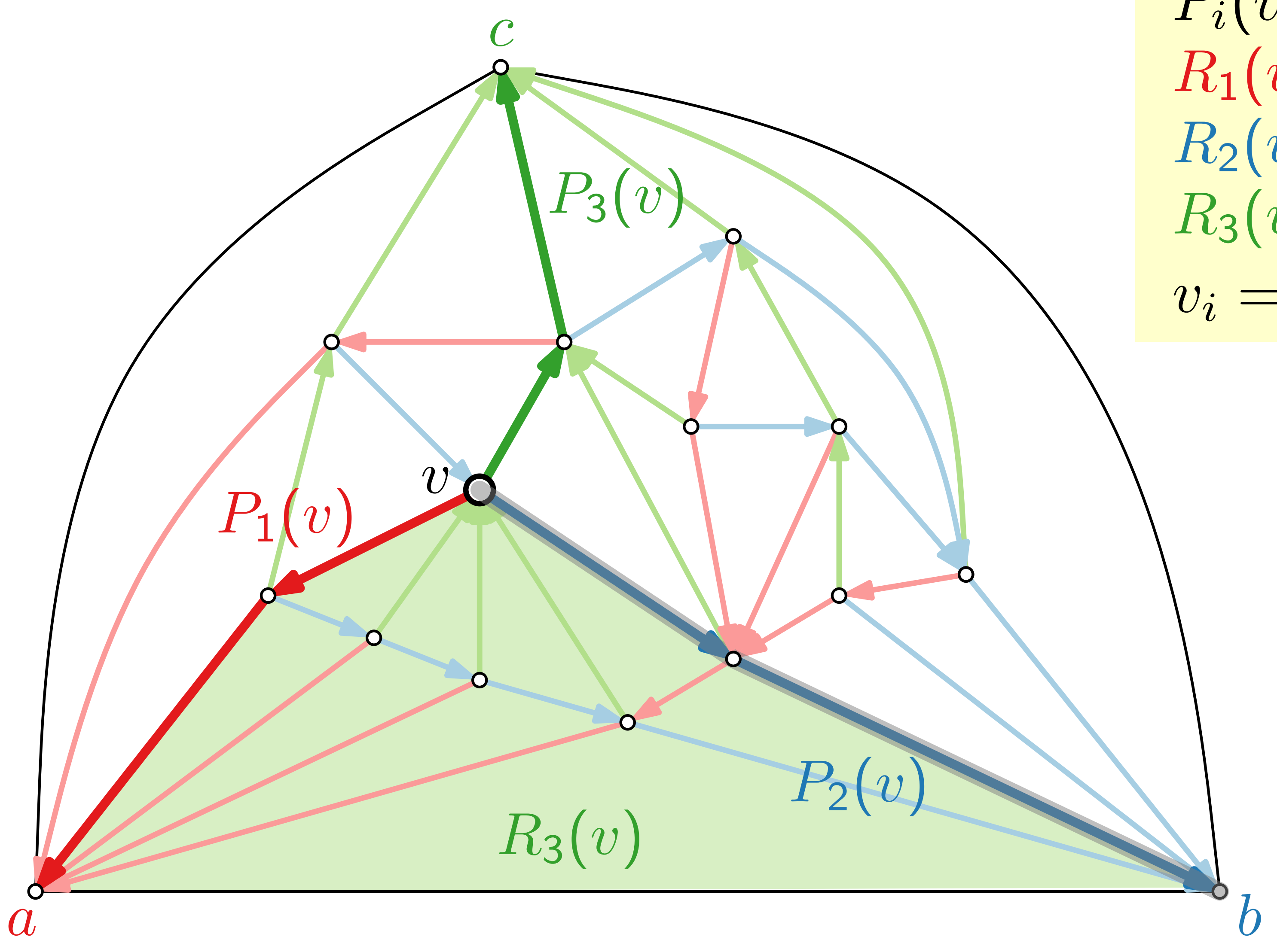

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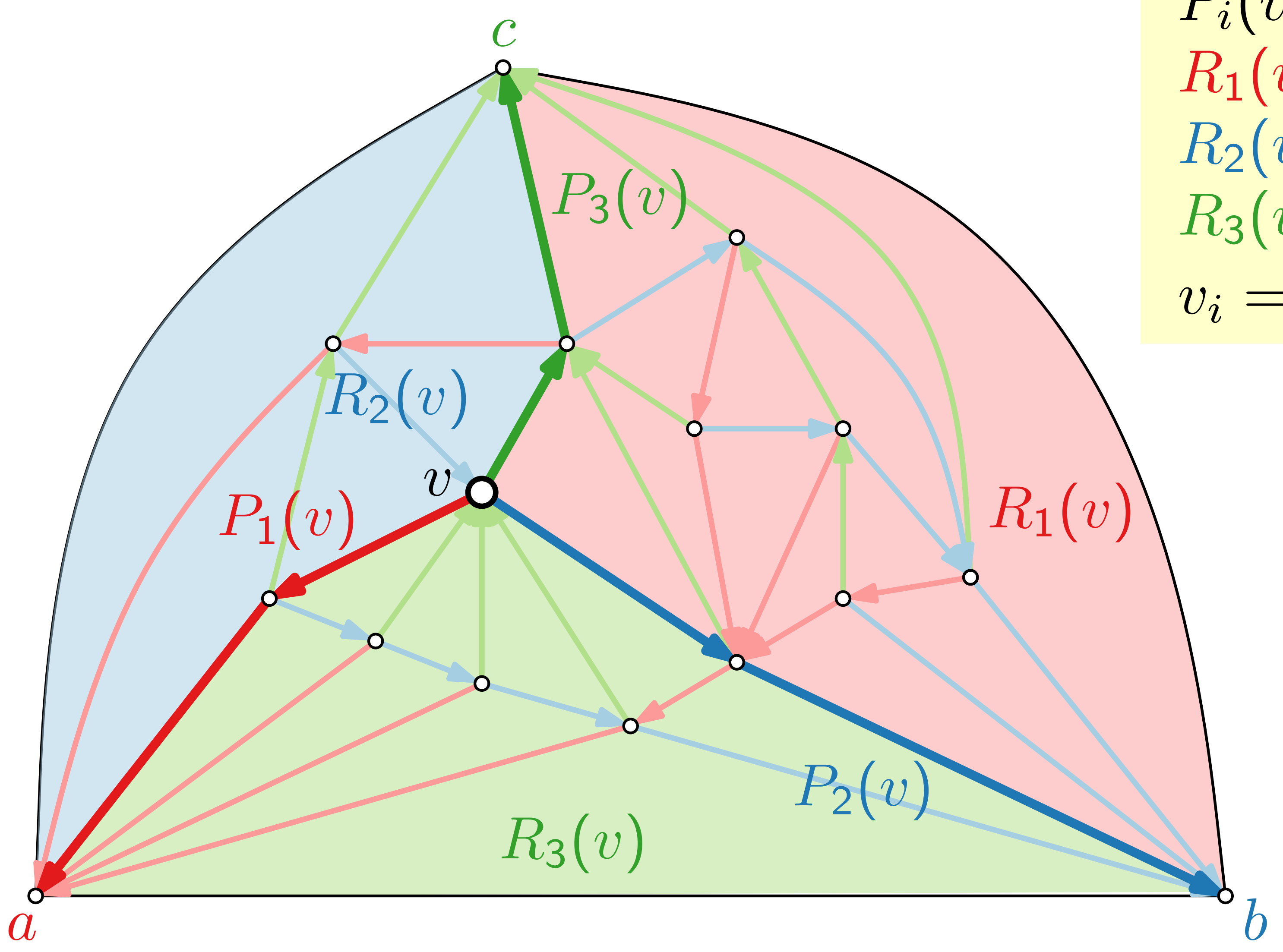

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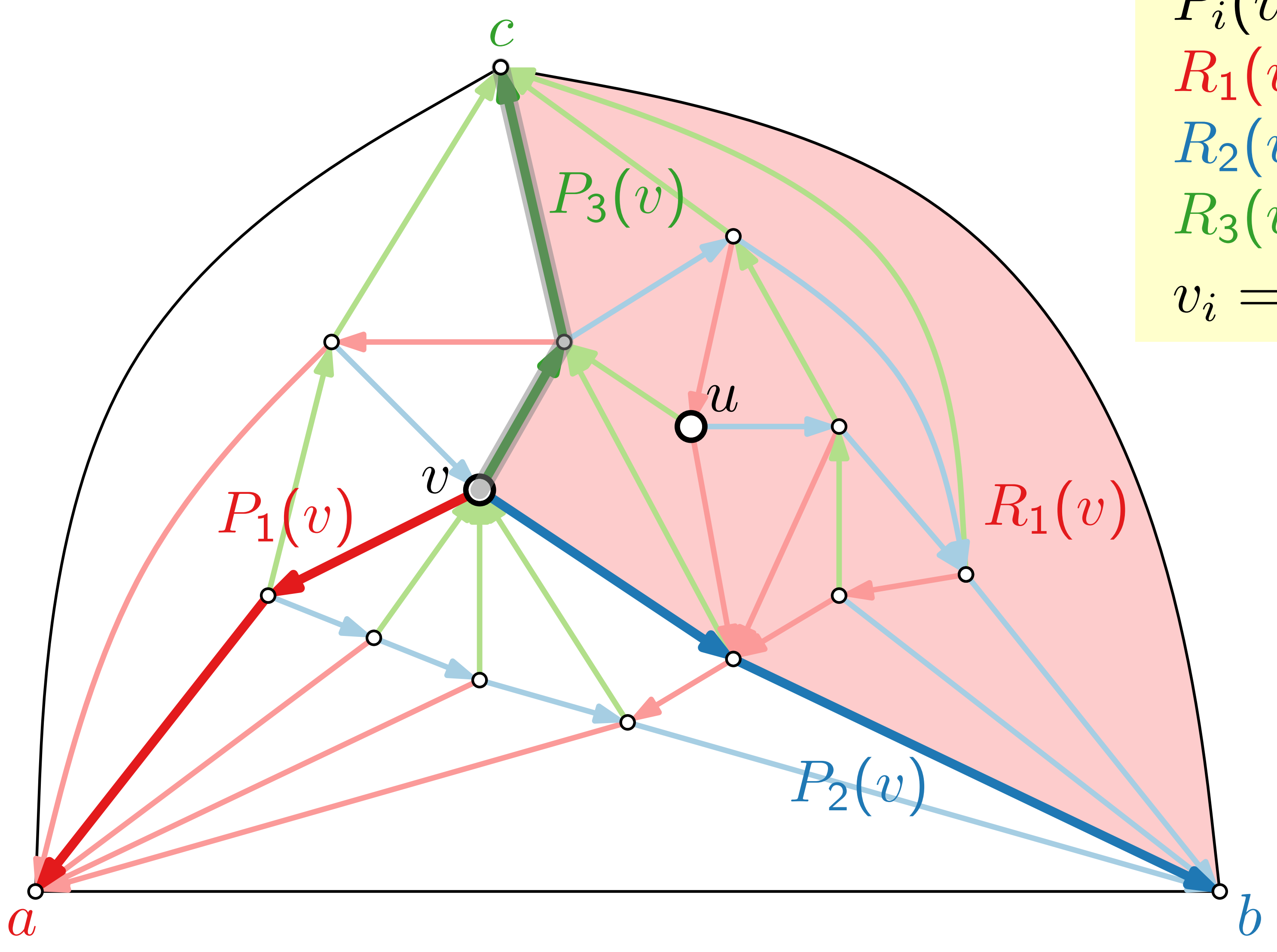

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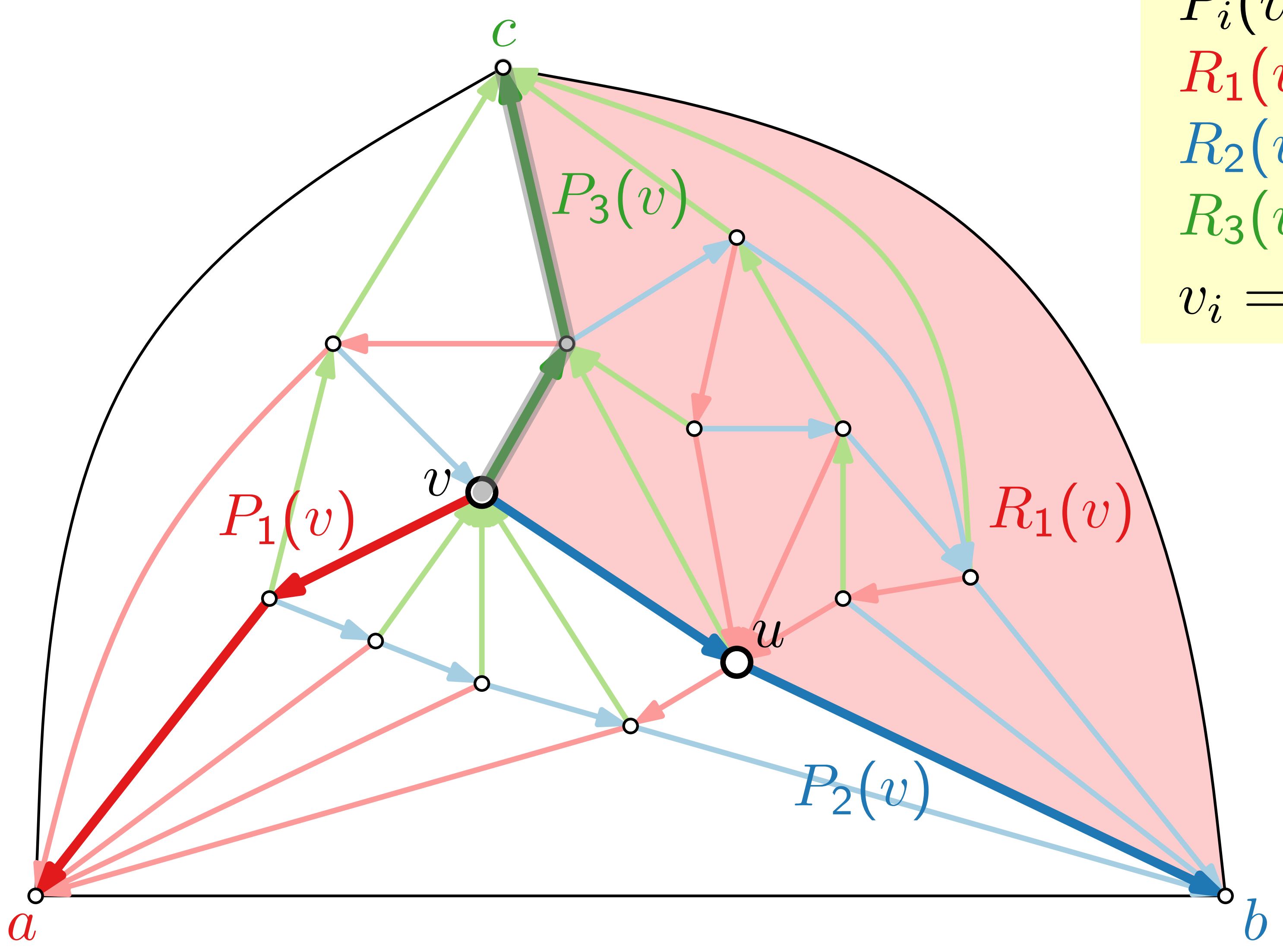

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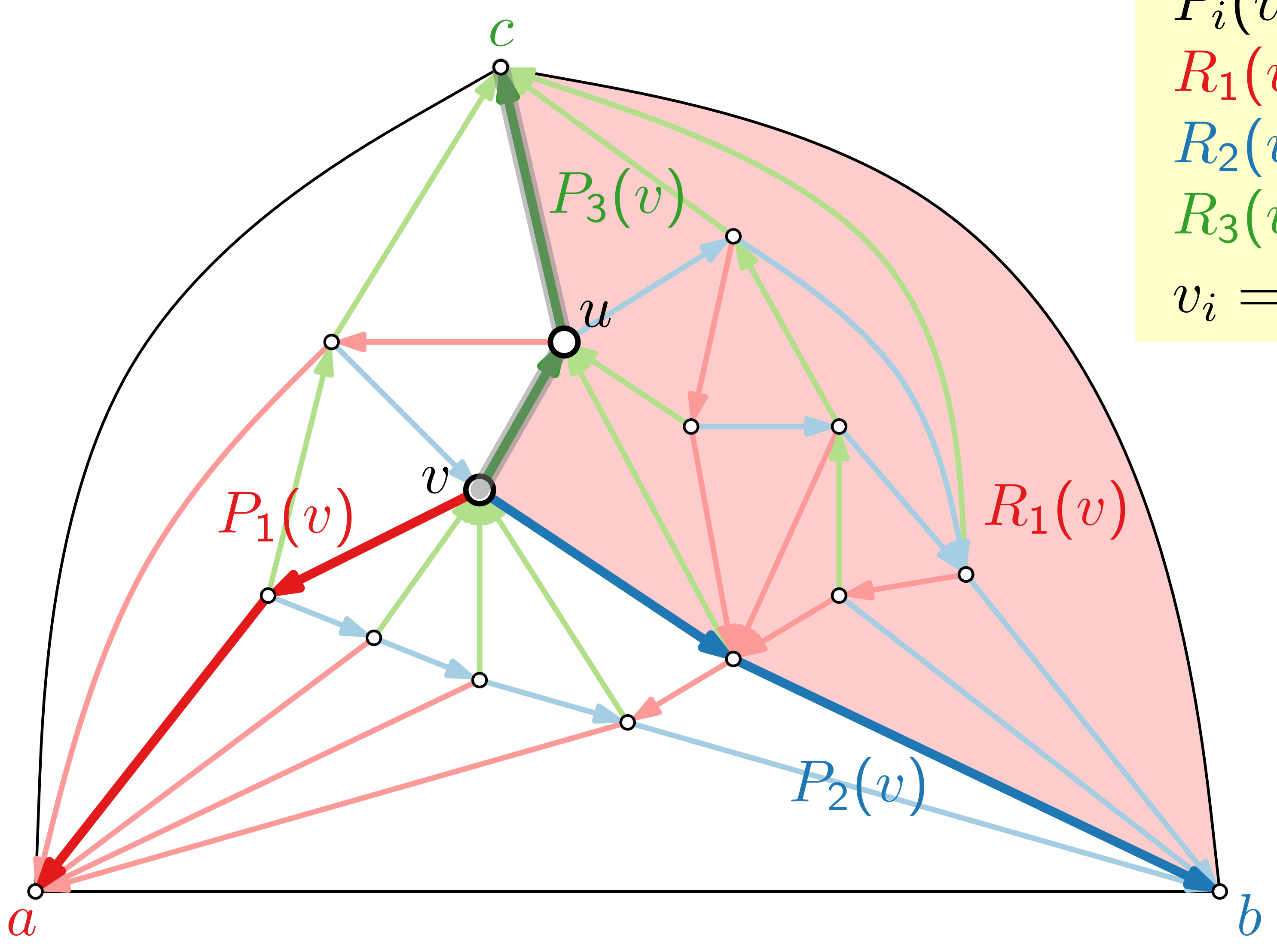

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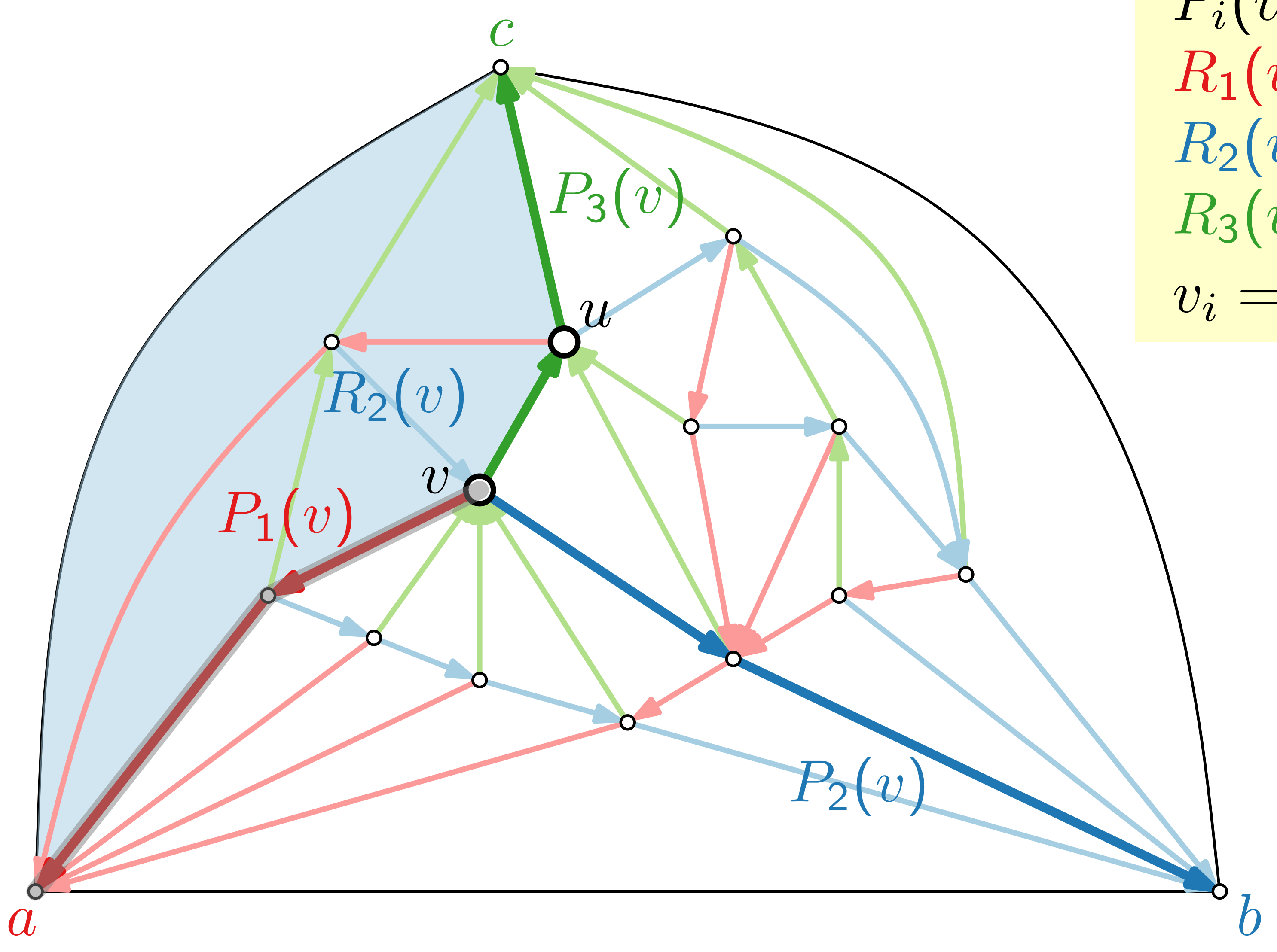

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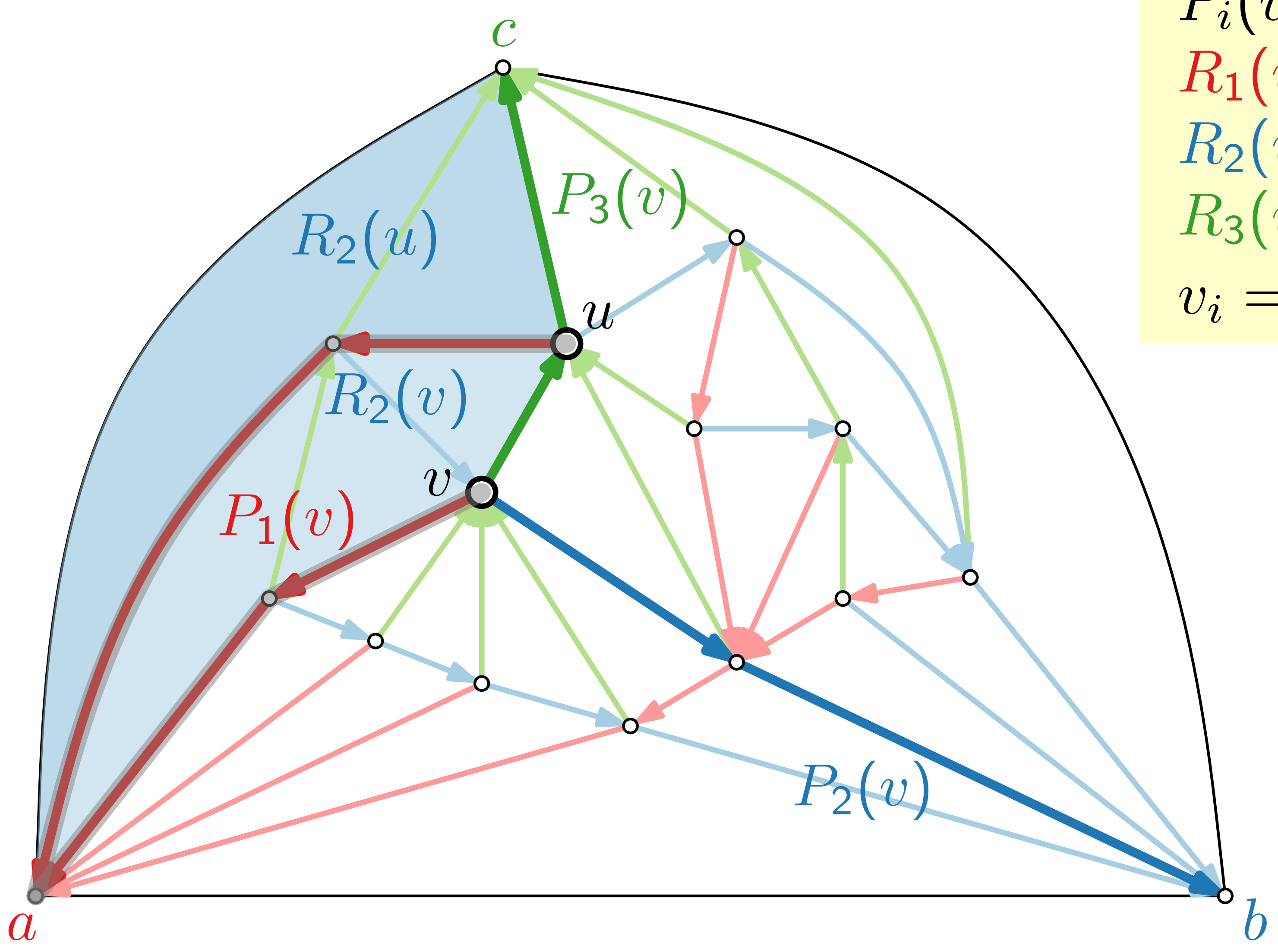

