Approximation Algorithms

Lecture 10: Minimum-Degree Spanning Tree via Local Search

Part I: Minimum-Degree Spanning Tree

Alexander Wolff

Winter 2022/23

Minimum-Degree Spanning Tree

Given: A connected graph G.

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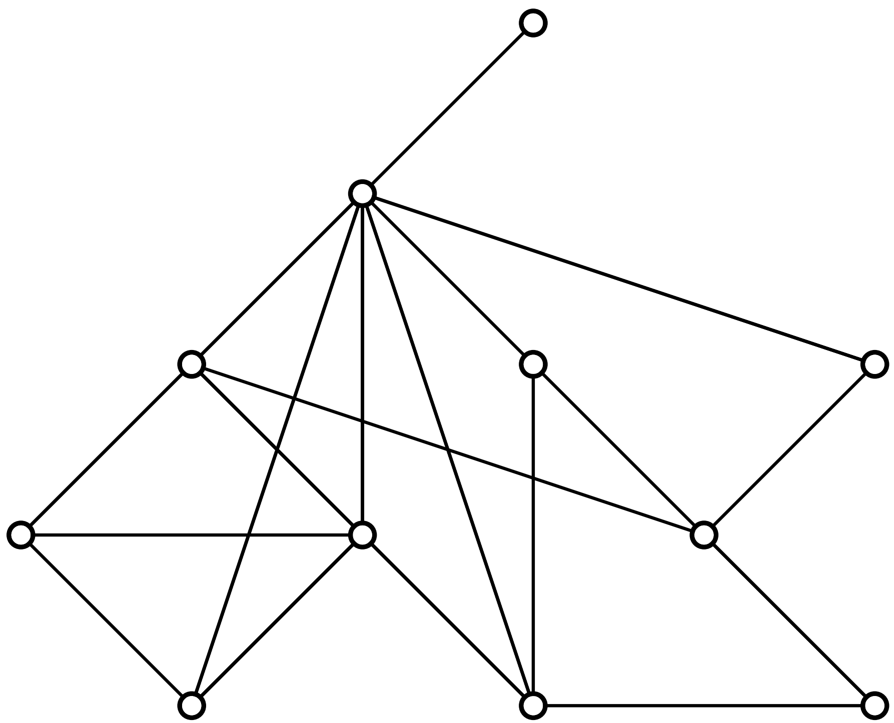
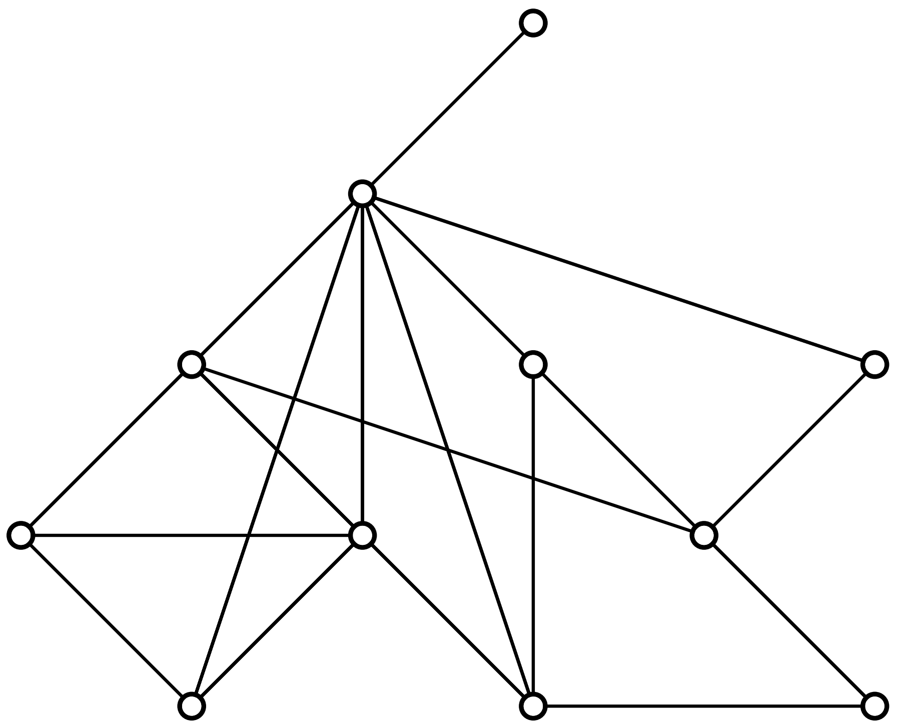


Minimum-Degree Spanning Tree

Given: A connected graph G.

Task: Find a spanning tree T that has the smallest maximum degree $\Delta(T)$ among all spanning trees of G.



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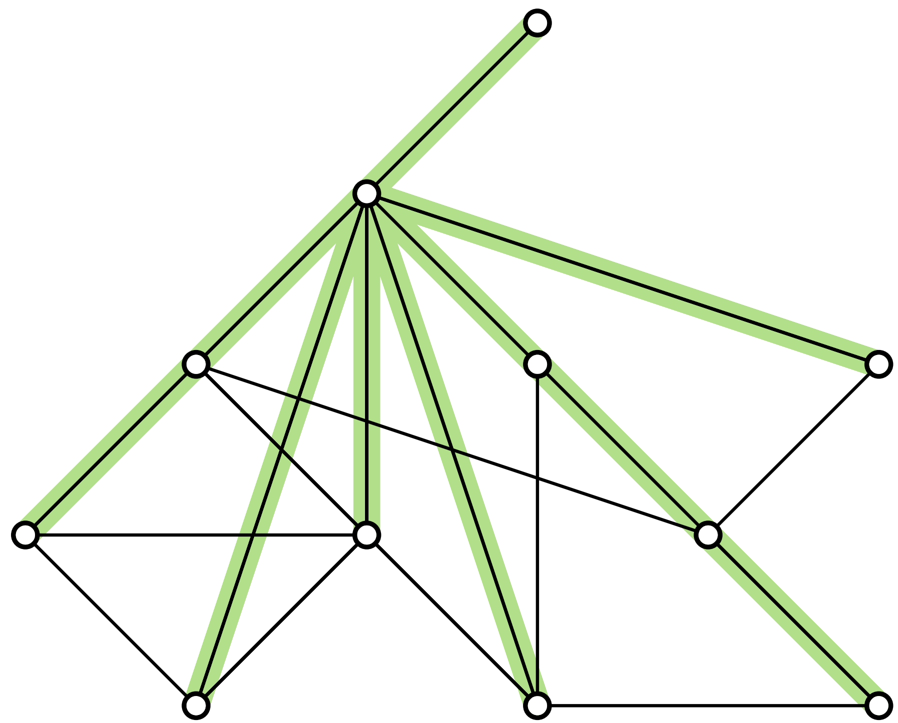

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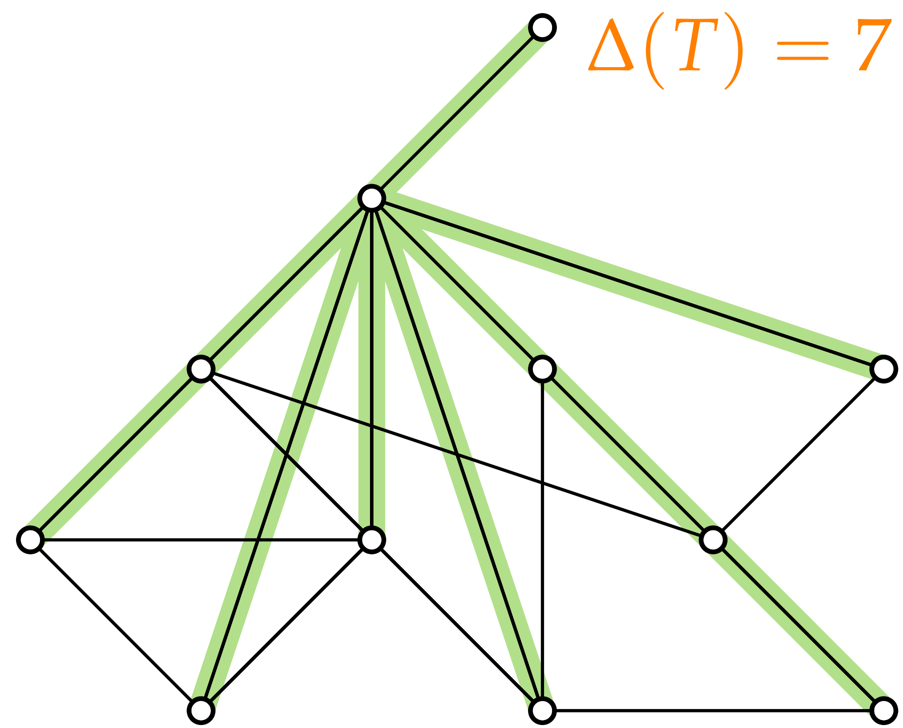

