Approximation Algorithms

Lecture 8: Approximation Schemes and the Knapsack Problem

Part I: Knapsack

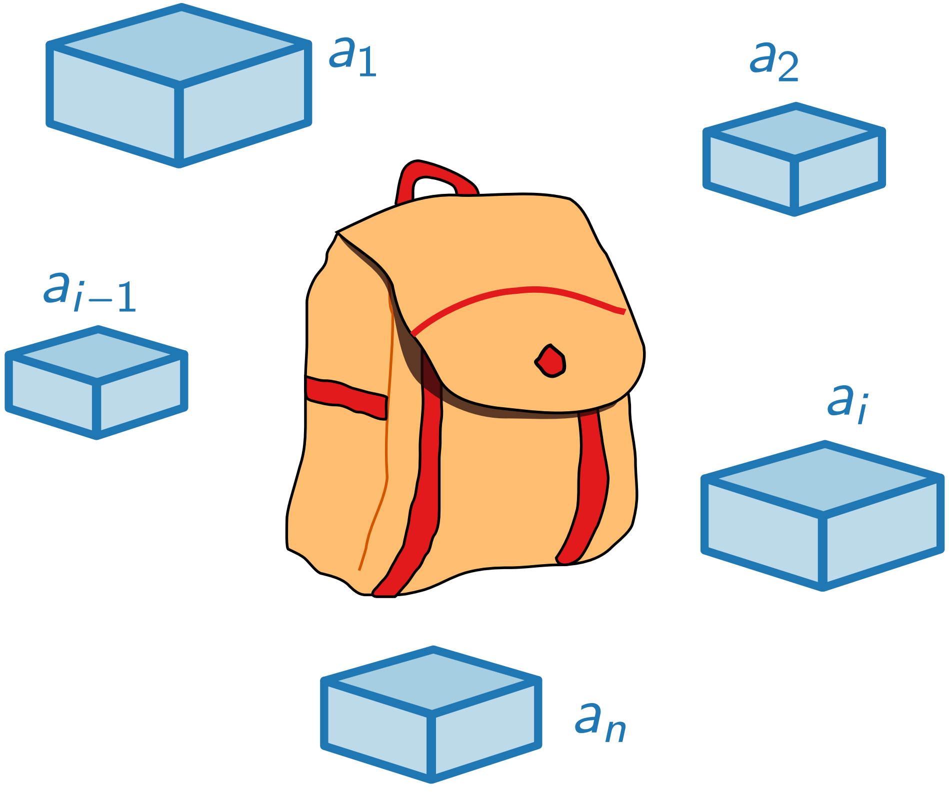
Alexander Wolff

Winter 2023/24

KNAPSACK

Given:

- A set $S = \{a_1, \ldots, a_n\}$ of **objects**.



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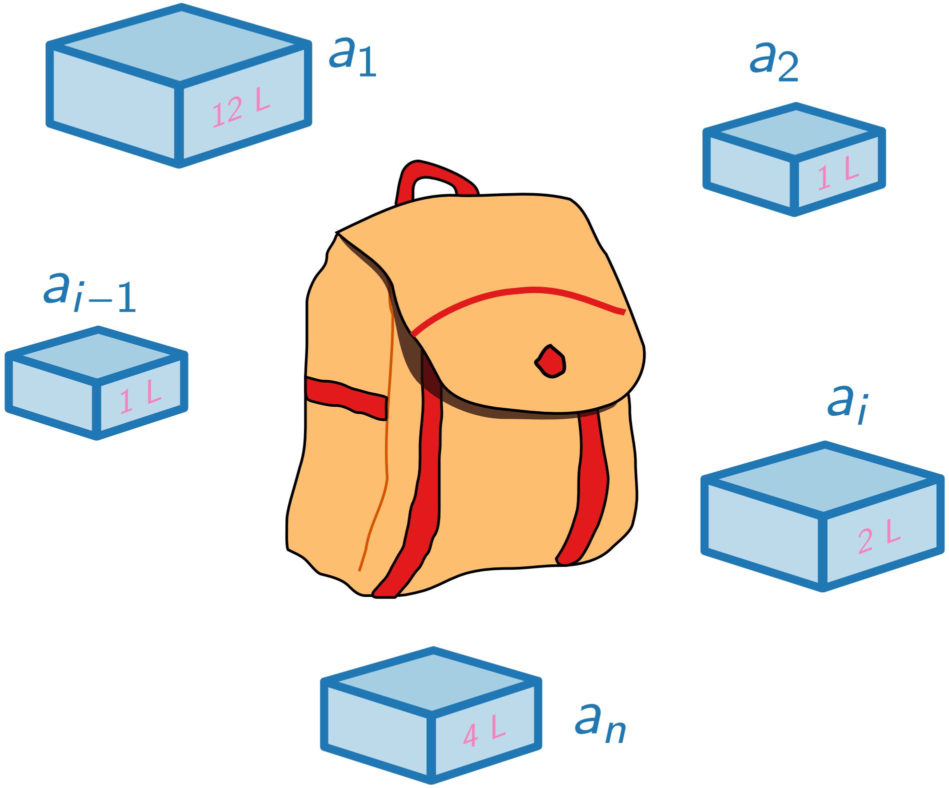

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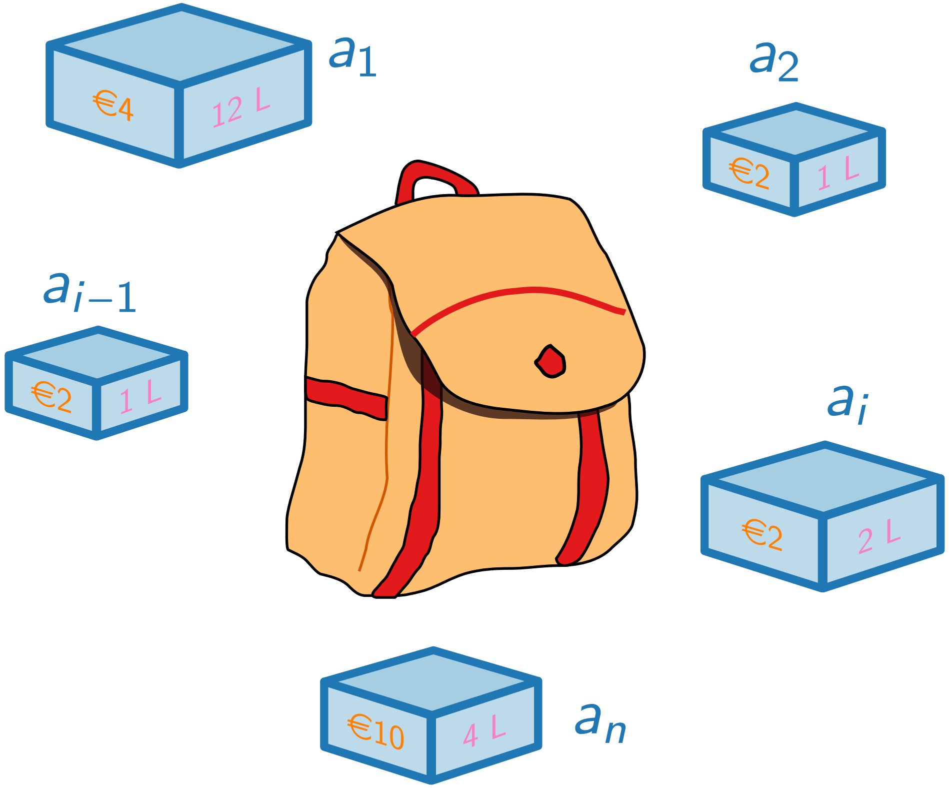

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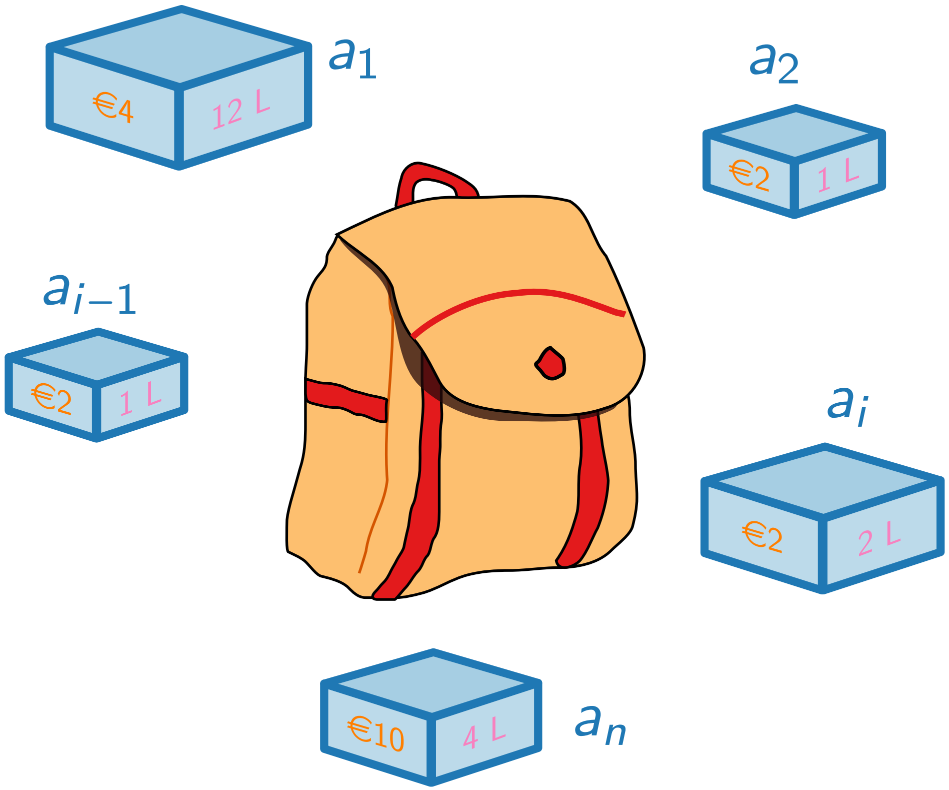

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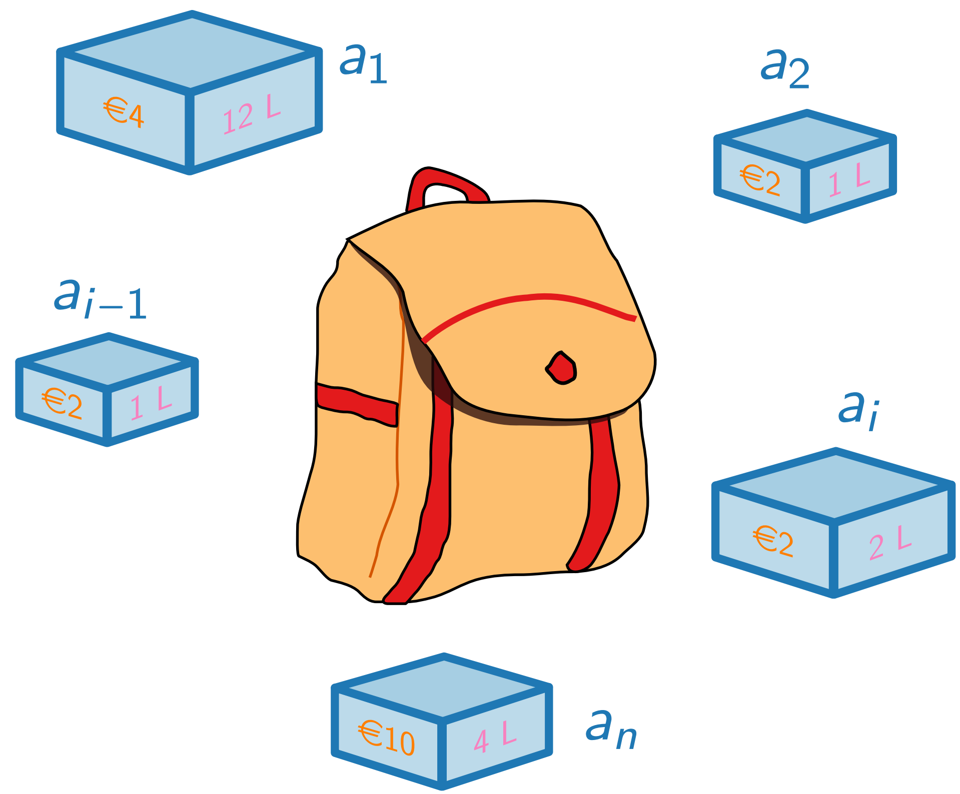
Task: Find a subset of objects whose **total size** is at most B and whose **total profit** is maximum.