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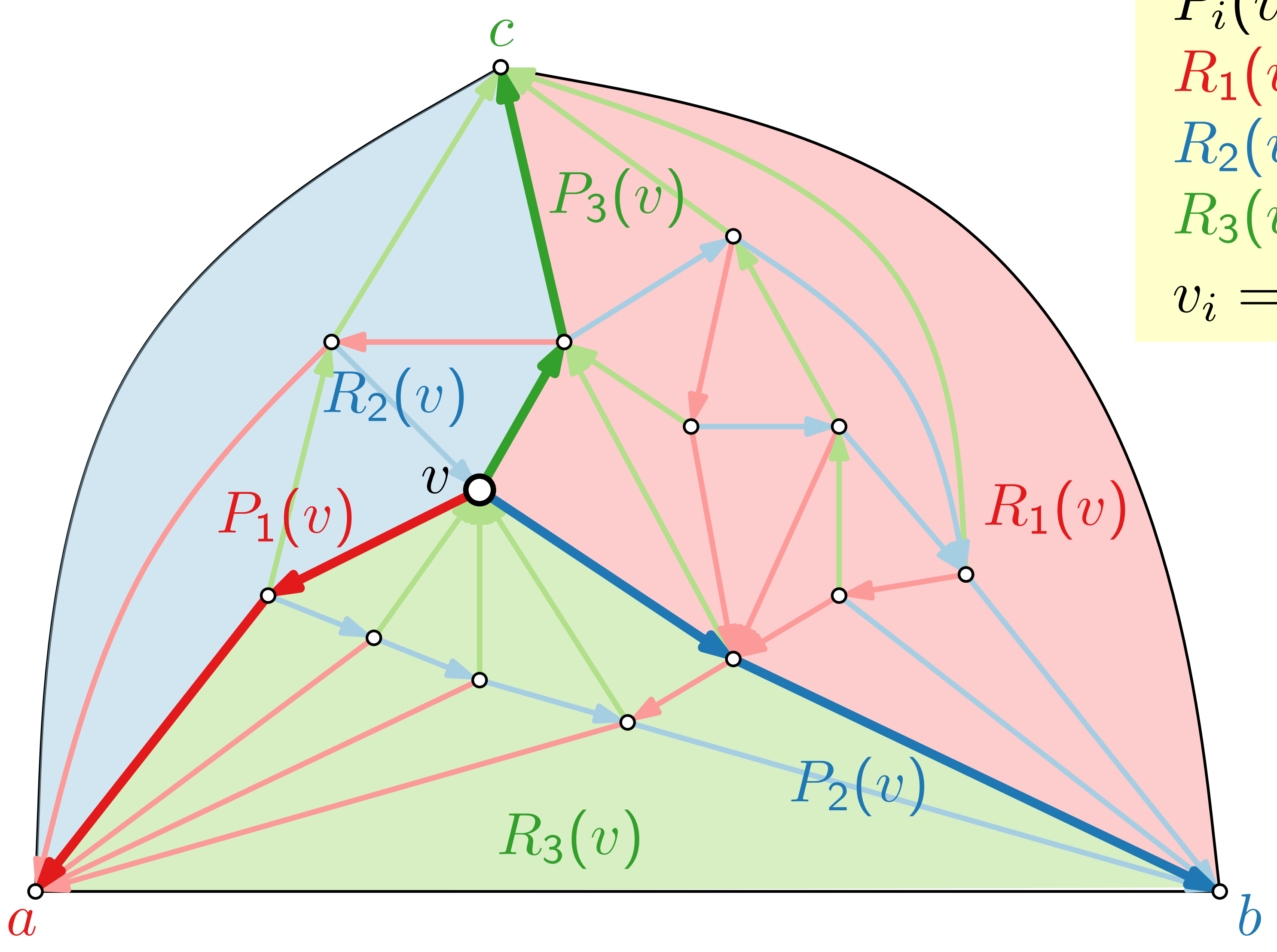

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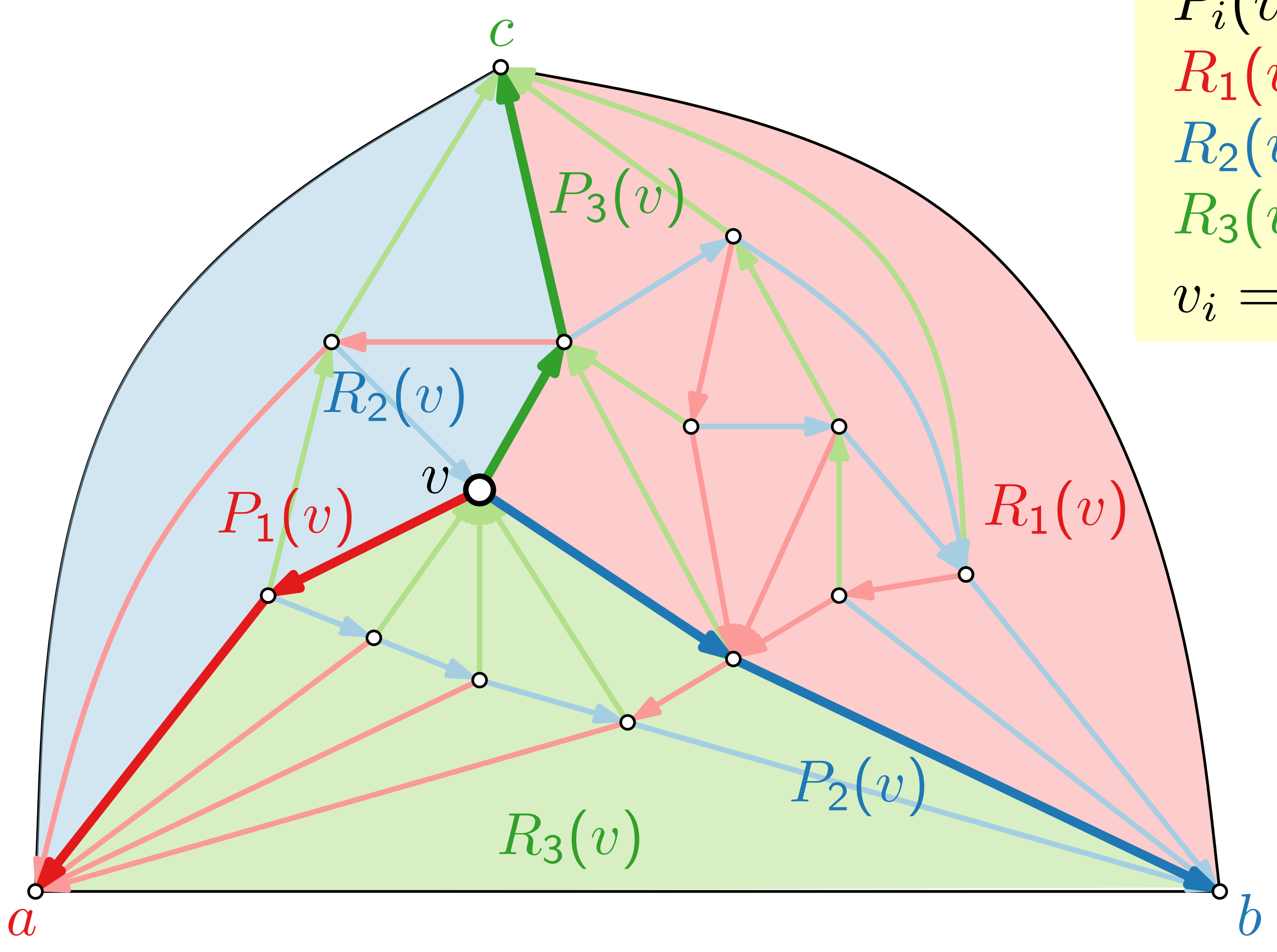

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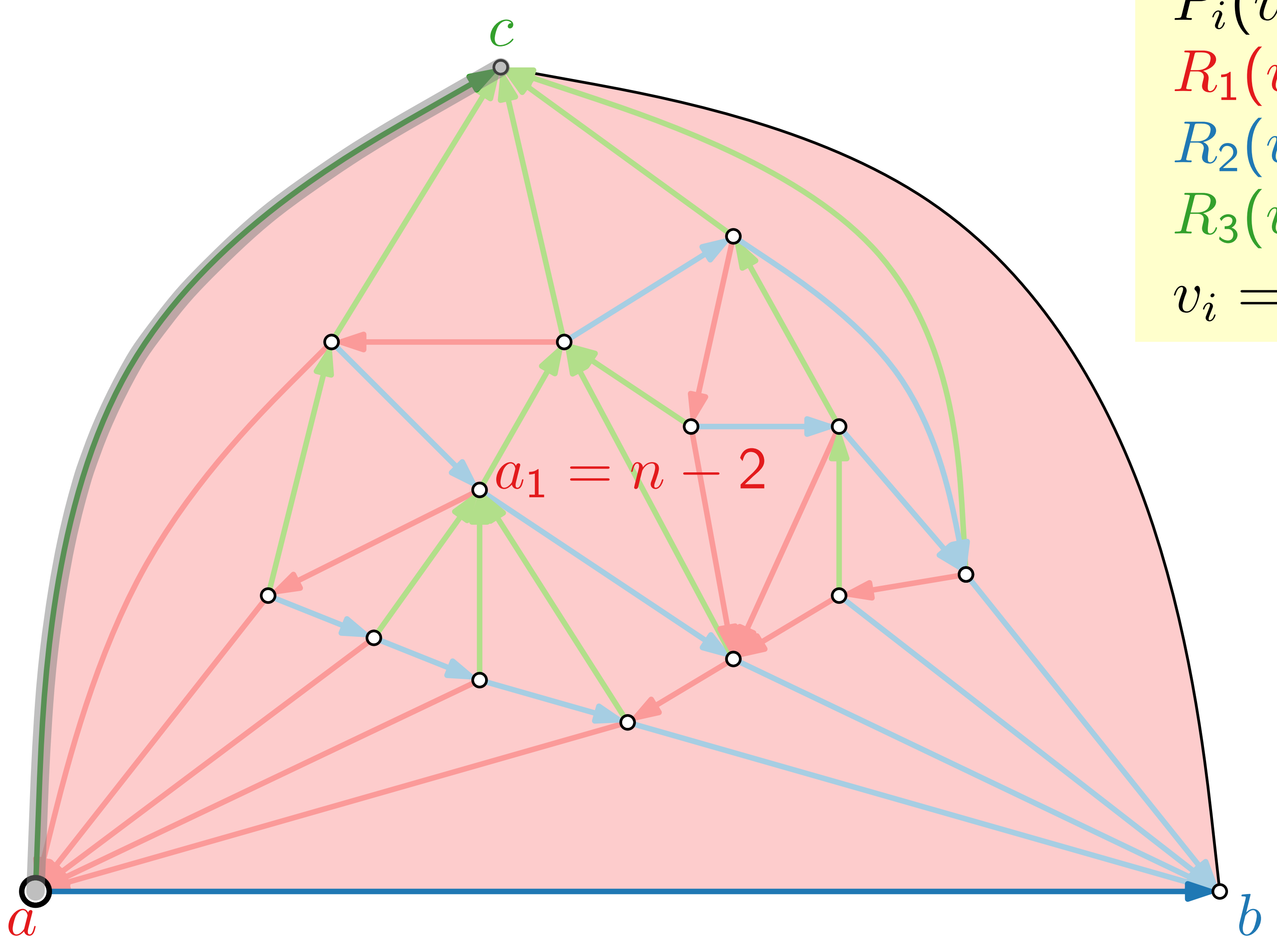
- For inner vertices $u \neq v$, it holds that $u \in R_i(v) \Rightarrow (u_i, u_{i+1}) <_{\text{lex}} (v_i, v_{i+1})$.
- $v_1 + v_2 + v_3 = n - 1$

Counting Vertices



$P_i(v)$: unique path from v to root of T_i

$R_1(v)$: subgraph bounded by $\langle P_2(v), bc, P_3(v)\rangle$

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$v_i = |V(R_i(v))| - |V(P_{i-1}(v))|$ (indices modulo 3)

$v_1 = 10 - 3 = 7$

$v_2 = 6 - 3 = 3$

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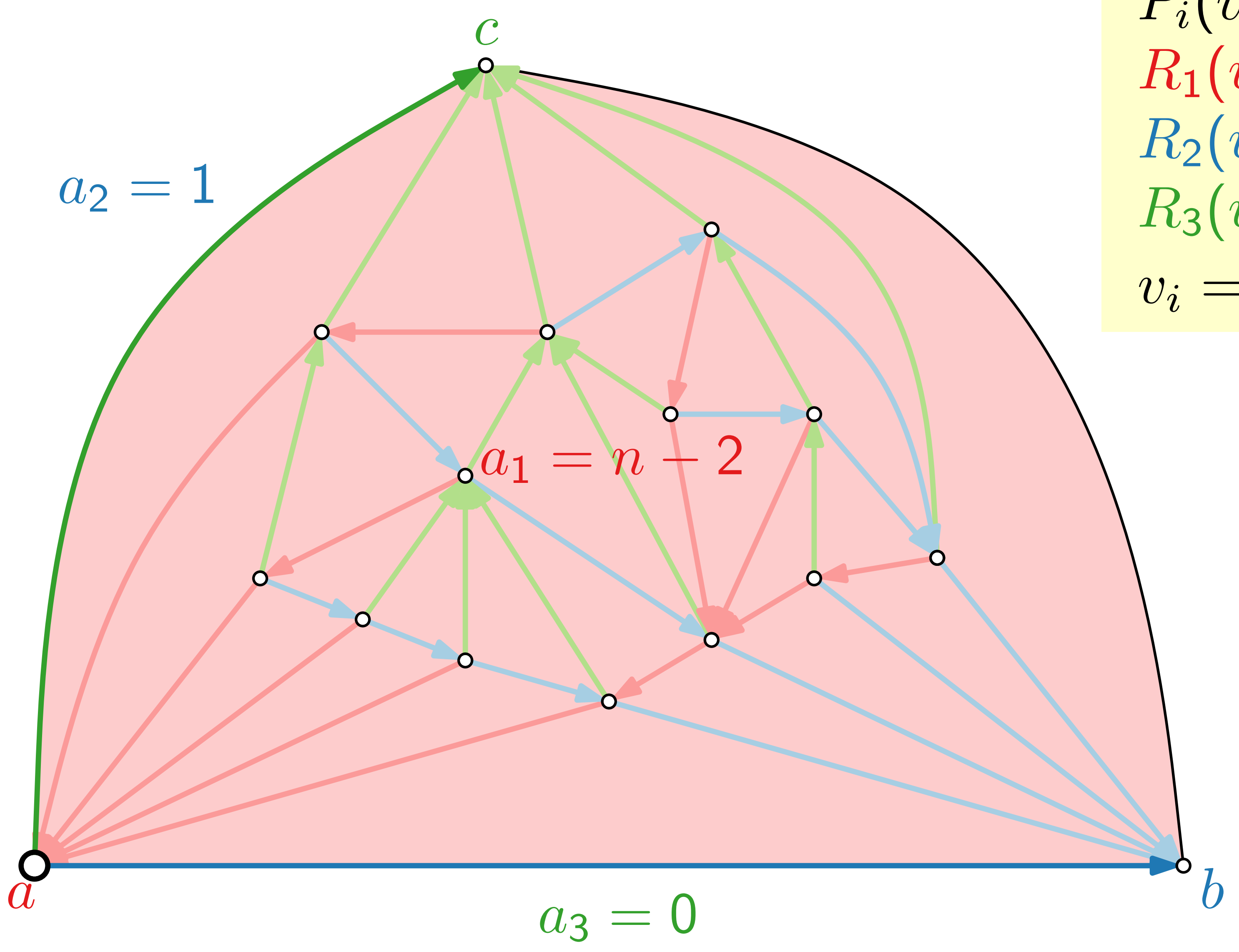

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Schnyder Drawing*

Set $A = (0,0)$, $B = (n-1, 0)$, and $C = (0, n-1)$.

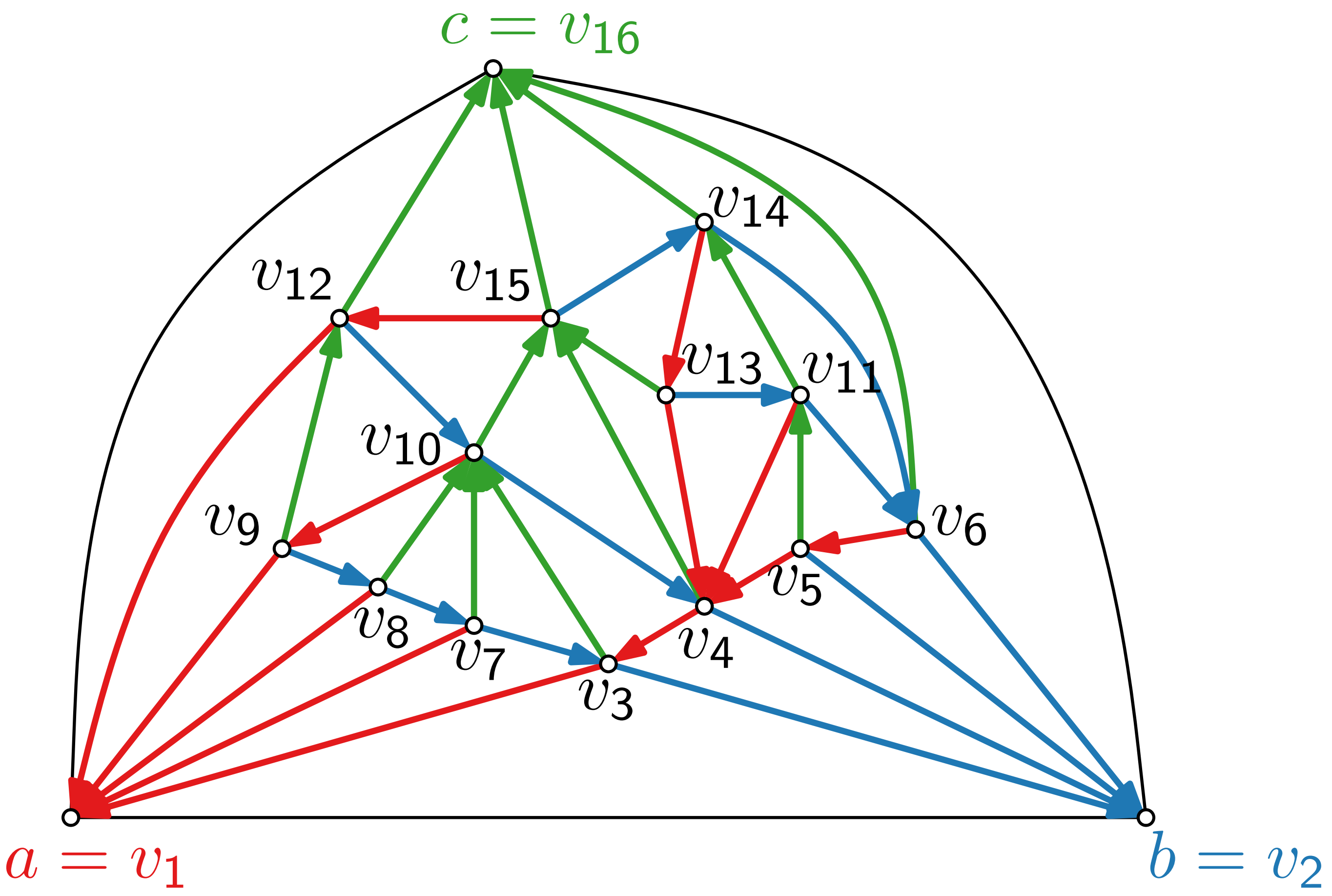
Theorem. [Schnyder '90]

For a plane triangulation G, the mapping

$$f\colon v \mapsto \frac{1}{n-1}(v_1, v_2, v_3)$$

is a weak barycentric representation of G and, thus, yields a planar straight-line drawing of G on the $(n-2) \times (n-2)$ grid.

Schnyder Drawing* – Example



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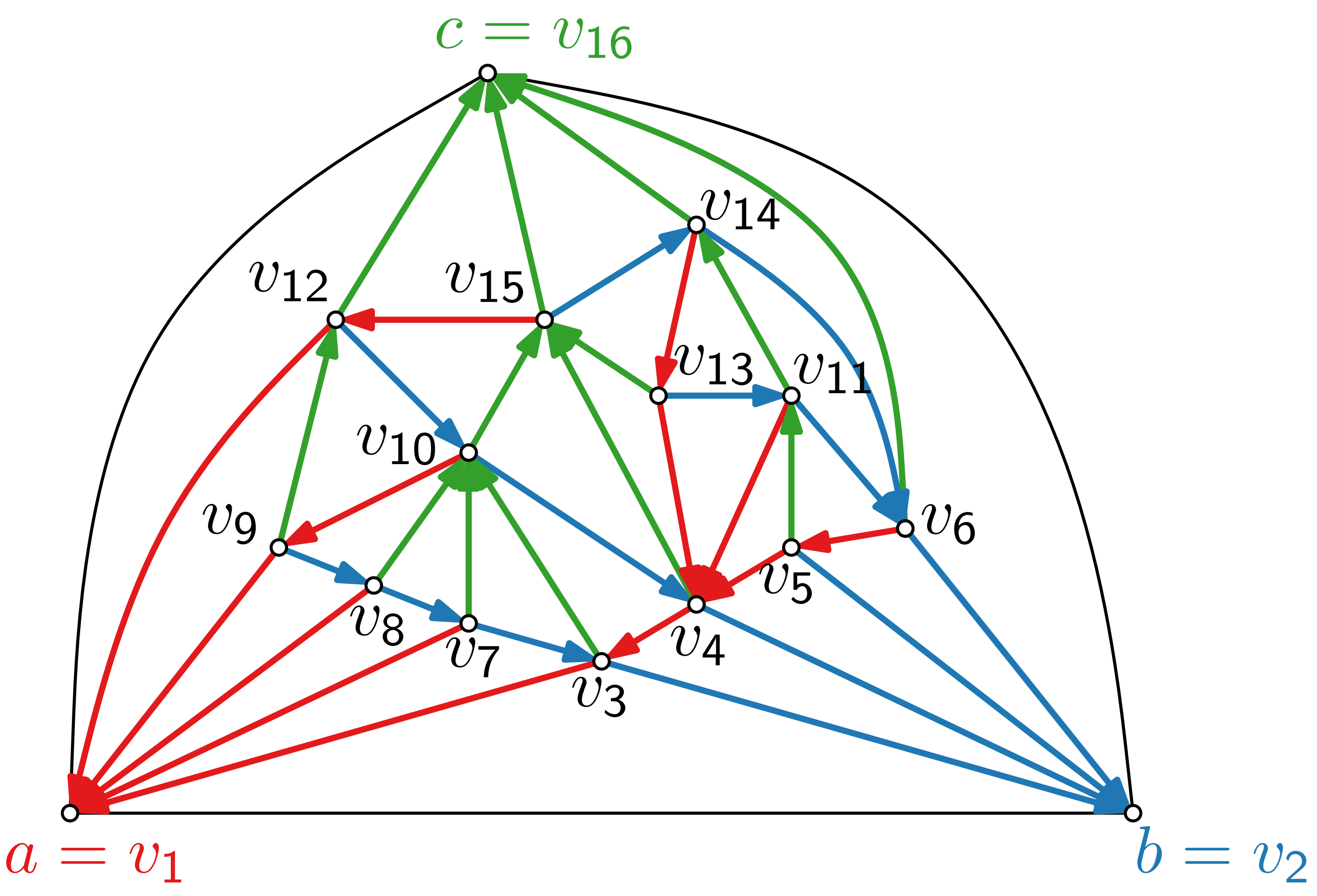

$n = 16$, $n - 2 = 14$

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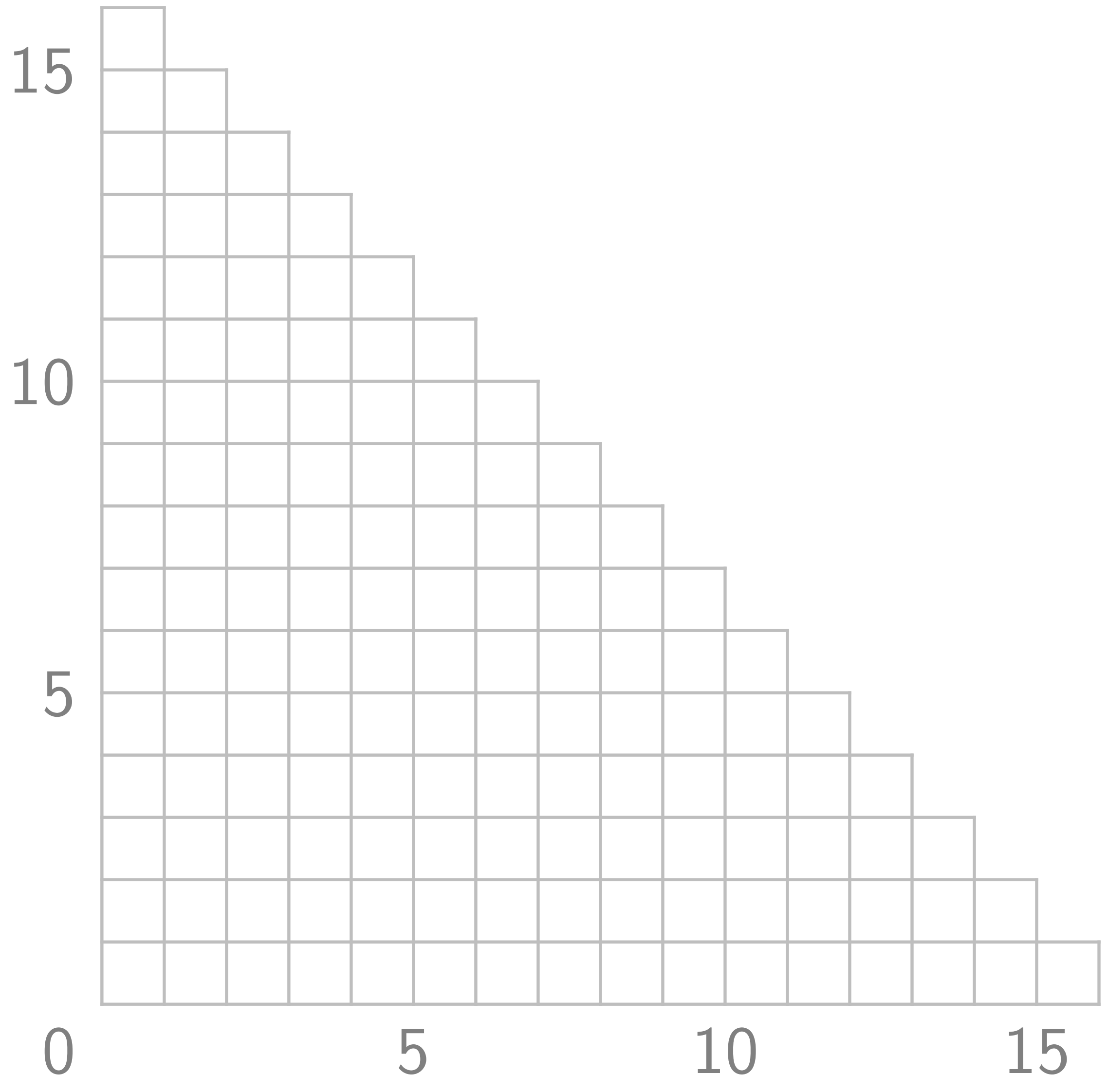

