Approximation Algorithms

Lecture 12:
SteinerForest via Primal–Dual

Part I:
SteinerForest

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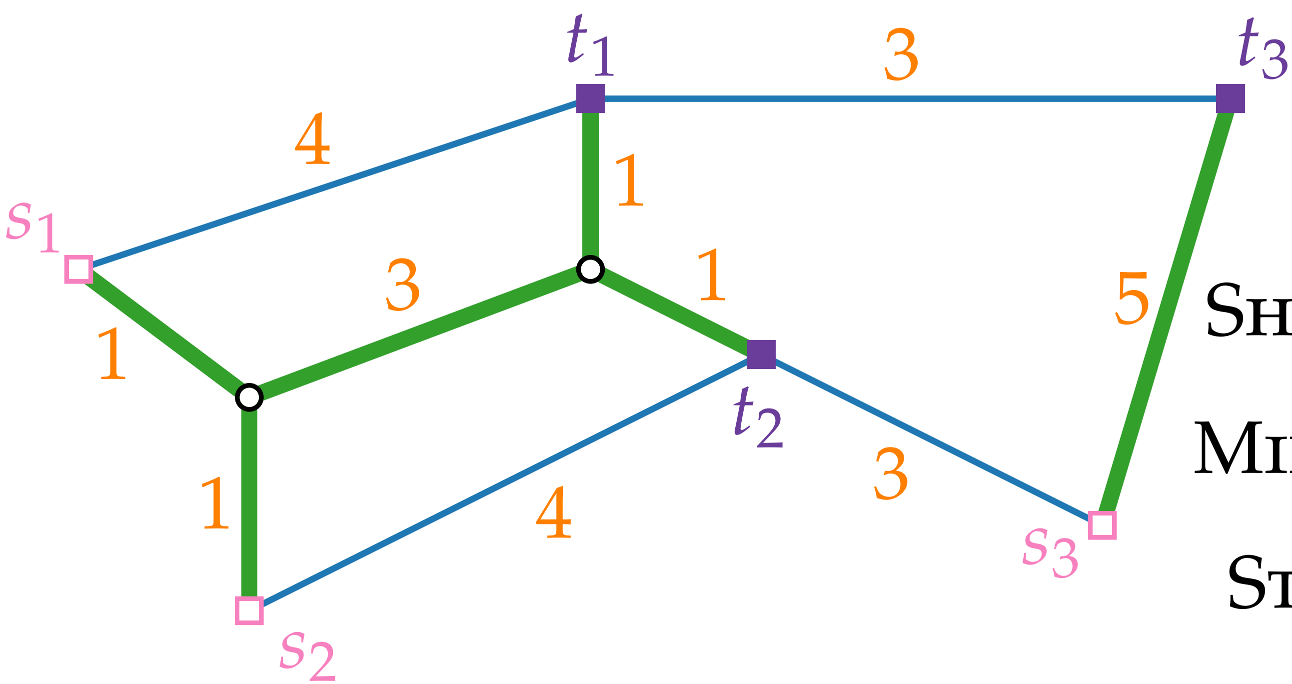
SteinerForest

Given: A graph $G = (V, E)$ with edge costs $c\colon E \to \mathbb{N}$ and a set $R = \{(s_1, t_1), \dots, (s_k, t_k)\}$ of k vertex pairs.

Task: Find an edge set $F \subseteq E$ of minimum total cost $c(F)$ such that the subgraph (V, F) connects every pair (s_i, t_i), $i = 1, \dots, k$.



Special cases?