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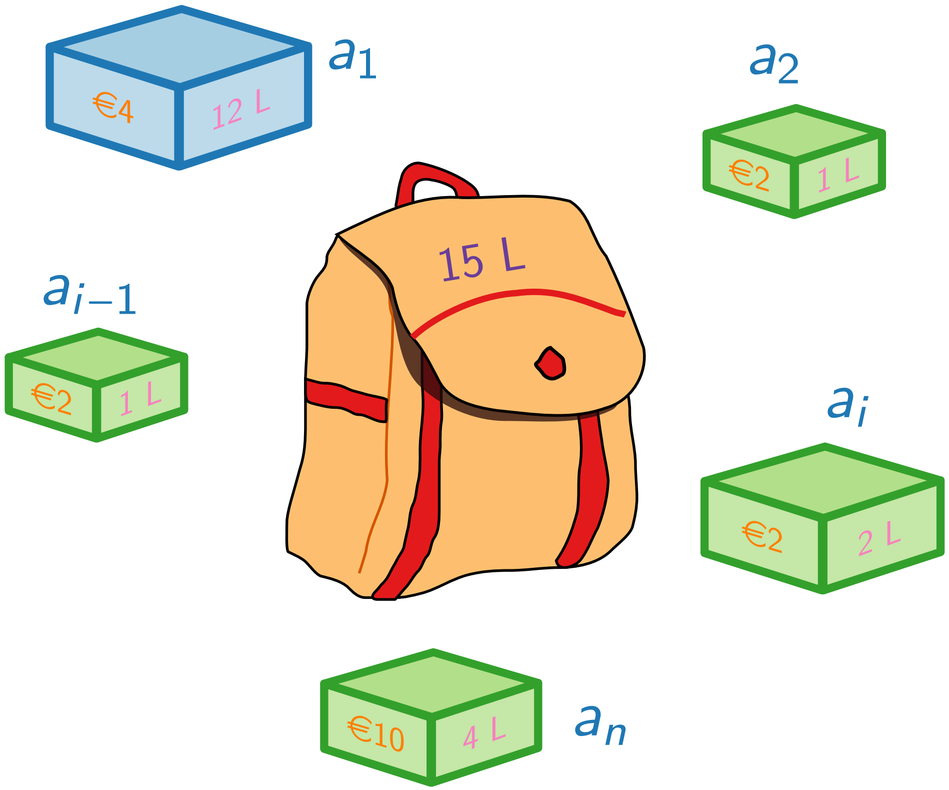

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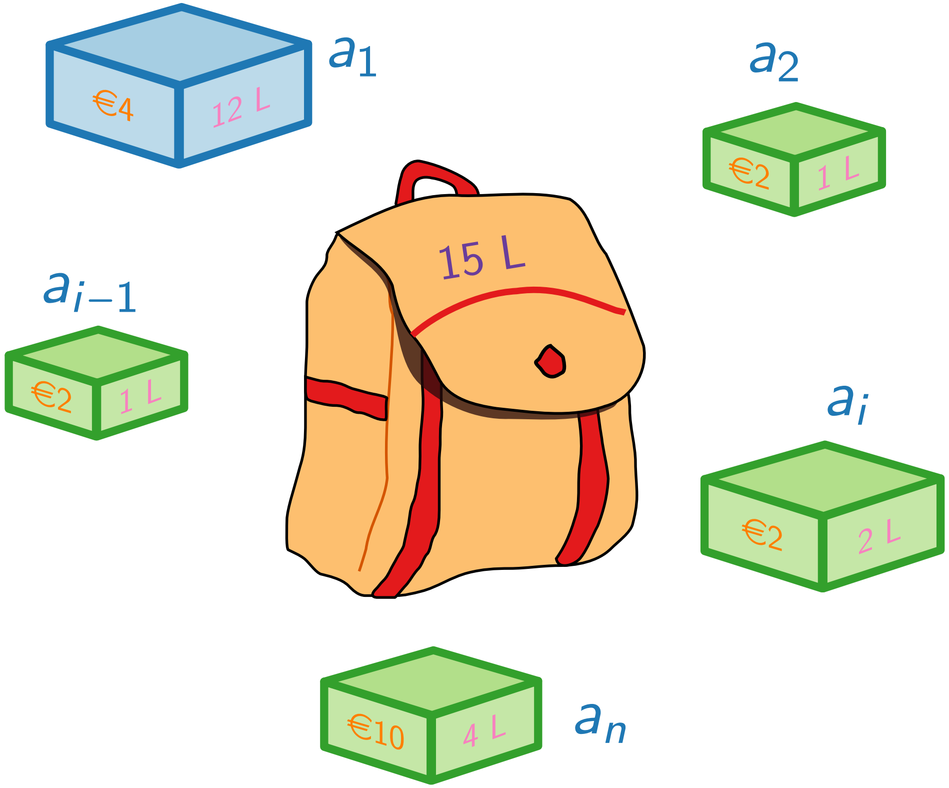

Approximation Algorithms

Lecture 8: Approximation Schemes and the KNAPSACK Problem

Part II: Pseudo-Polynomial Algorithms and Strong NP-Hardness

Pseudo-Polynomial Algorithms

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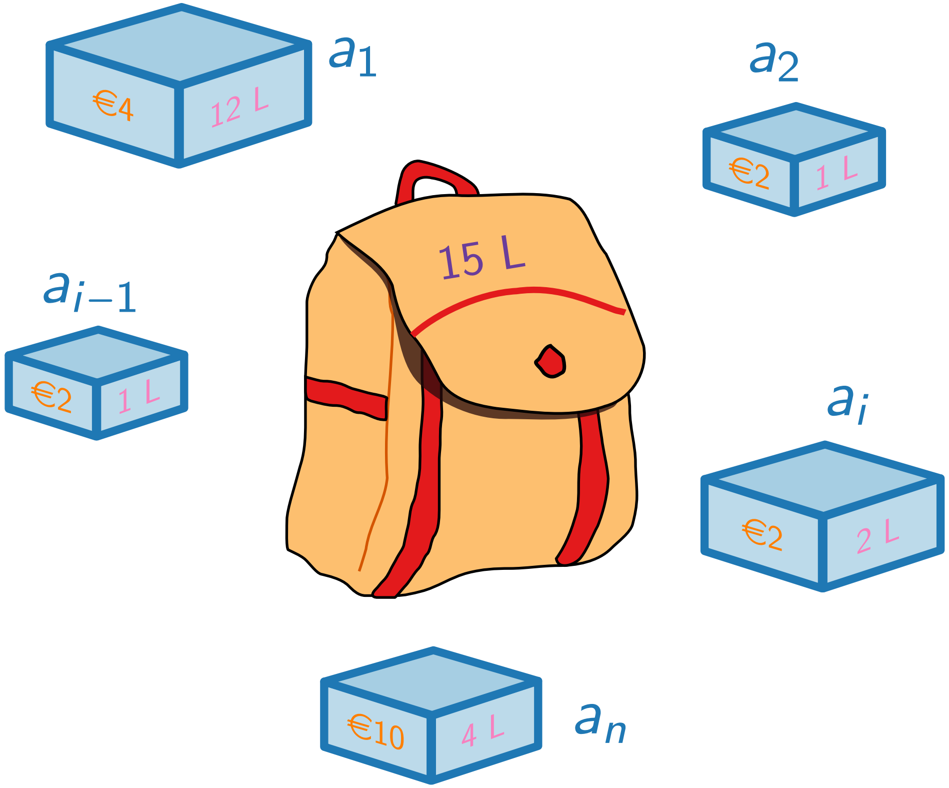

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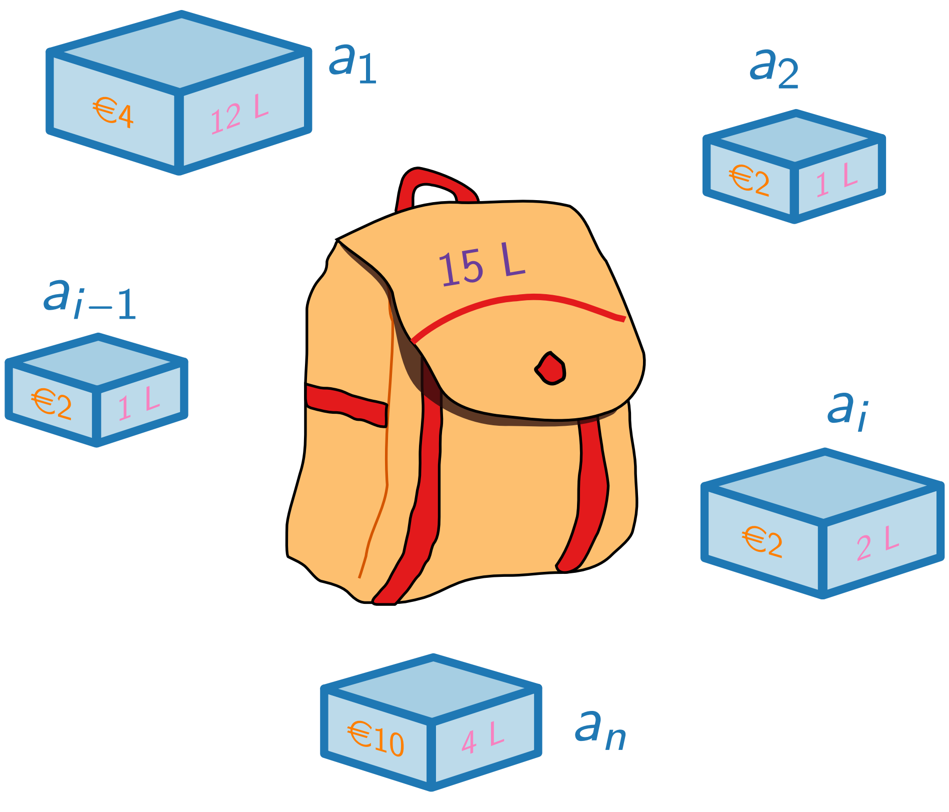

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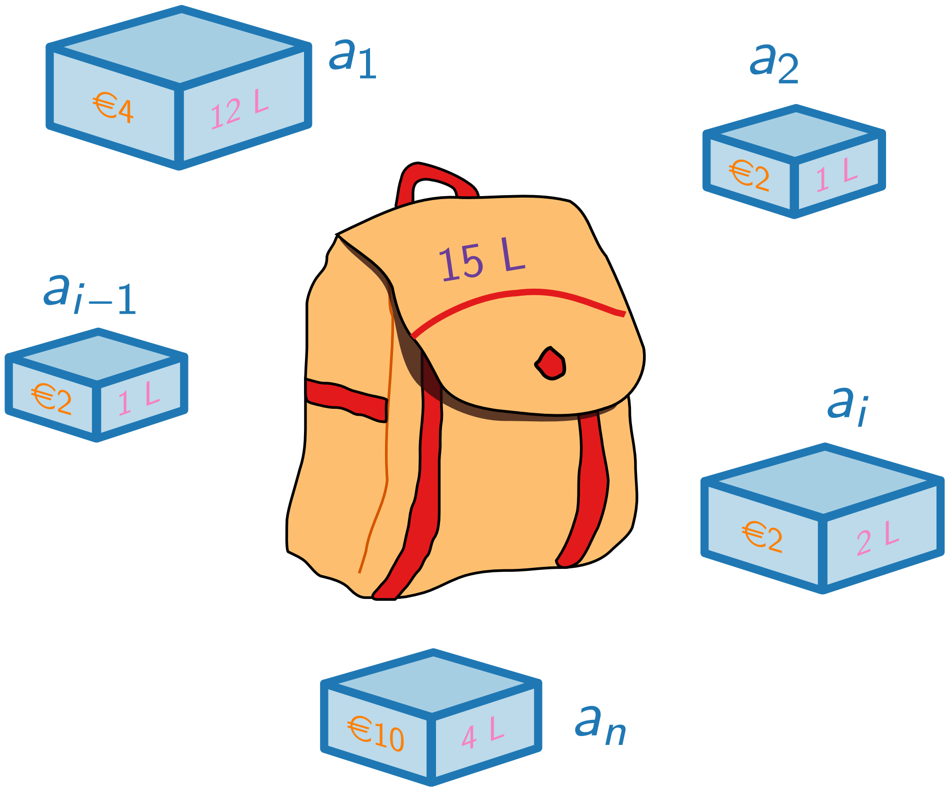