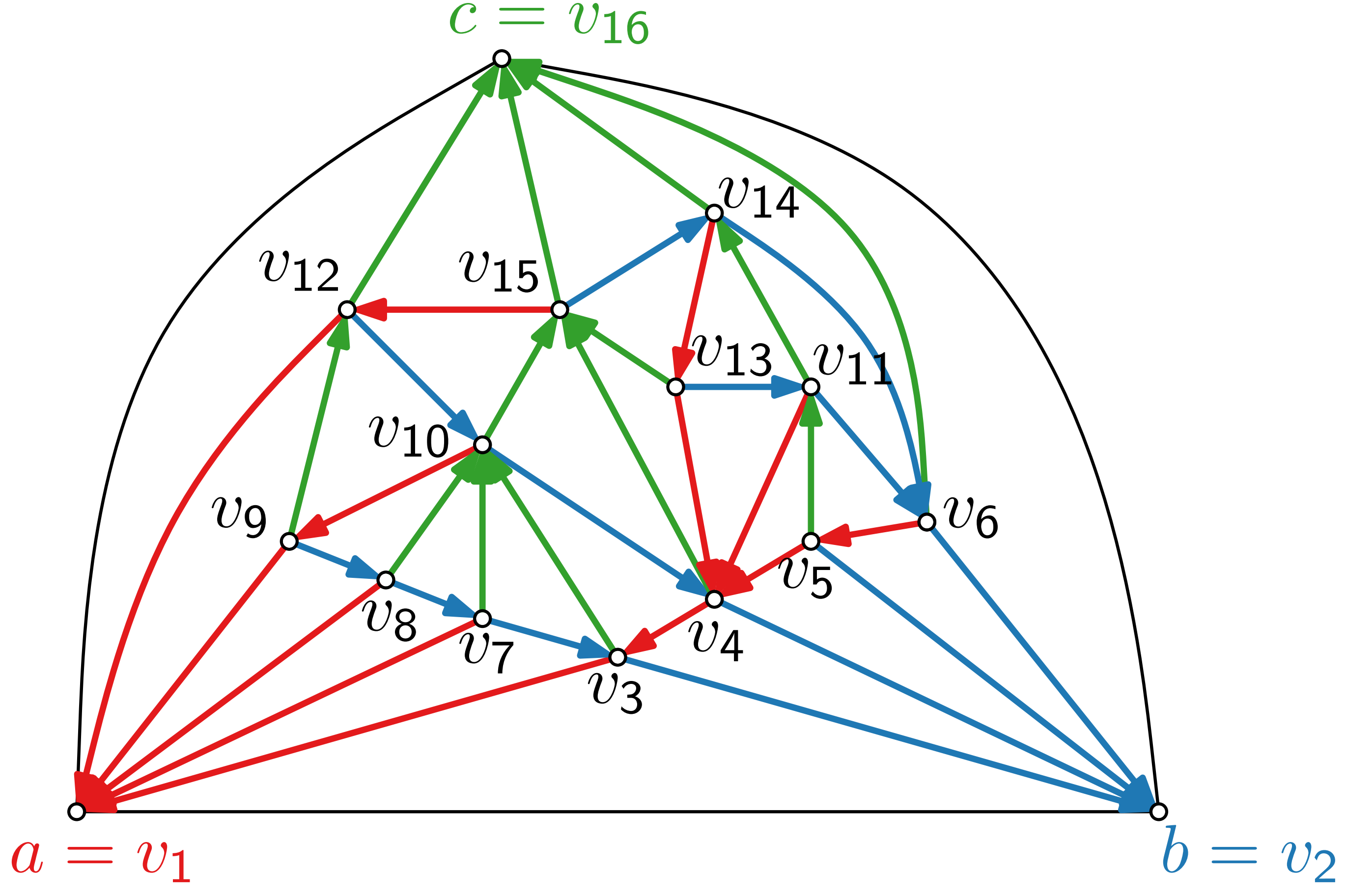

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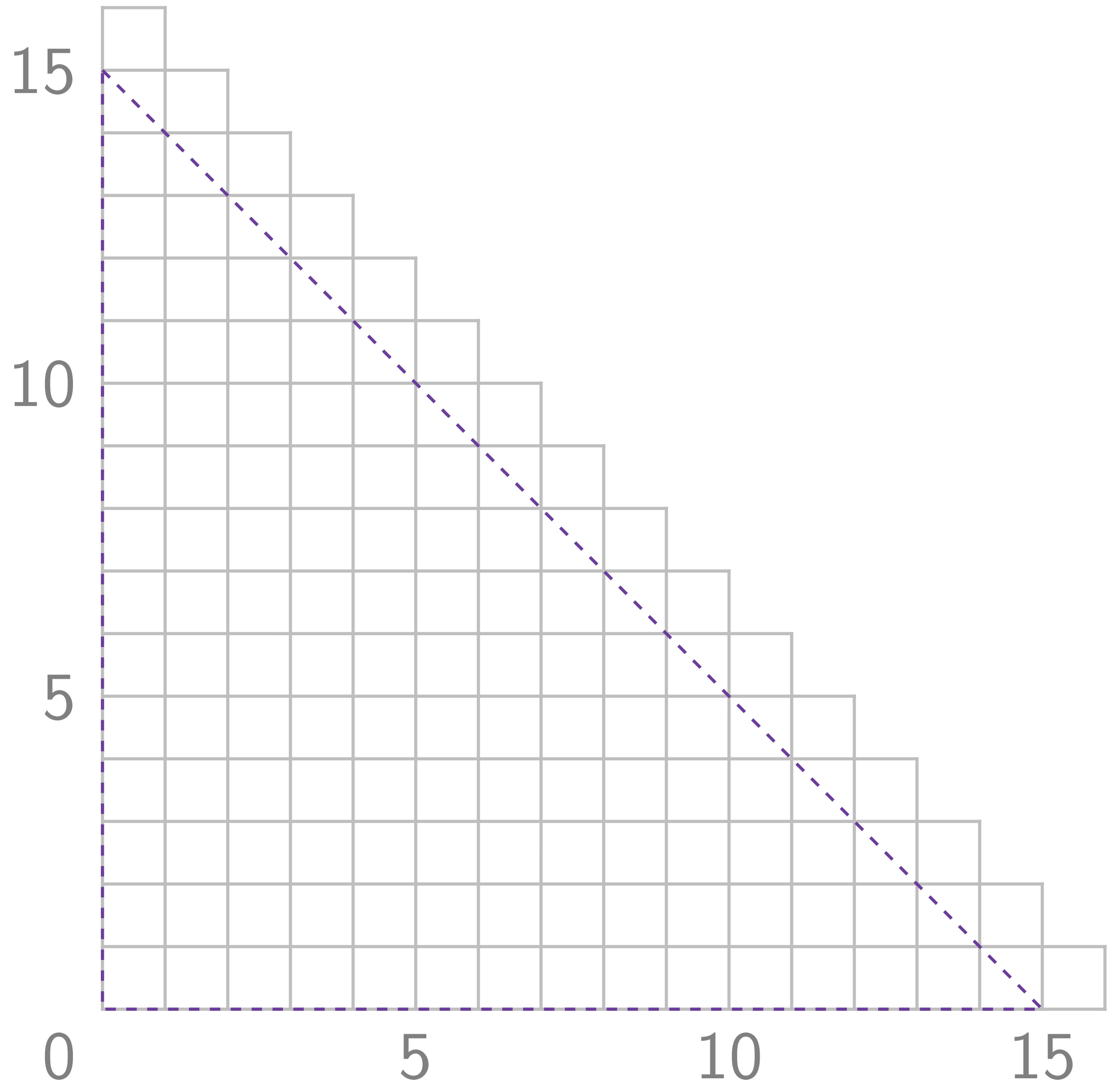

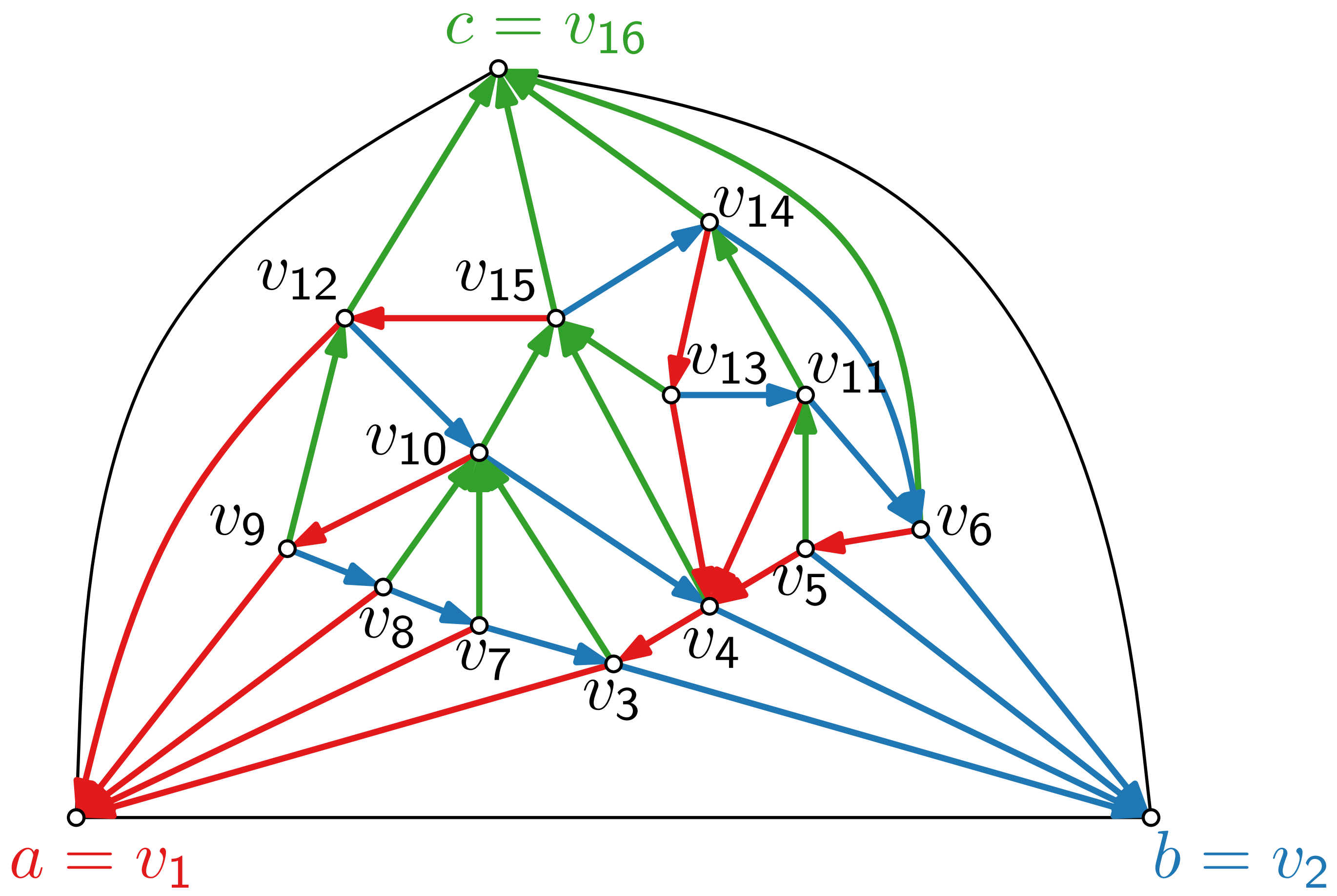

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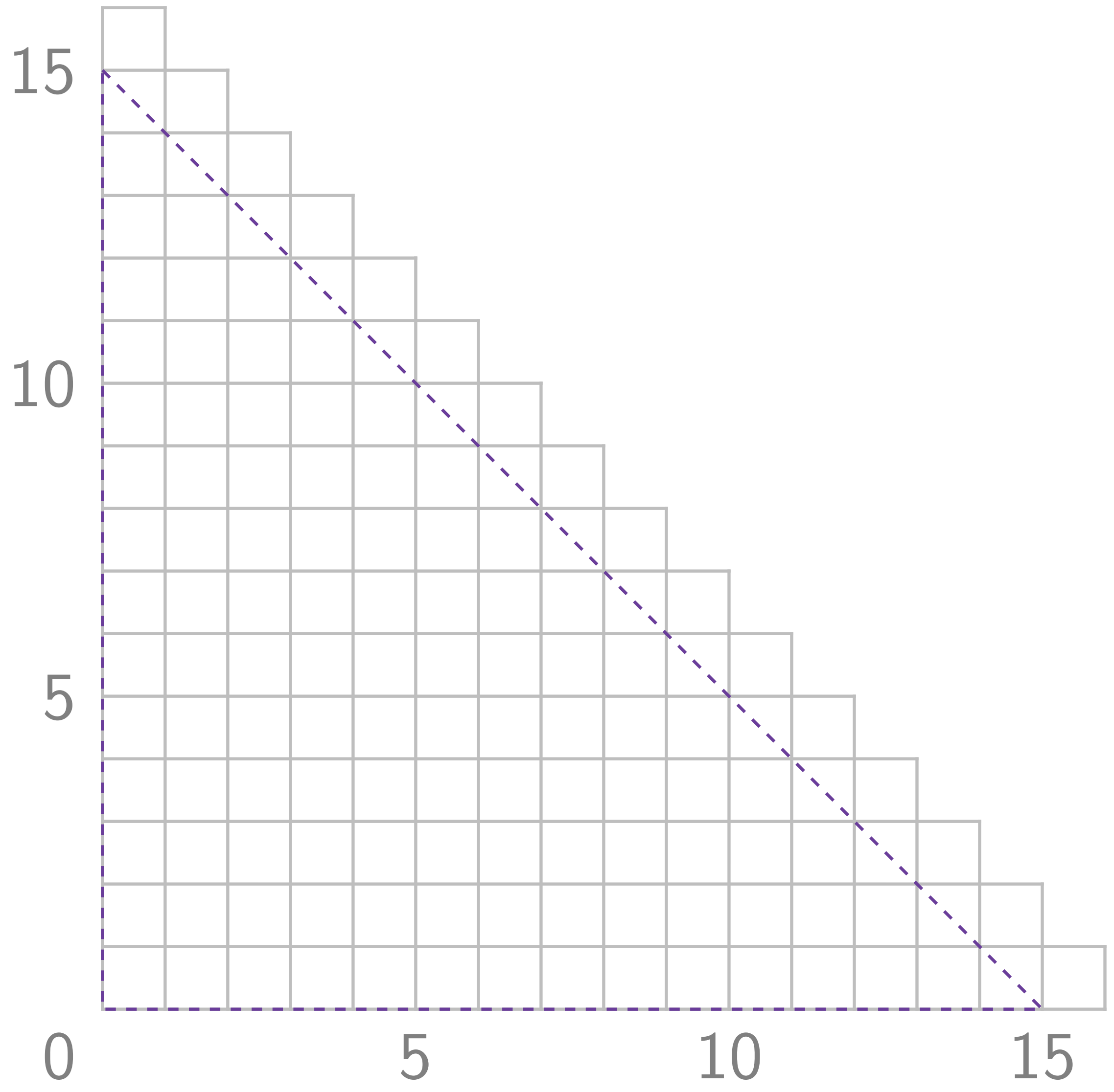

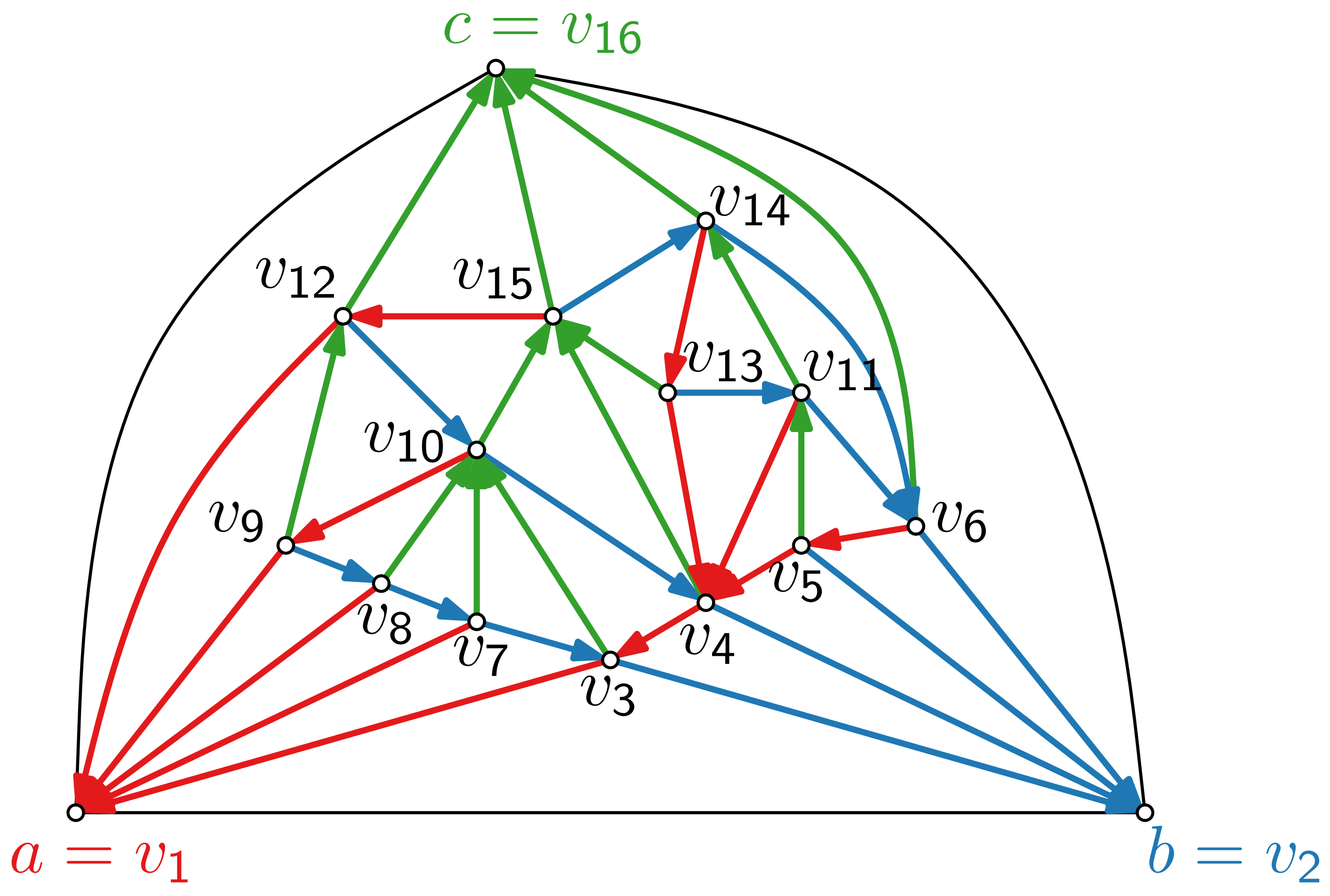

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$f(a) = (14, 1, 0)$

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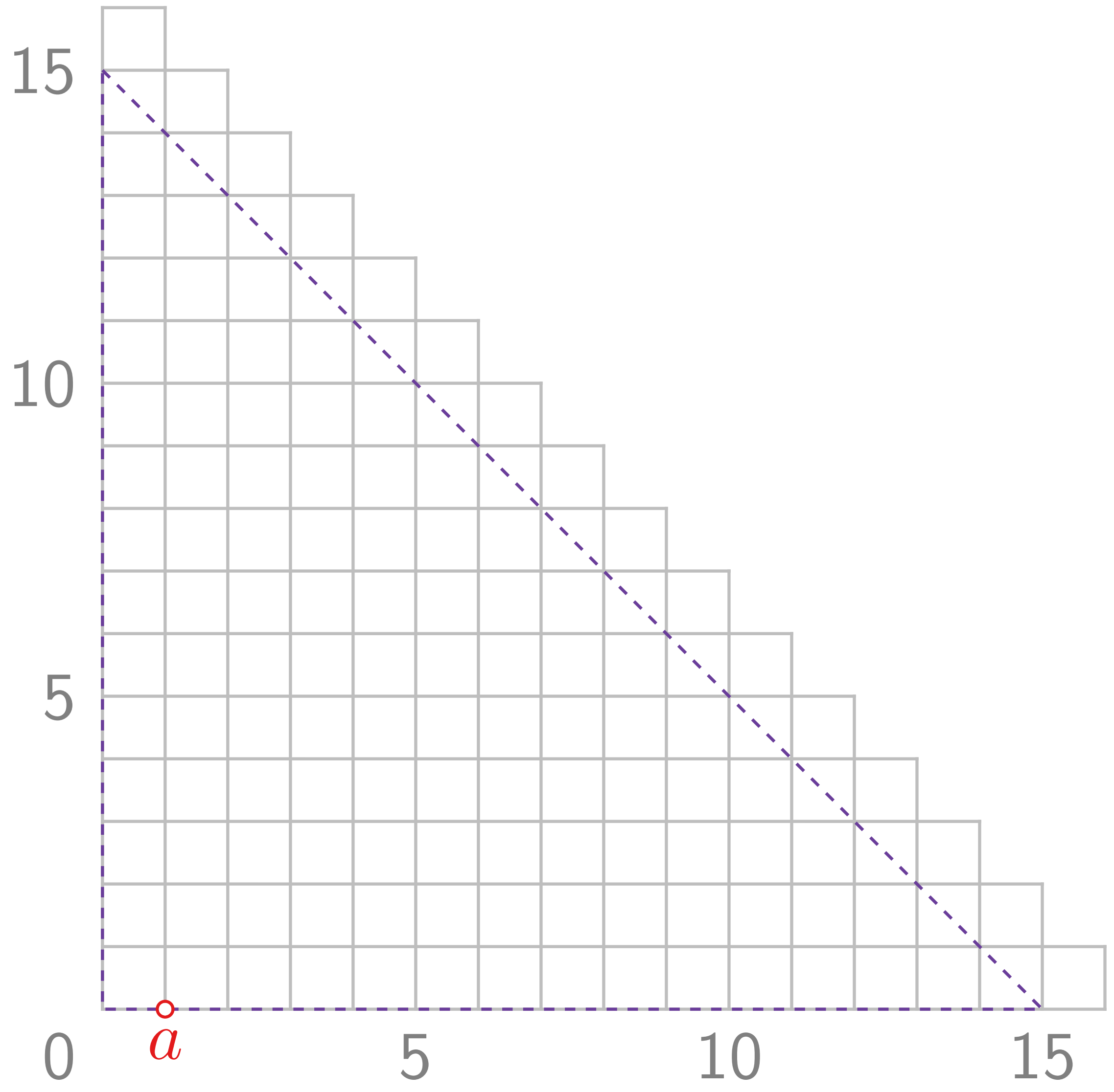

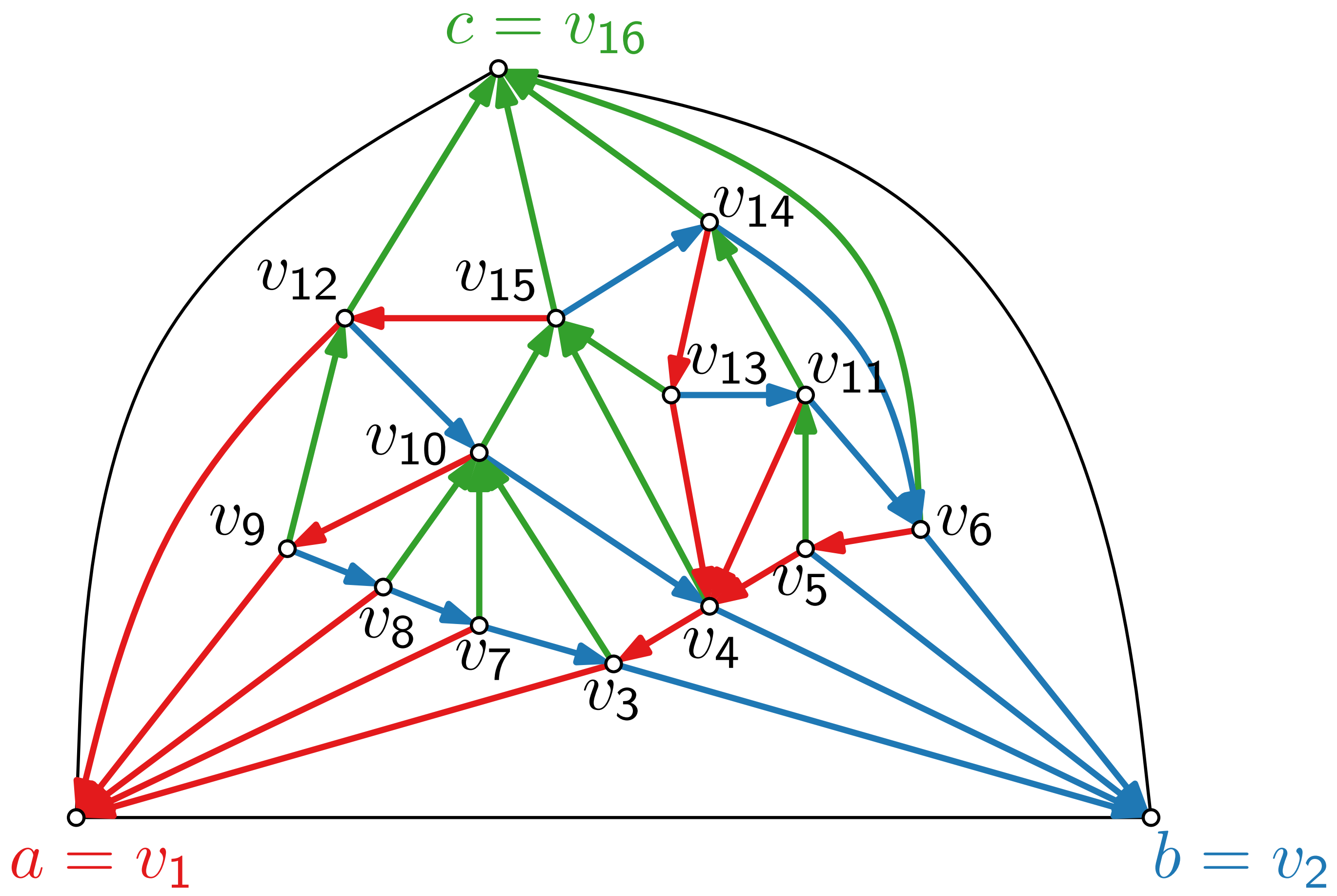

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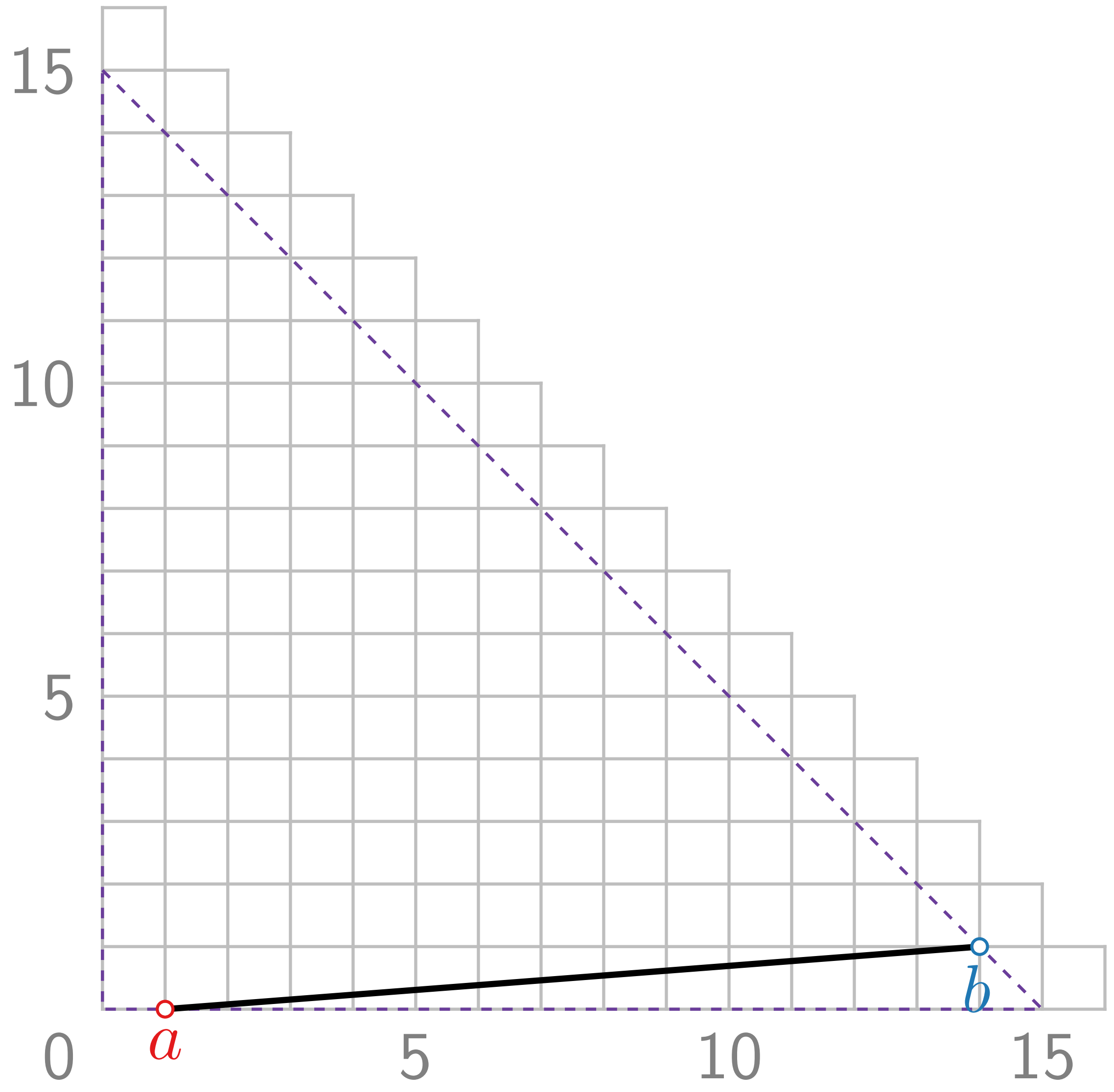