ShortestPath $(R = \{s, t\})$

MinSpanningTree $(R = E)$

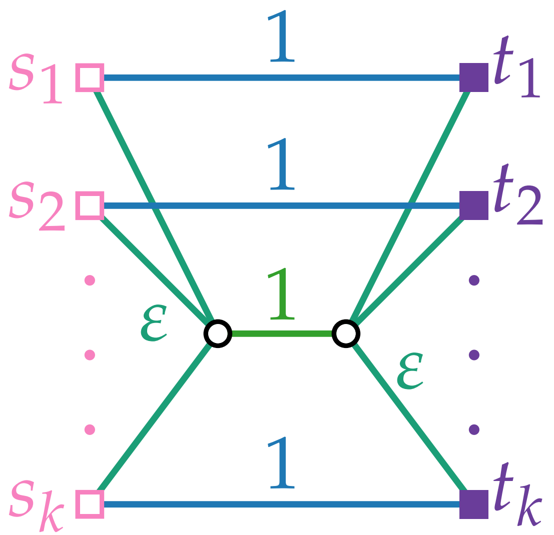
SteinerTree $(R = T \times T)$

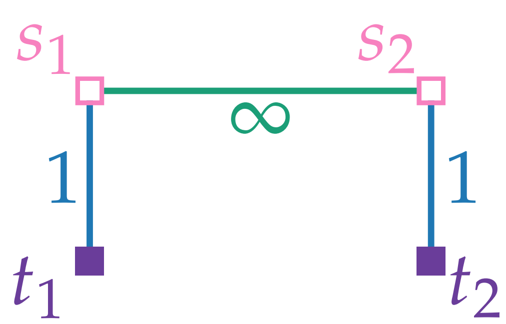
Approaches?

- Merge k shortest s_i–t_i paths
- SteinerTree on the set of terminals

Above approaches perform poorly :-(

Difficulty:
Which terminals belong to the same tree of the forest?





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Part II:
Primal and Dual LP

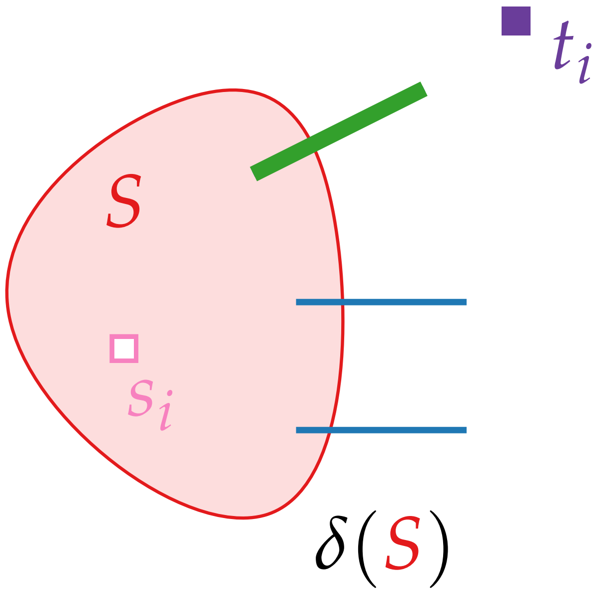
An ILP

$$\begin{array}{lll} \textbf{minimize} & \sum_{e \in E} c_e x_e & \\ \textbf{subject to} & \sum_{e \in \delta(S)} x_e \geq 1 & S \in \mathcal{S}_i,\ i \in \{1, \ldots, k\} \\ & x_e \in \{0,1\} & e \in E \end{array}$$

where $\mathcal{S}_i := \{S \subseteq V : |S \cap \{s_i, t_i\}| = 1\}$

and $\delta(S) := \{(u,v) \in E : u \in S \text{ and } v \notin S\}$

$\rightsquigarrow$ exponentially many constraints!