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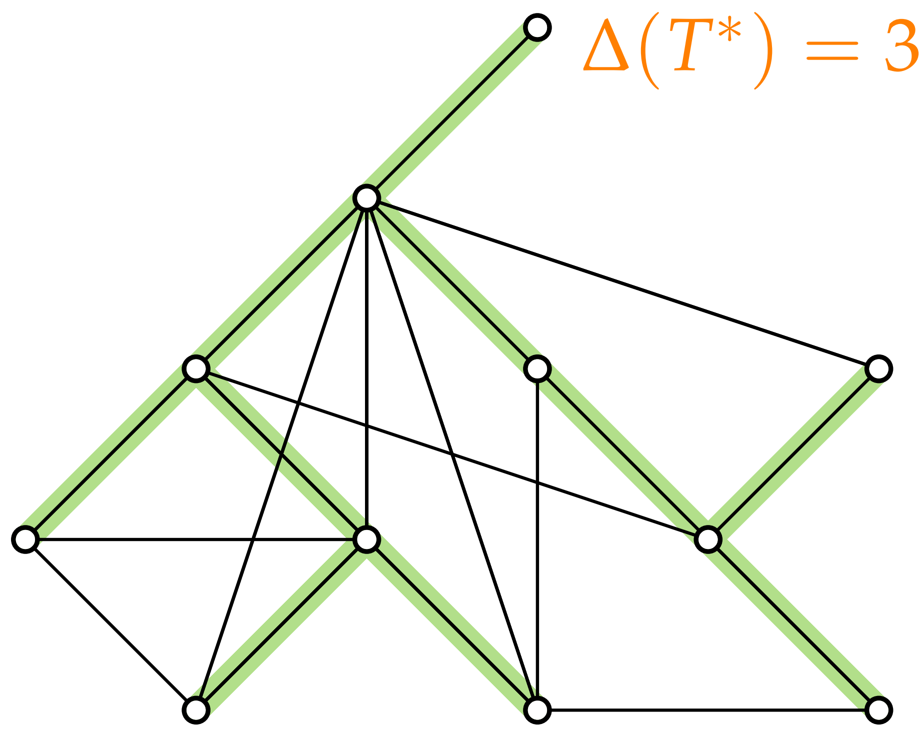

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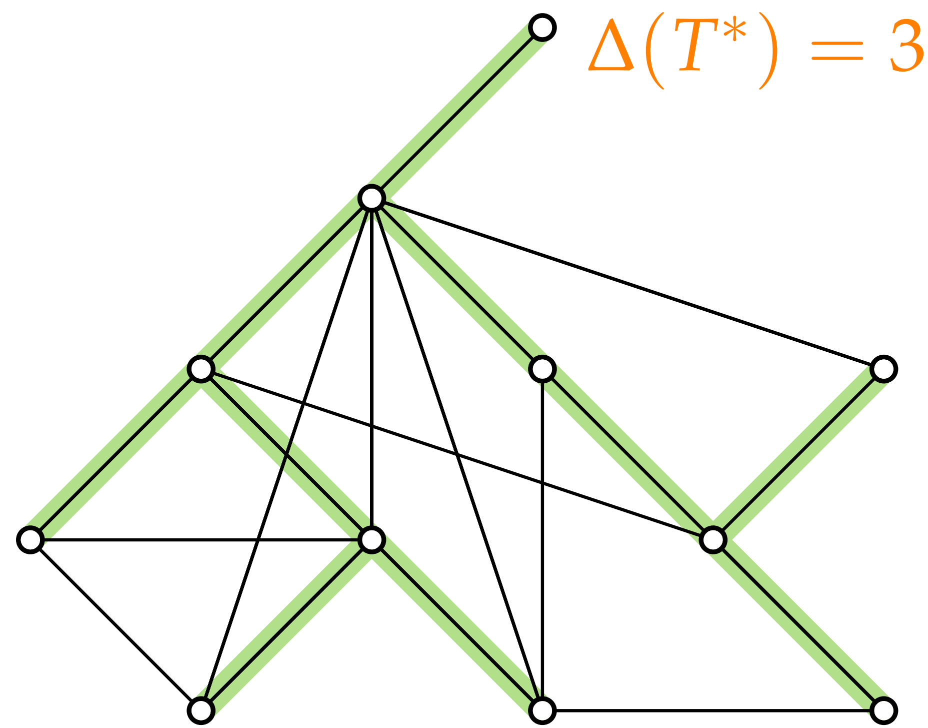
NP-hard. ☹



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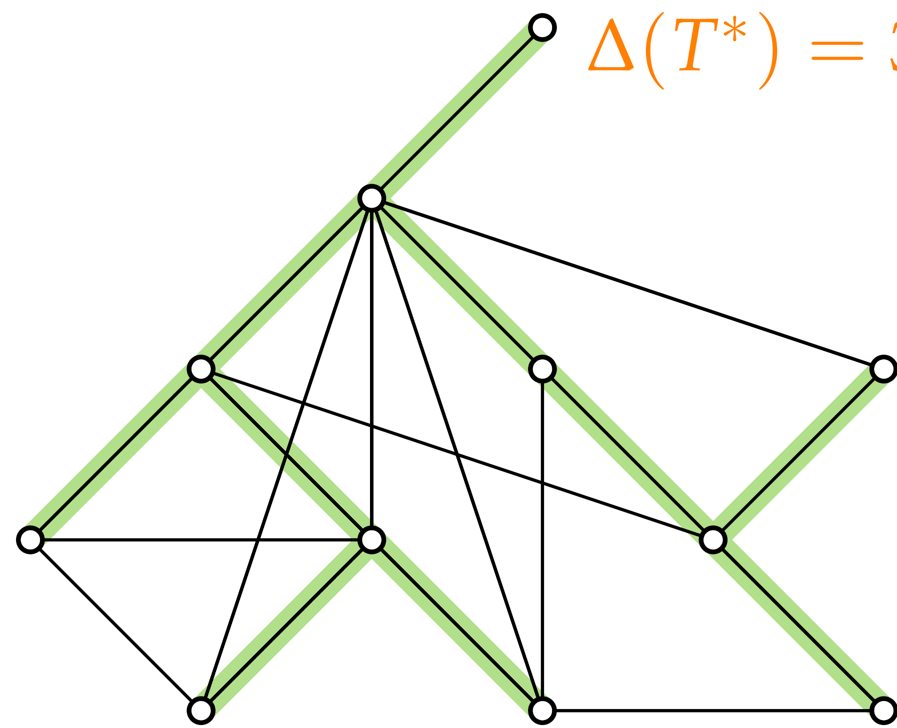
Why?



$\Delta(T^*) = 3$

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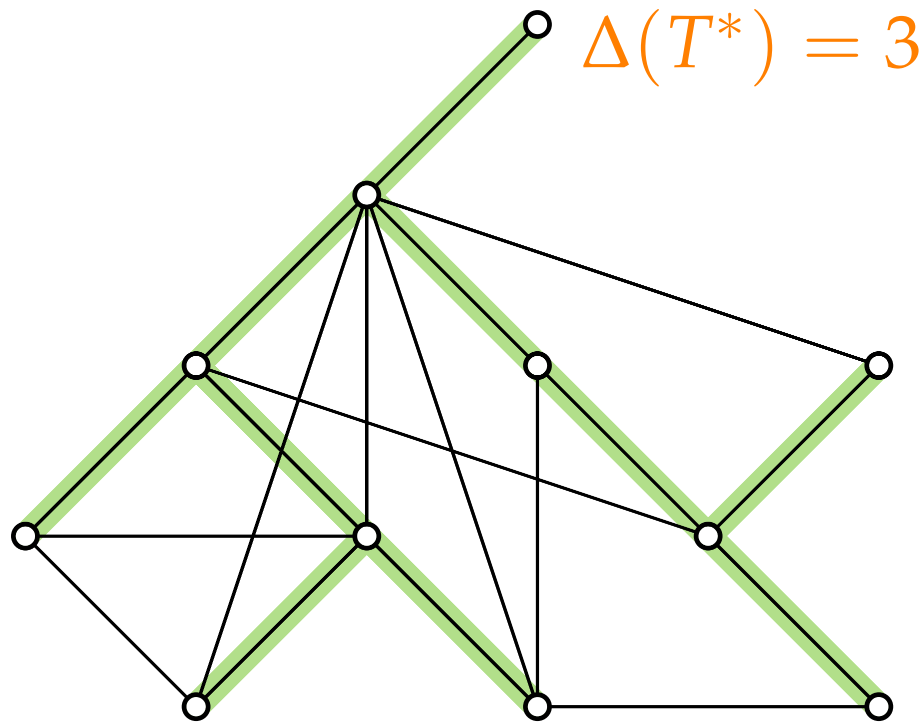
NP-hard. 