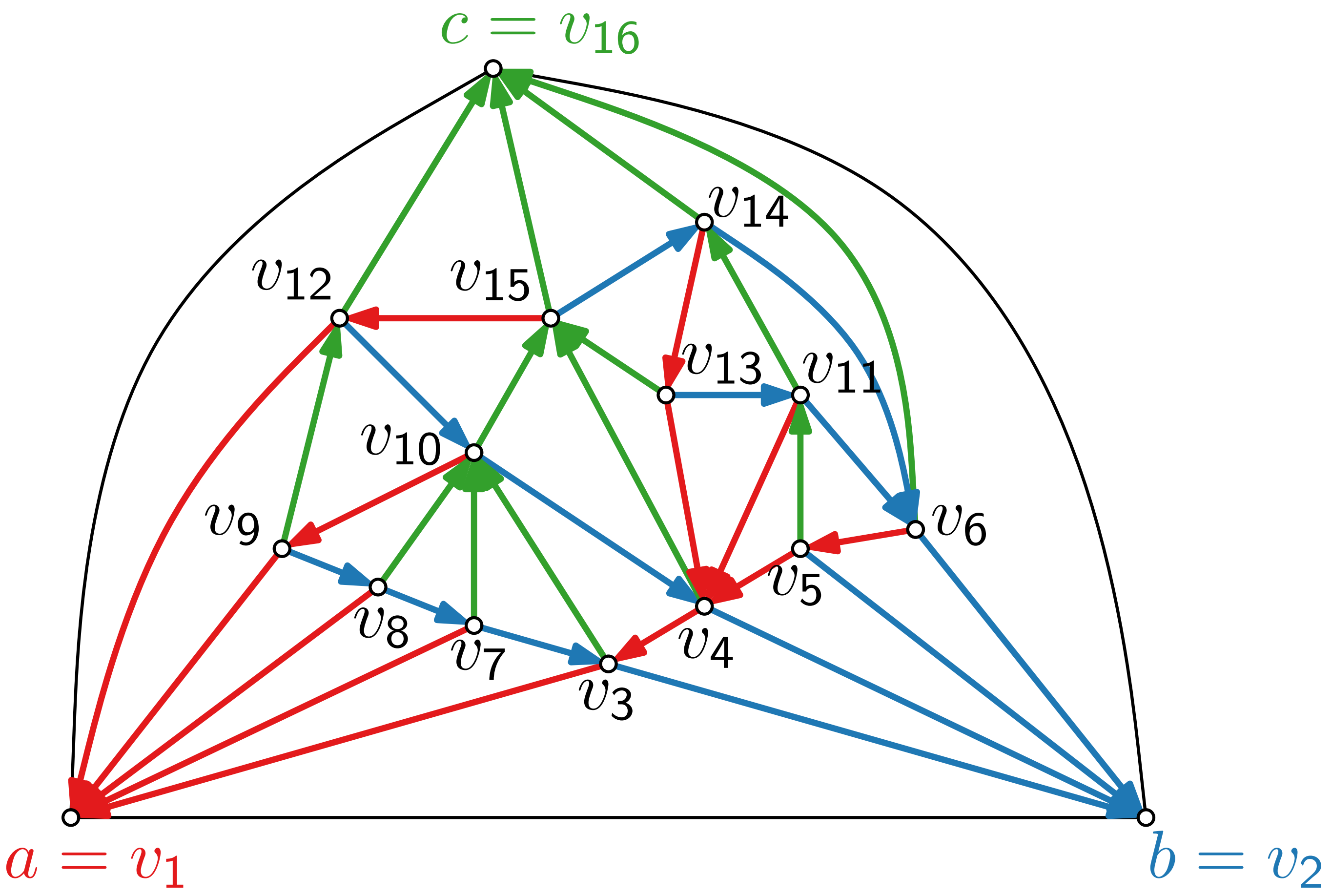

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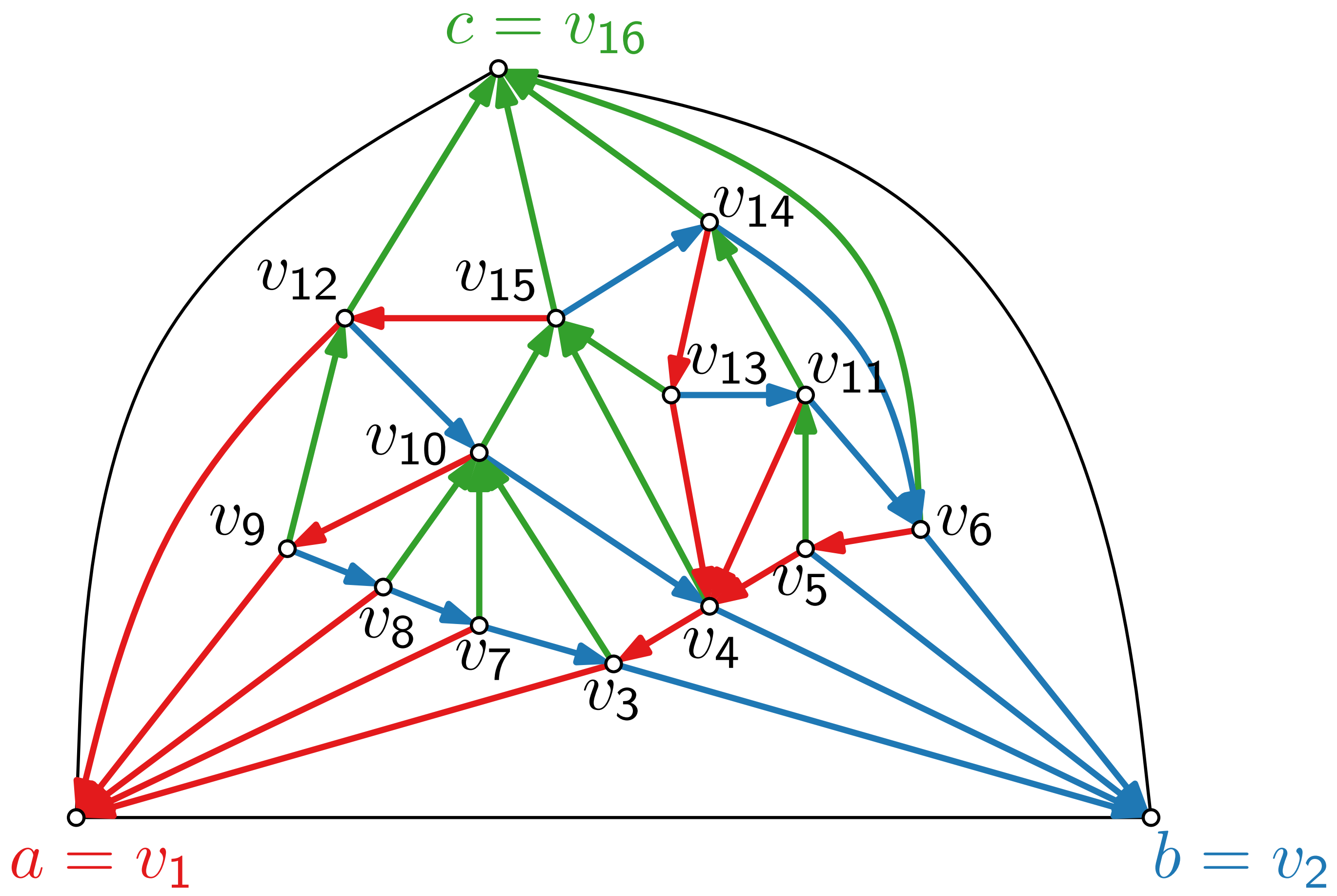

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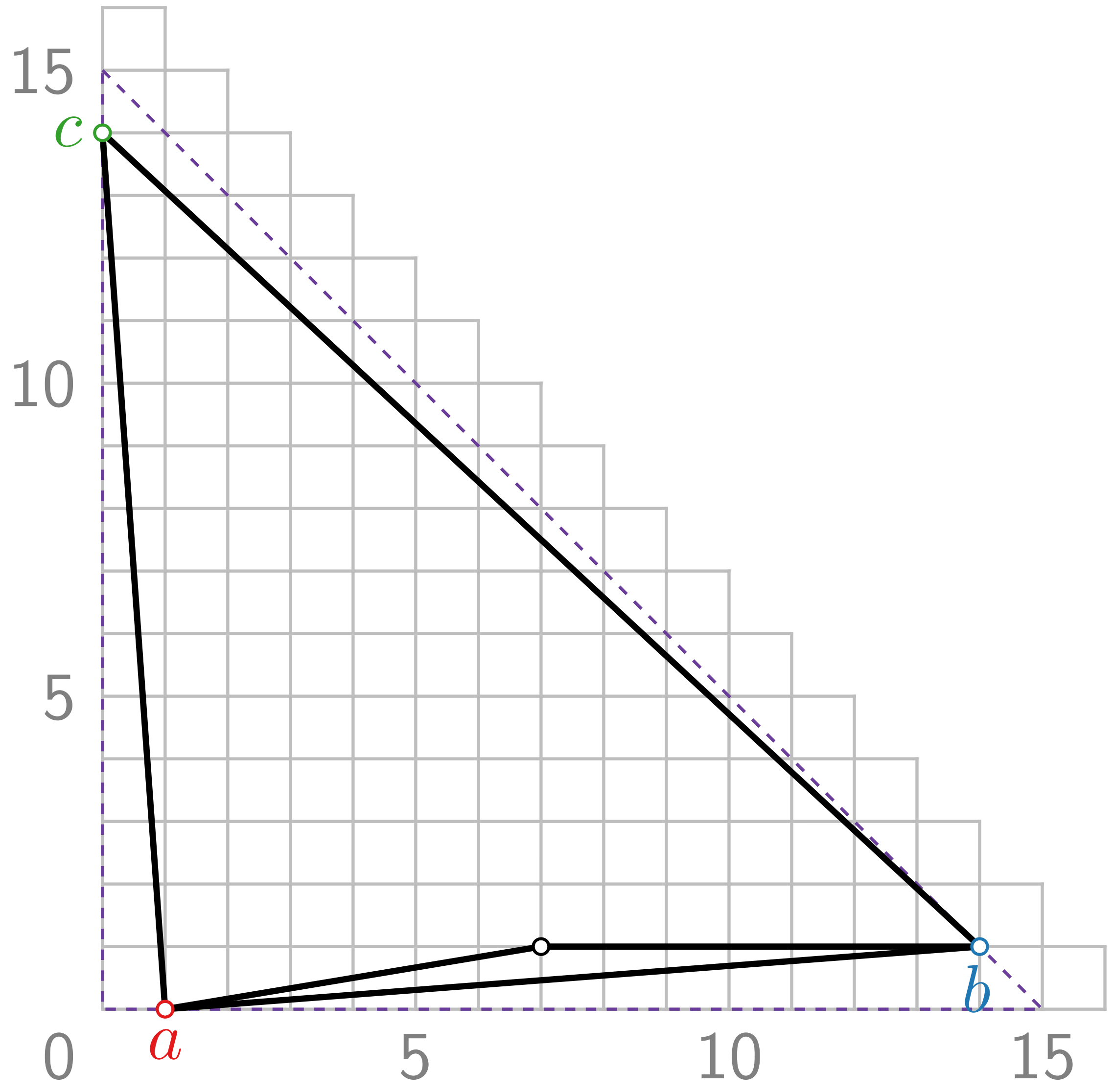

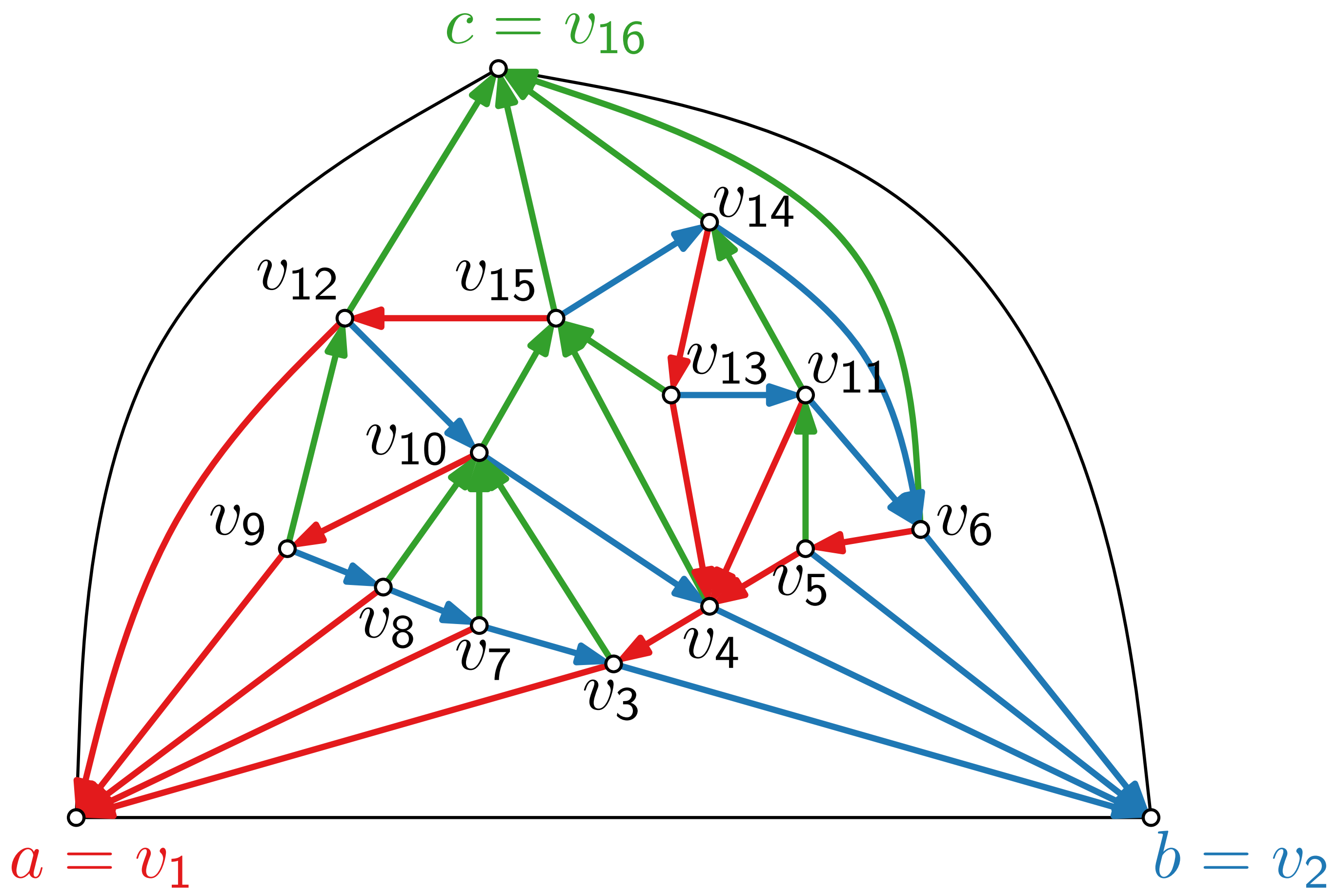

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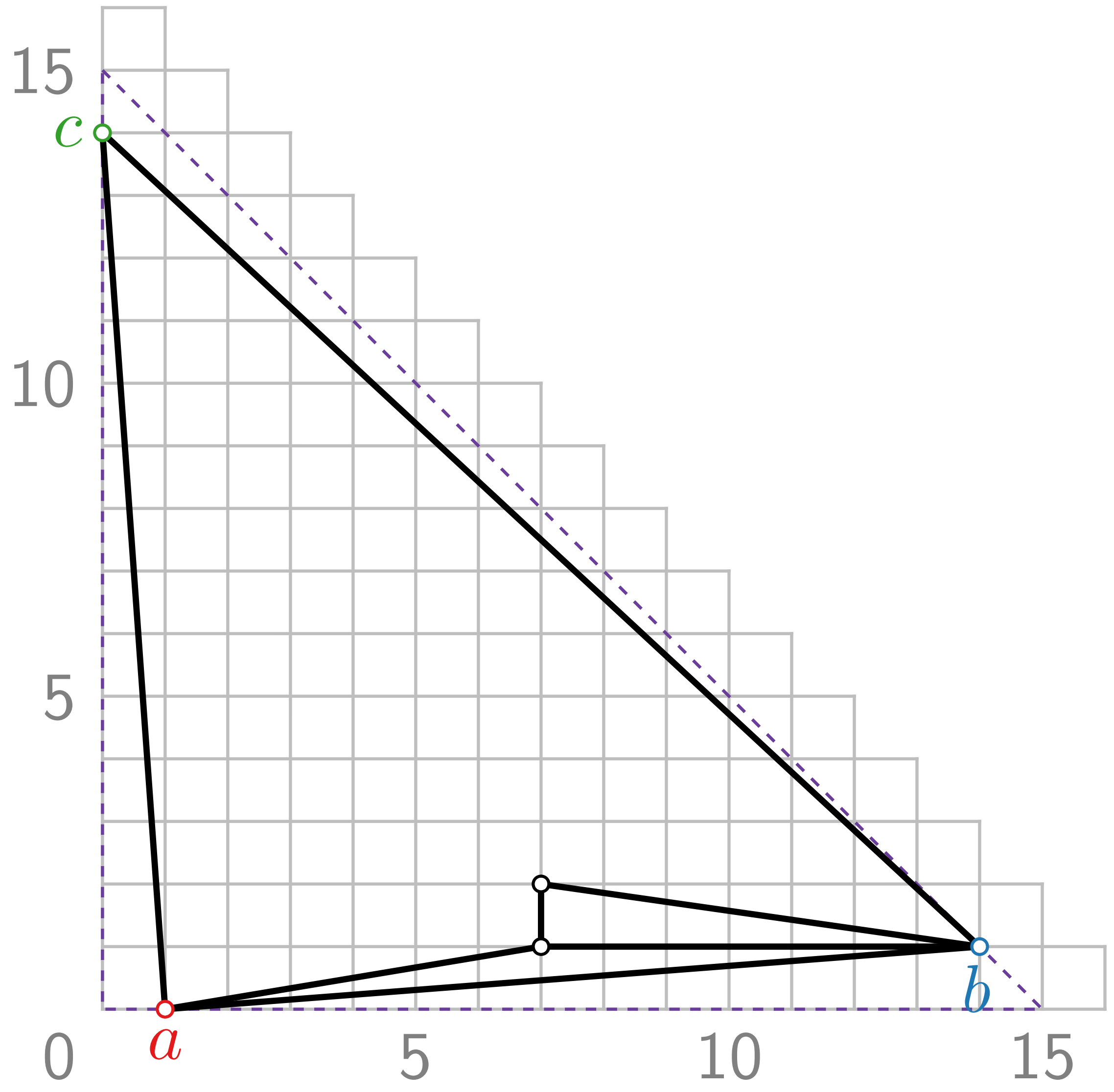

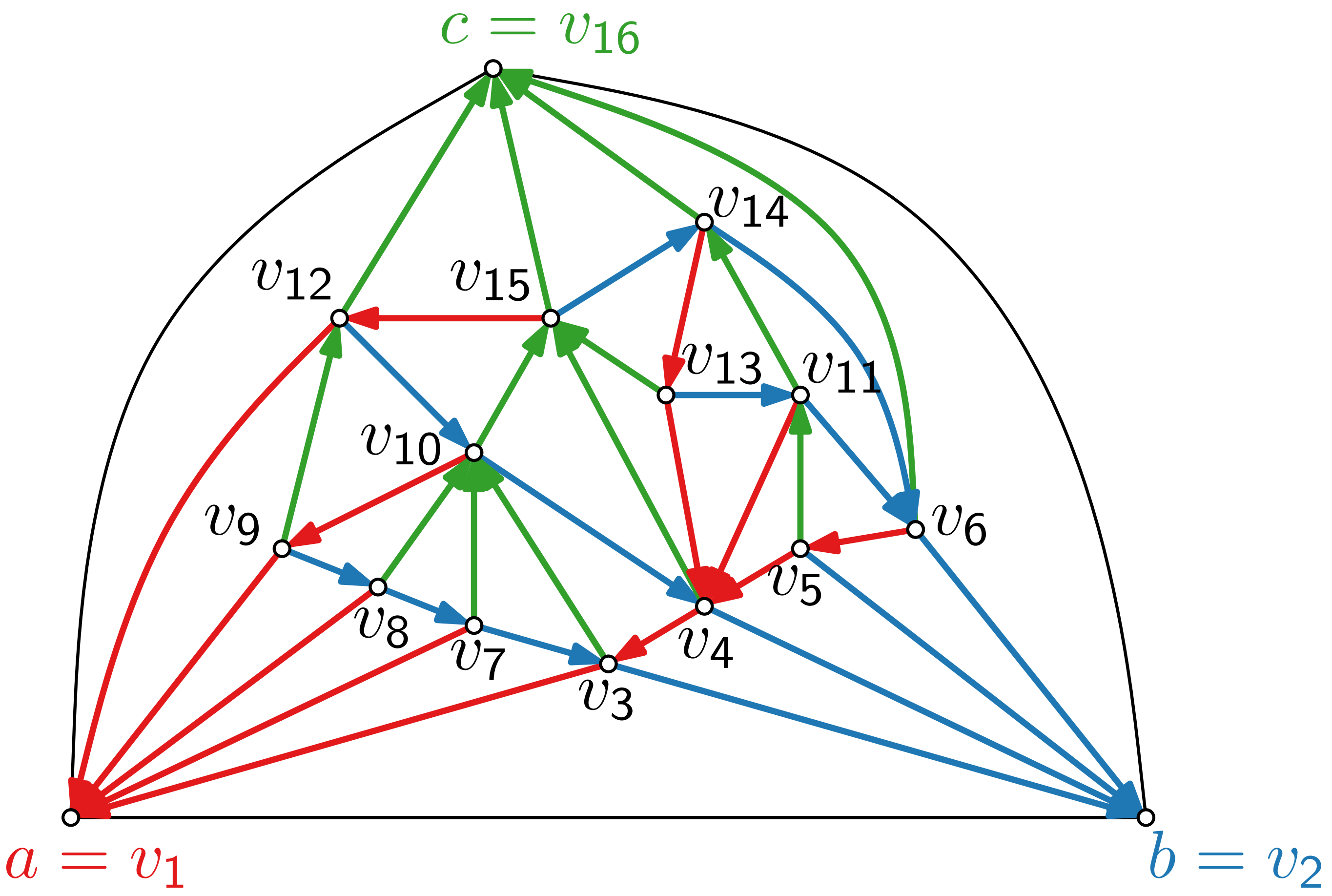

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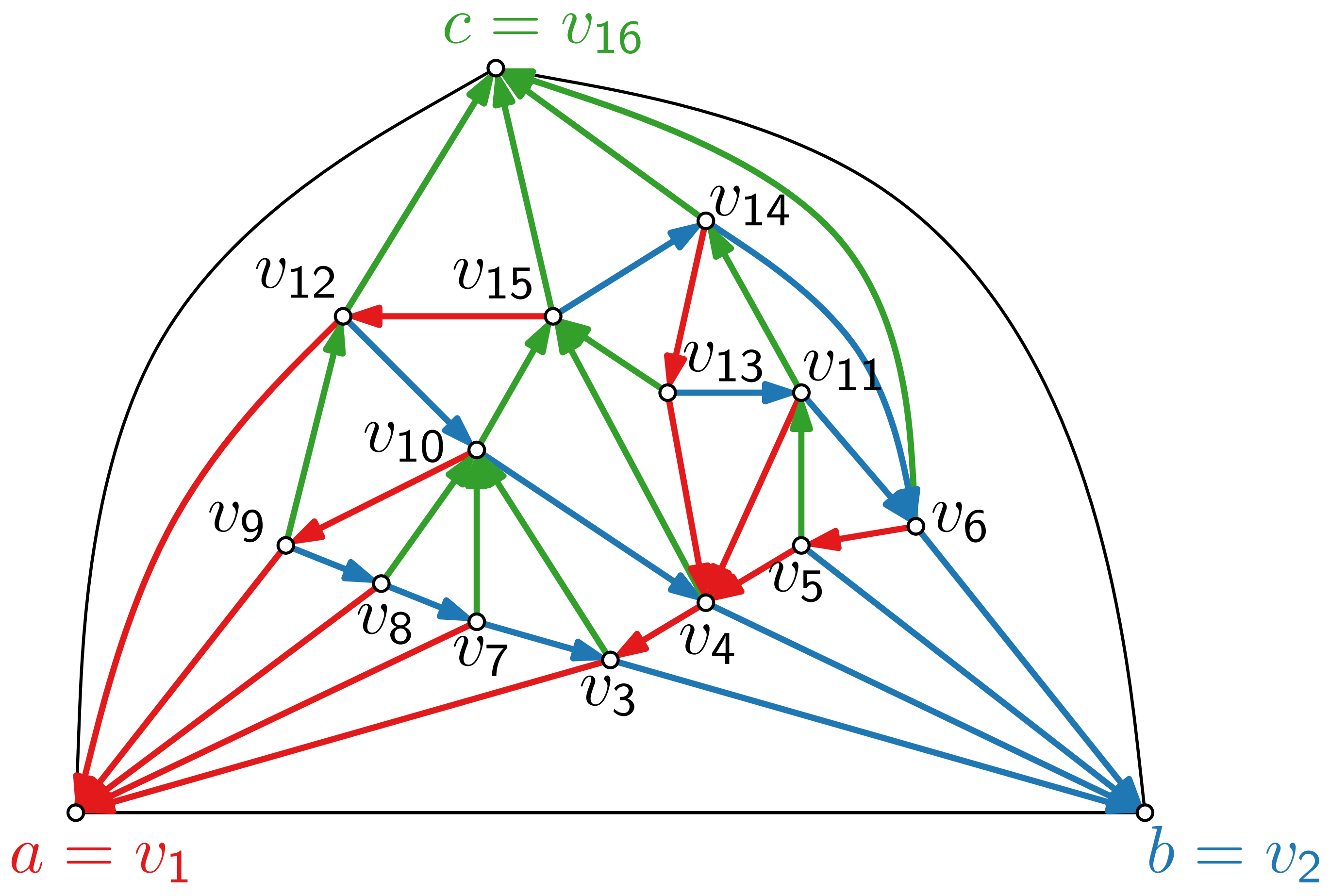