LP-Relaxation and Dual LP

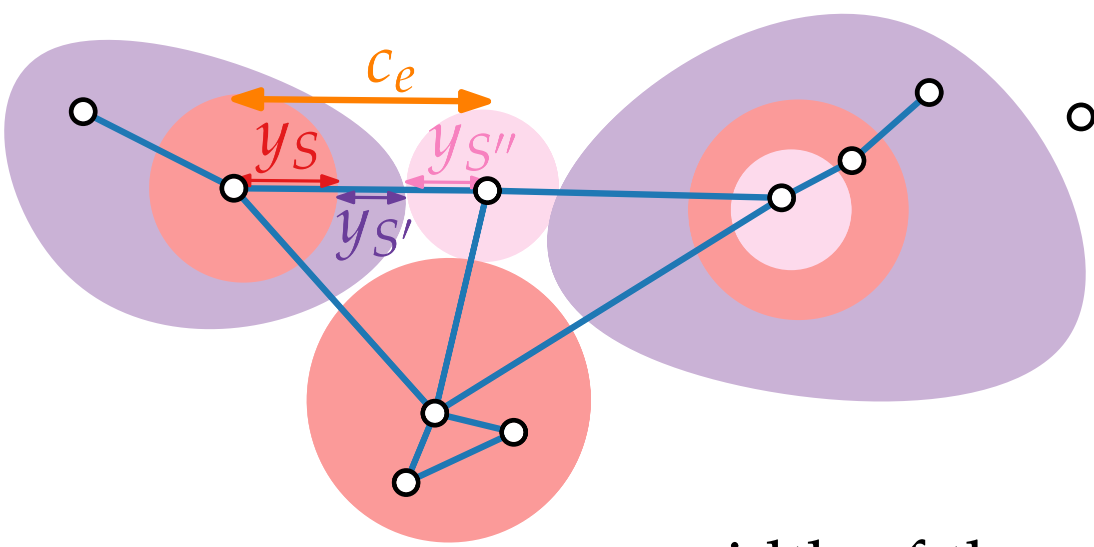
$$\begin{array}{lll} \textbf{minimize} & \displaystyle\sum_{e \in E} c_e x_e & \\ \textbf{subject to} & \displaystyle\sum_{e \in \delta(S)} x_e \geq 1 & S \in \mathcal{S}_i,\, i \in \{1, \ldots, k\} \;(y_S) \\ & x_e \geq 0 & e \in E \end{array}$$

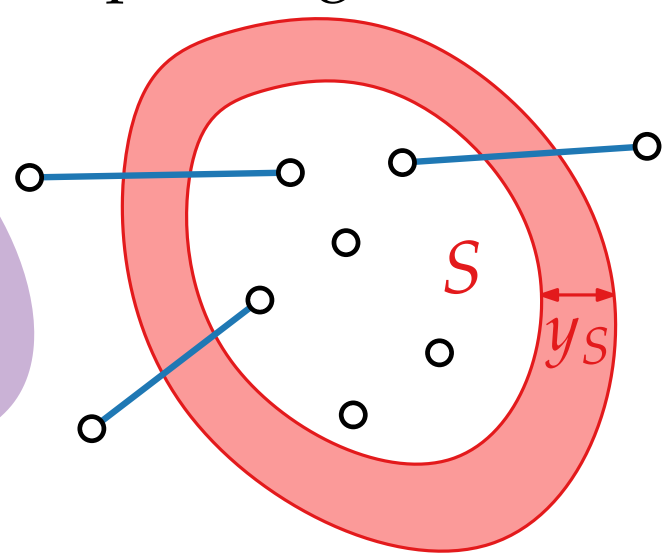
$$\begin{array}{lll} \textbf{maximize} & \displaystyle\sum_{\substack{S \in \mathcal{S}_i \\ i \in \{1,\ldots,k\}}} y_S & \\ \textbf{subject to} & \displaystyle\sum_{S:\, e \in \delta(S)} y_S \leq c_e & e \in E \\ & y_S \geq 0 & S \in \mathcal{S}_i,\, i \in \{1, \ldots, k\} \end{array}$$

Intuition for the Dual

$$\begin{array}{lll} \textbf{maximize} & \displaystyle\sum_{\substack{S \in \mathcal{S}_i \\ i \in \{1,\dots,k\}}} y_S & \\ \textbf{subject to} & \displaystyle\sum_{S:\, e \in \delta(S)} y_S \leq c_e & e \in E \\ & y_S \geq 0 & S \in \mathcal{S}_i,\, i \in \{1,\dots,k\} \end{array}$$

The graph is a network of **bridges**, spanning the **moats**.