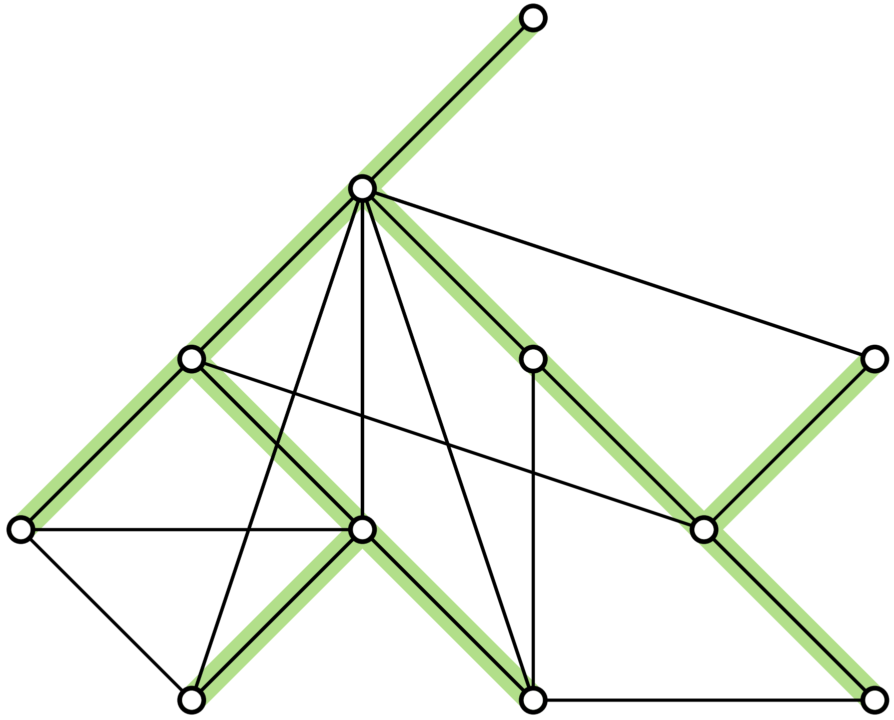
Why?

Special case of Hamiltonian Path!



Warm-up

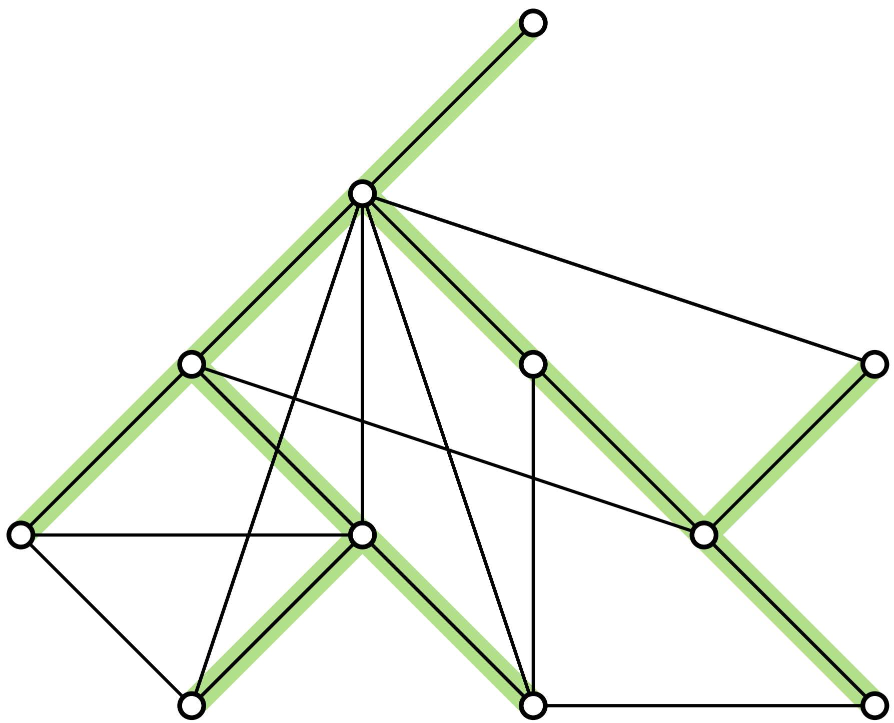
Obs. A spanning tree T has...



Warm-up

Obs. A spanning tree T has...

- n vertices and ? edges,



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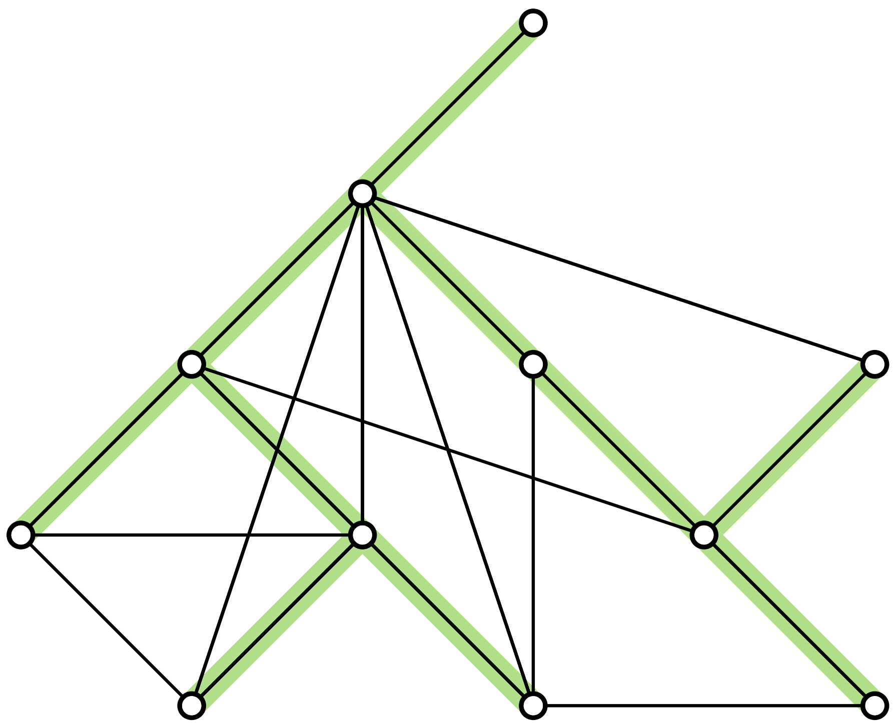
Obs. A spanning tree T has...

- n vertices and ? edges,
- sum of degrees $\sum_{v \in V} \deg_T(v) =$?



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Obs. A spanning tree T has...

- n vertices and ? edges,
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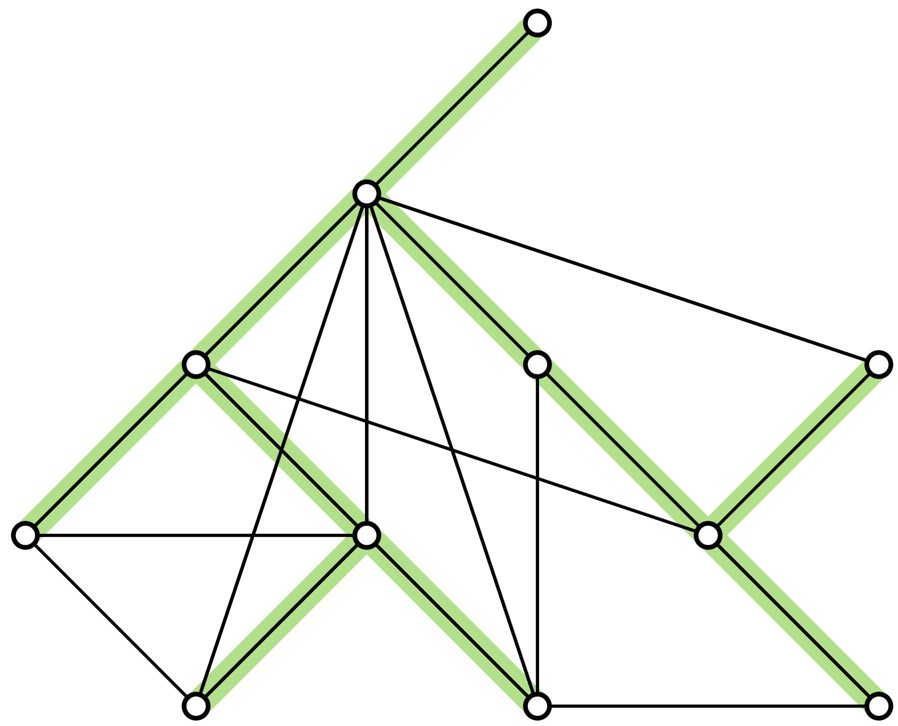
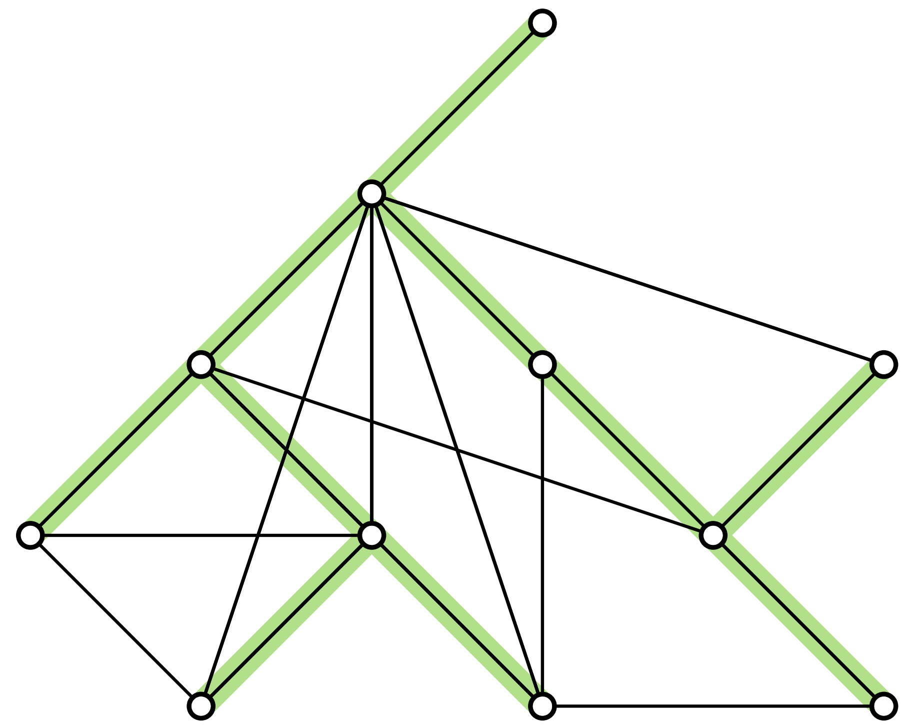
- average degree ?