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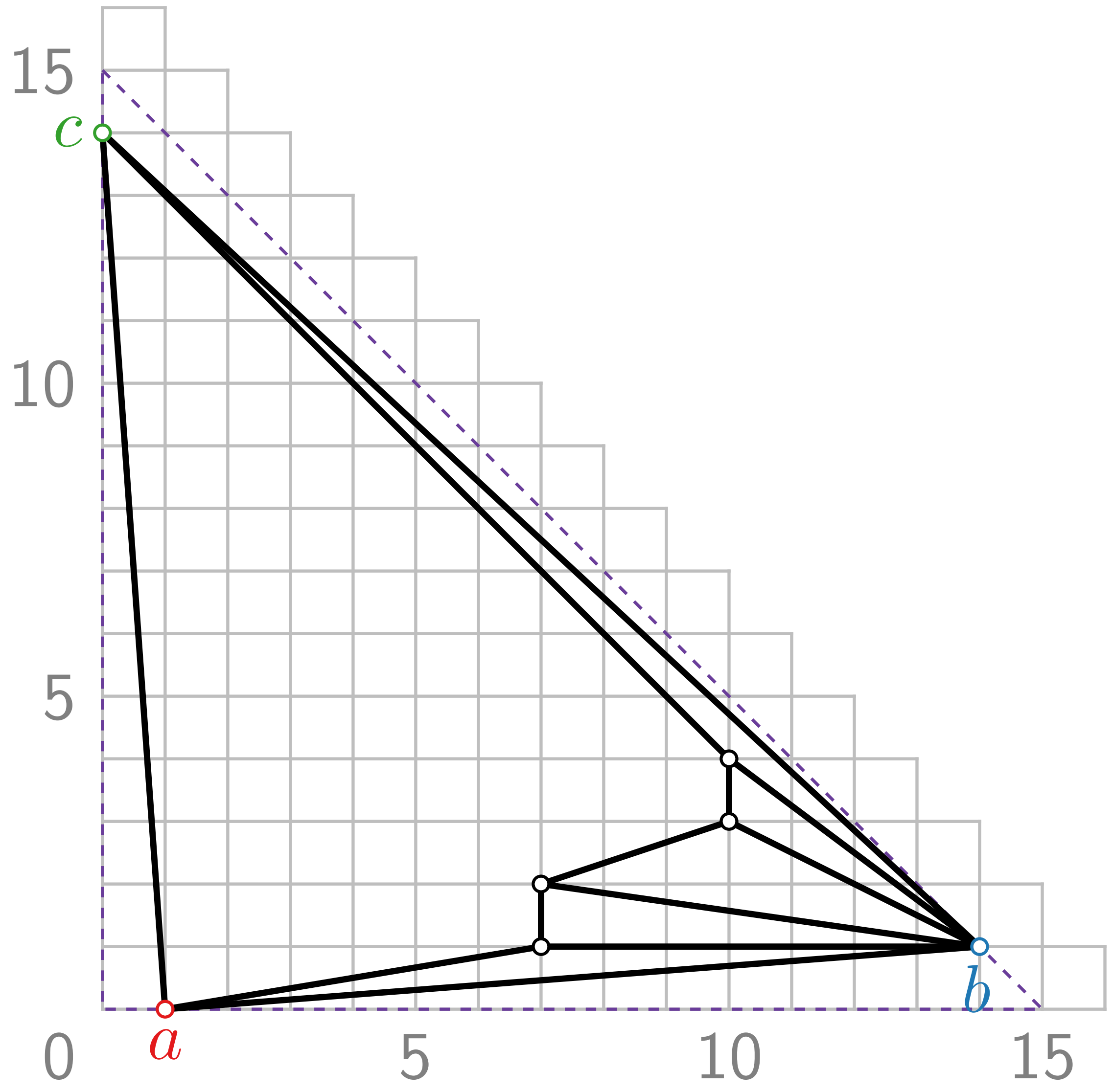

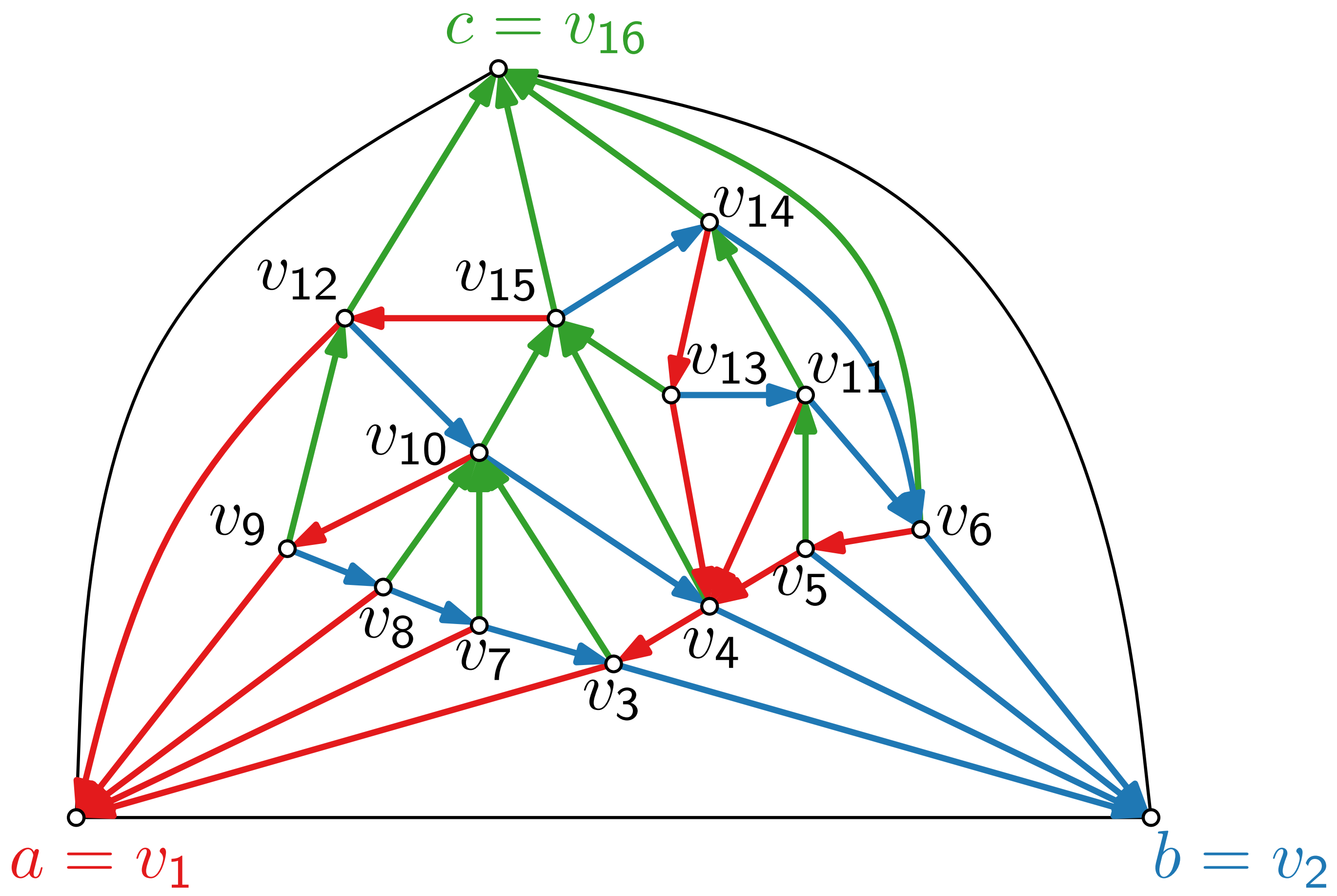

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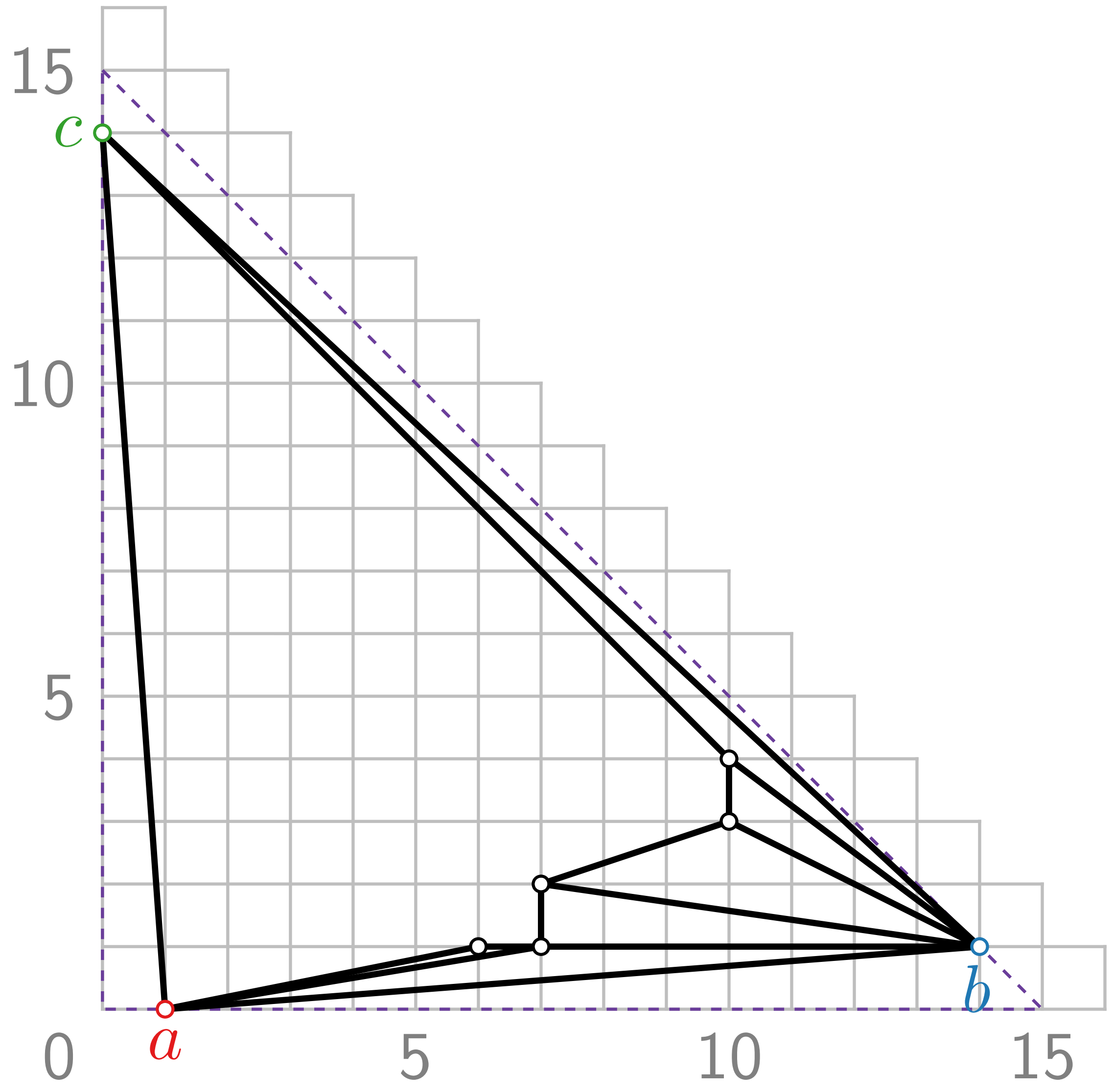

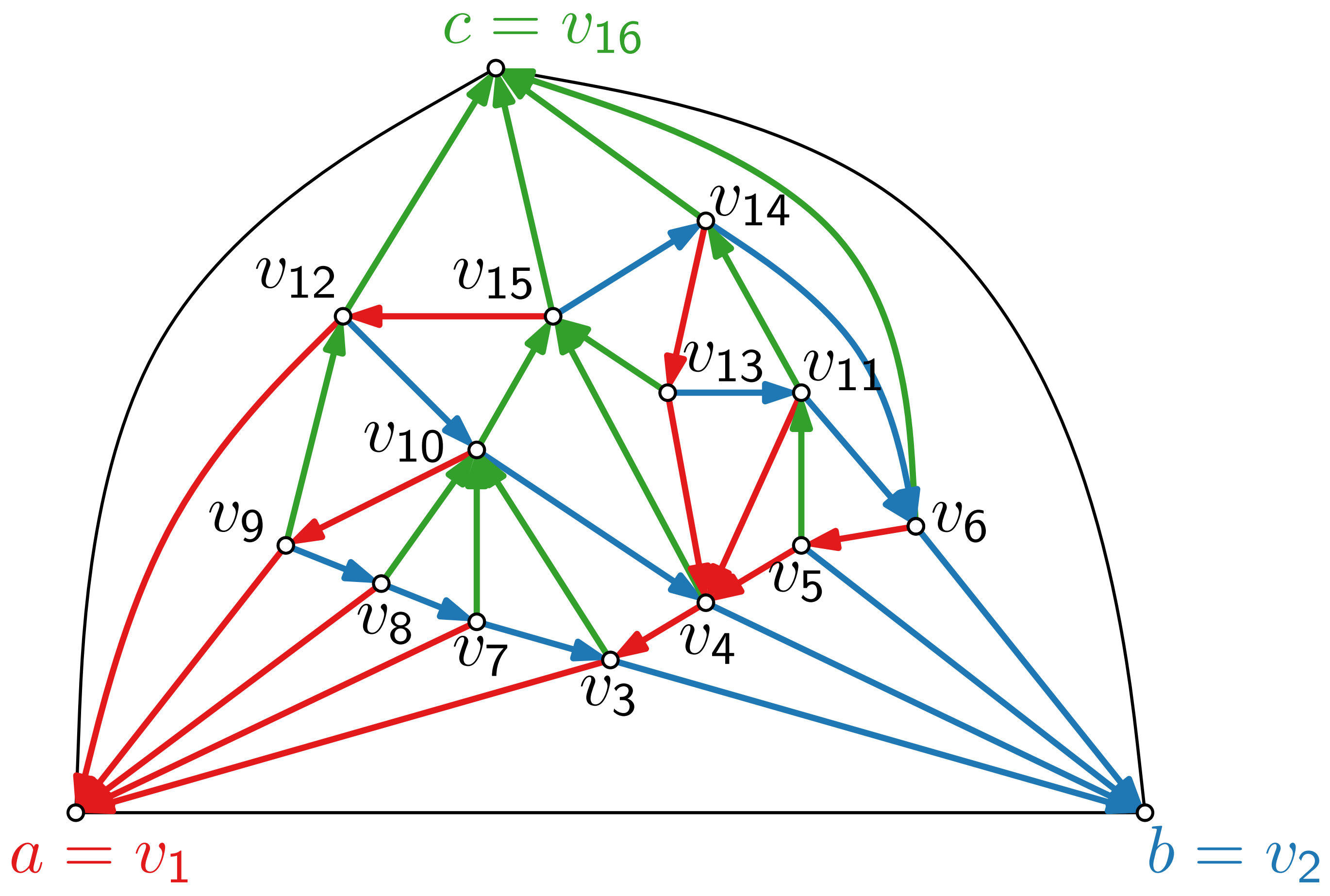

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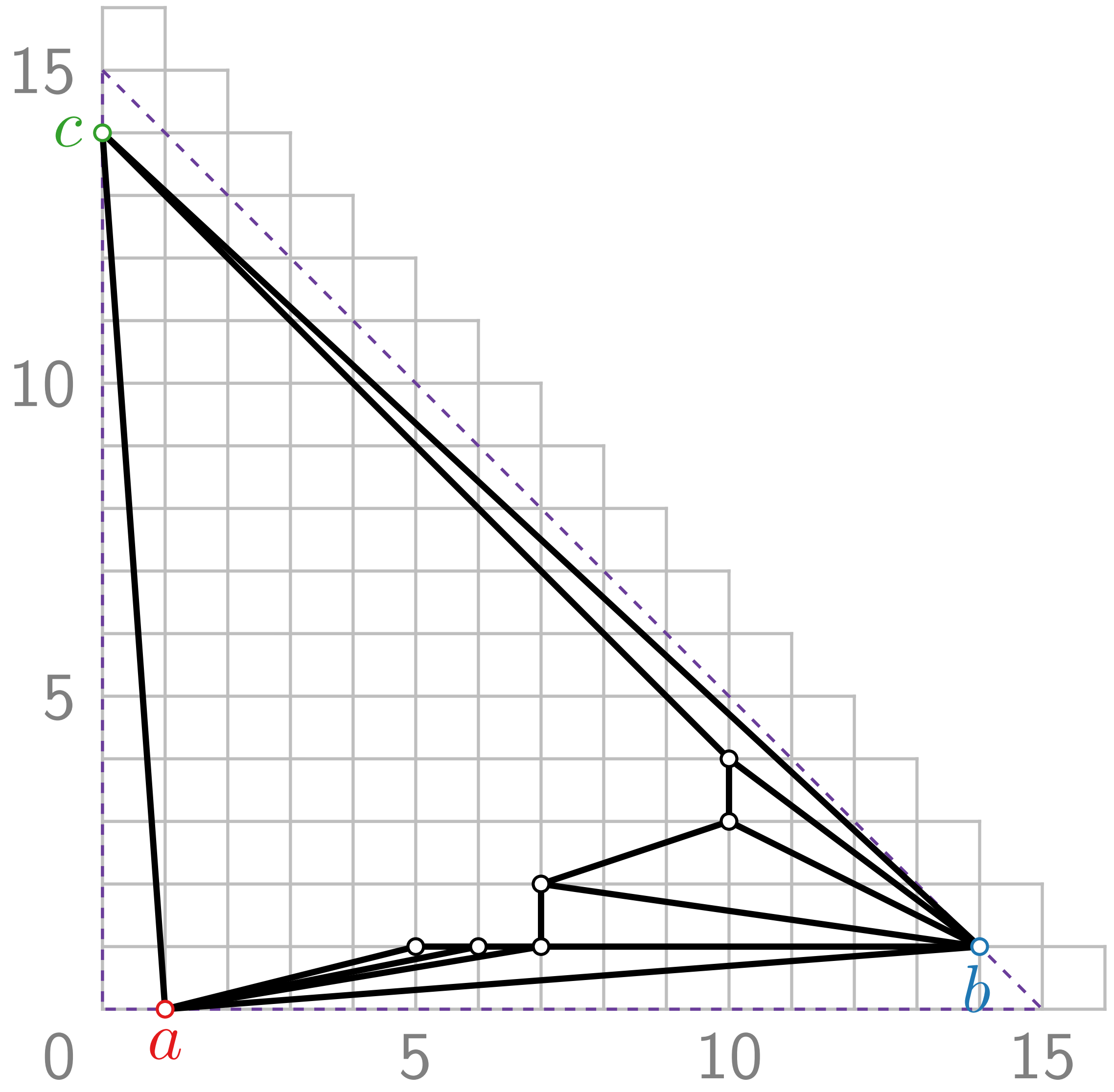

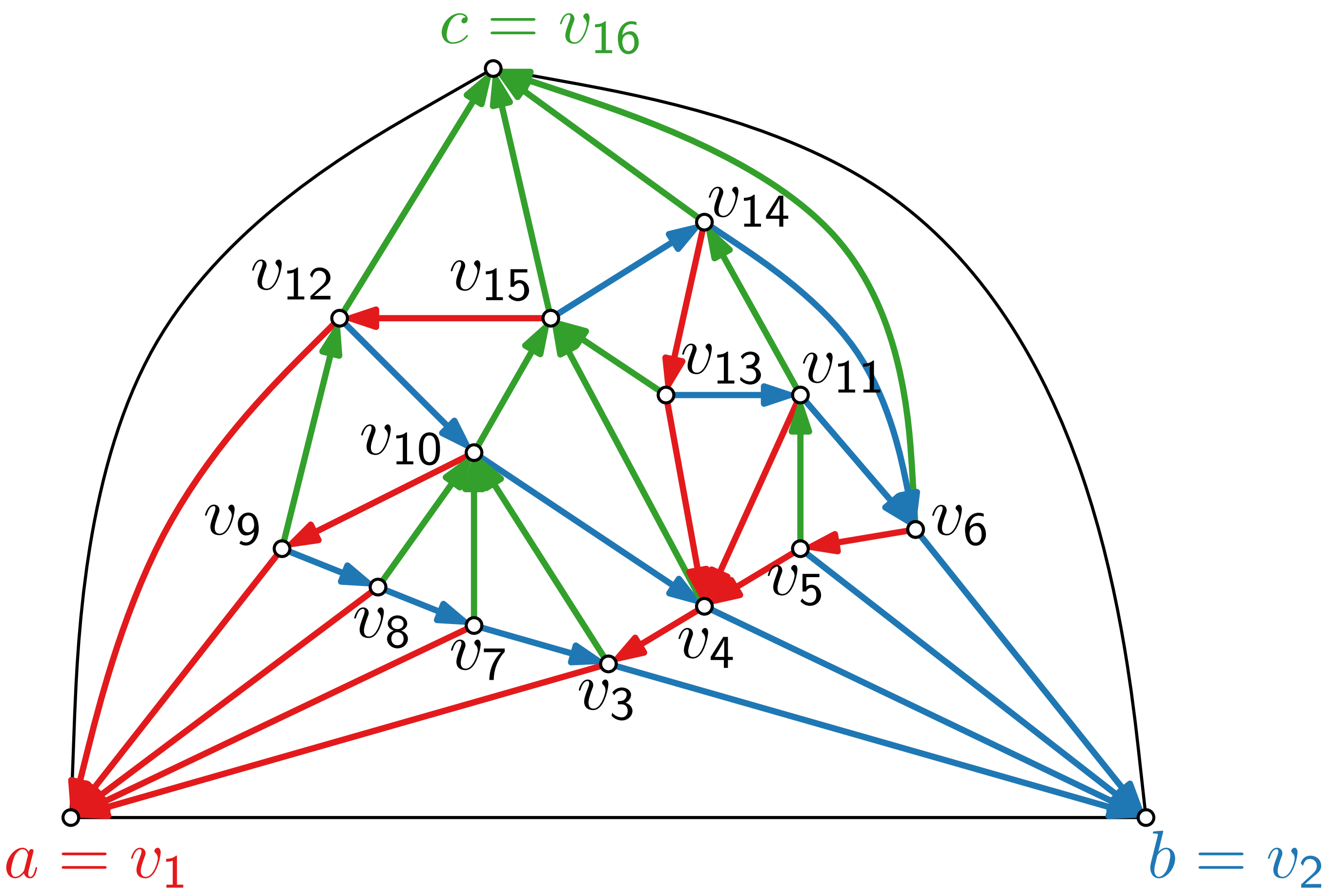

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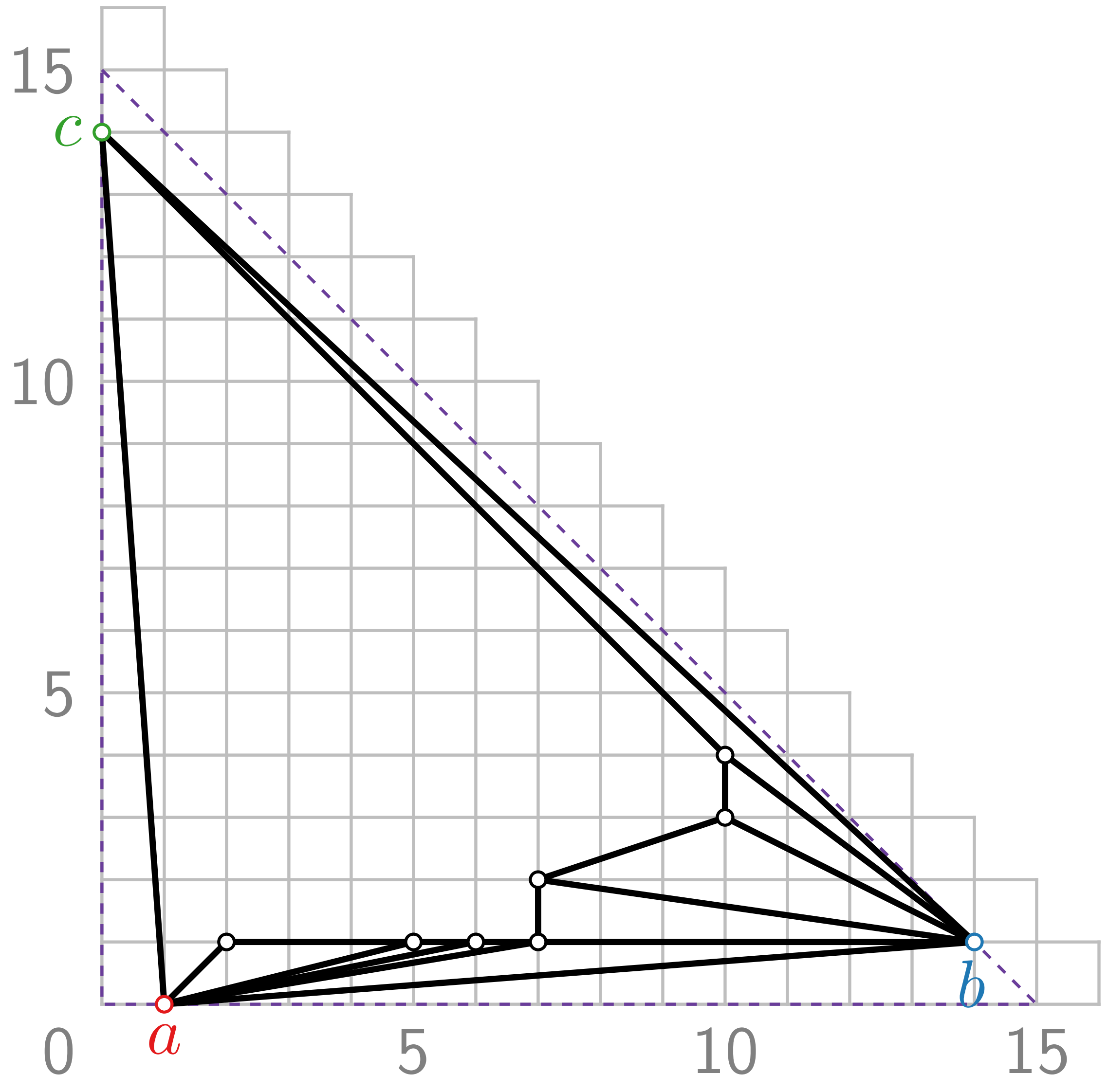