$\delta(S)$ = set of edges / bridges over the moat around S

y_S = width of the **moat** around S

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Part III:
A First Primal–Dual Approach

Complementary Slackness (Rep.)

$$\begin{array}{llll} \textbf{minimize} & c^\intercal x & & \\ \textbf{subject to} & Ax & \geq & b \\ & x & \geq & 0 \end{array}$$

$$\begin{array}{llll} \textbf{maximize} & b^\intercal y & & \\ \textbf{subject to} & A^\intercal y & \leq & c \\ & y & \geq & 0 \end{array}$$

Theorem. Let $x = (x_1, \ldots, x_n)$ and $y = (y_1, \ldots, y_m)$ be valid solutions for the primal and dual program (resp.). Then x and y are optimal if and only if the following conditions are met:

Primal CS:
For each $j = 1, \ldots, n$: either $x_j = 0$ or $\sum_{i=1}^{m} a_{ij} y_i = c_j$

Dual CS:
For each $i = 1, \ldots, m$: either $y_i = 0$ or $\sum_{j=1}^{n} a_{ij} x_j = b_i$

A First Primal–Dual Approach

Complementary slackness: $x_e > 0 \quad \Rightarrow \sum_{S:\, e \in \delta(S)} y_S = c_e$.

$\Rightarrow$ pick "critical" edges (and only those)

Idea: iteratively build a feasible integral primal solution.

How to find a violated primal constraint? $\left(\sum_{e \in \delta(S)} x_e < 1\right)$

$\rightsquigarrow$ Consider related connected component C!

How do we iteratively improve the dual solution?

$\rightsquigarrow$ Increase y_C (until some edge in $\delta(C)$ becomes critical)!

A First Primal–Dual Approach

PrimalDualSteinerForestNaive(G, c, R)

$y \leftarrow 0, F \leftarrow \emptyset$

while some $(s_i, t_i) \in R$ not connected in (V, F) **do**

  $C \leftarrow$ comp. in (V, F) with $|C \cap \{s_i, t_i\}| = 1$ for some i

  Increase y_C

    until $\sum_{S \colon e' \in \delta(S)} y_S = c_{e'}$ for some $e' \in \delta(C)$.

  $F \leftarrow F \cup \{e'\}$

return F

Running Time?

Trick: Handle all y_S with $y_S = 0$ implicitly

Analysis

The cost of the solution F can be written as

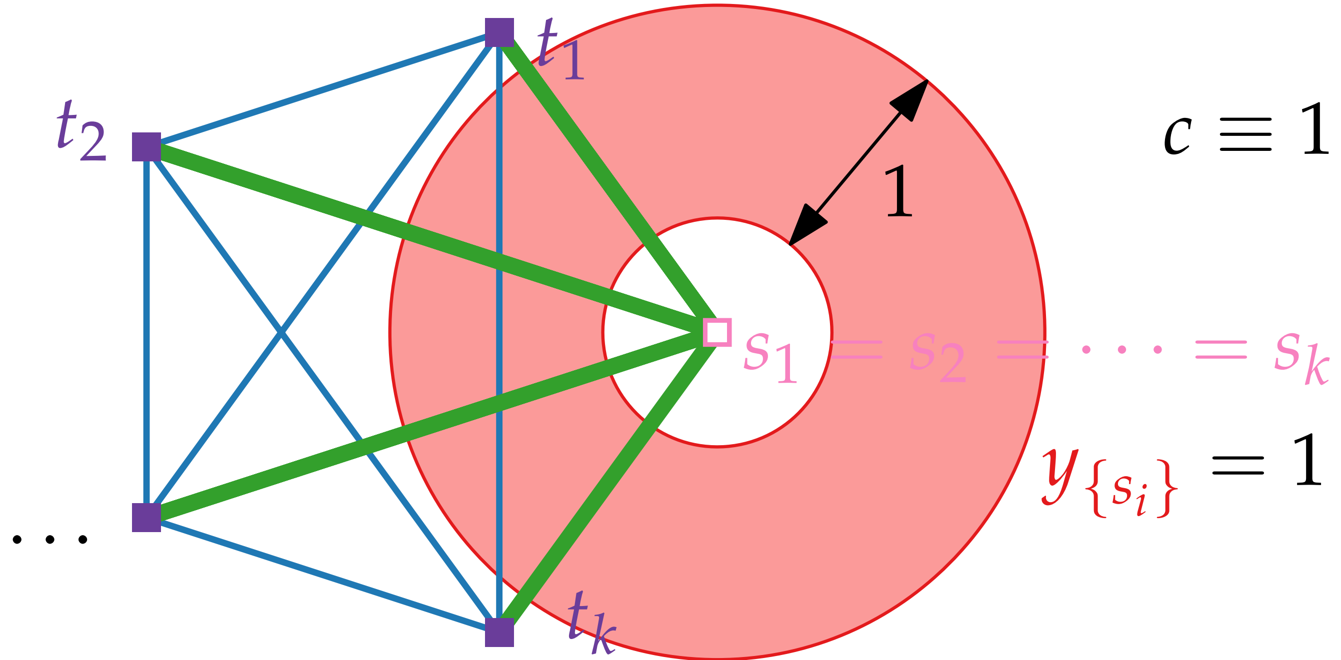
$$\sum_{e \in F} c_e \overset{\text{CS}}{=} \sum_{e \in F} \sum_{S\colon e \in \delta(S)} y_S = \sum_S |\delta(S) \cap F| \cdot y_S.$$