Warm-up

Obs. A spanning tree T has...

- n vertices and $n-1$ edges,
- sum of degrees $\sum_{v \in V} \deg_T(v) =$?
- average degree ?



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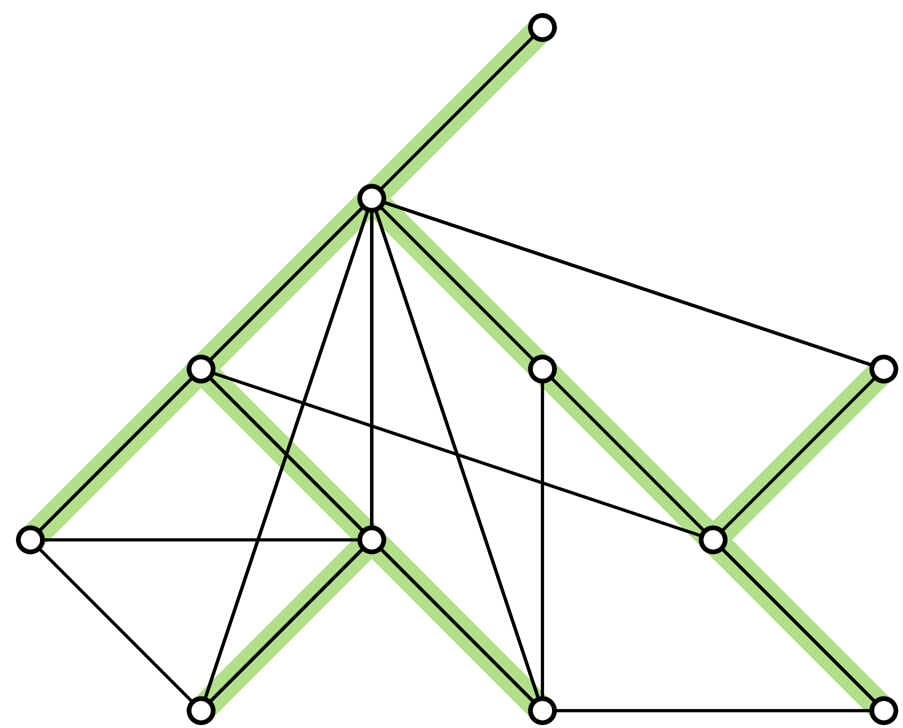
Obs. A spanning tree T has...

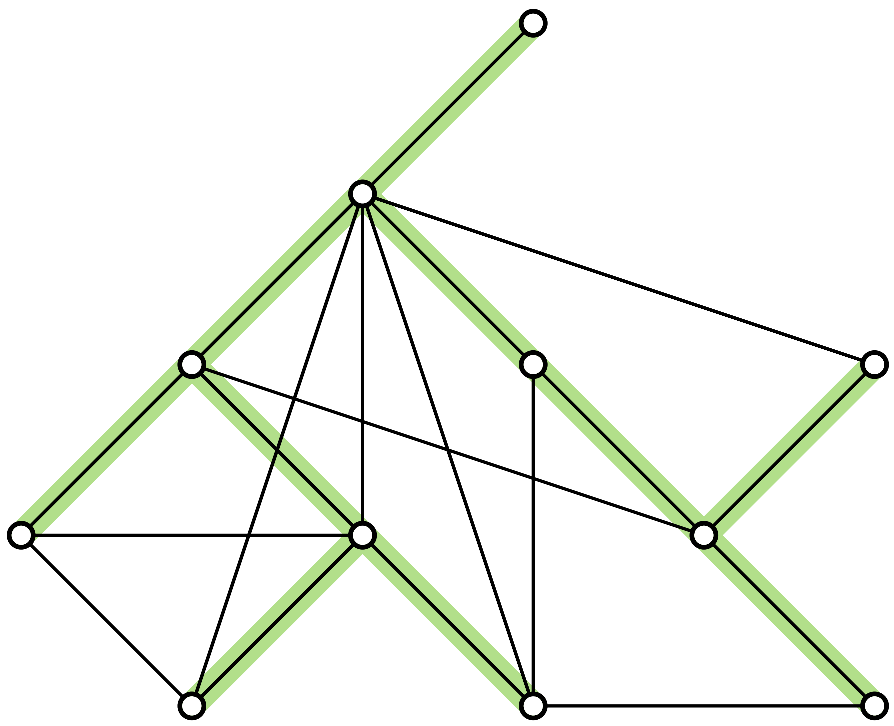
- n vertices and $n-1$ edges,
- sum of degrees $\sum_{v \in V} \deg_T(v) = 2n-2$,
- average degree ?



Warm-up

Obs. A spanning tree T has...

- n vertices and $n-1$ edges,
- sum of degrees $\sum_{v \in V} \deg_T(v) = 2n-2$,
- average degree < 2.



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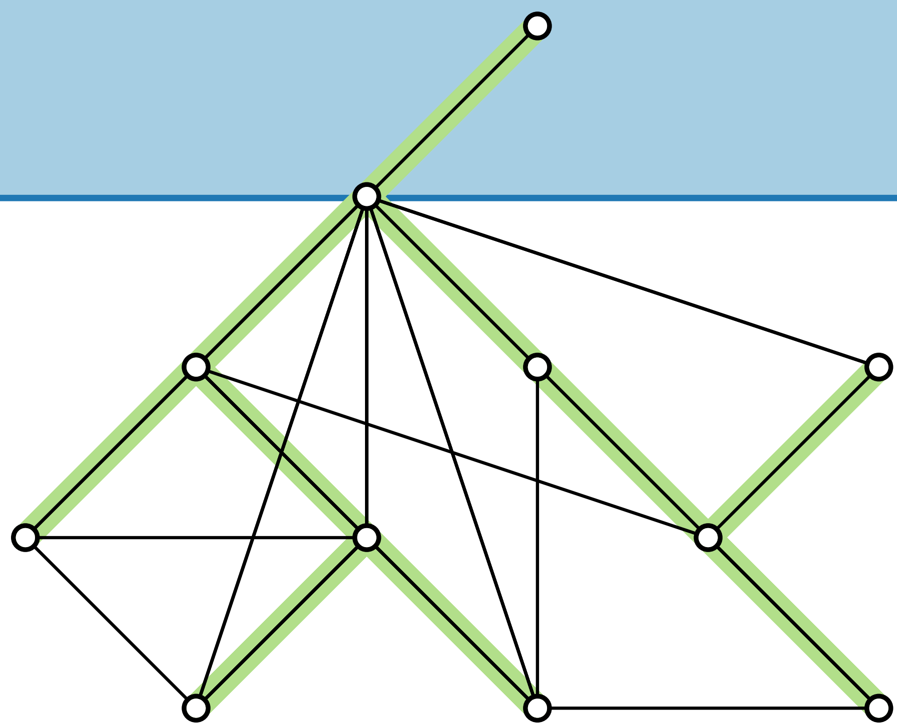
Obs. Let $V' \subseteq V(G)$.
Then $\Delta(G) \geq$?



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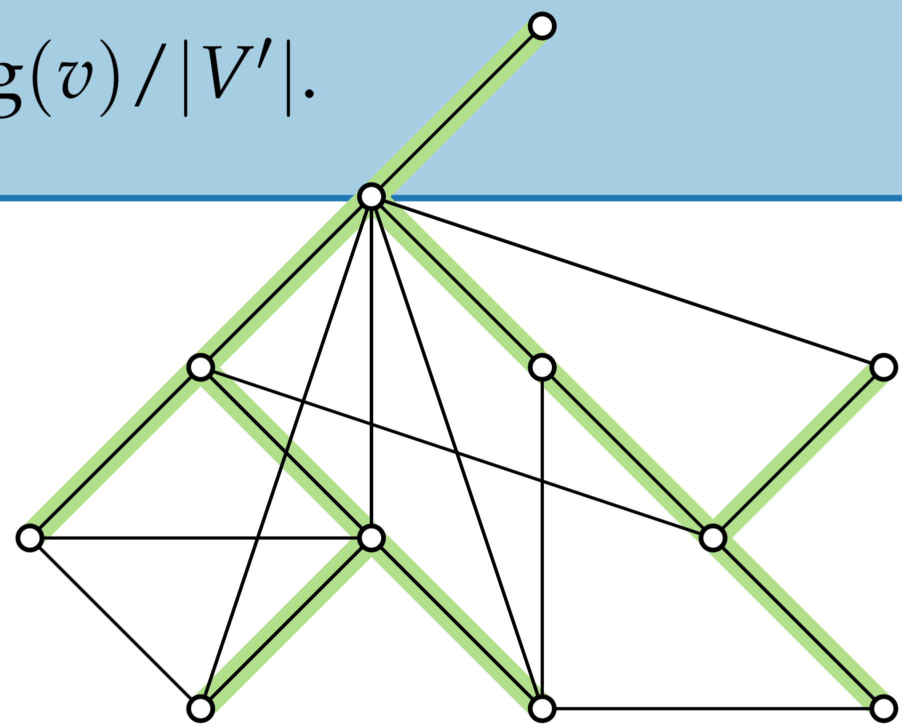
Obs. Let $V' \subseteq V(G)$.
Then $\Delta(G) \geq \sum_{v \in V'} \deg(v) / |V'|$.