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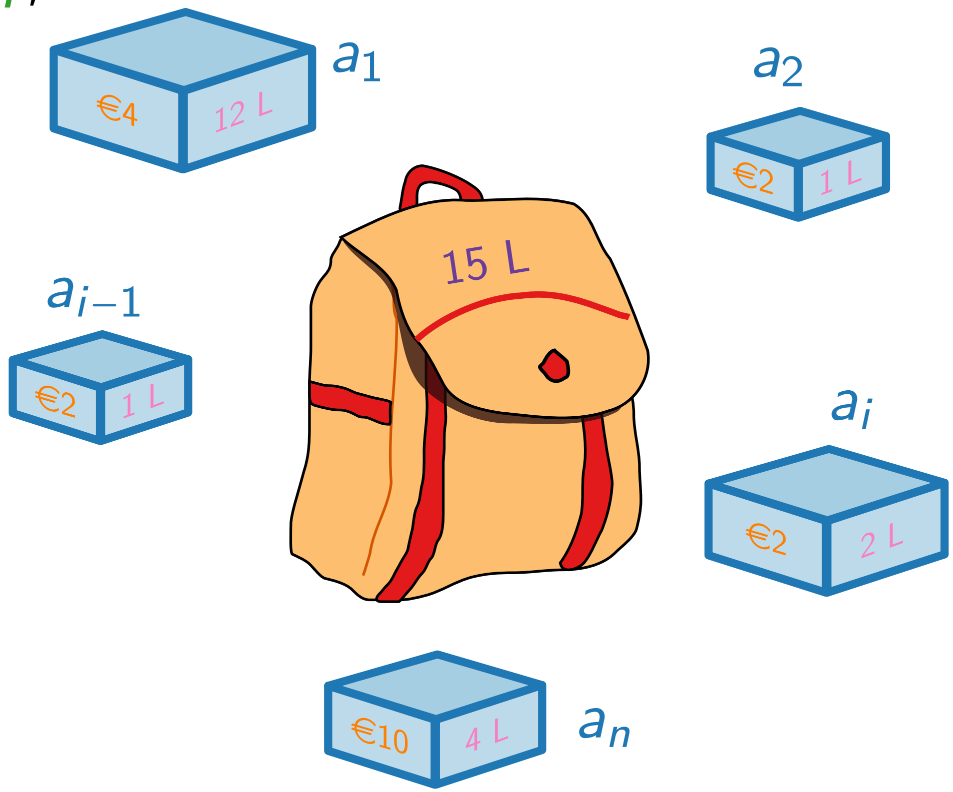

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The running time of a **pseudo-polynomial algorithm** is polynomial in $|I|_u$.

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The running time of a polynomial algorithm for Π is polynomial in $|I|$.

The running time of a **pseudo-polynomial algorithm** is polynomial in $|I|_u$.

The running time of a pseudo-polynomial algorithm may not be polynomial in $|I|$.

Strong NP-Hardness

An optimization problem is called **strongly NP-hard** if it remains NP-hard under unary encoding.

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Theorem. A strongly NP-hard problem has no pseudo-polynomial algorithm unless P = NP.

Approximation Algorithms

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Approximation Schemes and the KNAPSACK Problem

Part III:
Pseudo-Polynomial Algorithm for KNAPSACK

Pseudo-Polynomial Alg. for KNAPSACK

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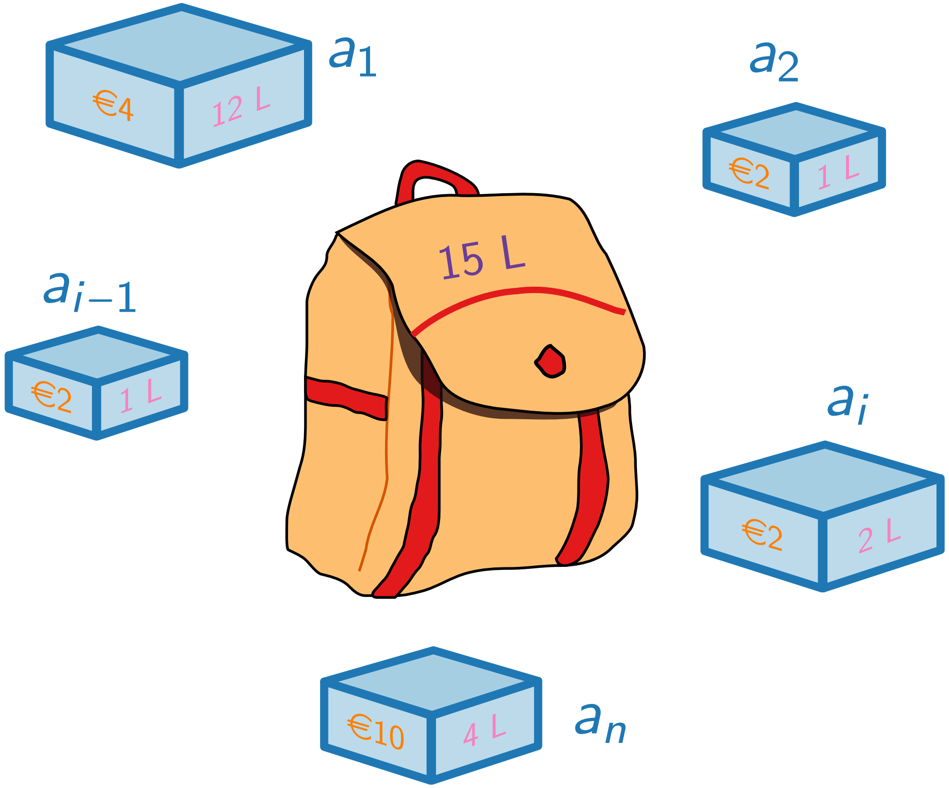

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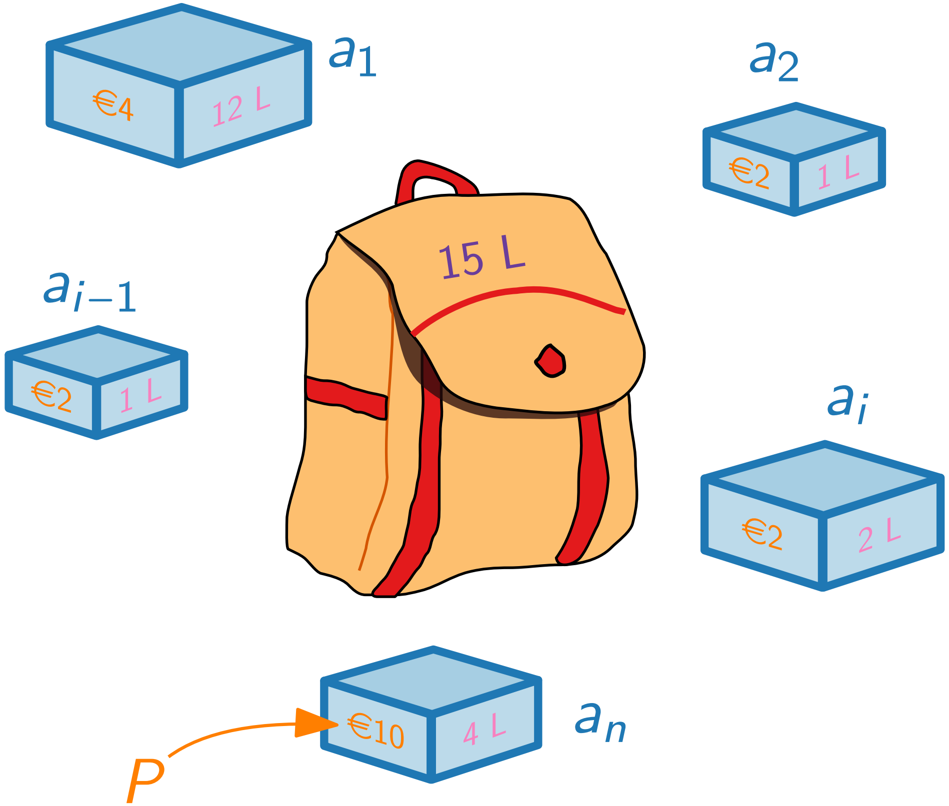

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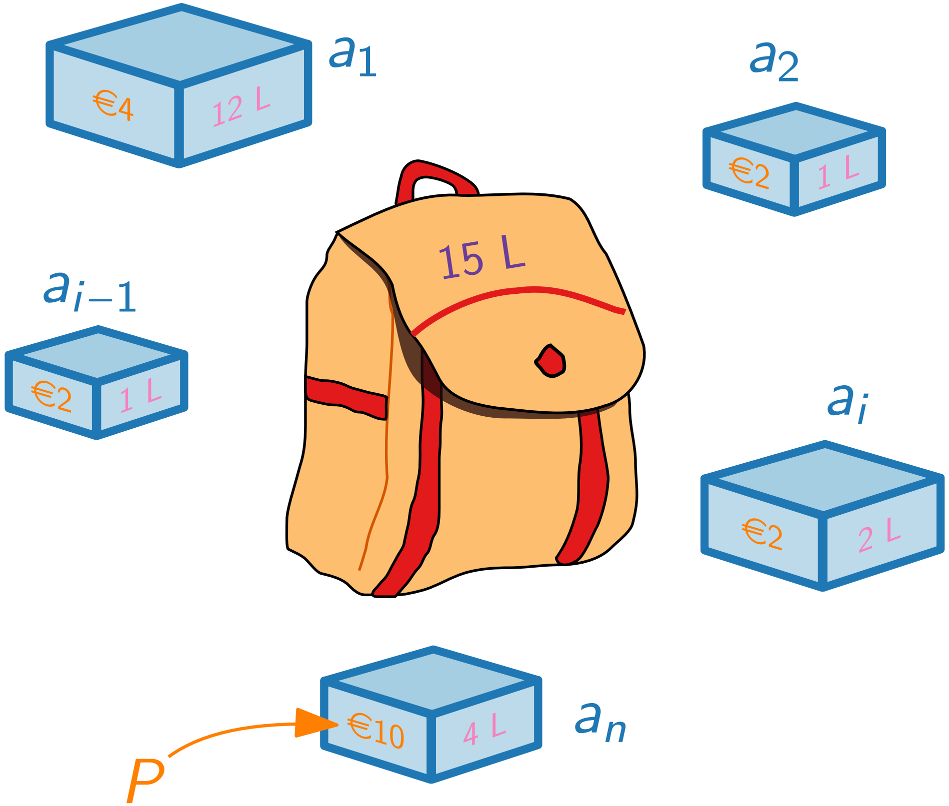

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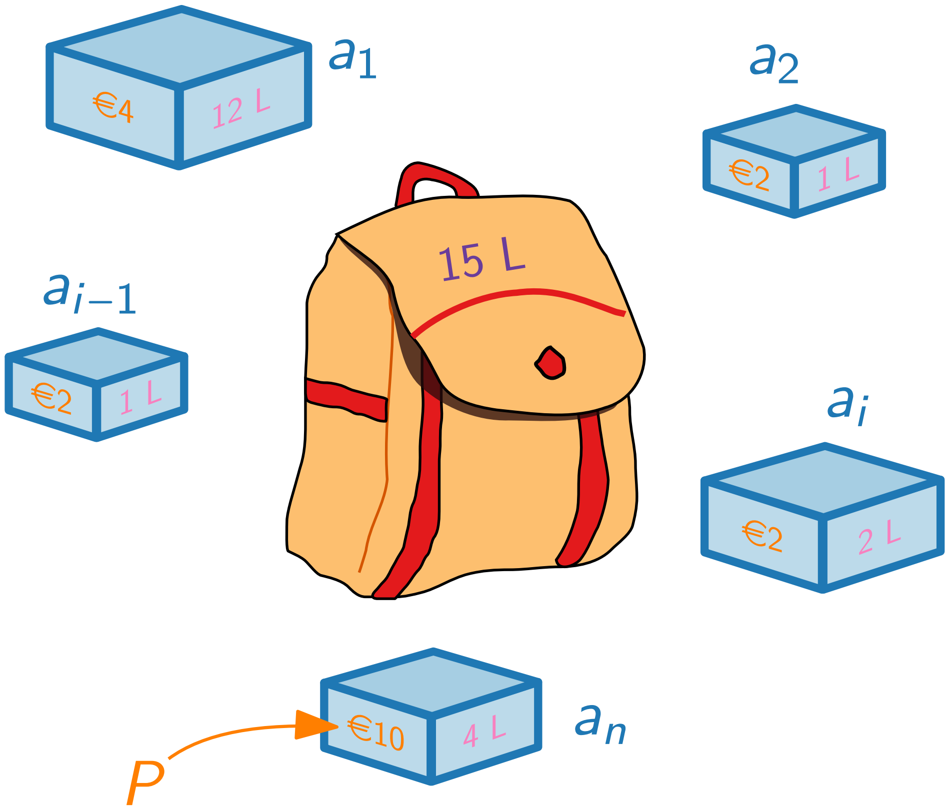