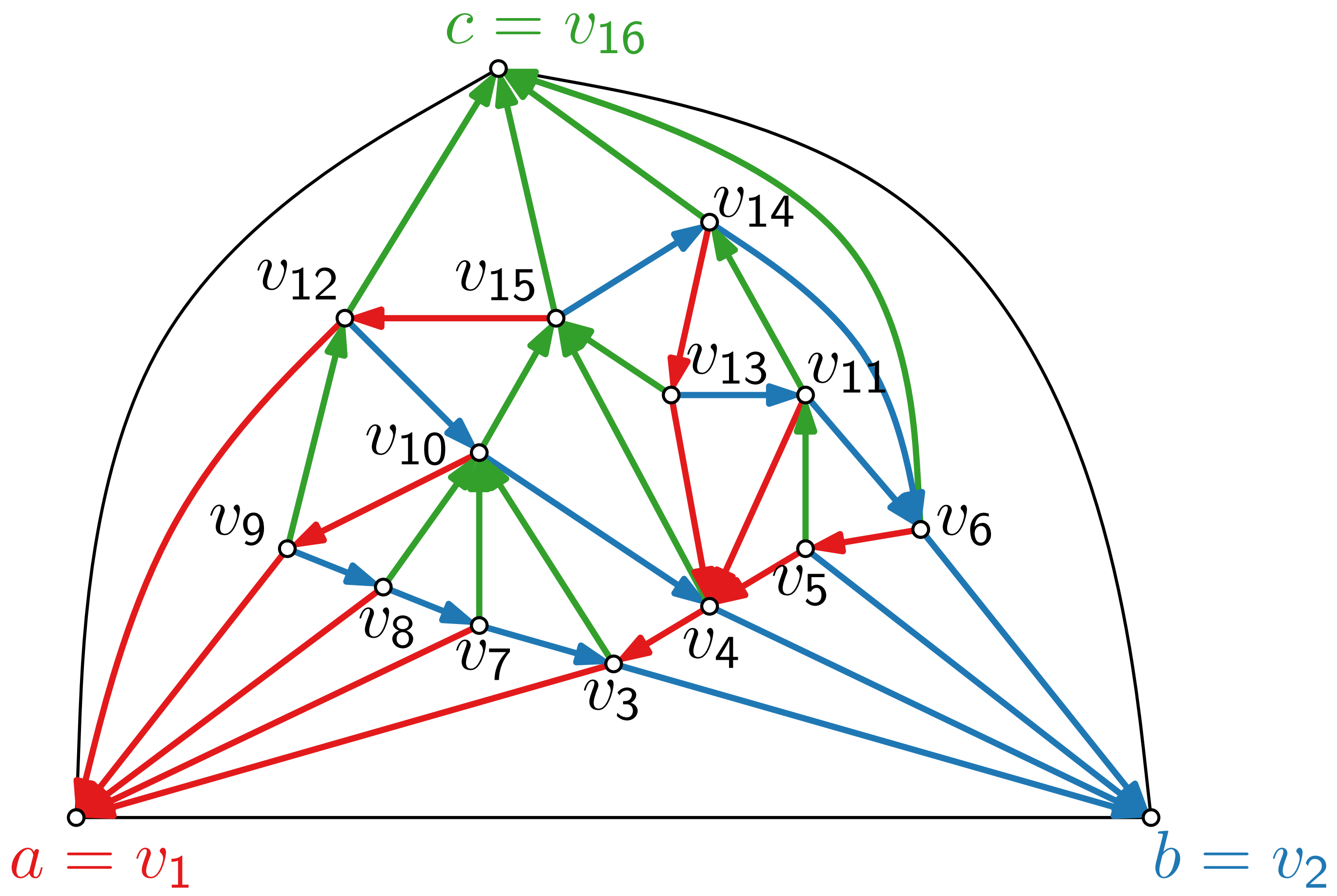

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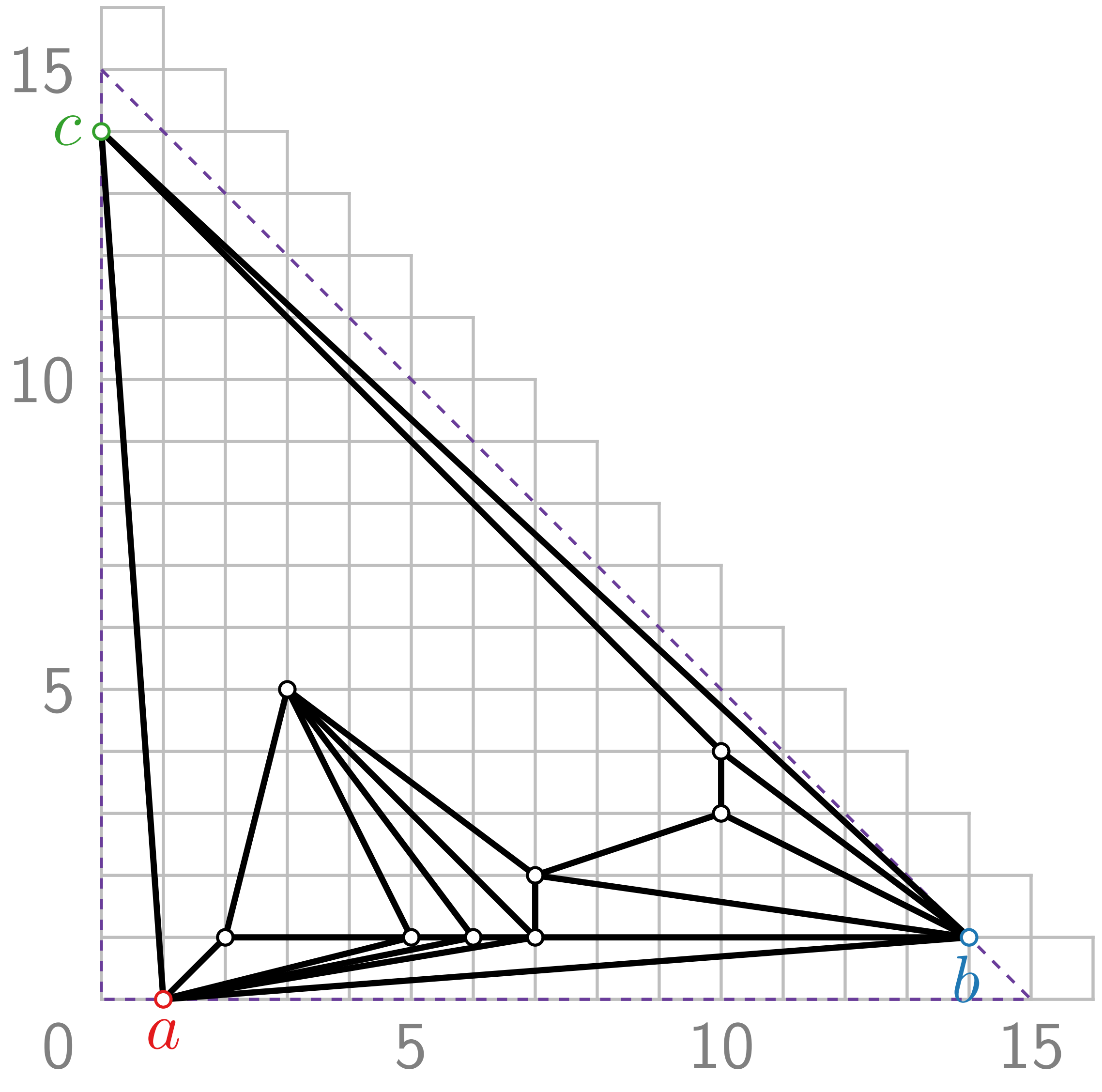

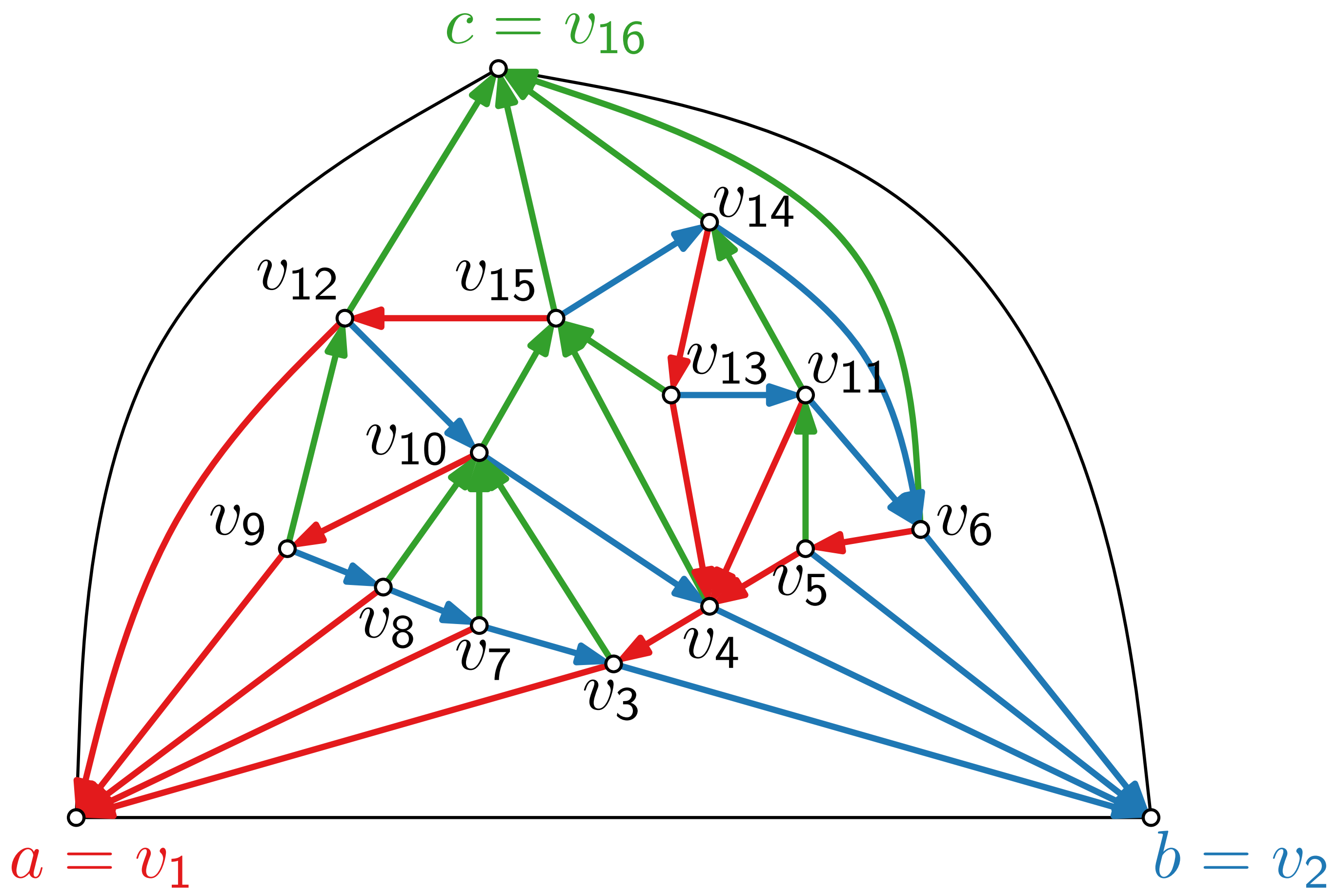

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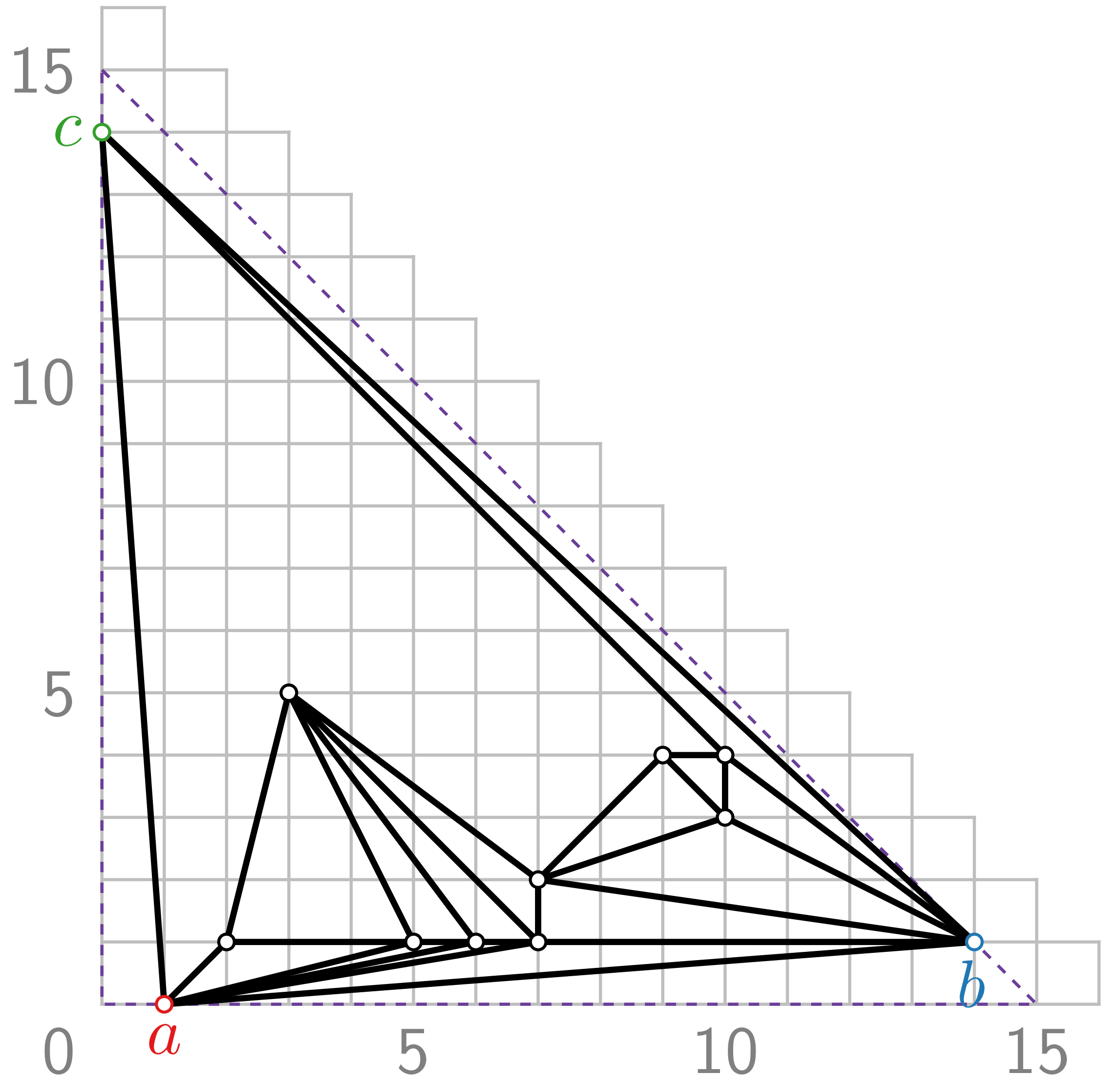

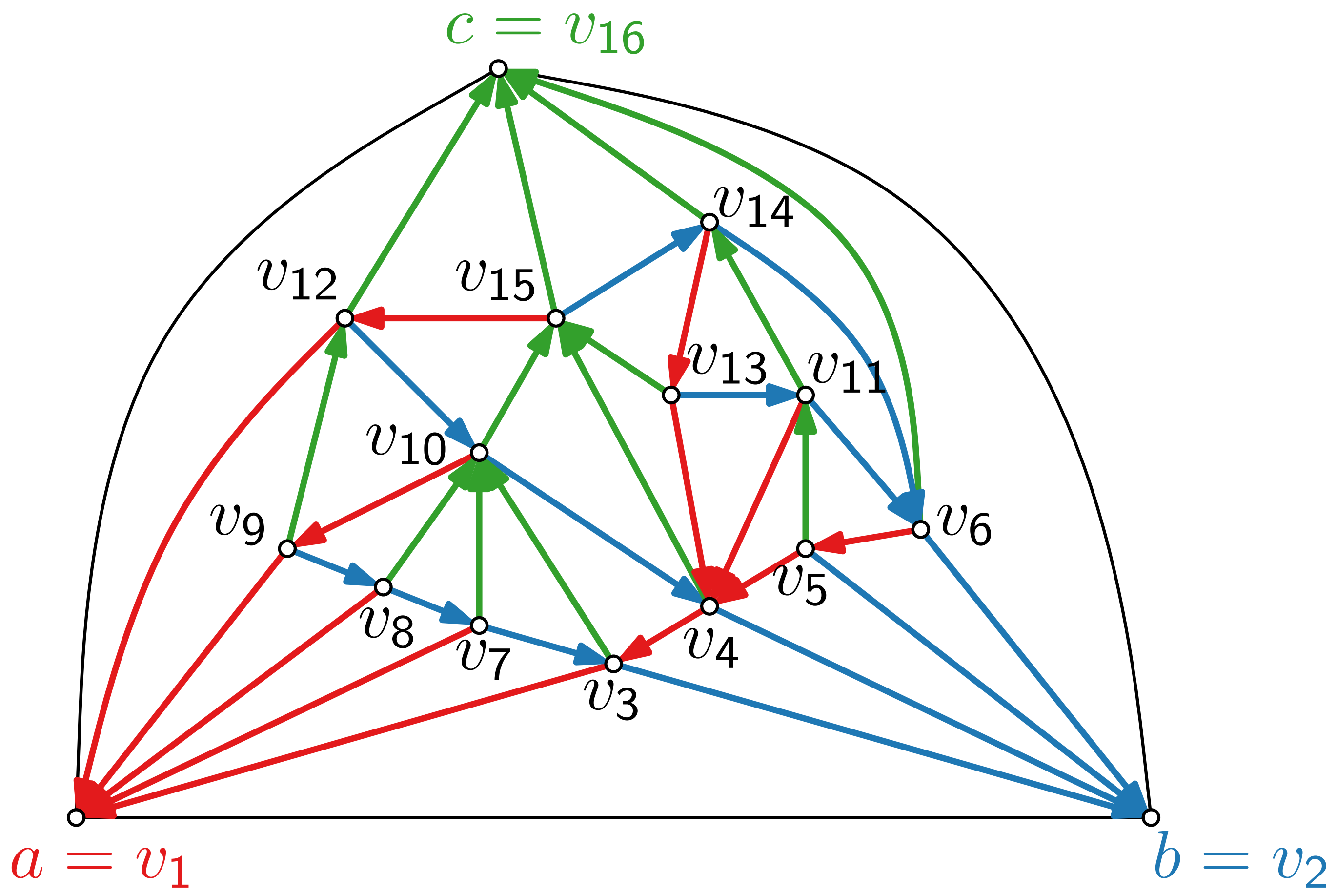

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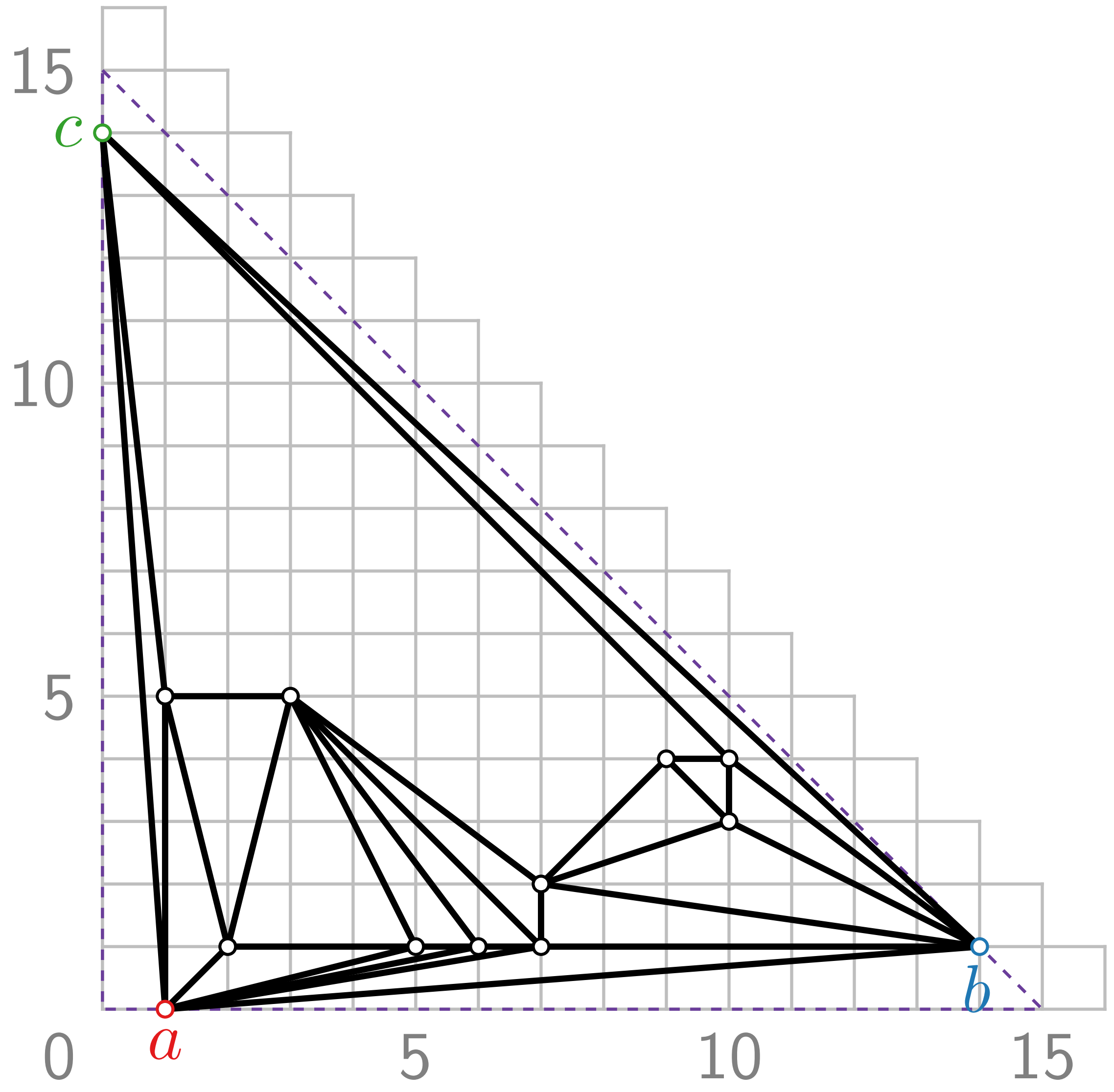

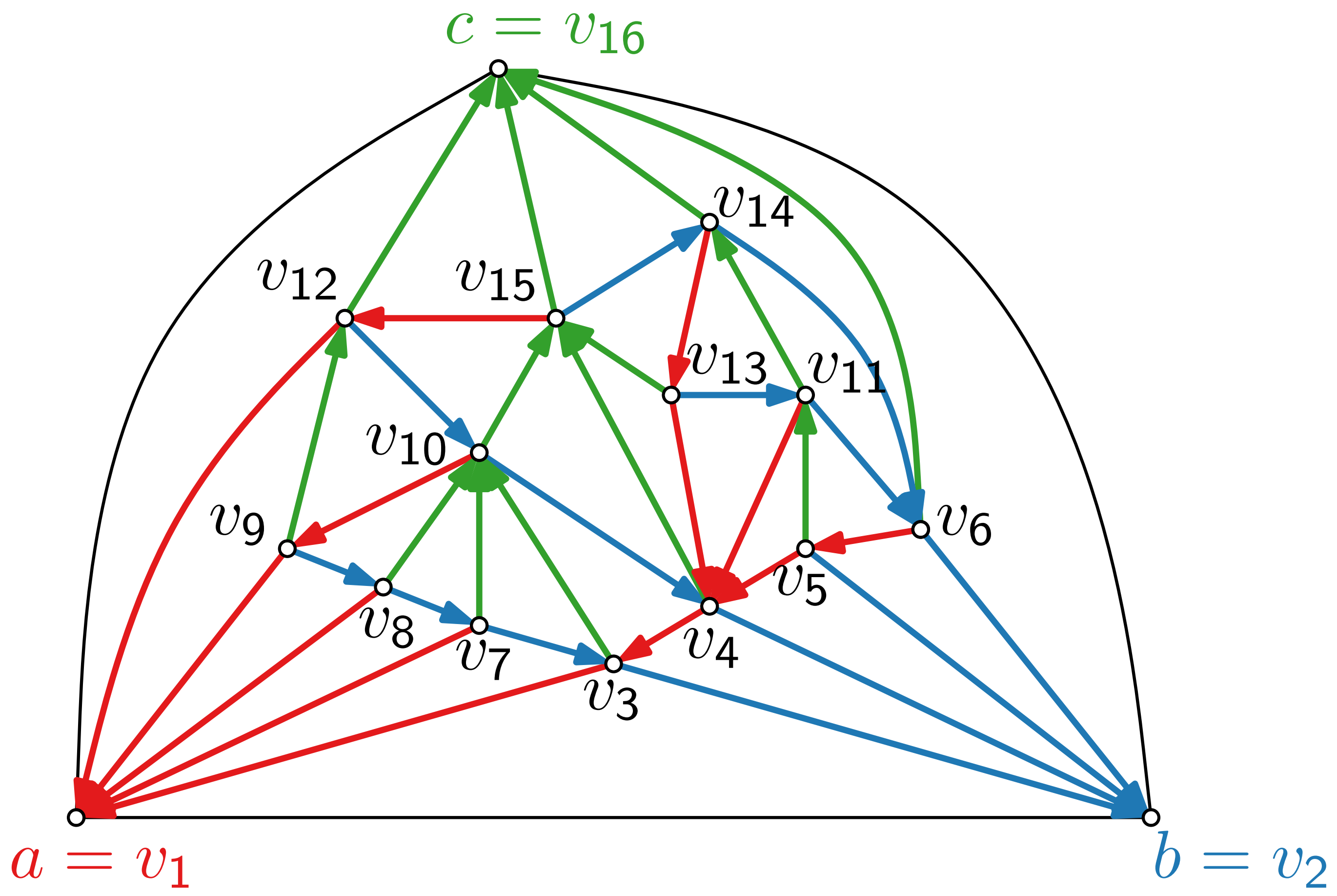

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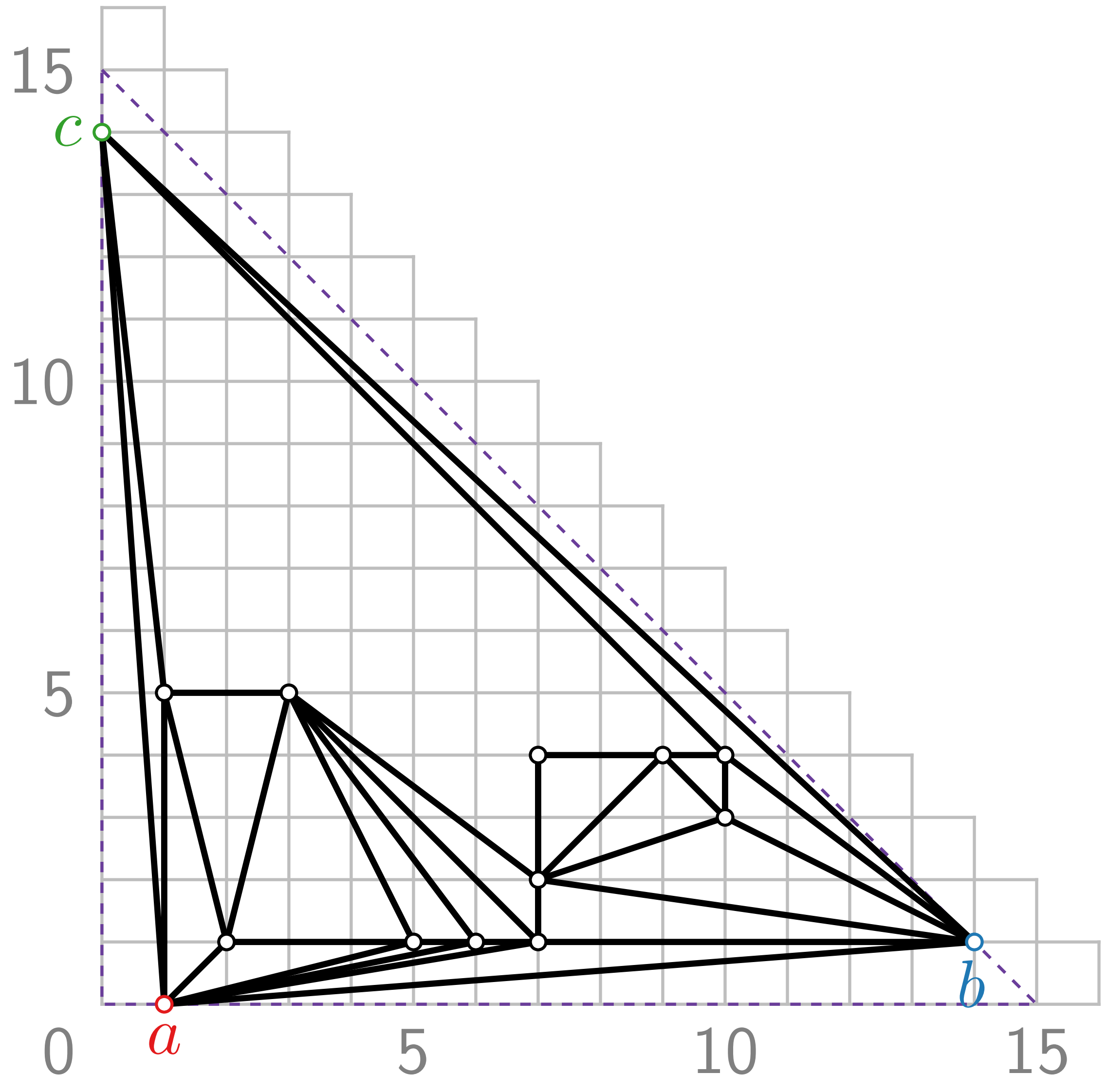