Compare to the value of the dual objective function $\sum_S y_S$

There are examples with $|\delta(S) \cap F| = k$ for each $y_S > 0$:

But: Average degree of component is 2!

$\Rightarrow$ Increase y_C for all components C simultaneously!



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Part IV:
Primal–Dual with Synchronized Increases

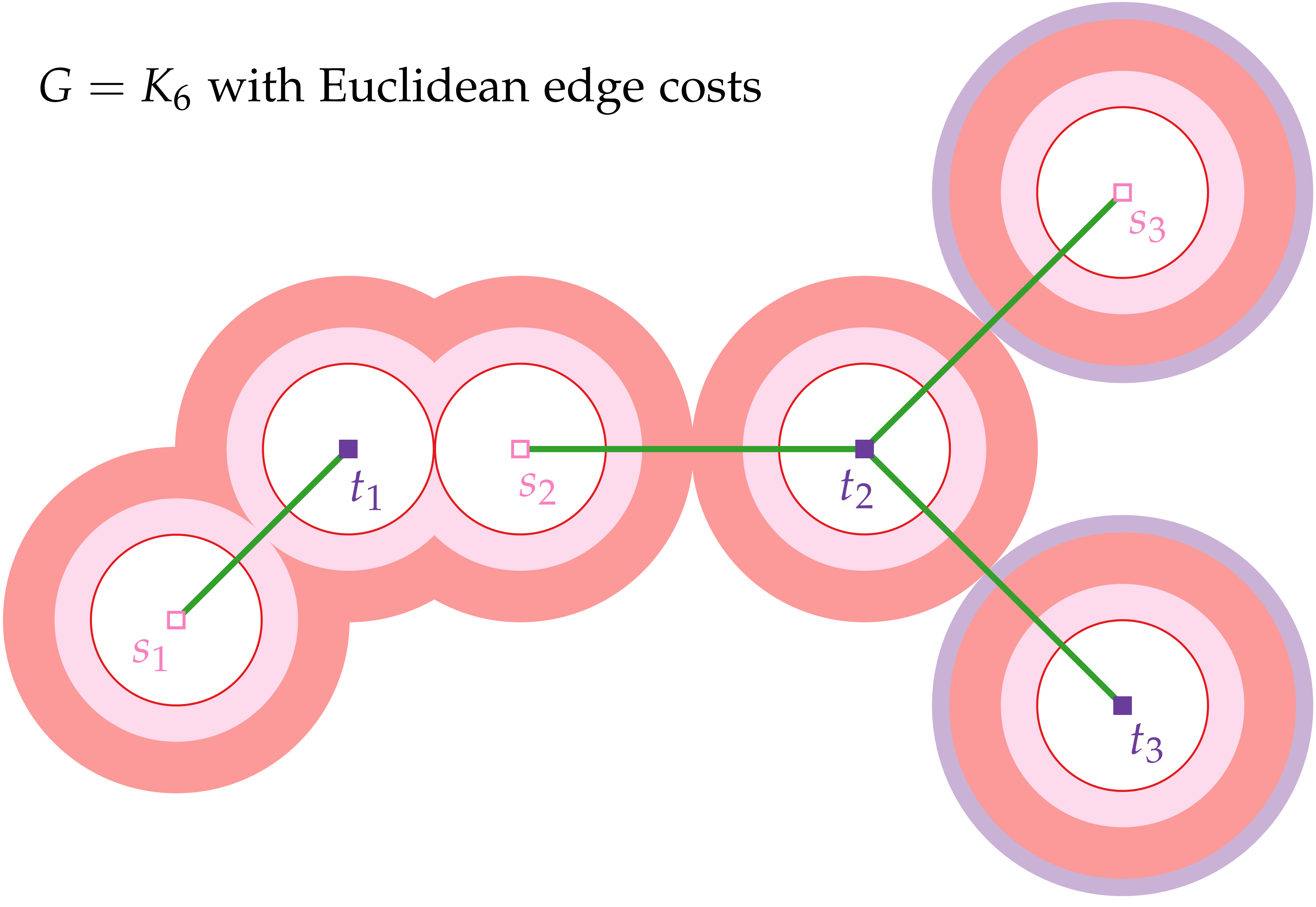
Primal–Dual with Synchronized Increases

PrimalDualSteinerForest(G, c, R)

$y \leftarrow 0, F \leftarrow \emptyset, \ell \leftarrow 0$
while some $(s_i, t_i) \in R$ not connected in (V, F) **do**
    $\ell \leftarrow \ell + 1$
    $\mathcal{C} \leftarrow \{$comp. C in (V, F) with $|C \cap \{s_i, t_i\}| = 1$ for some $i\}$