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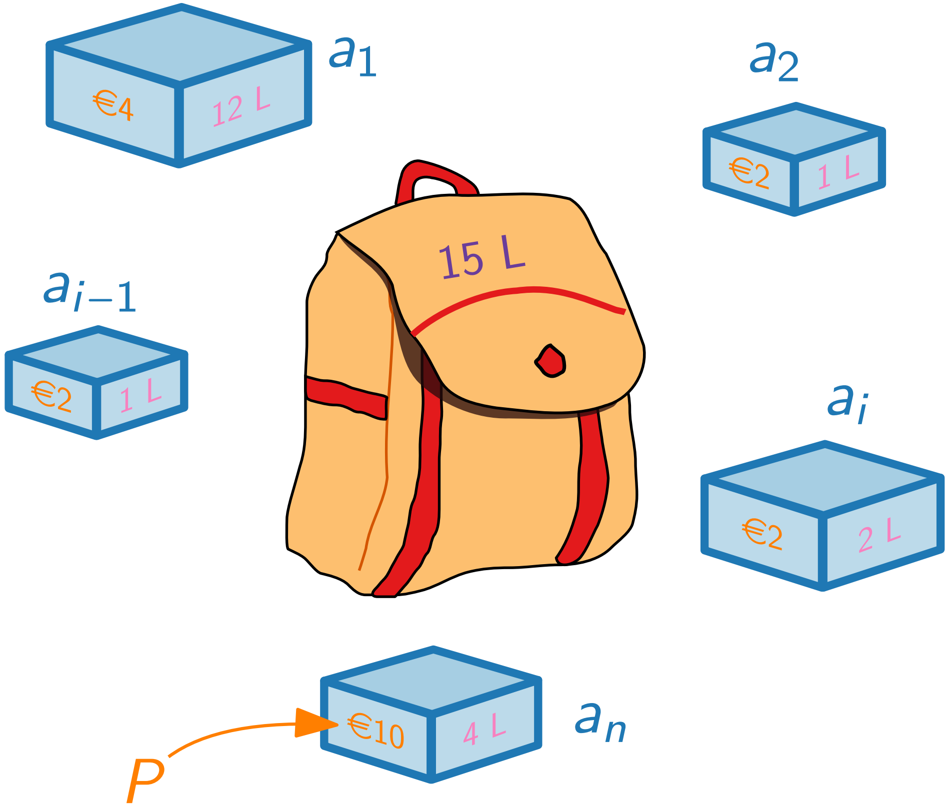

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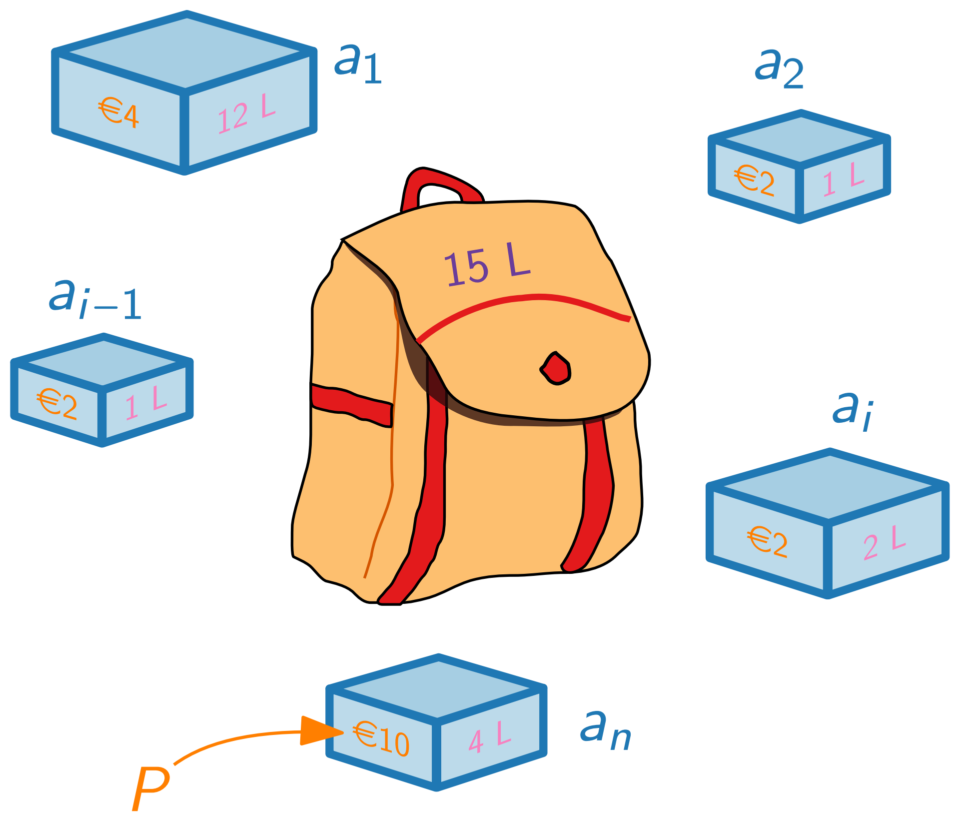

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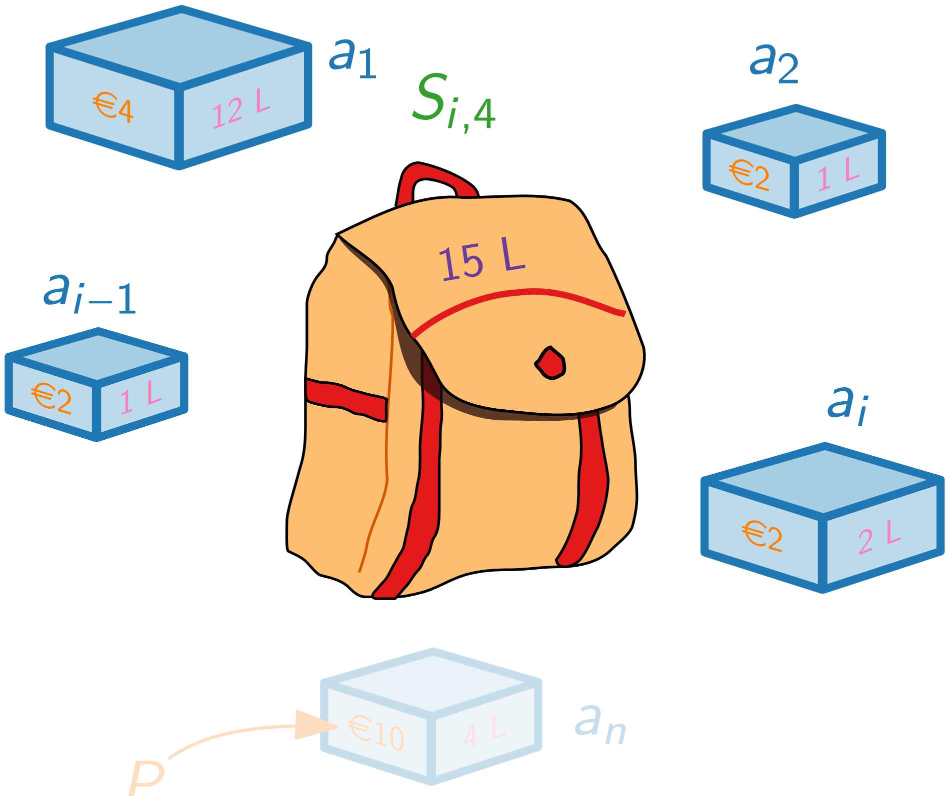

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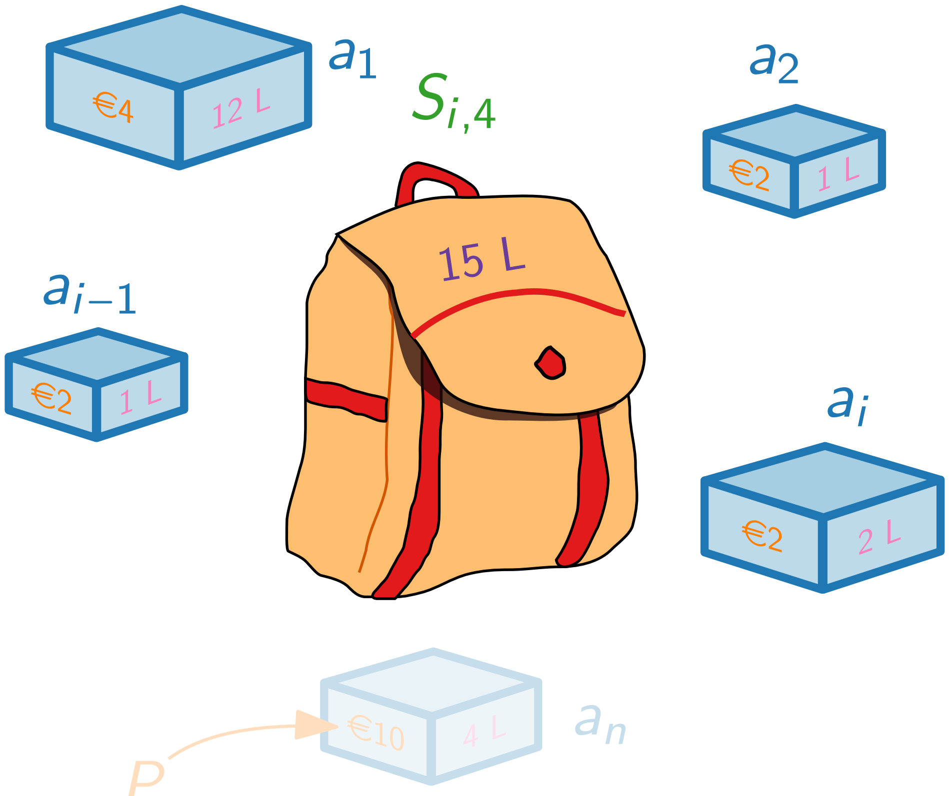

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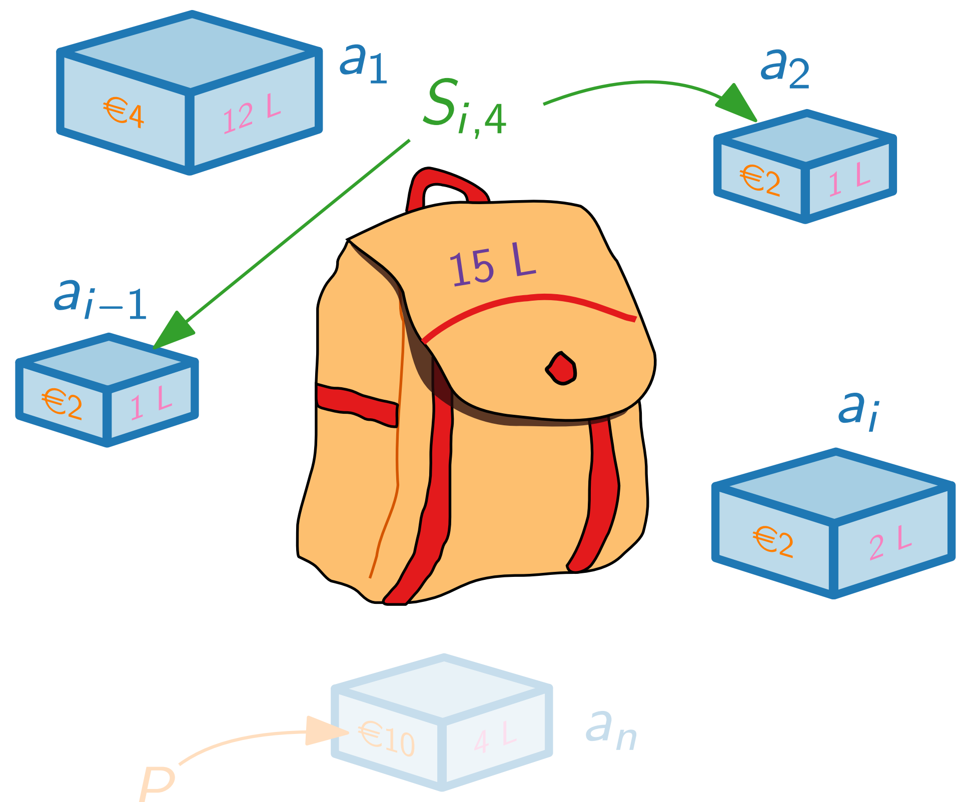