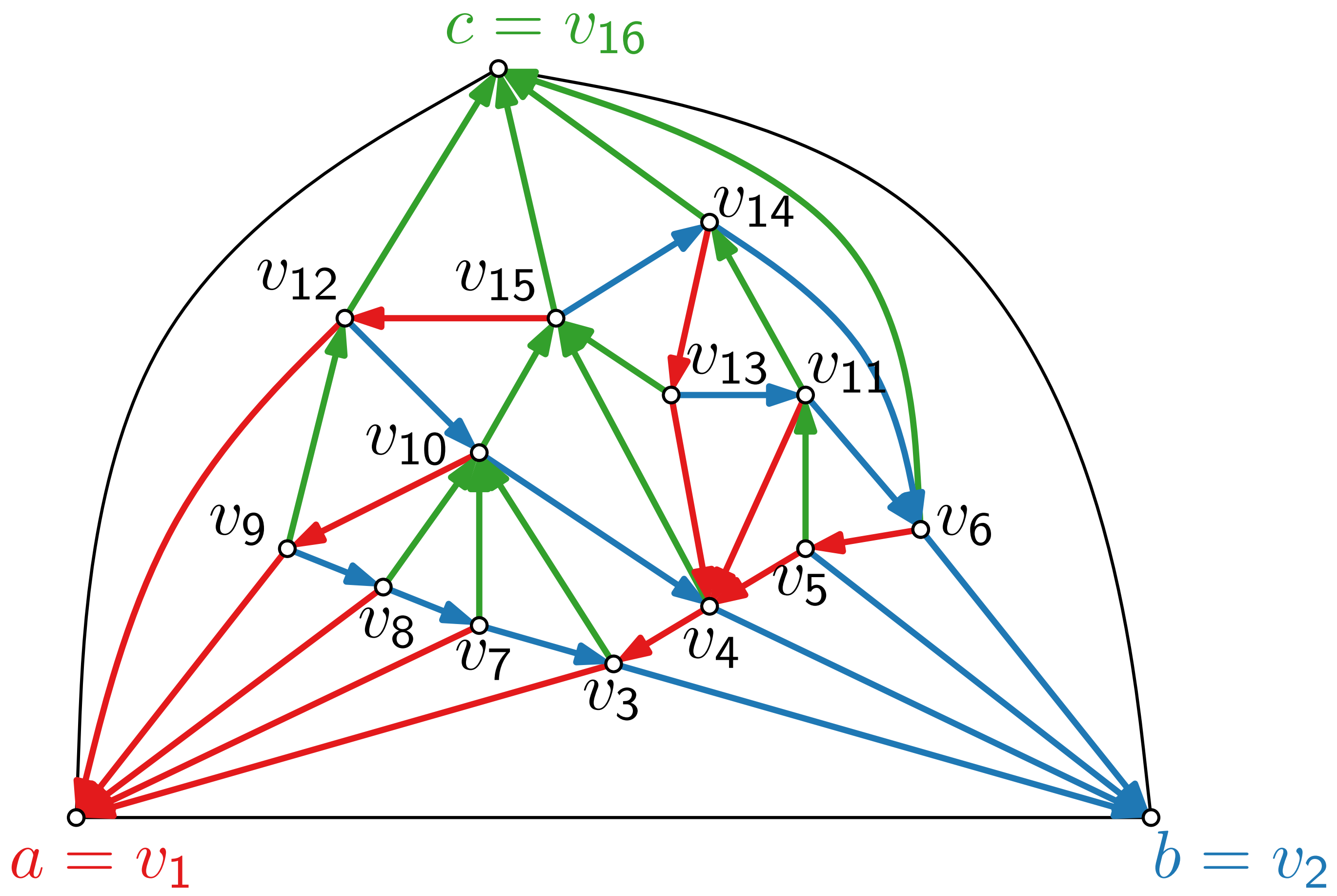

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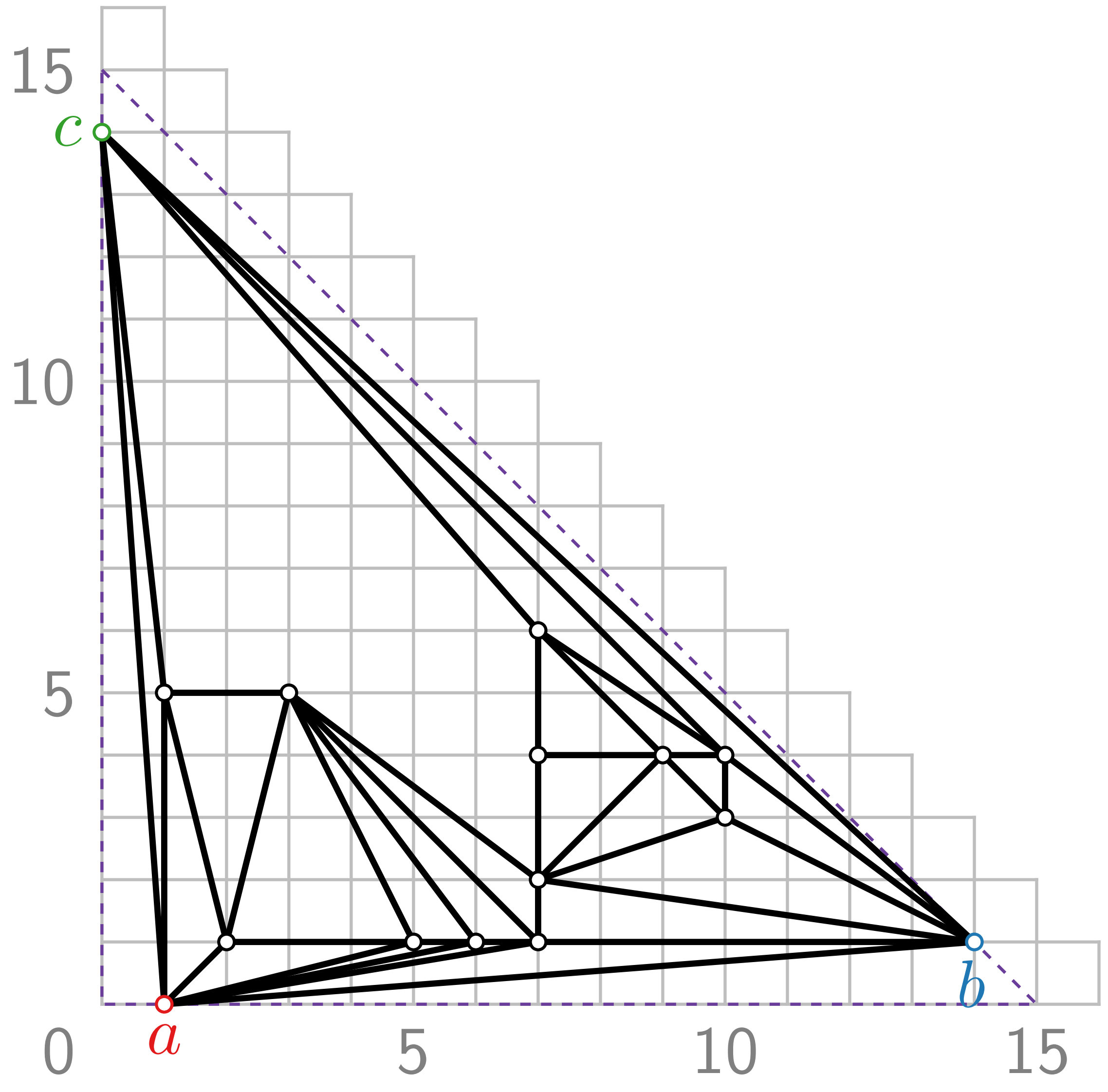

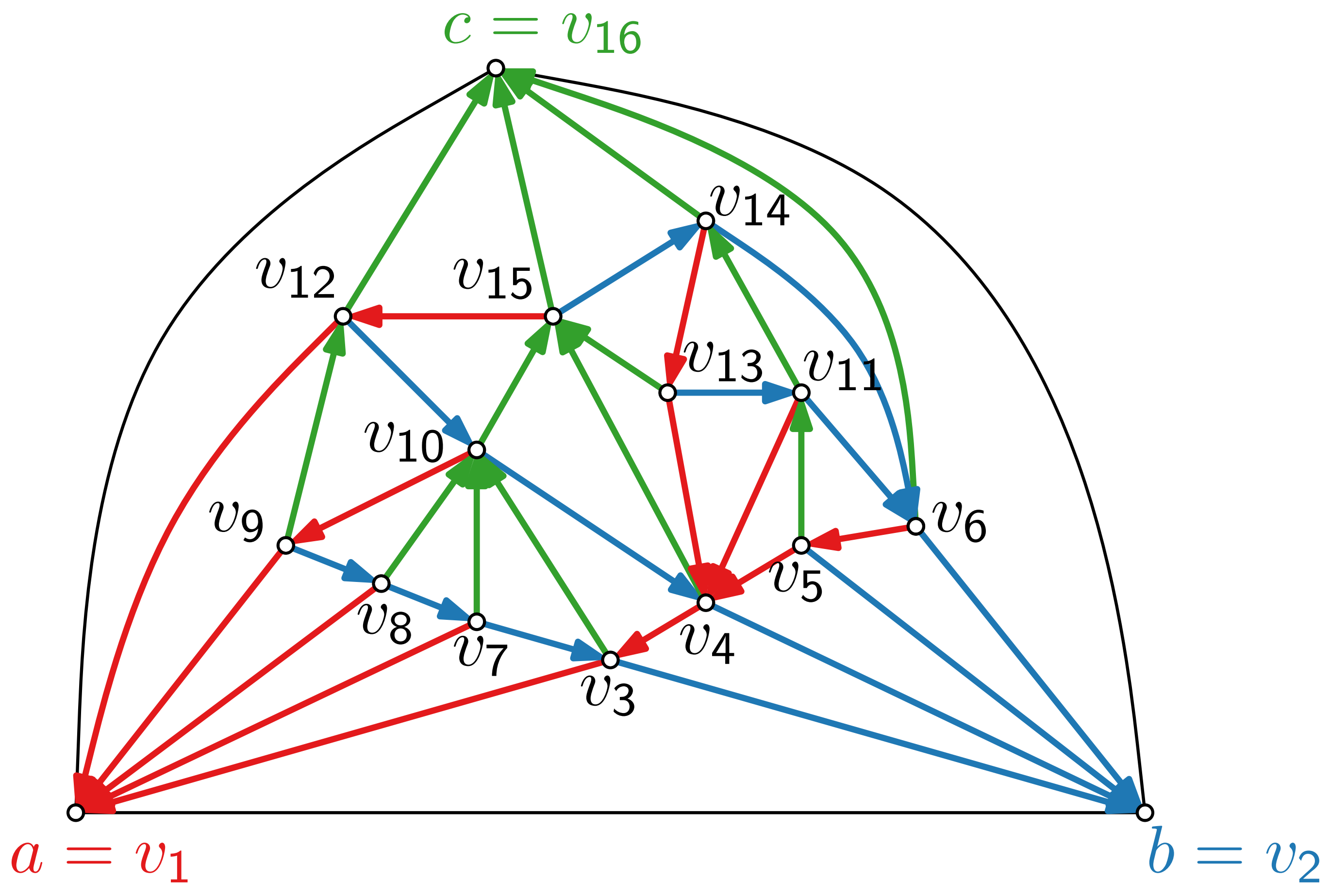

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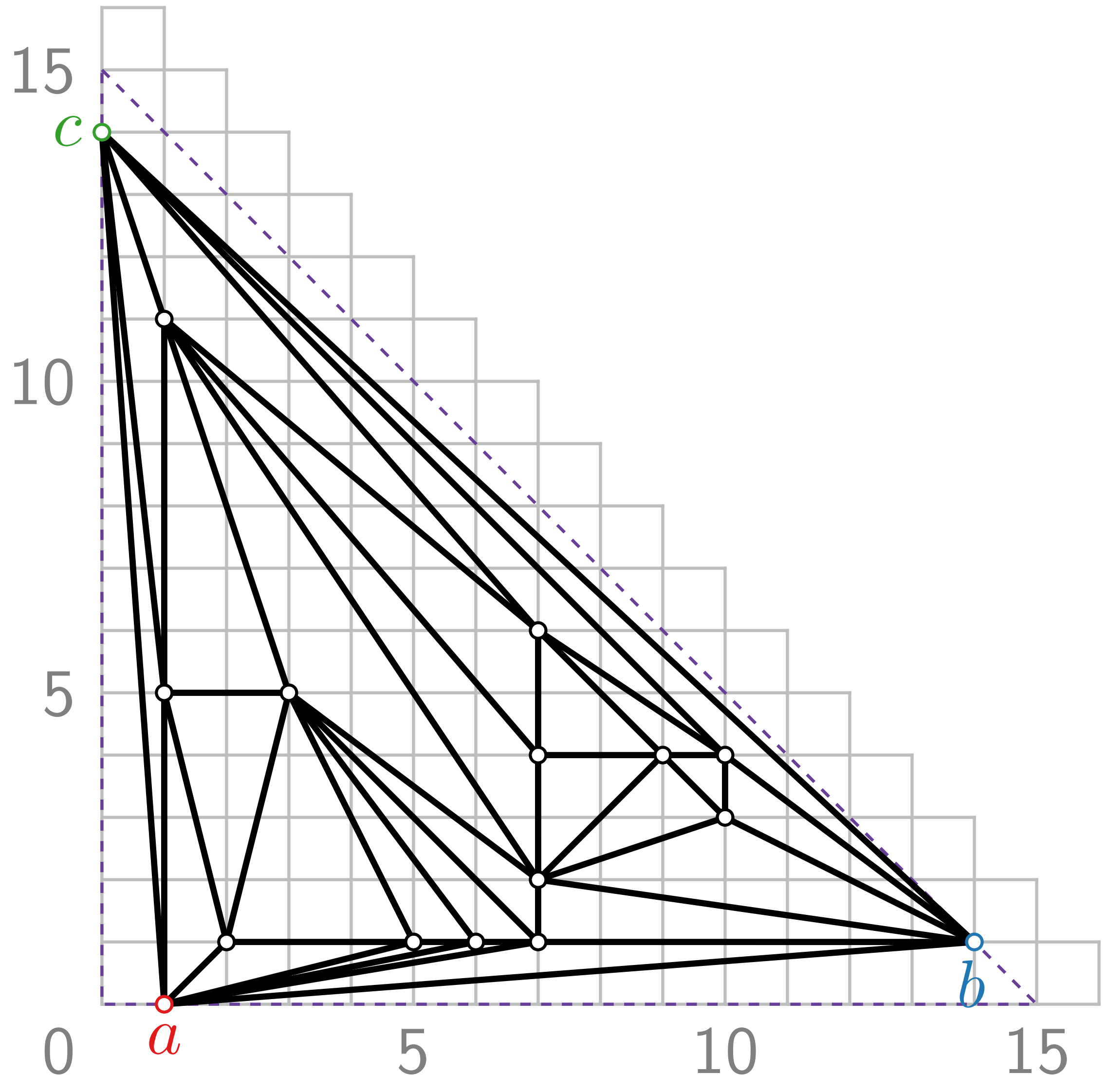

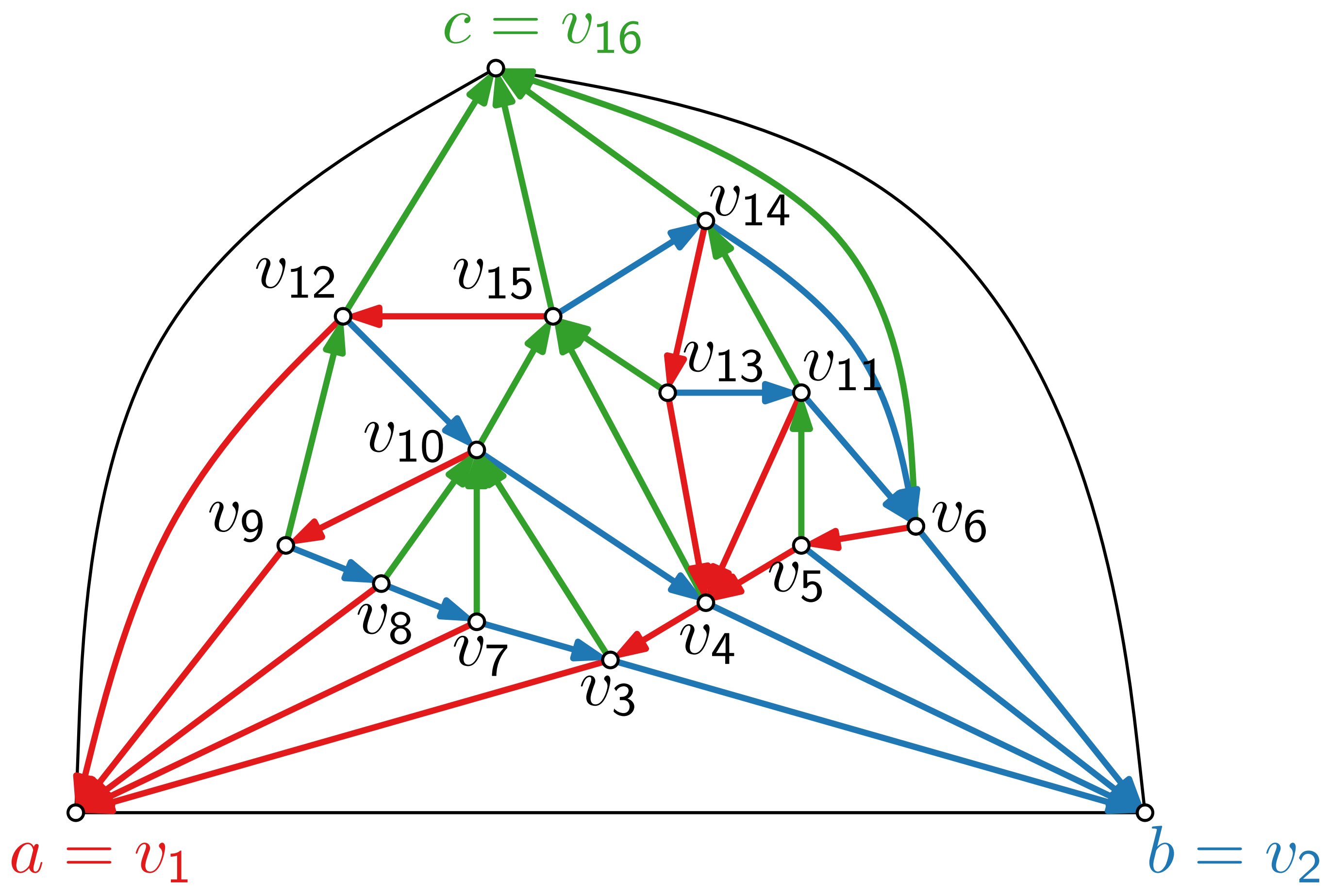

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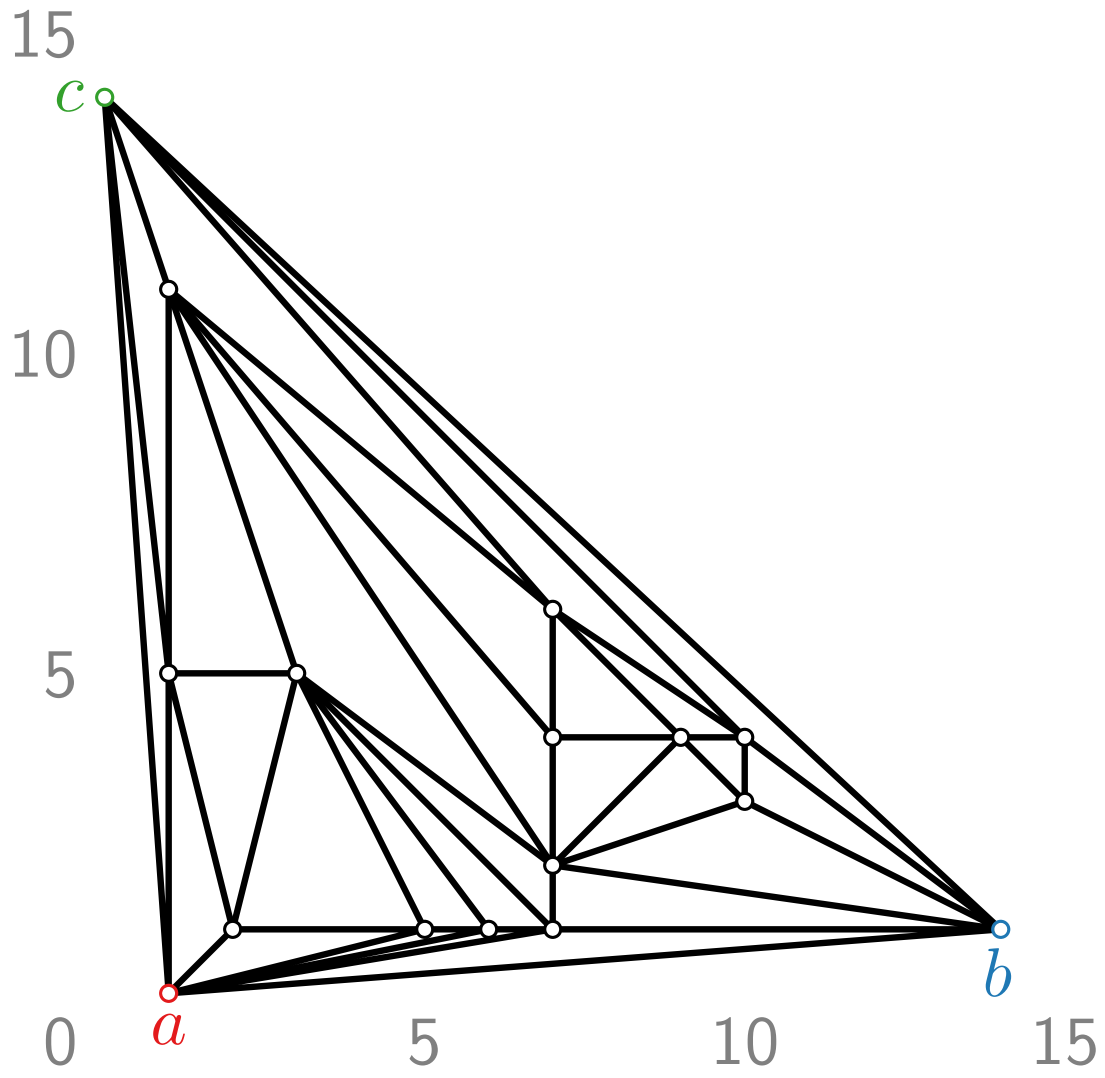

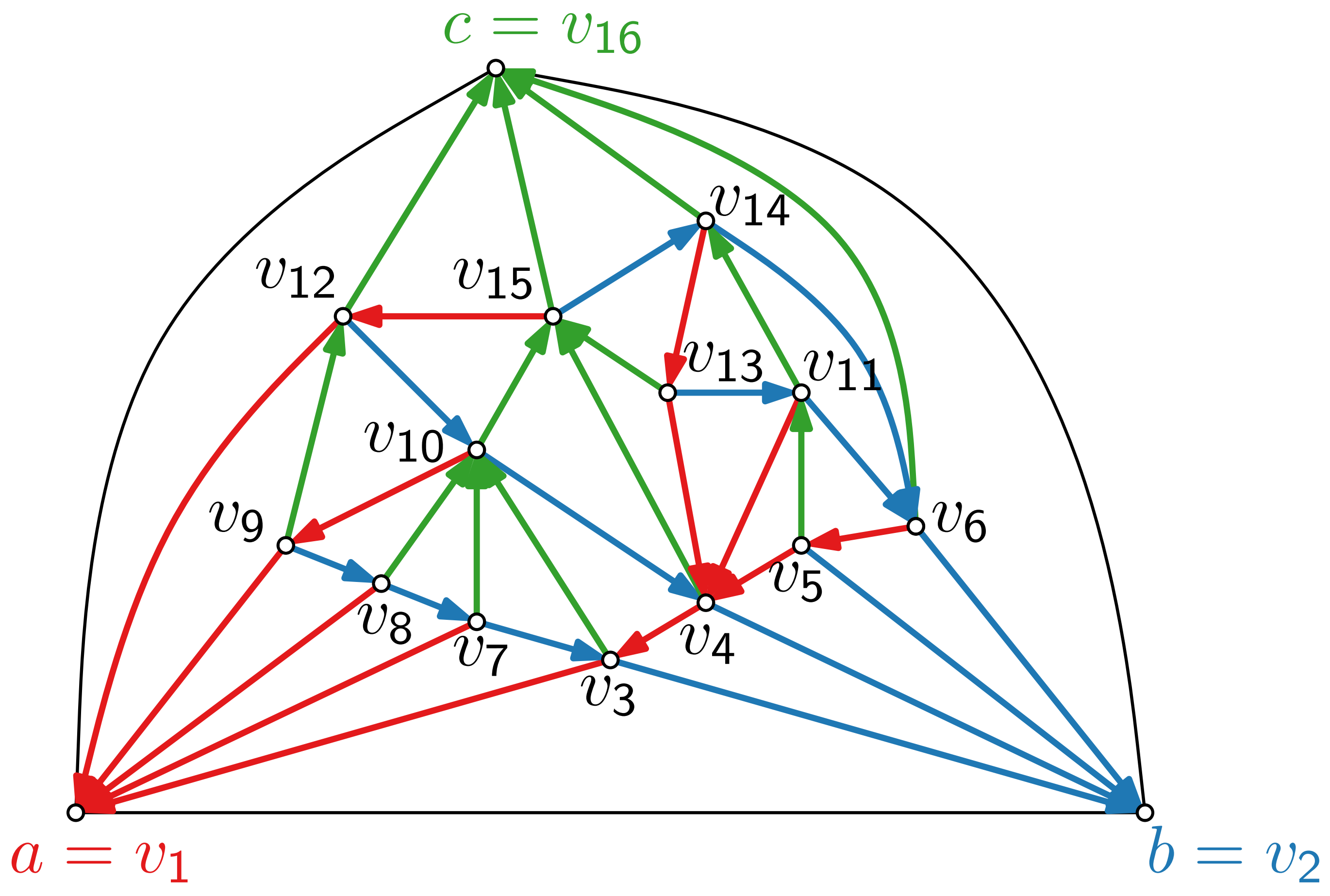

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Results & Variations

Theorem. [De Fraysseix, Pach, Pollack '90]
Every n-vertex planar graph has a planar straight-line drawing of size $(2n-4) \times (n-2)$. Such a drawing can be computed in $O(n)$ time.