    Increase y_C for all $C \in \mathcal{C}$ simultaneously
        until $\sum_{S\colon e_\ell \in \delta(S)} y_S = c_{e_\ell}$ for some $e_\ell \in \delta(C)$, $C \in \mathcal{C}$.
    $F \leftarrow F \cup \{e_\ell\}$
$F' \leftarrow F$
// Pruning
for $j \leftarrow \ell$ **down to** 1 **do**
    if $F' \setminus \{e_j\}$ is feasible solution **then**
        $F' \leftarrow F' \setminus \{e_j\}$
return F'

Illustration

$G = K_6$ with Euclidean edge costs



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Part V:
Structure Lemma

Structure Lemma

Lemma. For the set $\mathcal{C}$ in any iteration of the algorithm:

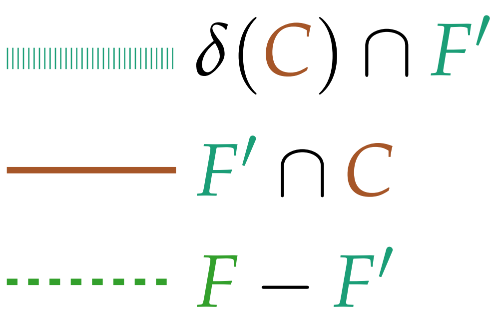
$$\sum_{C \in \mathcal{C}} |\delta(C) \cap F'| \leq 2|\mathcal{C}|\,.$$