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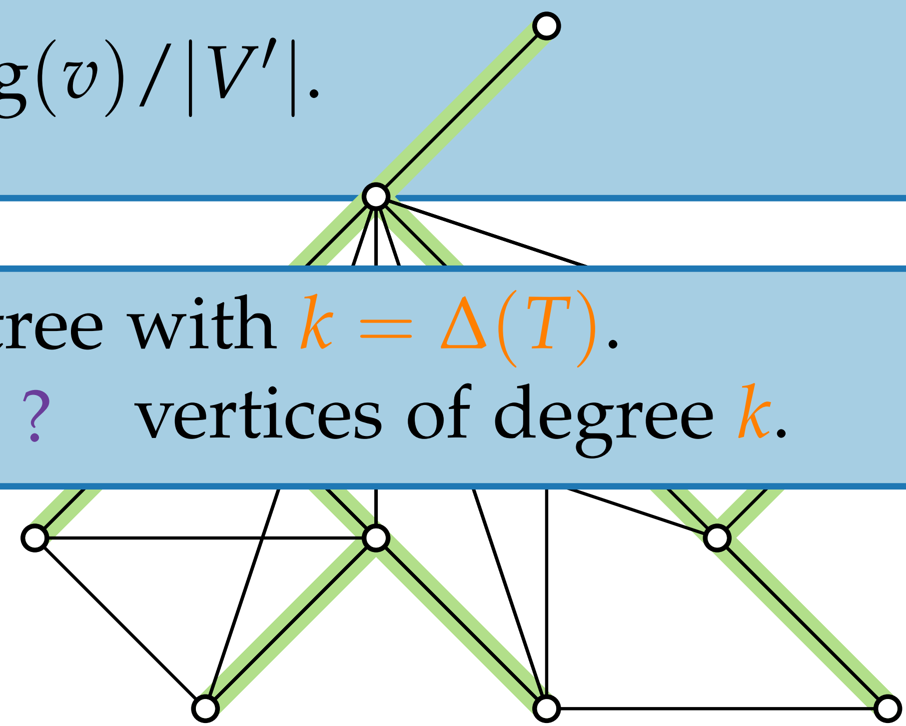
Obs. Let T be a spanning tree with $k = \Delta(T)$.
Then T has at most ? vertices of degree k.



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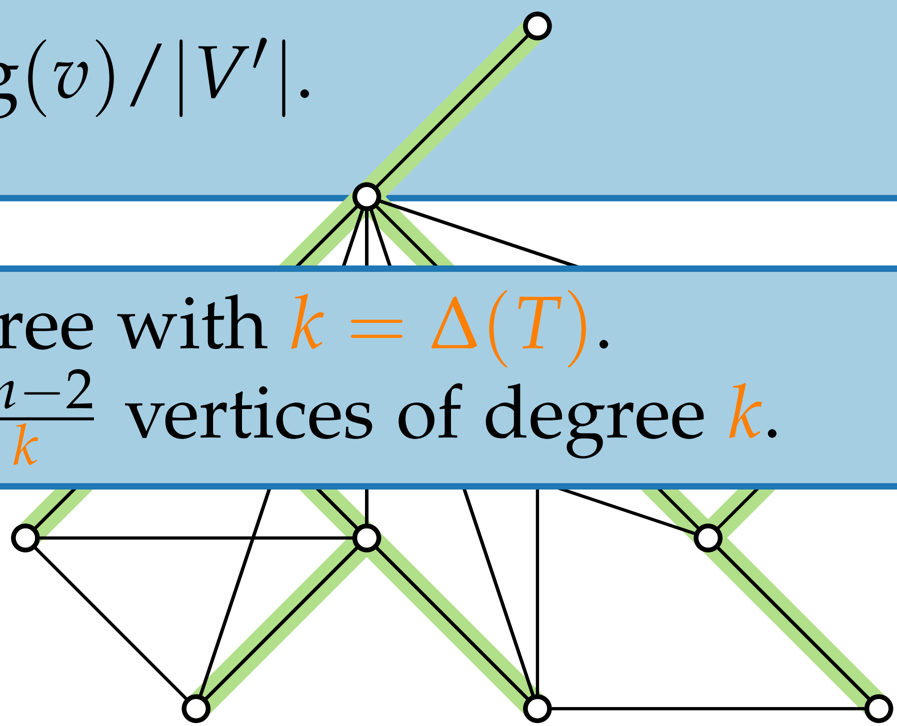
Then $\Delta(G) \geq \sum_{v \in V'} \deg(v) / |V'|$.

Obs. Let T be a spanning tree with $k = \Delta(T)$.
Then T has at most $\frac{2n-2}{k}$ vertices of degree k.



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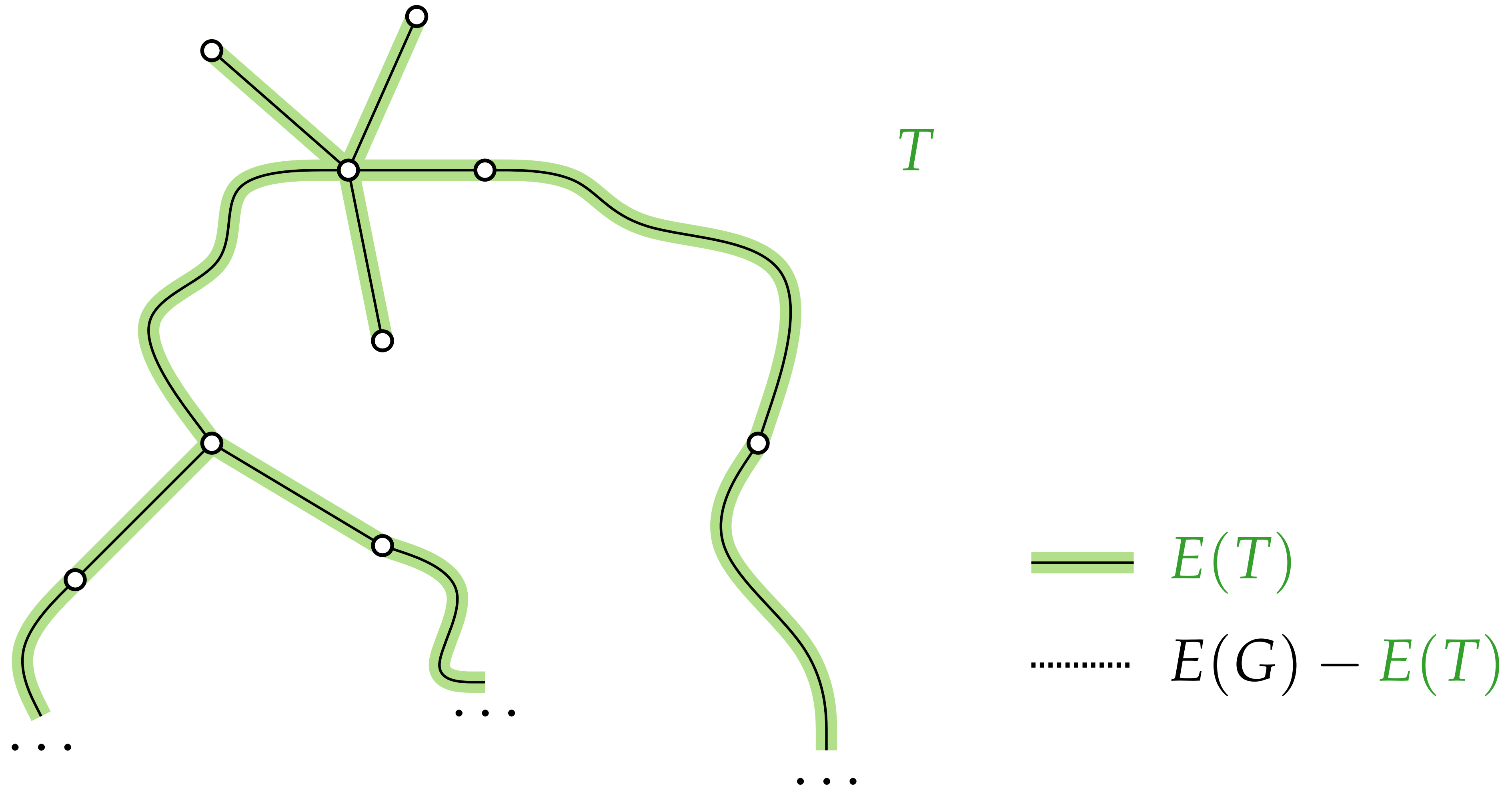
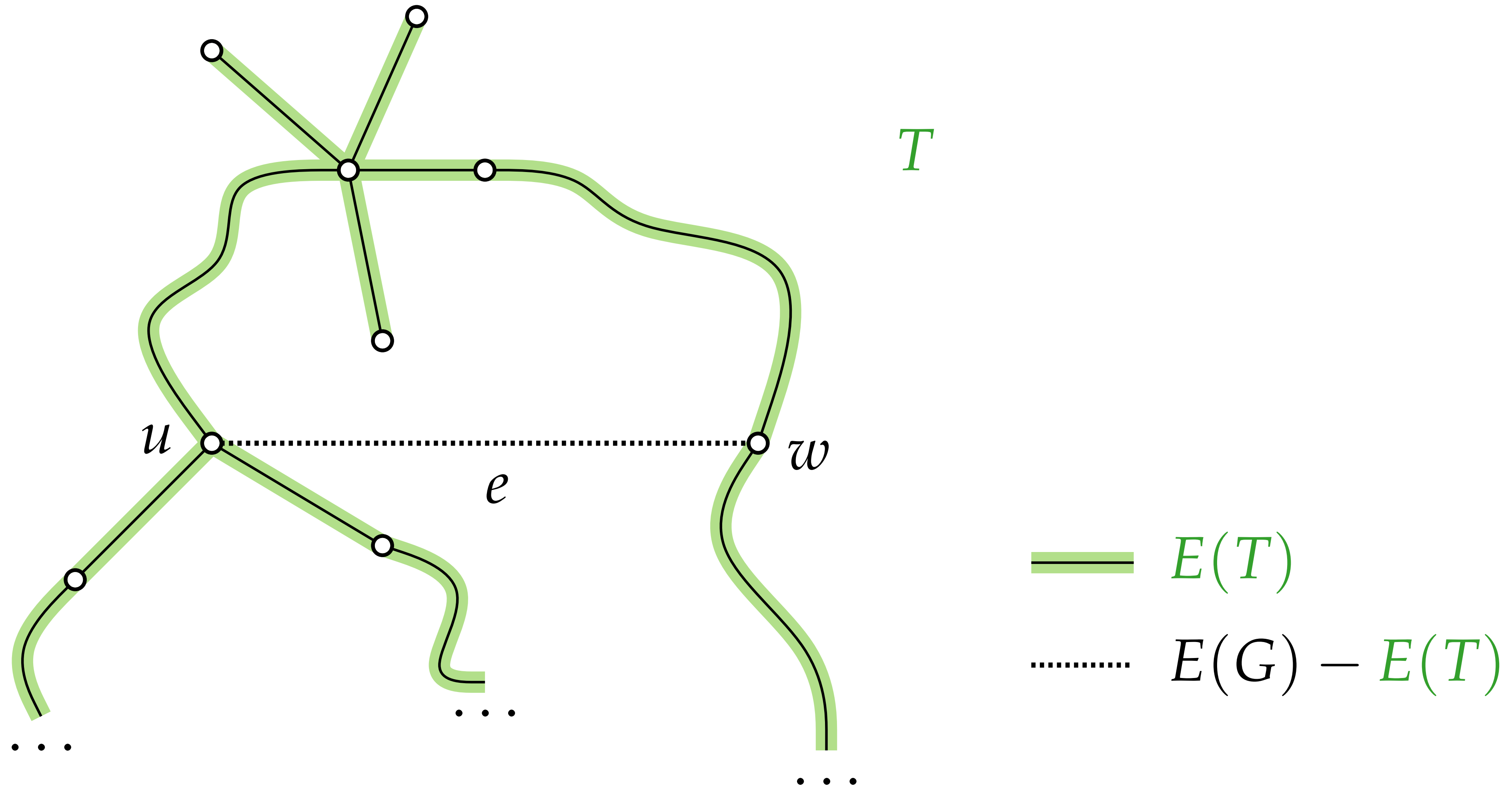
via Local Search