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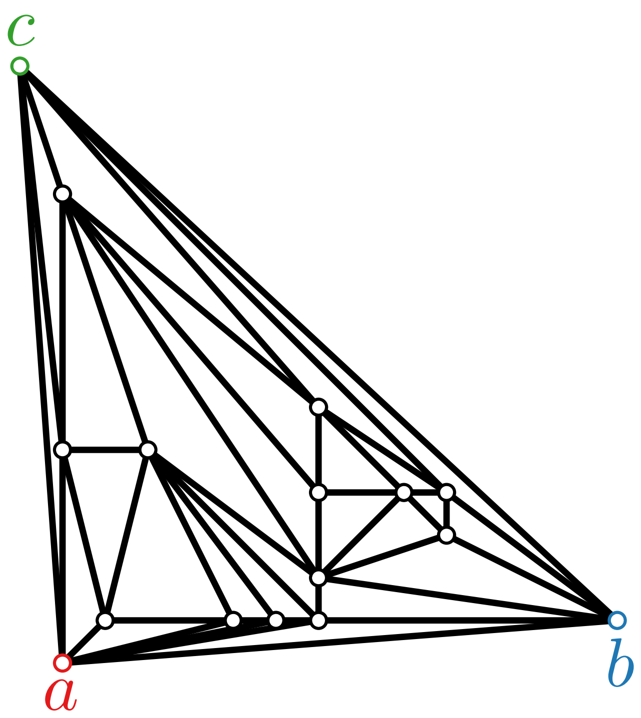

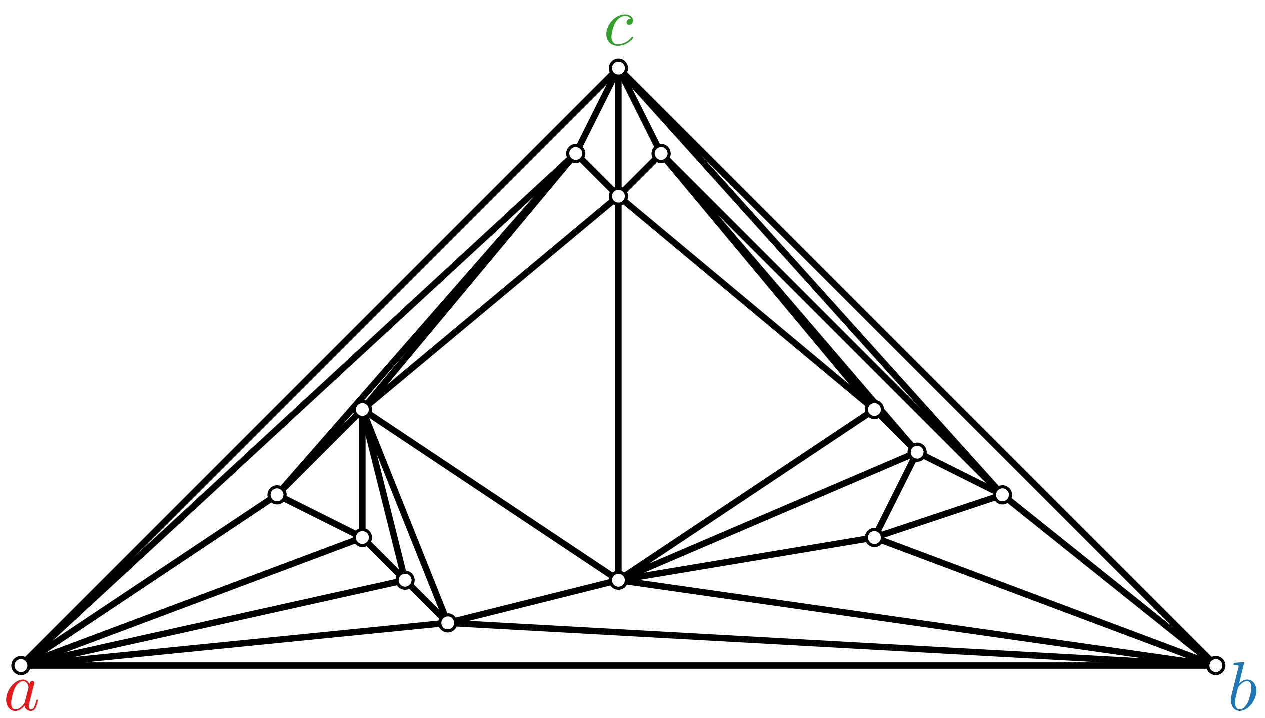

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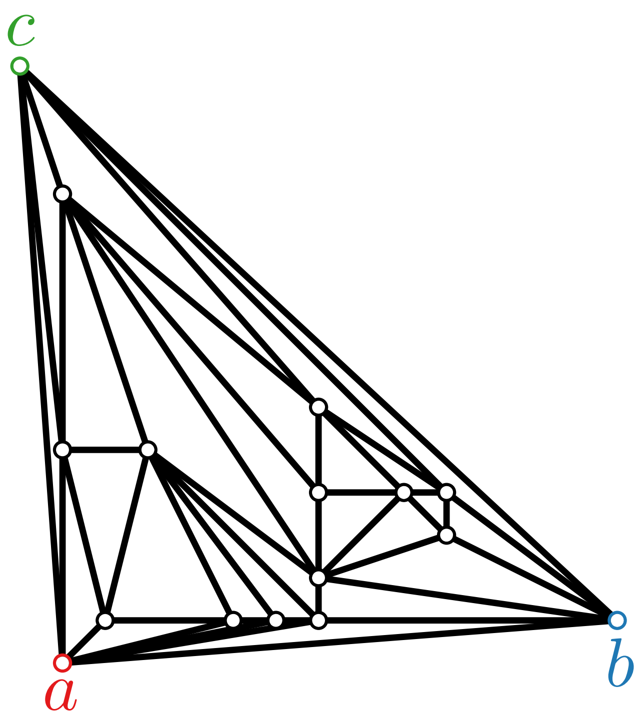

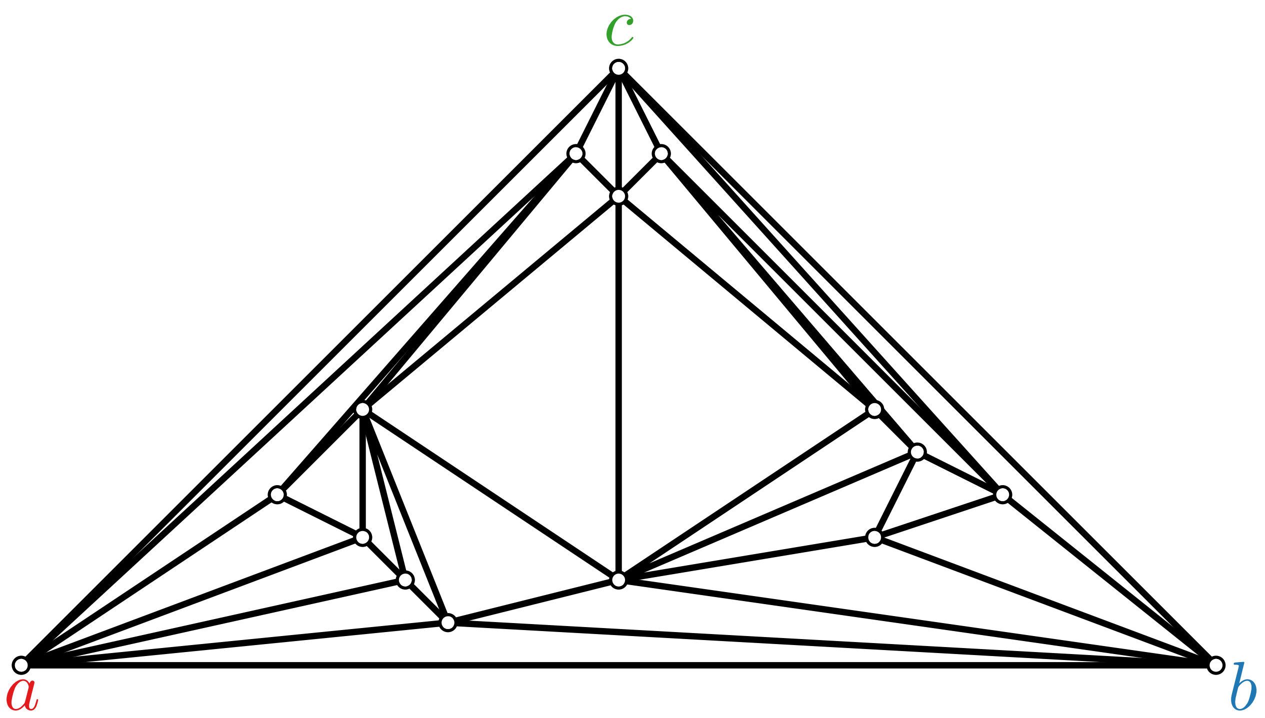

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Theorem. [Dolev, Leighton, Trickey '84]
There exist n-vertex plane graphs such that any planar straight-line drawing of them has an area of at least $(2n/3-1) \times (2n/3-1)$.

Results & Variations

Theorem. [De Fraysseix, Pach, Pollack '90]
Every n-vertex planar graph has a planar straight-line drawing of size $(2n-4) \times (n-2)$. Such a drawing can be computed in $O(n)$ time.

Theorem. [Schnyder '90]
Every n-vertex planar graph has a planar straight-line drawing of size $(n-2) \times (n-2)$. Such a drawing can be computed in $O(n)$ time. ← *Exercise!*

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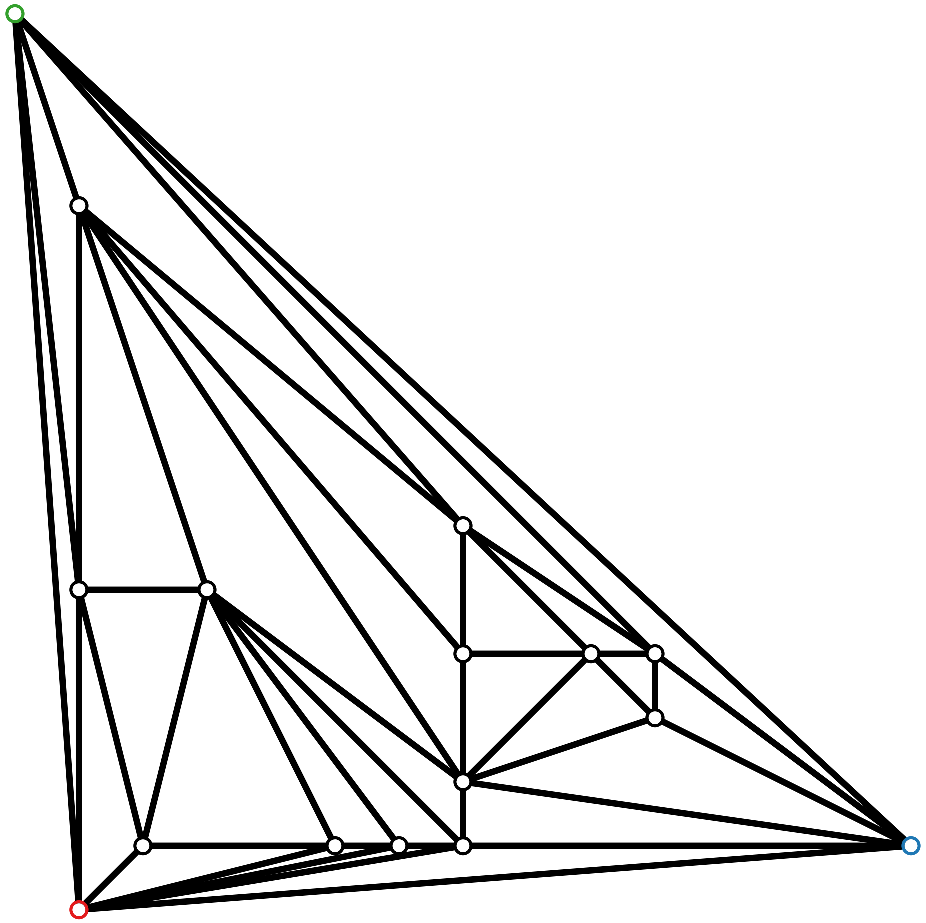
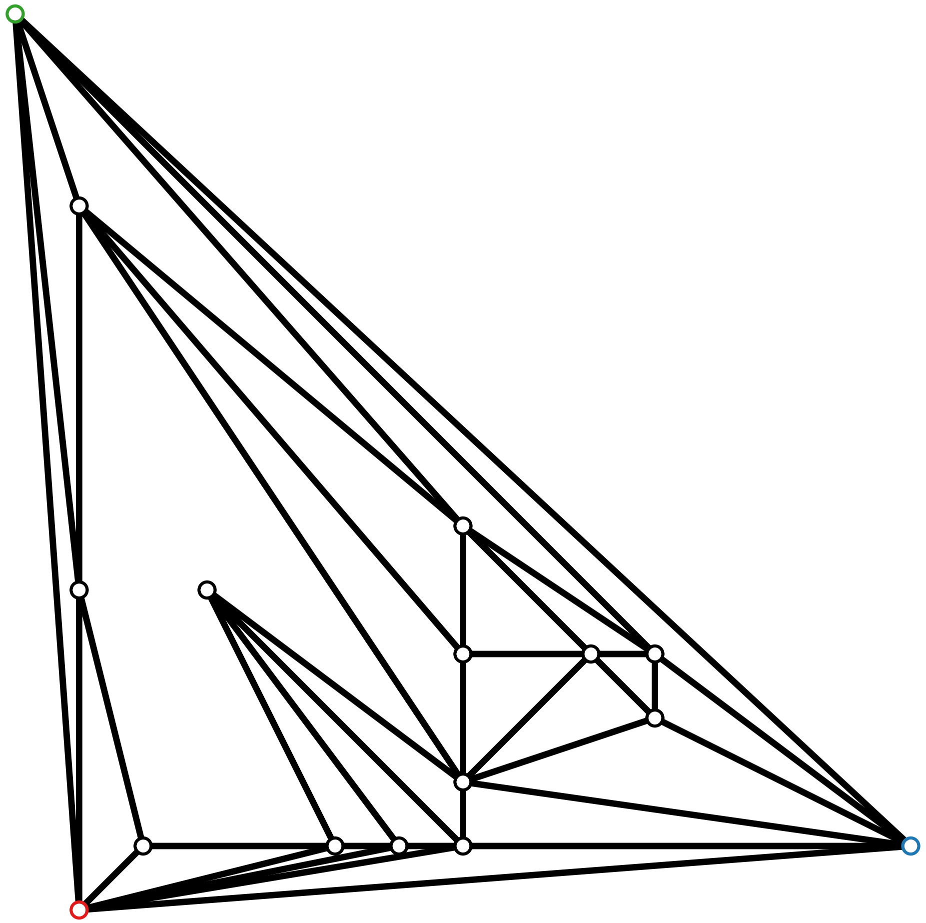
[Frati, Patrignani '07] Area at least $n^2/9 + \Omega(n)$ in the variable-embedding setting.

Results & Variations



Results & Variations



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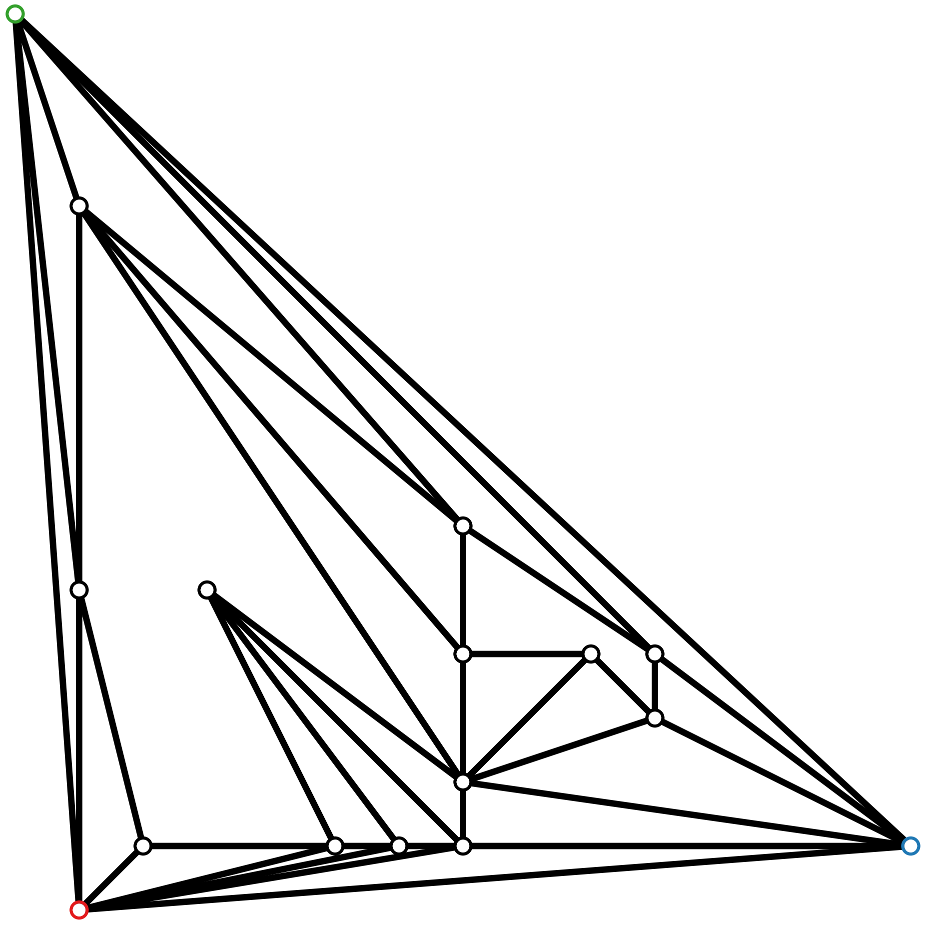

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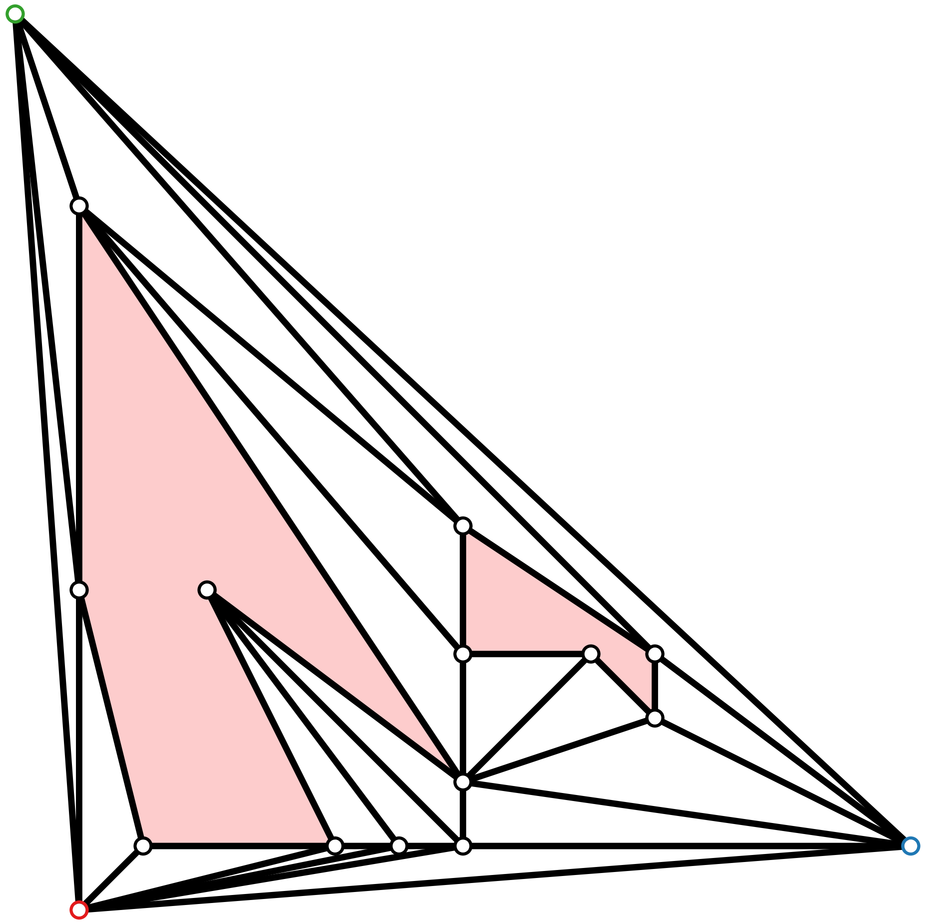

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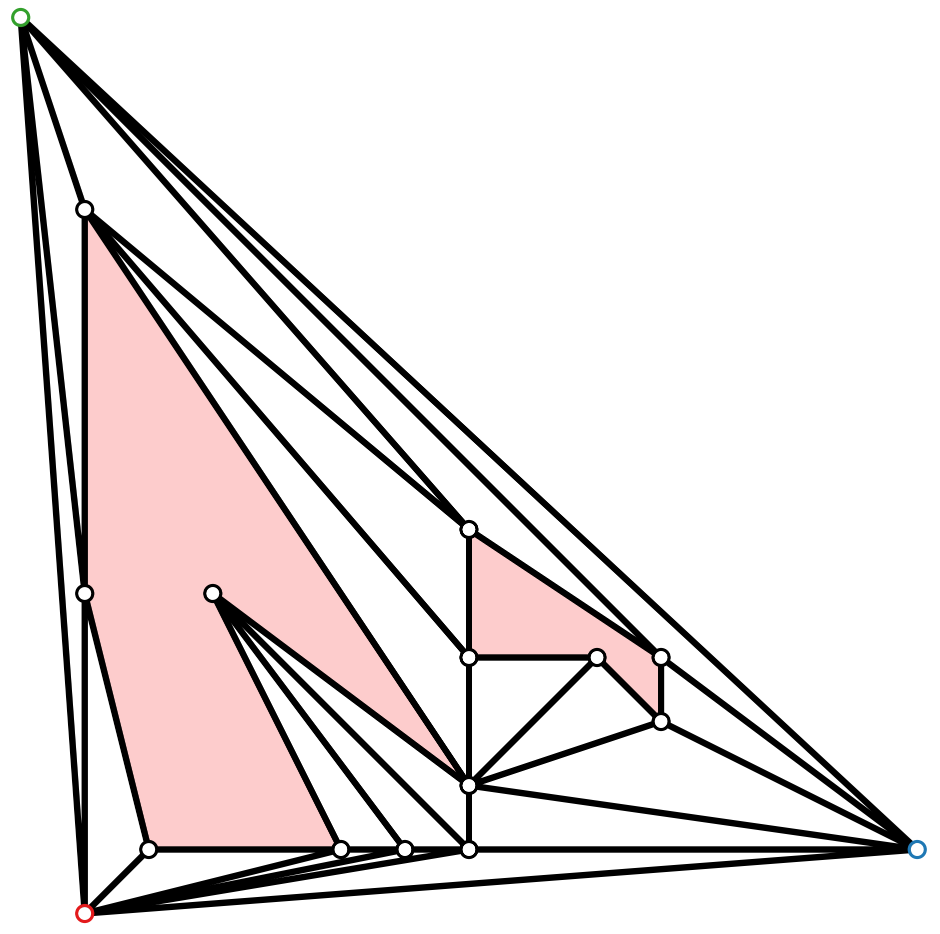

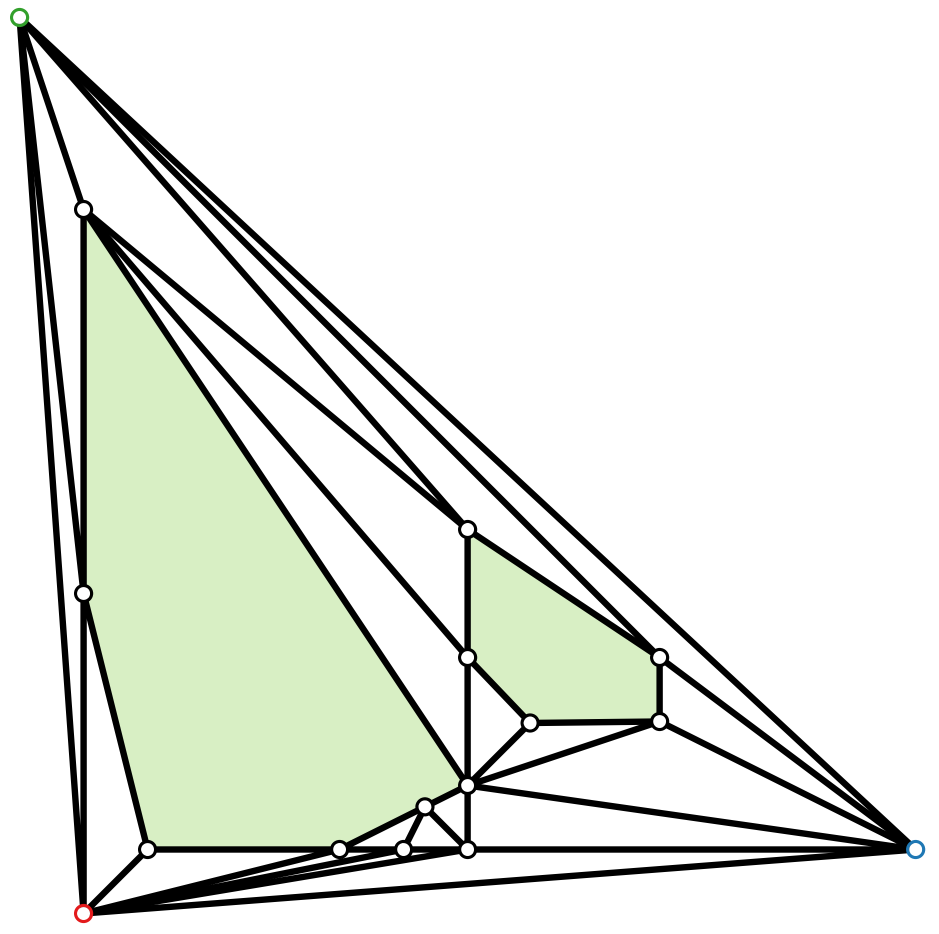

Results & Variations

Theorem. [Kant '96]
Every n-vertex 3-connected planar graph has a planar straight-line drawing of size $(2n-4) \times (n-2)$ where all faces are drawn convex.
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Theorem. [Felsner '01]
Every 3-connected planar graph with f faces has a planar straight-line drawing of size $(f-1) \times (f-1)$ where all faces are drawn convex.
Such a drawing can be computed in $O(n)$ time.

Literature

- [PGD Ch. 4.3] for detailed explanation of Schnyder woods etc.
- [Sch90] "Embedding planar graphs on the grid", Walter Schnyder, SoCG 1990 – original paper on Schnyder realizer method.