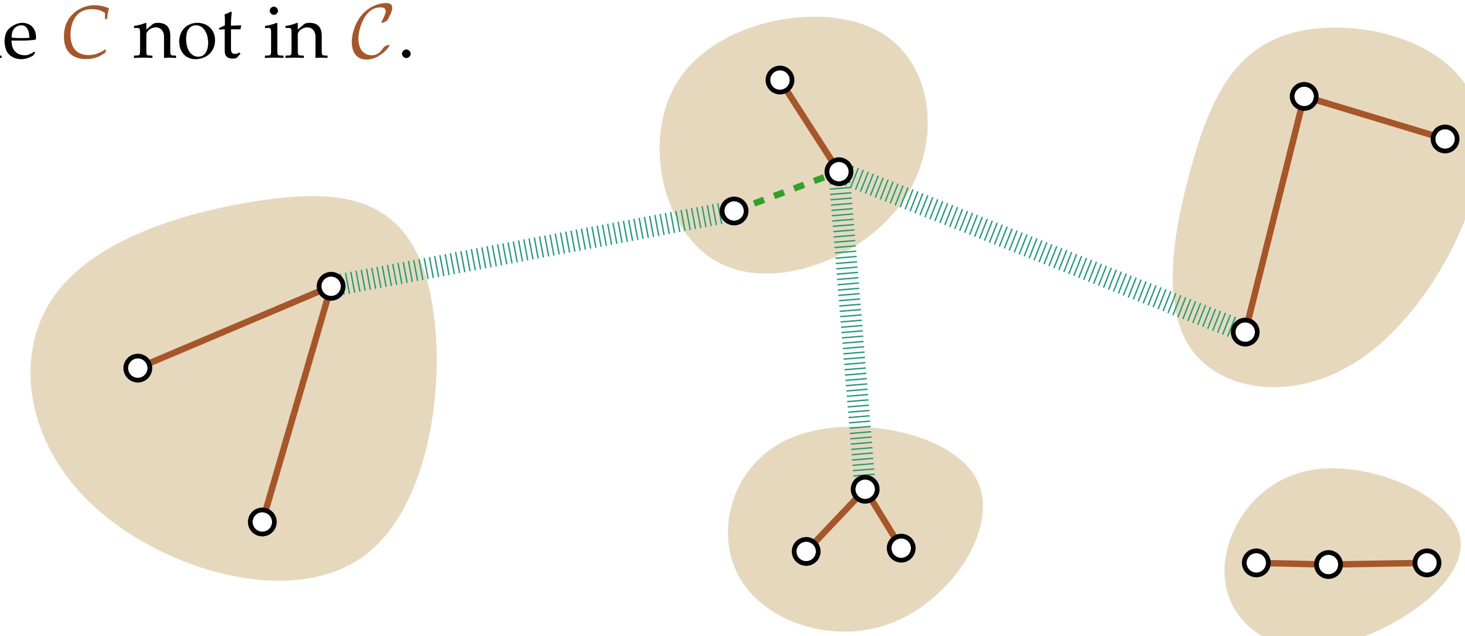
Proof. First the intuition...

Every connected component C of F is a forest in F'.

$\rightsquigarrow$ average degree ≤ 2

Difficulty: Some C not in $\mathcal{C}$.





Proof of the Structure Lemma

Lemma. For the set $\mathcal{C}$ in any iteration of the algorithm:

$$\sum_{C \in \mathcal{C}} |\delta(C) \cap F'| \leq 2|\mathcal{C}|\,.$$

Proof.

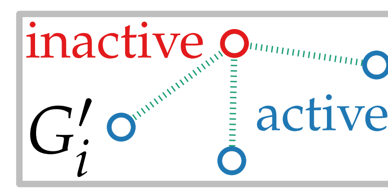
For $i = 1, \ldots, \ell$, consider i-th iteration (when e_i was added to F).

Let $F_i = \{e_1, \ldots, e_i\}$, $G_i = (V, F_i)$, and $G_i^* = (V, F_i \cup F')$.

Contract every component C of G_i in G_i^* to a single vertex $\rightsquigarrow G_i'$.

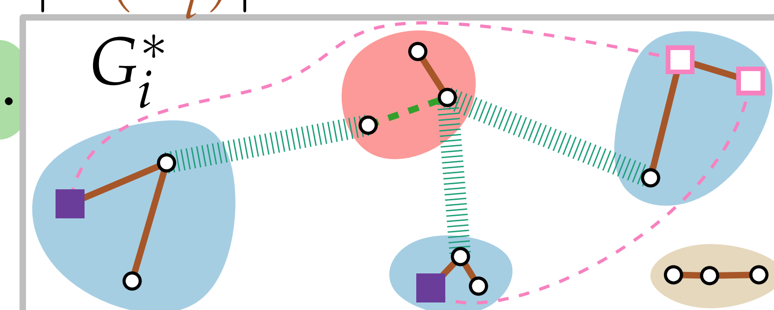
(Ignore components C with $\delta(C) \cap F' = \emptyset$.)