Part II:
Edge Flips and Local Search

Edge Flips



Edge Flips



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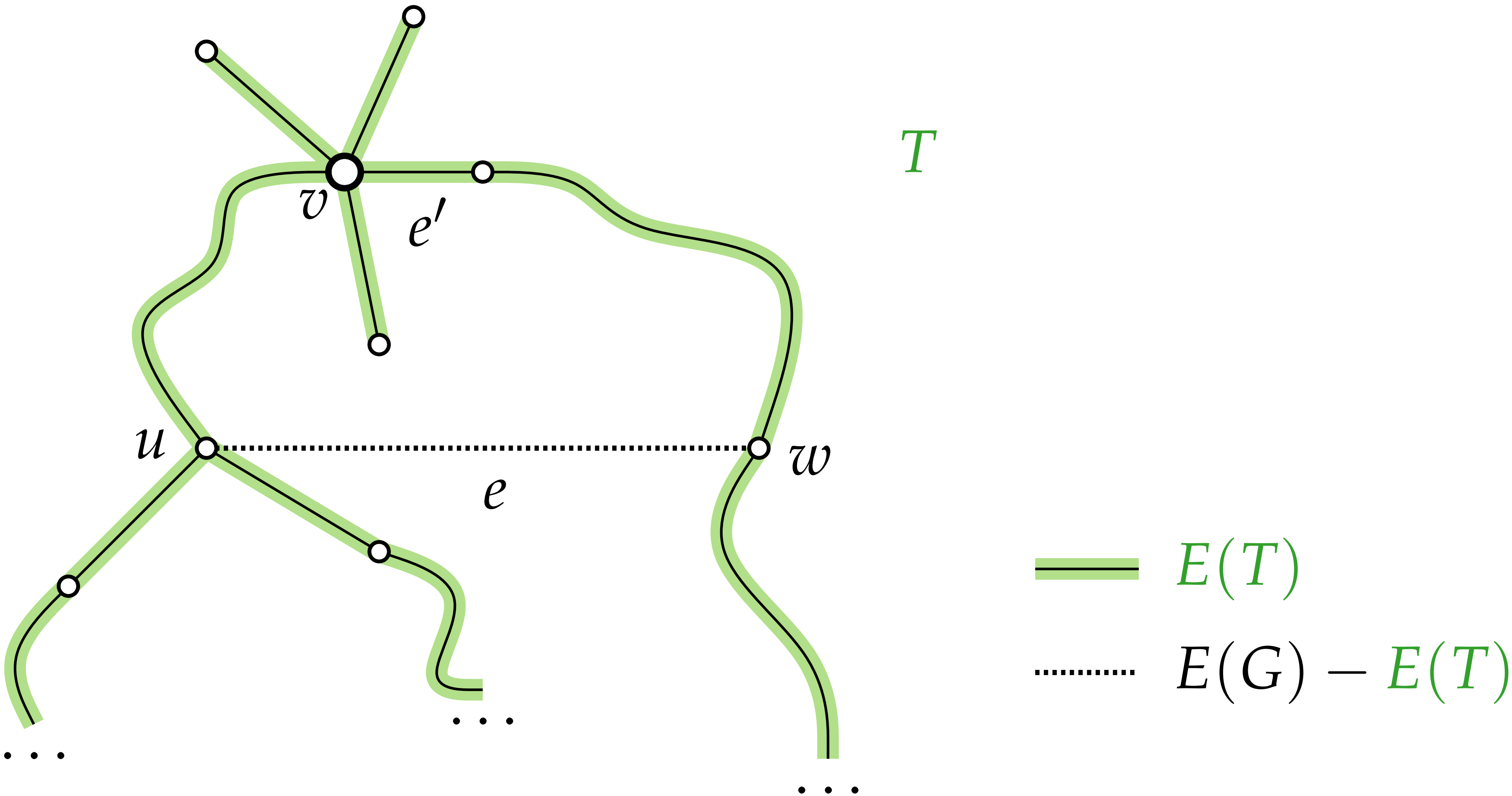