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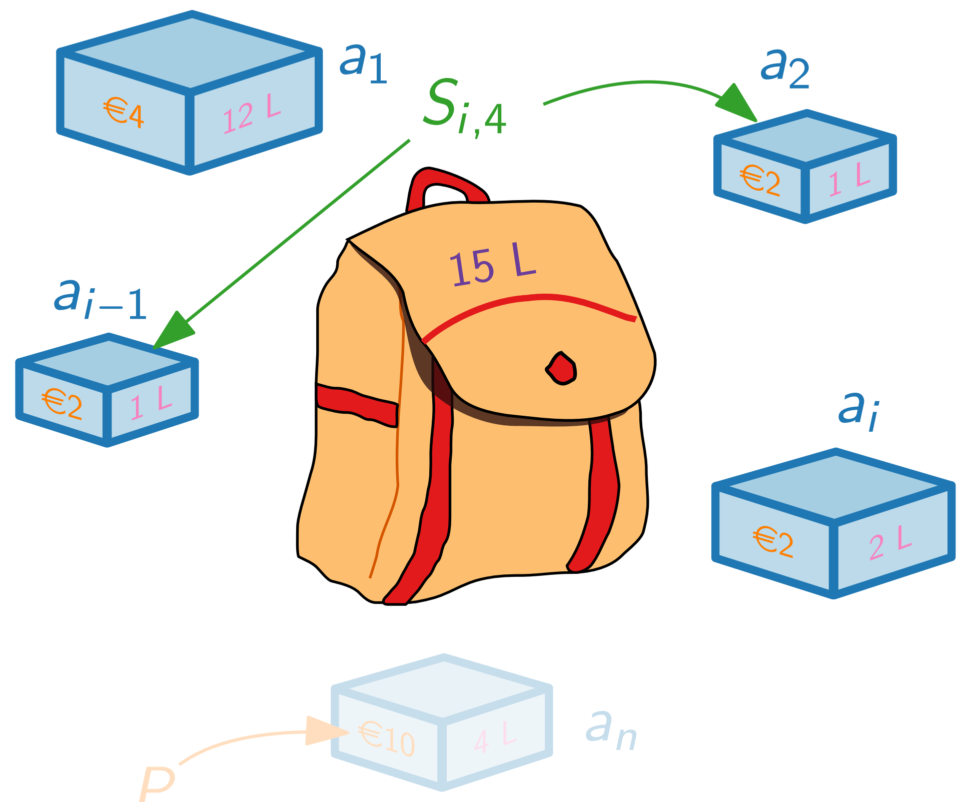

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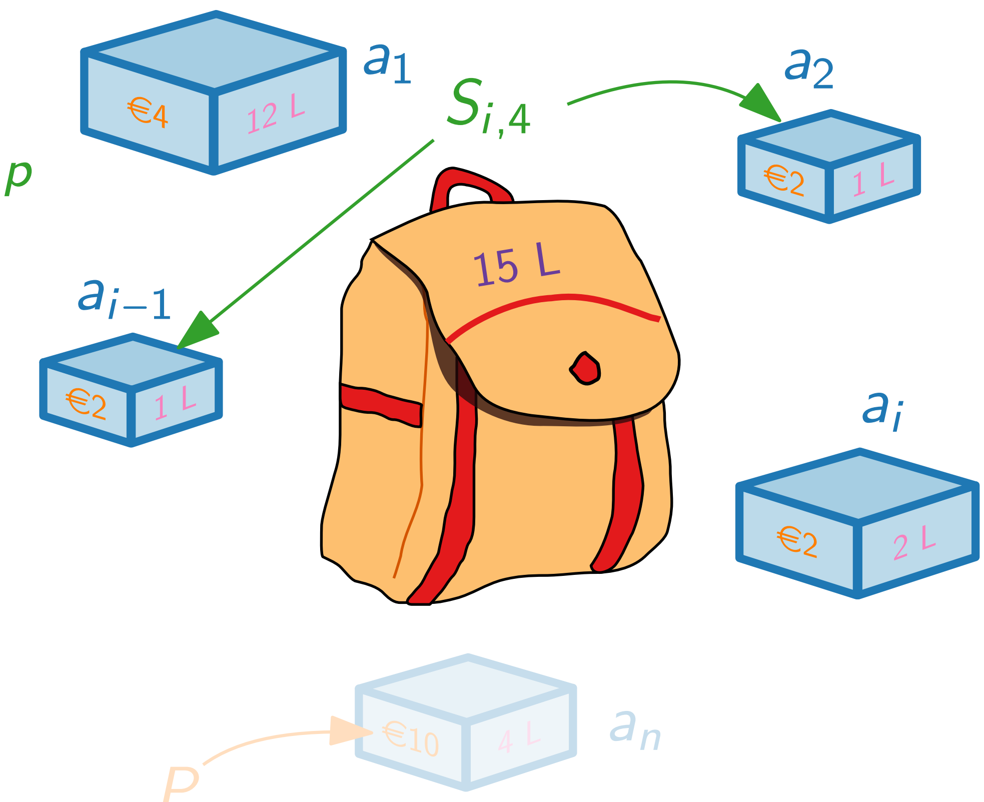

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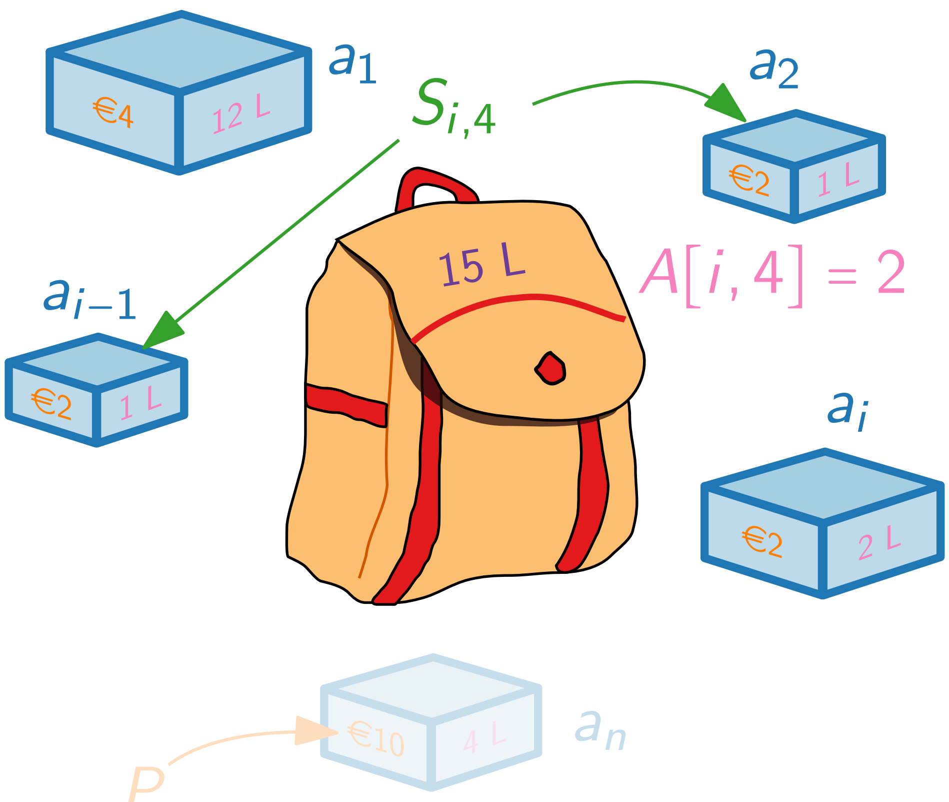

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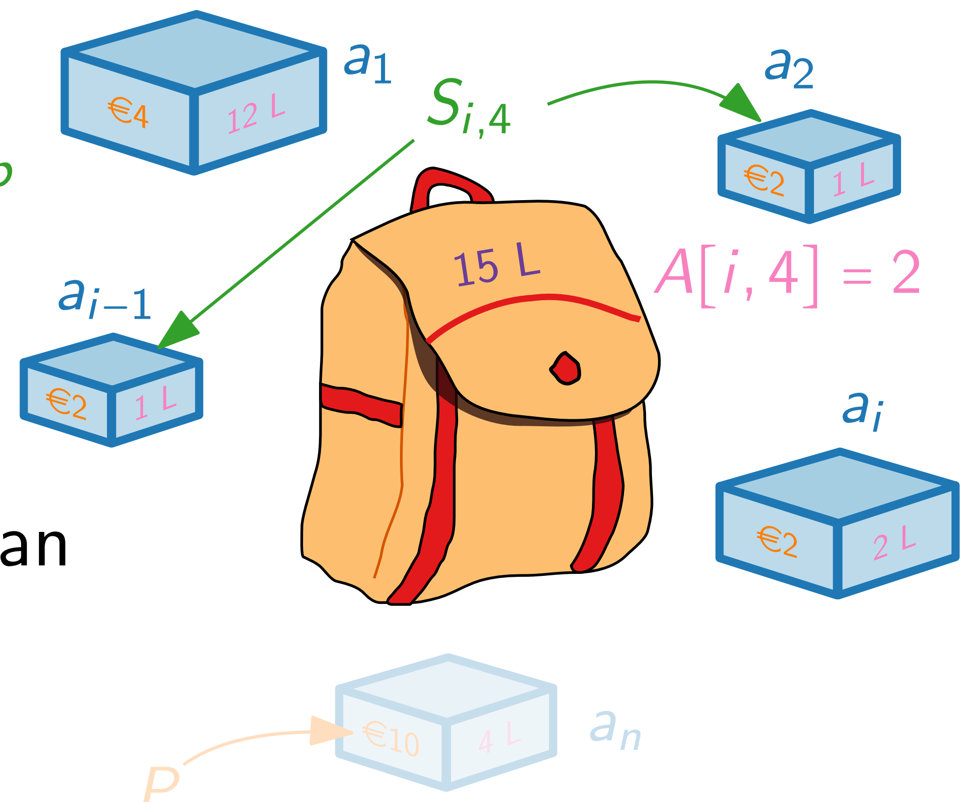

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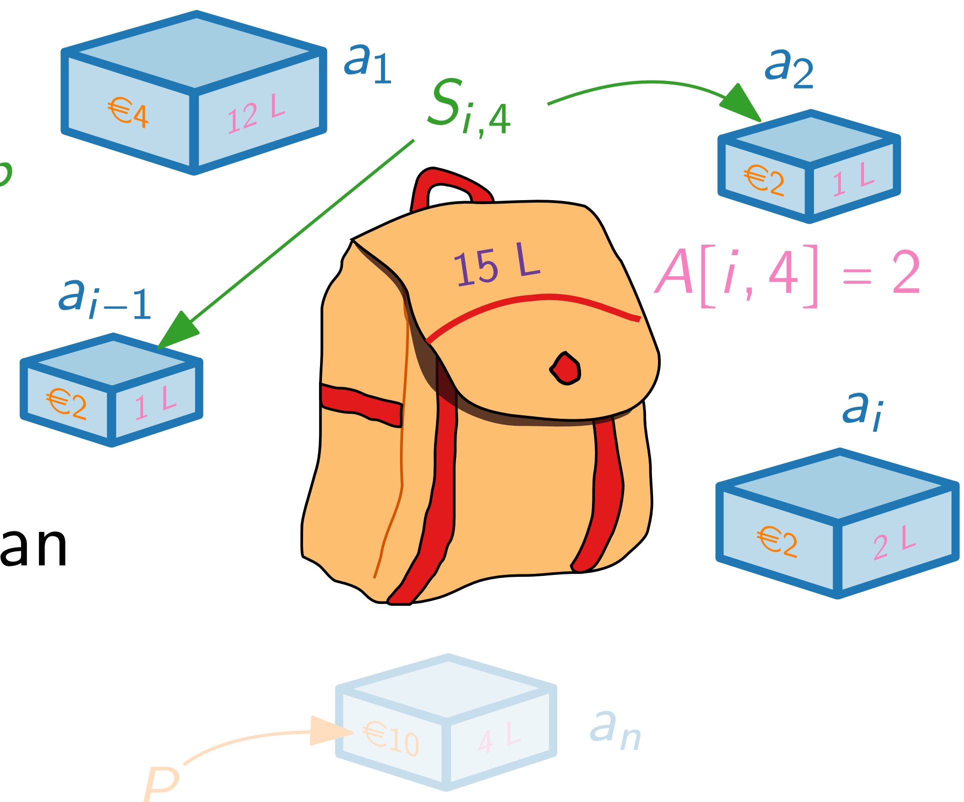
If all $A[i, p]$ are known, then we can compute $\text{OPT} = \max\{\, p \mid A[n, p] \leq B \,\}$.