Claim. G_i' is a forest.



Note: $\sum_{C \text{ comp.}} |\delta(C) \cap F'| = \sum_{v \in V(G_i')} \deg_{G'}(v)$
$= 2|E(G_i')| \leq 2|V(G_i')|$

Claim. Inactive vertices have degree ≥ 2.



$\Rightarrow \sum_{v \text{ active}} \deg_{G'}(v) \leq$
$2 \cdot |V(G')| - 2 \cdot \#(\text{inactive}) = 2|\mathcal{C}|$. $\square$

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Part VI:
Analysis

Analysis

Theorem. The Primal–Dual algorithm with synchronized increases yields a 2-approximation for SteinerForest.

Proof.

As mentioned before,

$$\sum_{e \in F'} c_e \overset{\text{CS}}{=} \sum_{e \in F'} \sum_{S:\, e \in \delta(S)} y_S = \sum_S |\delta(S) \cap F'| \cdot y_S.$$

We prove by induction over the number of iterations of the algorithm that

$$\sum_S |\delta(S) \cap F'| \cdot y_S \leq 2 \sum_S y_S. \qquad (*)$$

From that, the claim of the theorem follows.

Analysis

Theorem. The Primal–Dual algorithm with synchronized increases yields a 2-approximation for SteinerForest.

Proof.

$$\sum_S |\delta(S) \cap F'| \cdot y_S \leq 2 \sum_S y_S. \qquad (*)$$

Base case trivial since we start with $y_S = 0$ for every S.

Assume that $(*)$ holds at the start of the current iteration.

In the current iteration, we increase y_C for every $C \in \mathcal{C}$ by the same amount, say $\varepsilon \geq 0$.

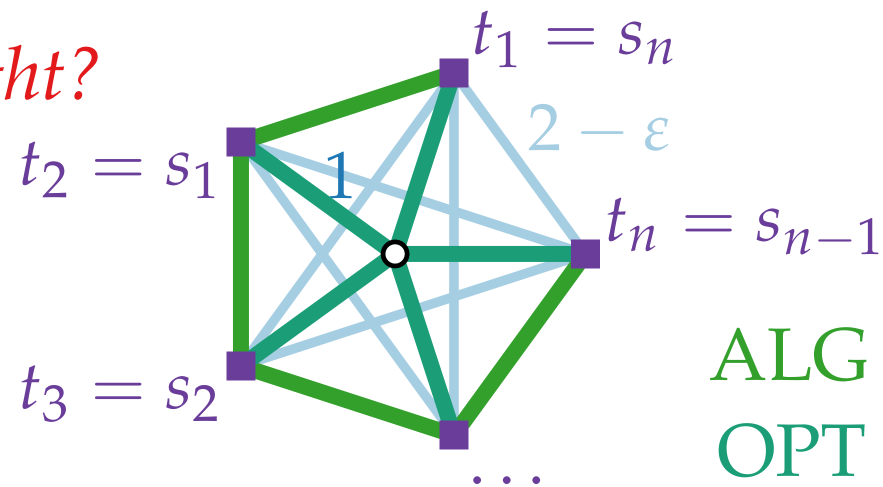
This increases the left side of $(*)$ by $\varepsilon \cdot \sum_{C \in \mathcal{C}} |\delta(C) \cap F'|$ and the right side by $\varepsilon \cdot 2|\mathcal{C}|$.

Structure lemma $\Rightarrow (*)$ also holds after the current iteration. □

Summary

Theorem. The Primal–Dual algorithm with synchronized increases gives a 2-approximation for SteinerForest.

Is our analysis tight?



$\text{ALG} = (2 - \varepsilon)(n - 1)$

$\text{OPT} = n$

Can we do better?

No better approximation factor is known. :-(

The integrality gap is $2 - 1/n$.

SteinerForest (as SteinerTree) cannot be approximated within factor $\frac{96}{95} \approx 1.0105$ (unless P = NP). [Chlebík, Chlebíková '08]