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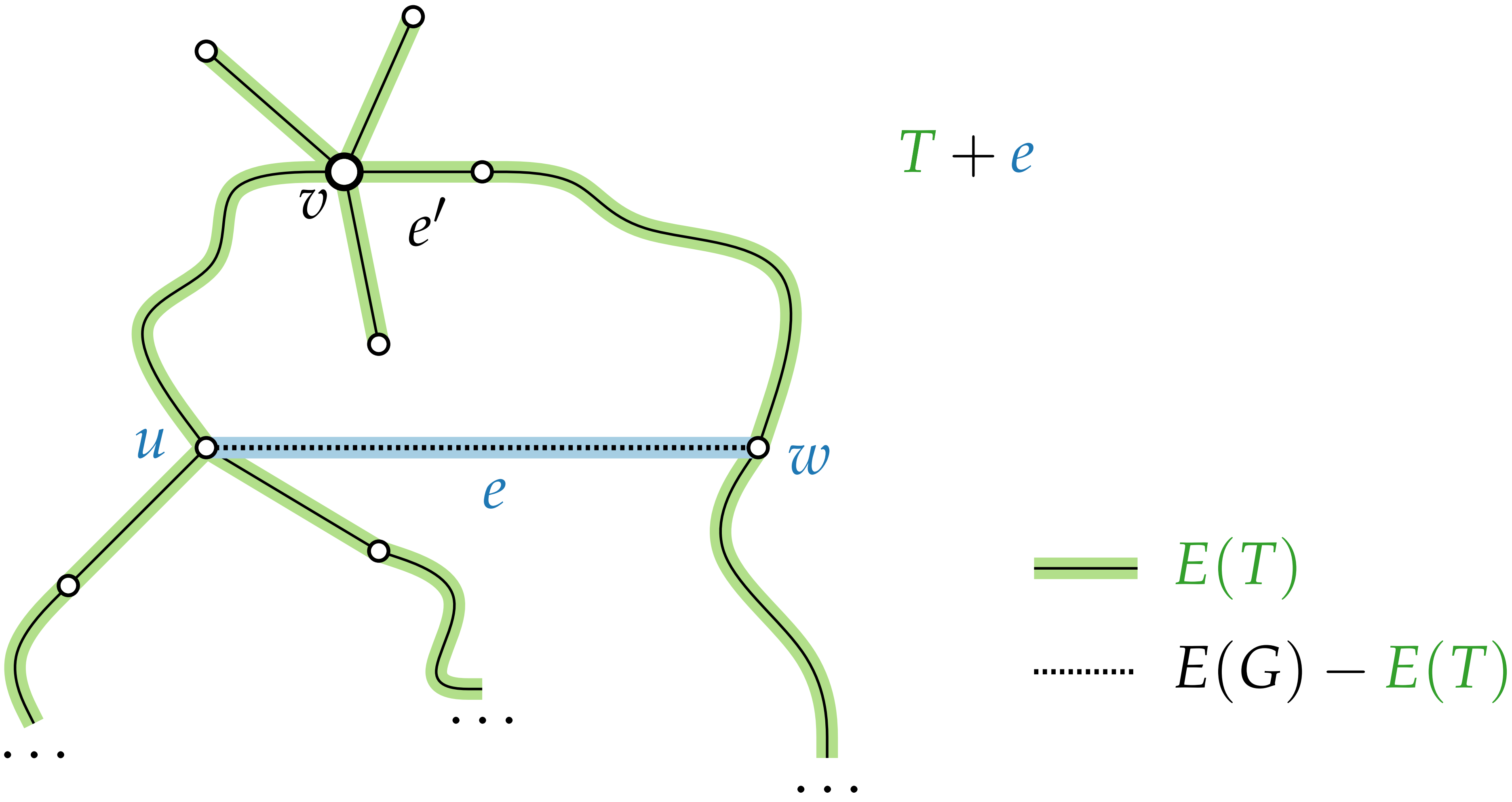

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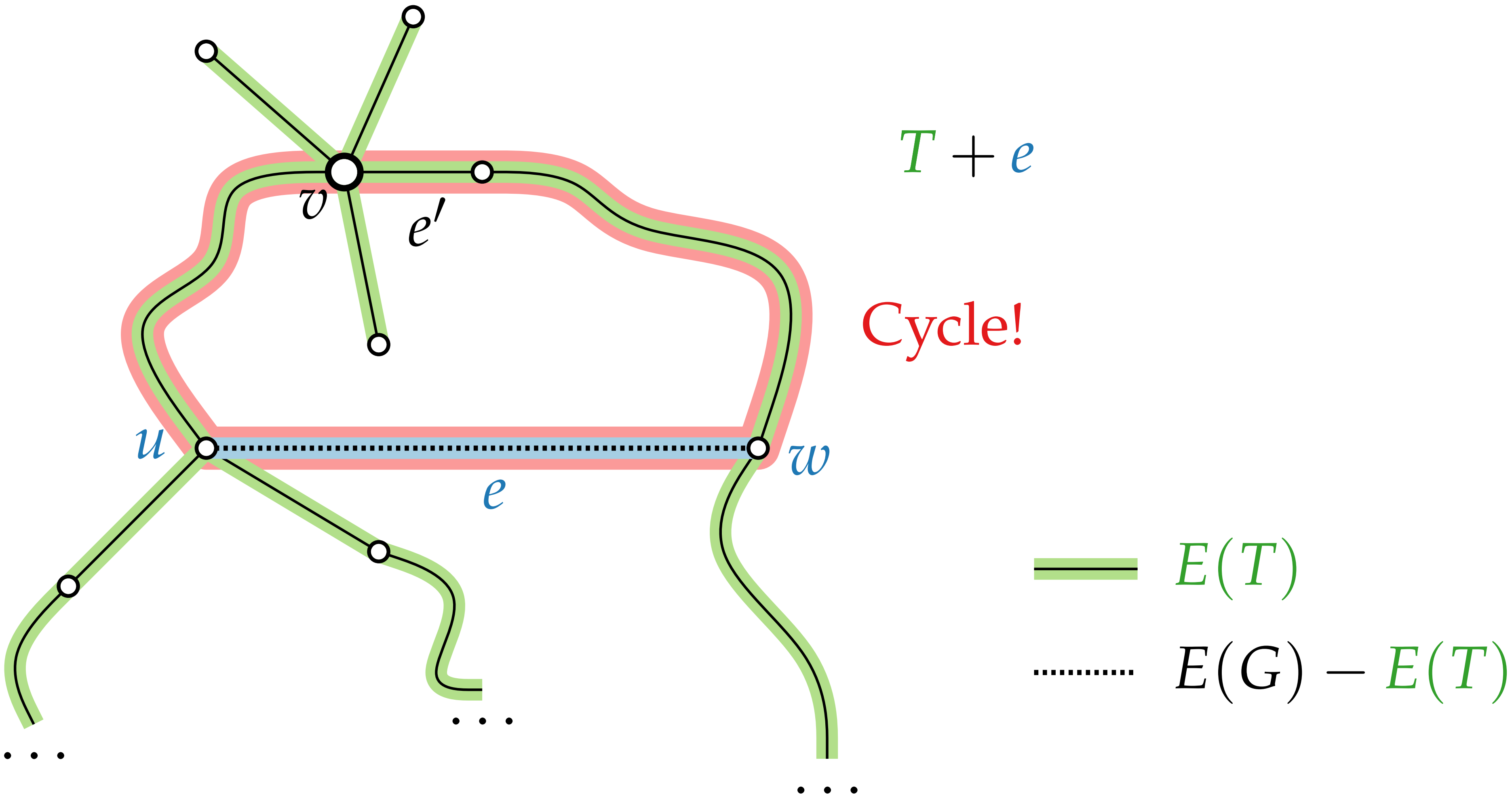

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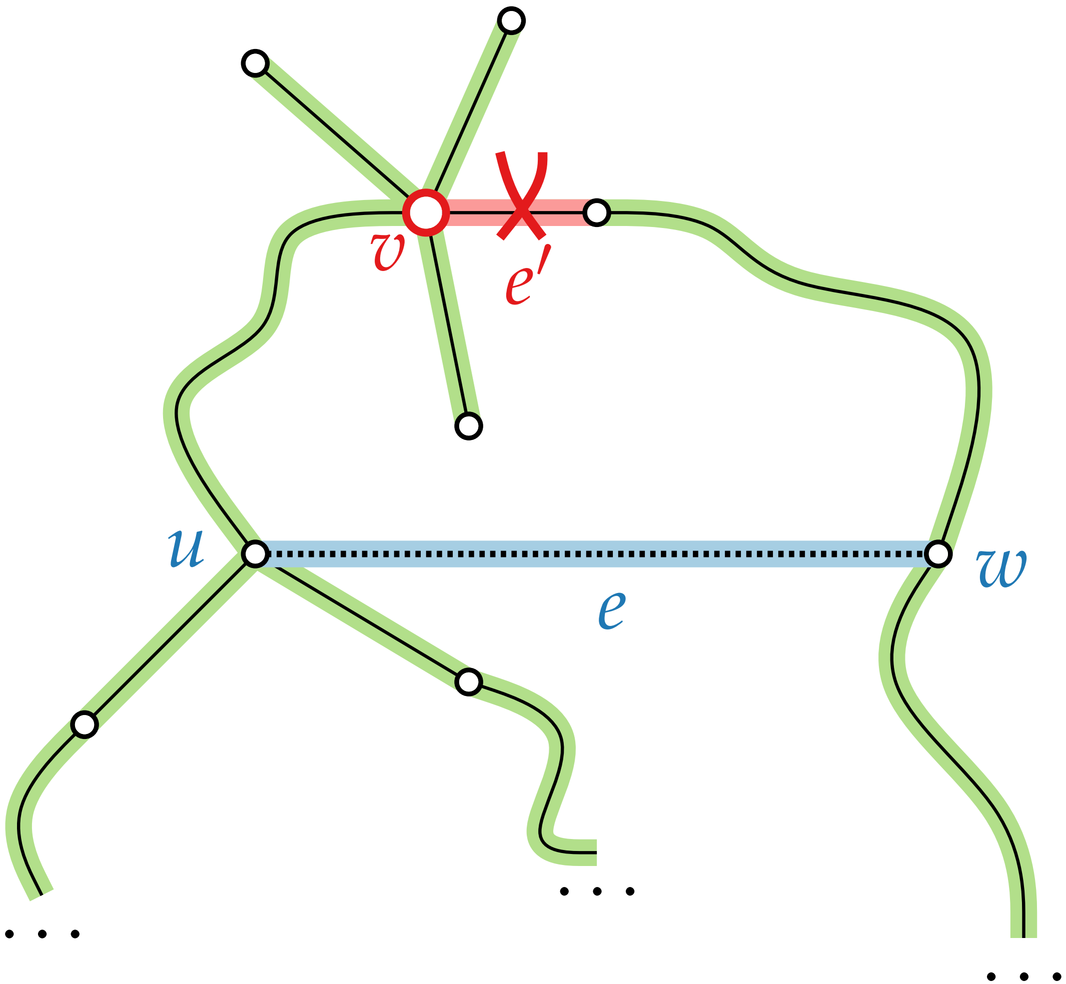