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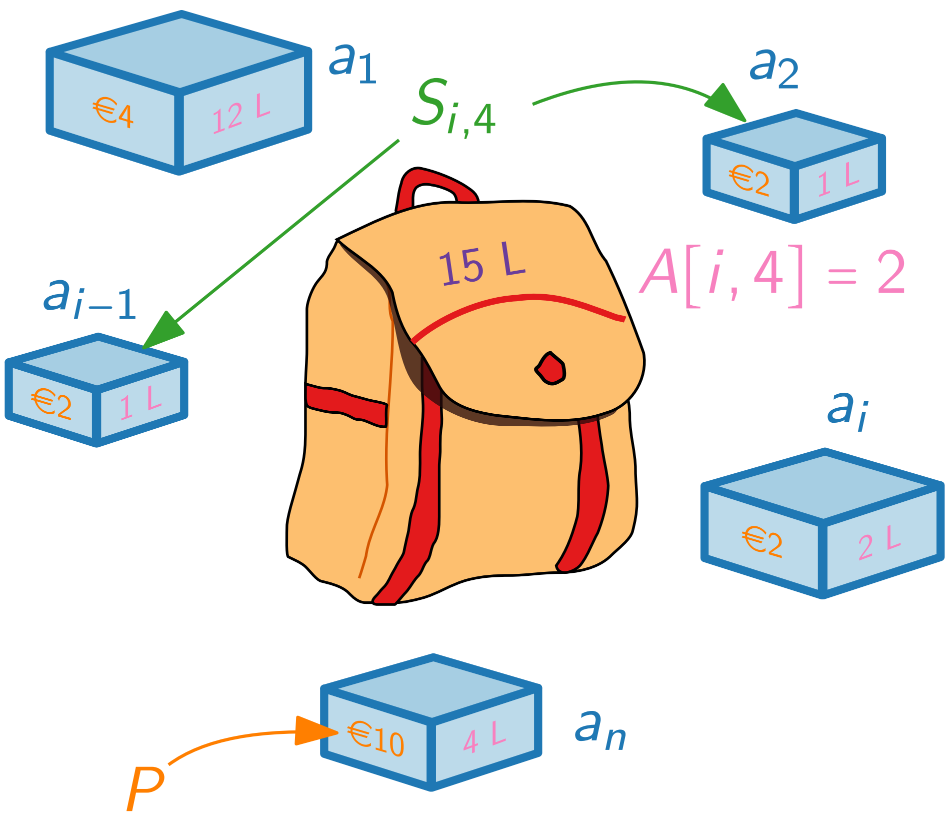

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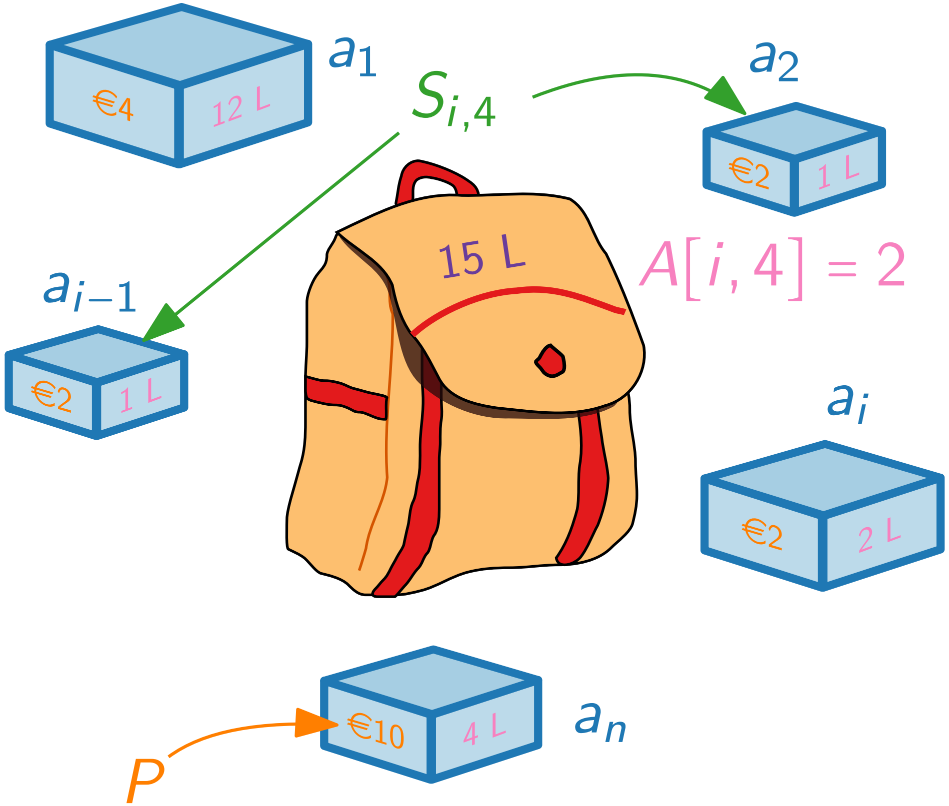

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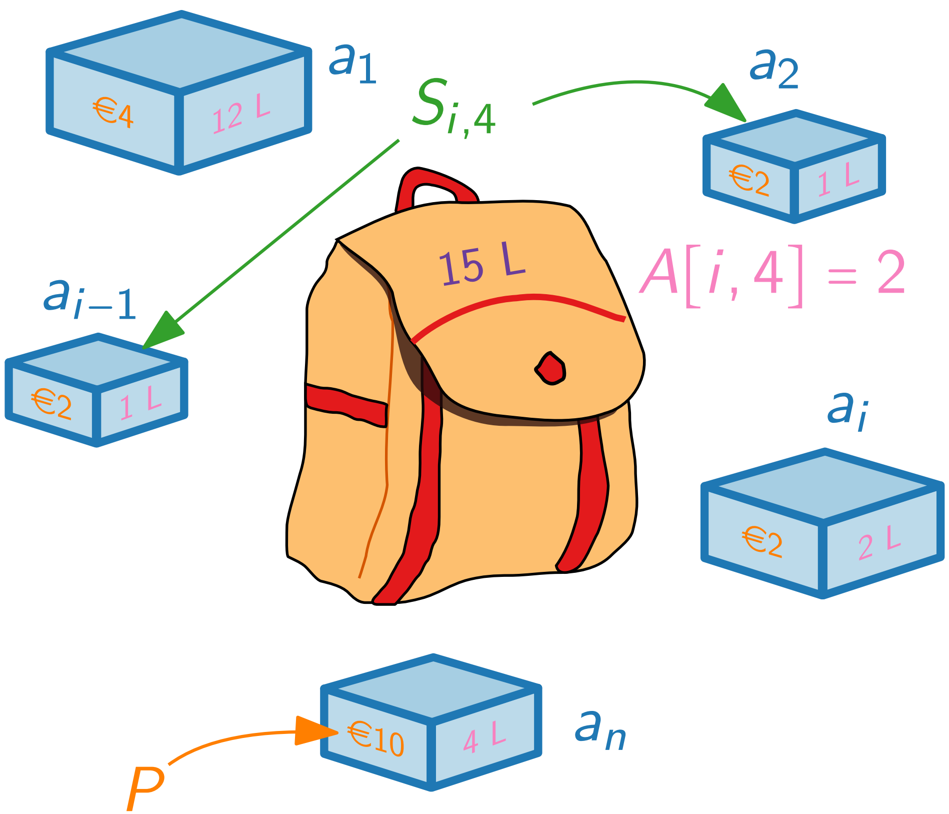

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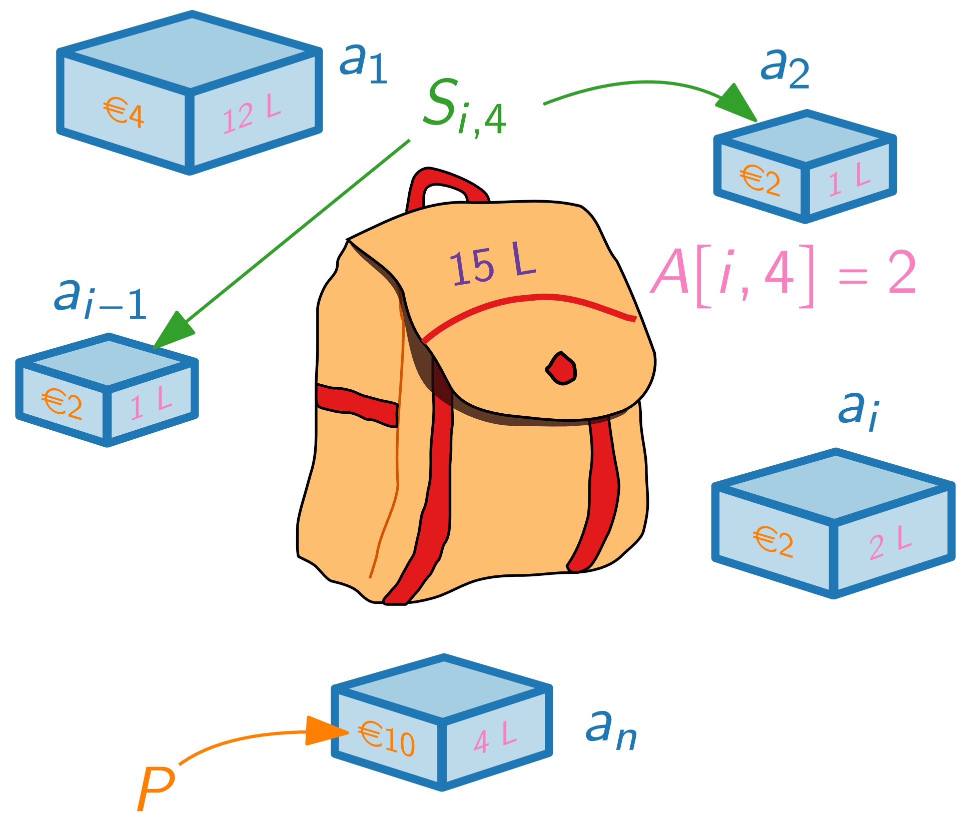

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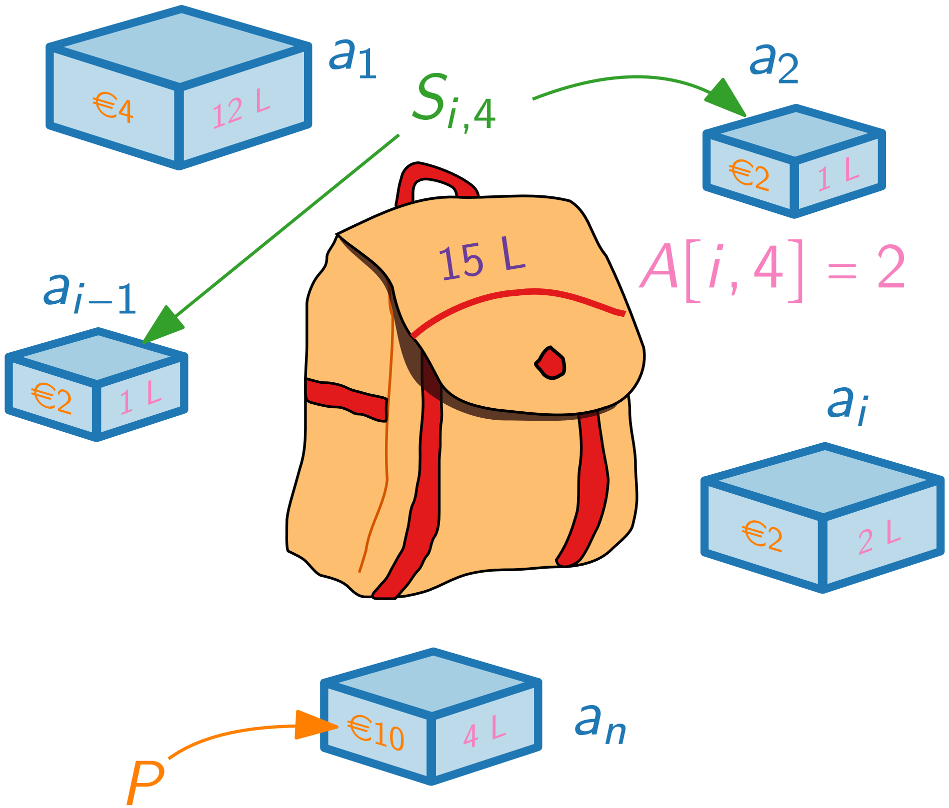

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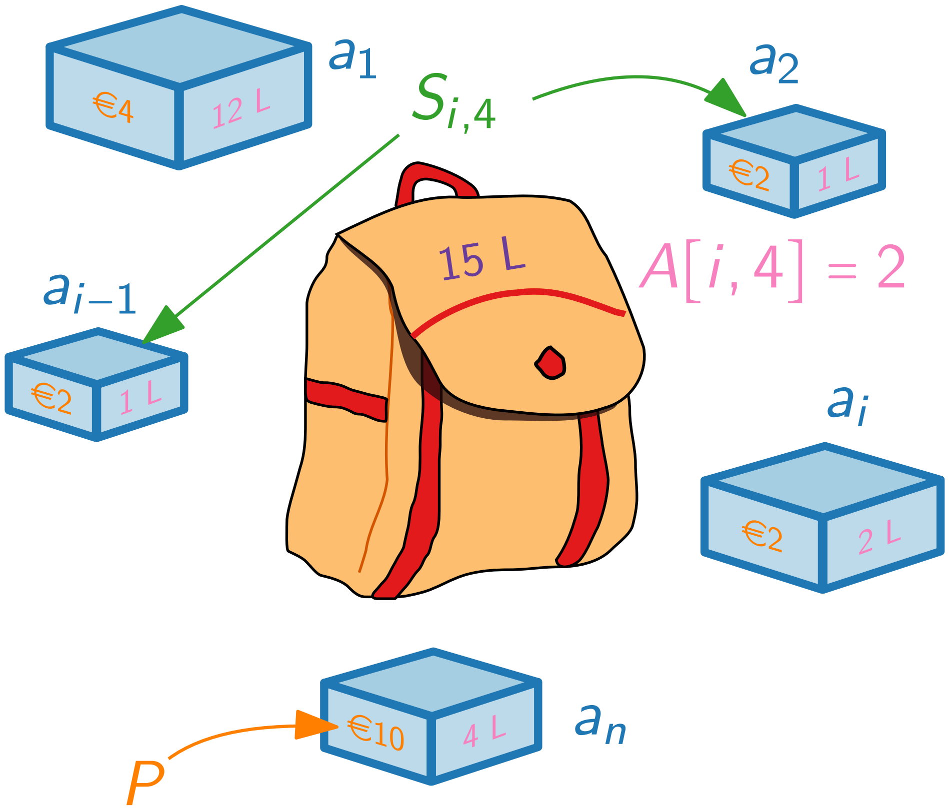

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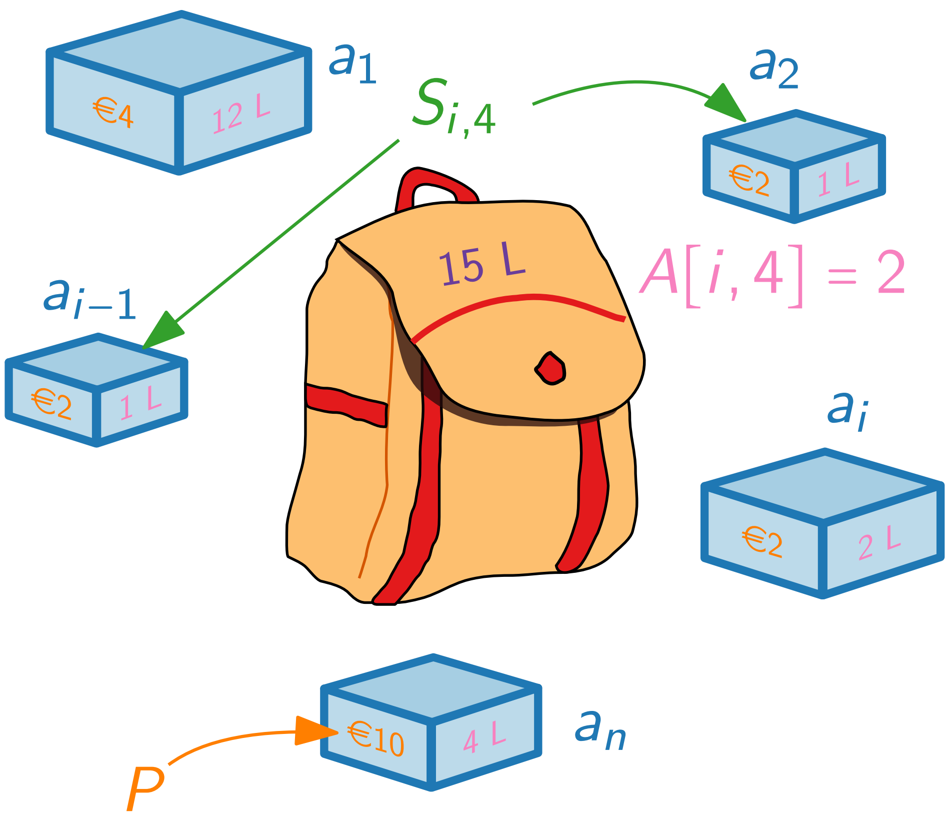

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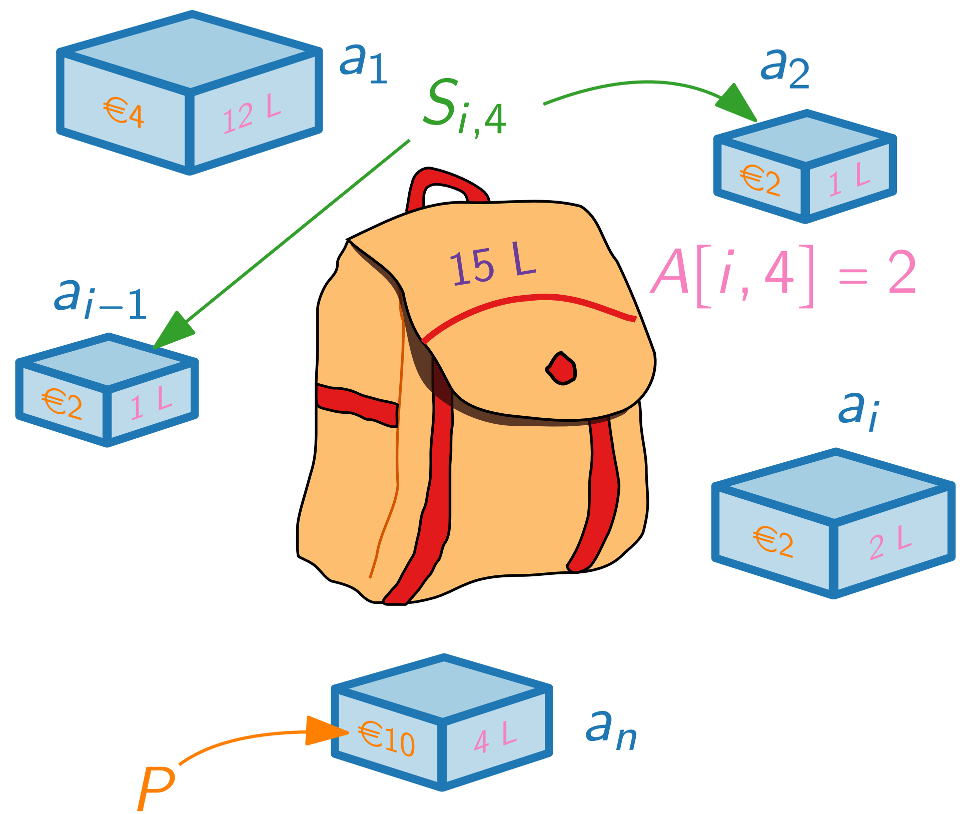

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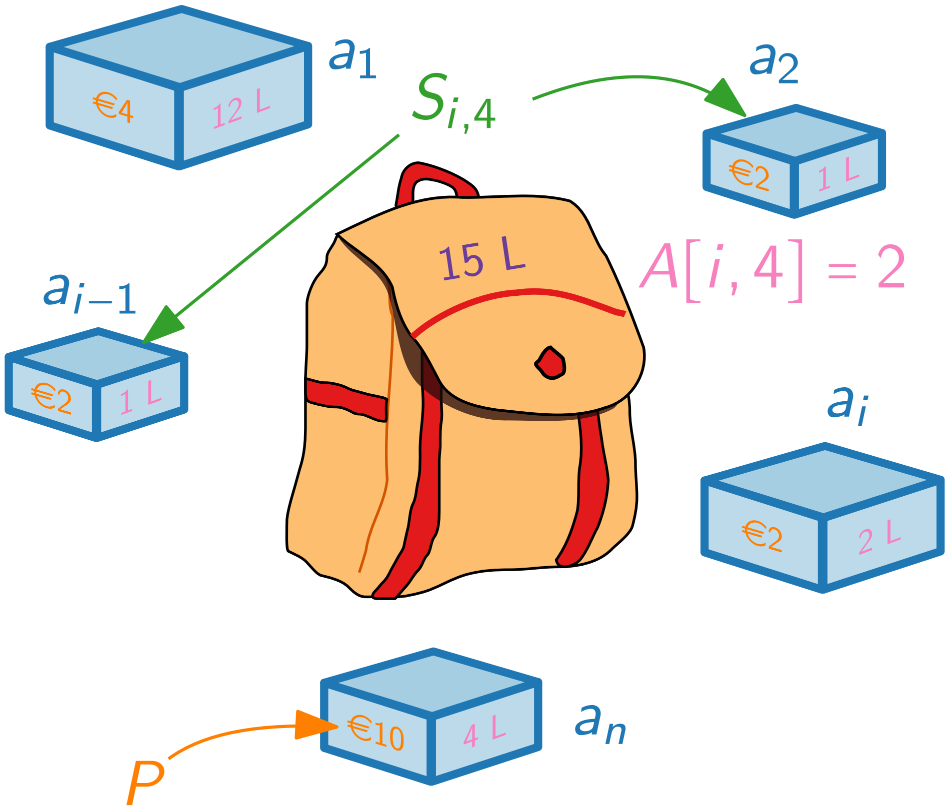
$\Rightarrow$ All values $A[i, p]$ can be computed in total time $O(n^2 P)$.



Pseudo-Polynomial Alg. for KNAPSACK

$A[1, p]$ can be computed for every $p \in \{0, \ldots, nP\}$.

Set $A[i, p] := \infty$ for $p < 0$ (for convenience).

$A[i+1, p] = \min\{A[i, p], \text{size}(a_{i+1}) + A[i, p - \text{profit}(a_{i+1})]\}$

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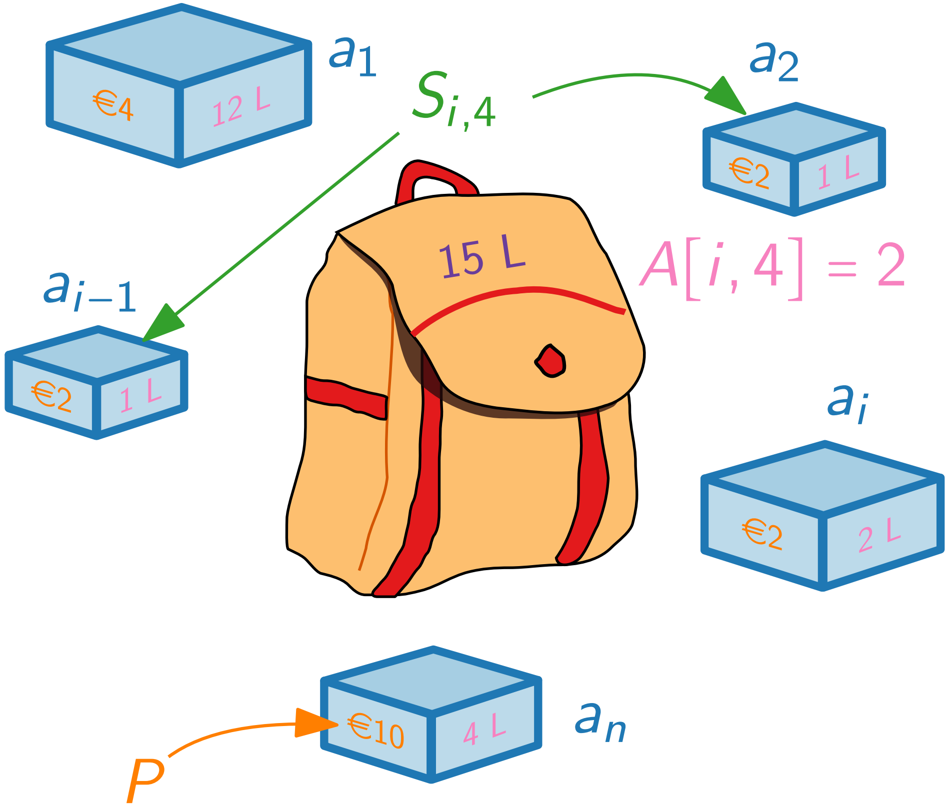

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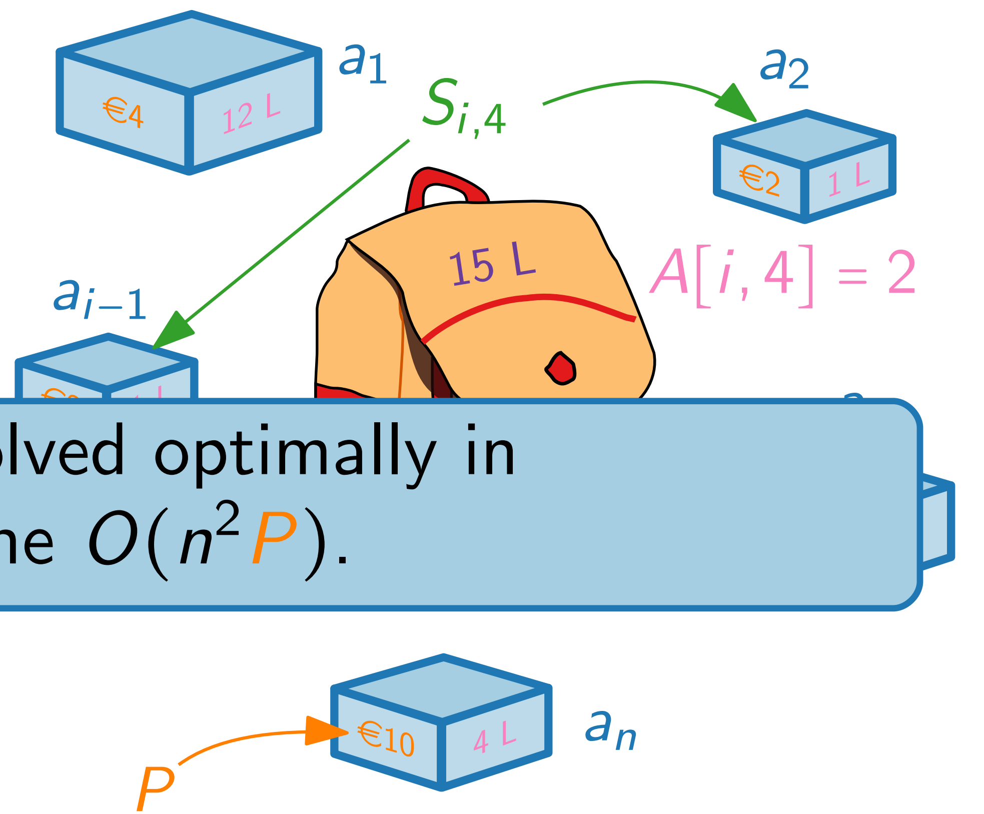

Theorem. KNAPSACK can be solved optimally in pseudo-polynomial time $O(n^2 P)$.

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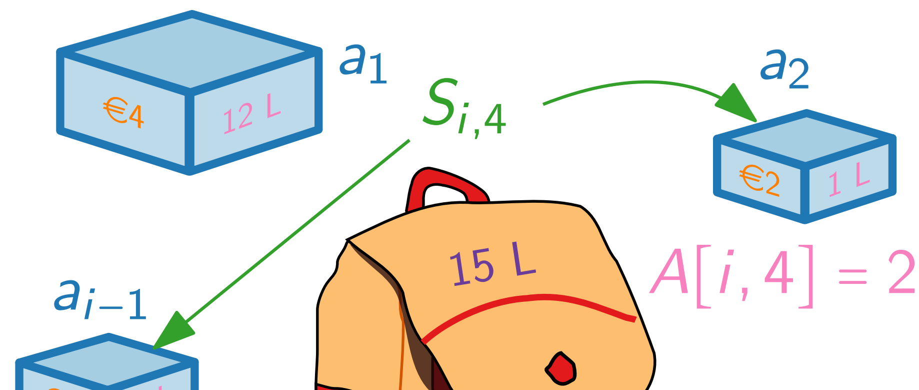
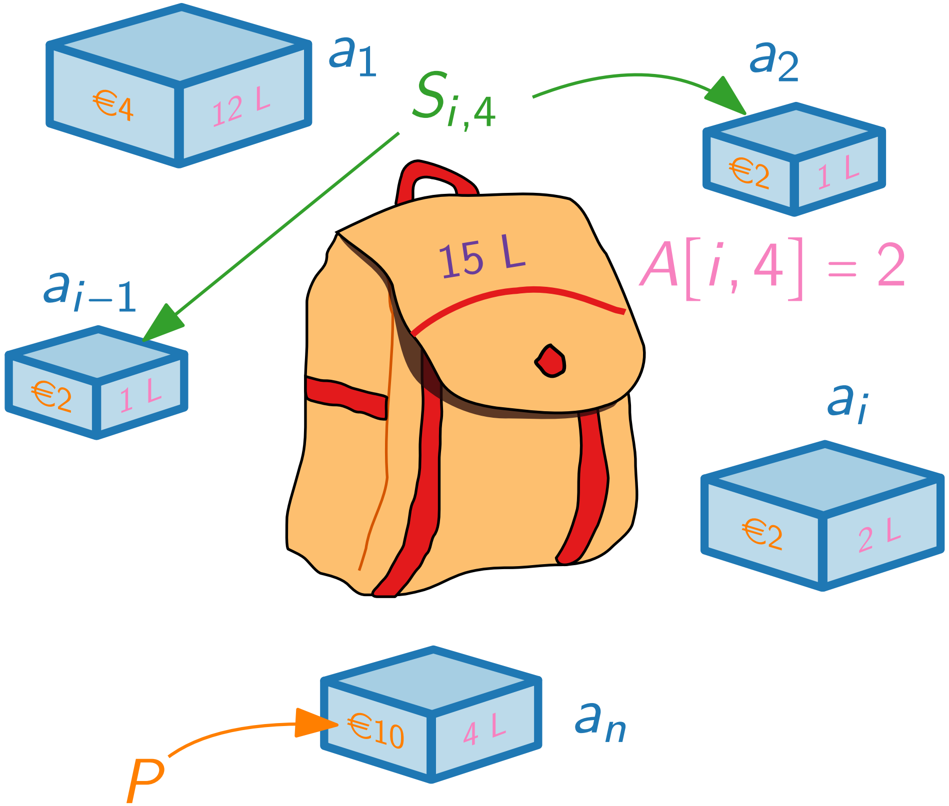


Theorem. KNAPSACK can be solved optimally in pseudo-polynomial time $O(n^2 P)$.

Corollary. KNAPSACK is weakly NP-hard.

Pseudo-Polynomial Alg. for Knapsack

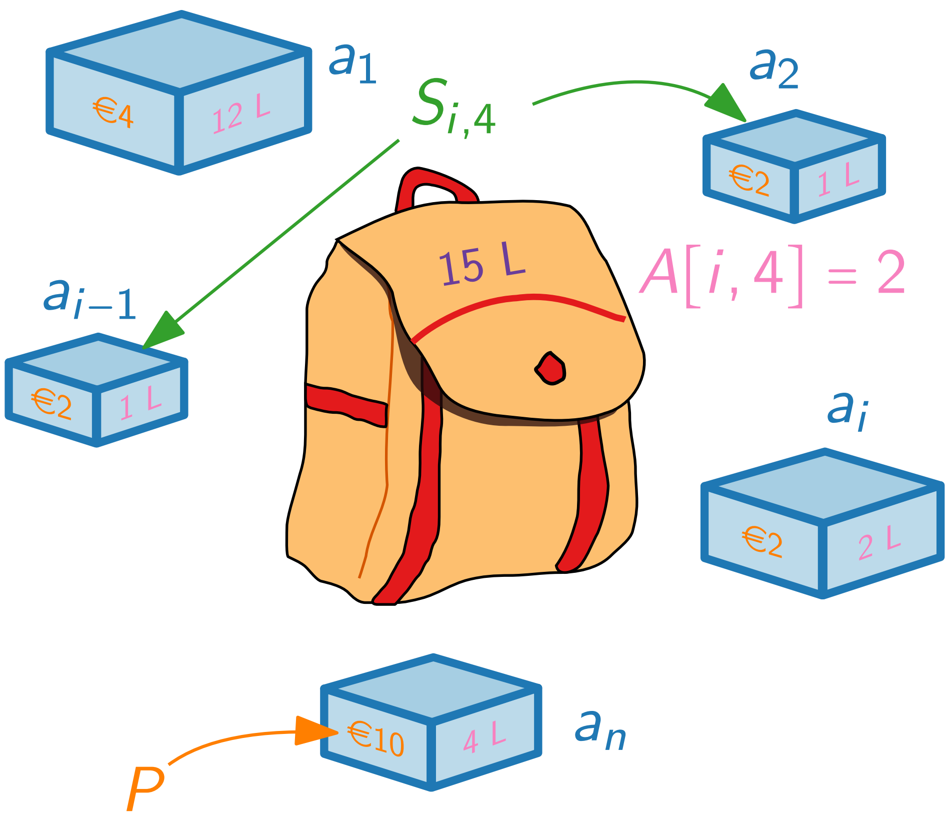
Theorem. Knapsack can be solved optimally in pseudo-polynomial time $O(n^2 P)$.



Pseudo-Polynomial Alg. for KNAPSACK

Theorem. KNAPSACK can be solved optimally in pseudo-polynomial time $O(n^2 P)$.

Observe. The running time $O(n^2 P)$ is polynomial in n if P is polynomial in n.



Approximation Algorithms

Lecture 8:
Approximation Schemes and the KNAPSACK Problem

Part IV:
Approximation Schemes

Approximation Schemes

Let Π be an optimization problem.

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Lecture 8:
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Part V:
FPTAS for KNAPSACK

An FPTAS for KNAPSACK via Scaling

FPTAS idea: **Scale** profits to polynomial size (as required by the error parameter ε)...

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$\Rightarrow K \cdot \sum_i \text{profit}'(o_i) \geq \text{OPT} - \ell K \geq \text{OPT} - nK = \text{OPT} - \varepsilon P$.

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FPTAS idea: **Scale** profits to polynomial size (as required by the error parameter ε)...

An FPTAS for KNAPSACK via Scaling

KnapsackScaling (I, ε)

$K = \varepsilon P / n$ // scaling factor

$\text{profit}'(a_i) = \lfloor \text{profit}(a_i)/K \rfloor$

Compute optimal solution S' for I w.r.t. $\text{profit}'(\cdot)$.

return S'

Lemma. $\text{profit}(S') \geq (1 - \varepsilon) \cdot \text{OPT}$.

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Approximation Algorithms

Lecture 8: Approximation Schemes and the KNAPSACK Problem

Part VI: Connections Between the Concepts

FPTAS and Pseudo-Poly. Algorithms

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FPTAS and Pseudo-Poly. Algorithms

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FPTAS and Strong NP-Hardness

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Corollary. Let Π be an NP-hard optimization problem that fulfills the restrictions above.
If Π is strongly NP-hard, then there is no FPTAS for Π (unless P = NP).