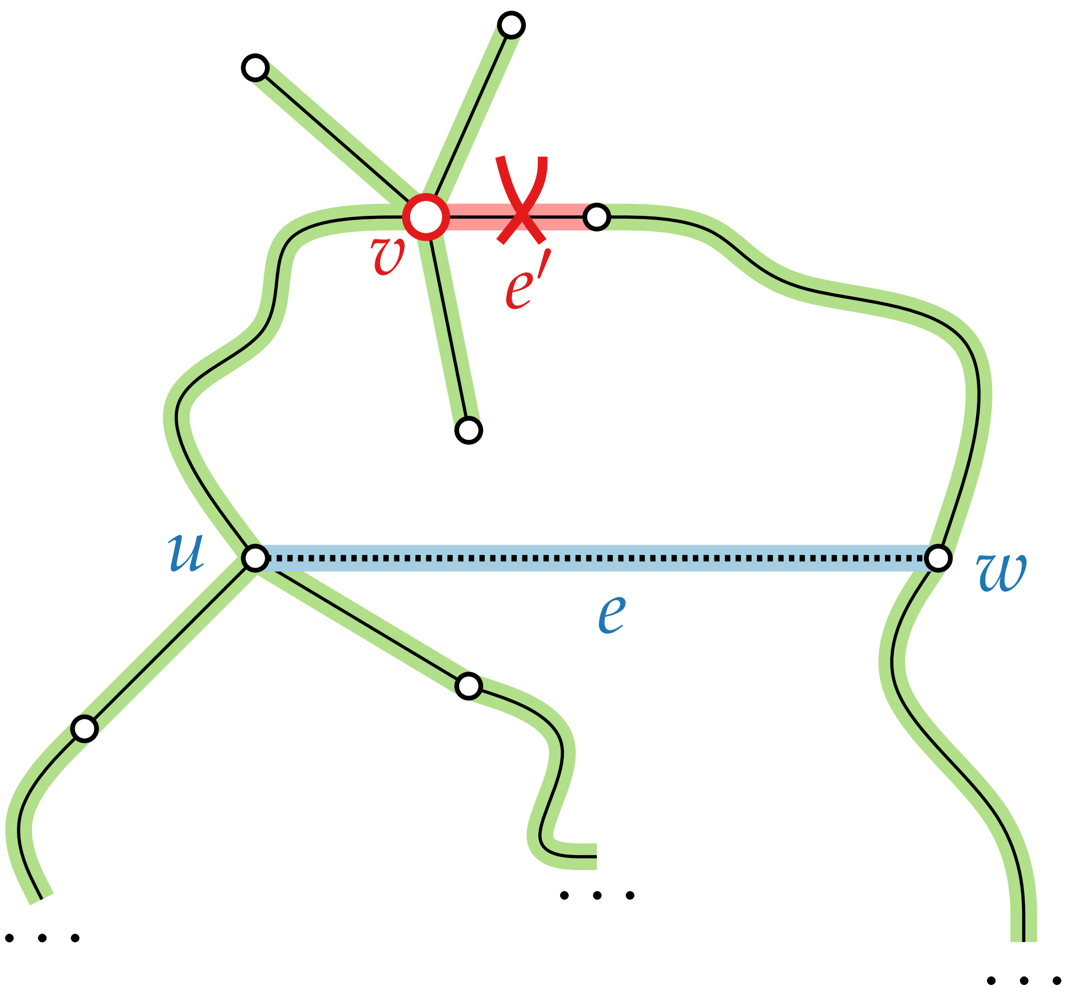
$T + e - e'$
is a new spanning tree

$E(T)$

$E(G) - E(T)$

Edge Flips

Def. An **improving flip** in T for a vertex v and an edge $uw \in E(G) \setminus E(T)$ is a flip with $\deg_T(v) >$



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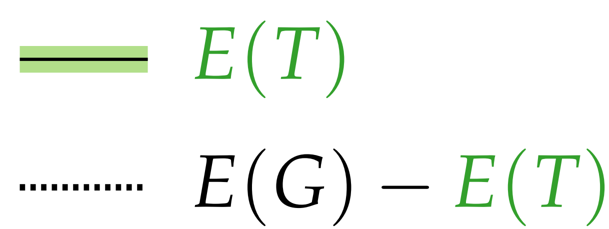
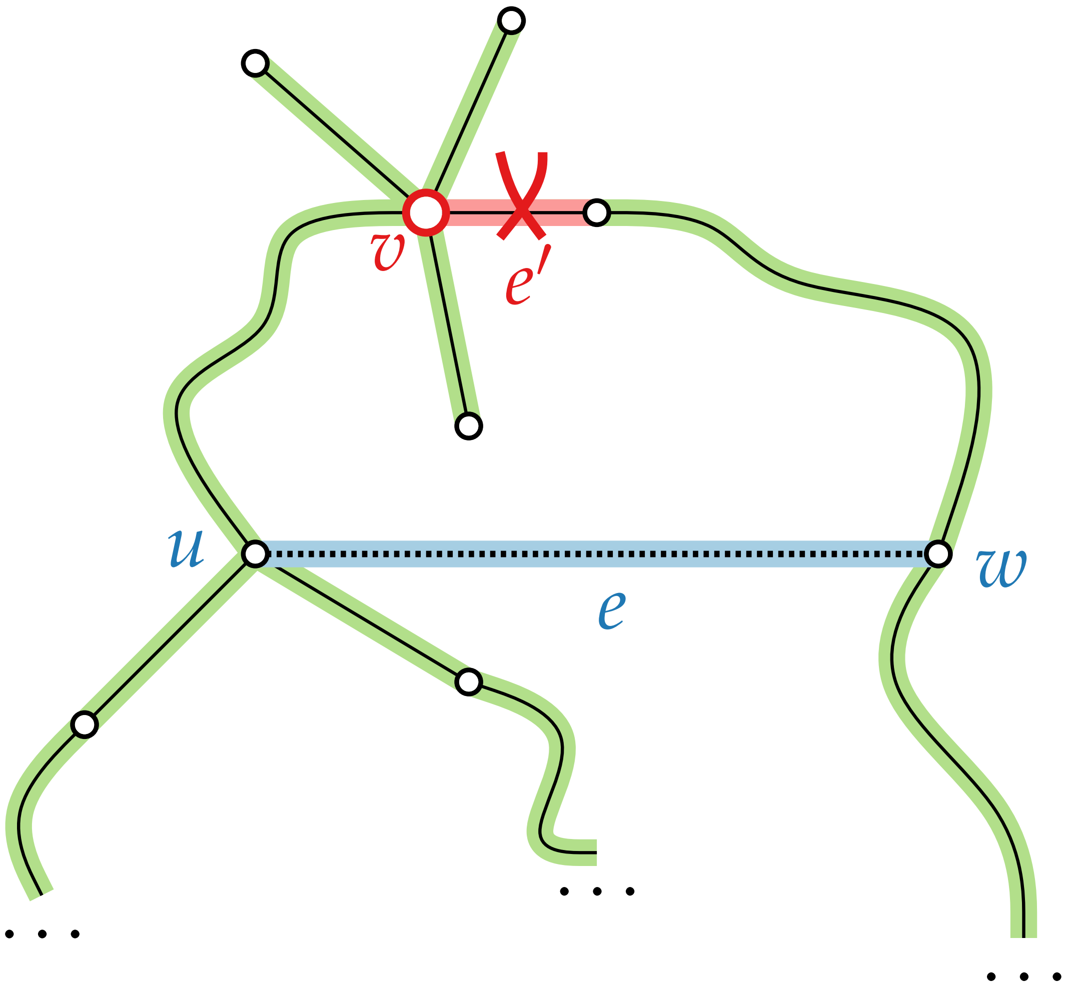


Edge Flips

Def. An **improving flip** in T for a vertex v and an edge $uw \in E(G) \setminus E(T)$ is a flip with $\deg_T(v) > \max\{\deg_T(u), \deg_T(w)\} + 1$.



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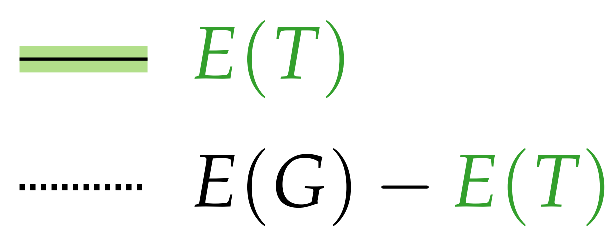

Local Search

MinDegSpanningTreeLocalSearch(graph G)
  $T \leftarrow$ any spanning tree of G
  while $\exists$ improving flip in T for a vertex v
      with $\deg_T(v) \geq \Delta(T) - \ell$ **do**
    do the improving flip
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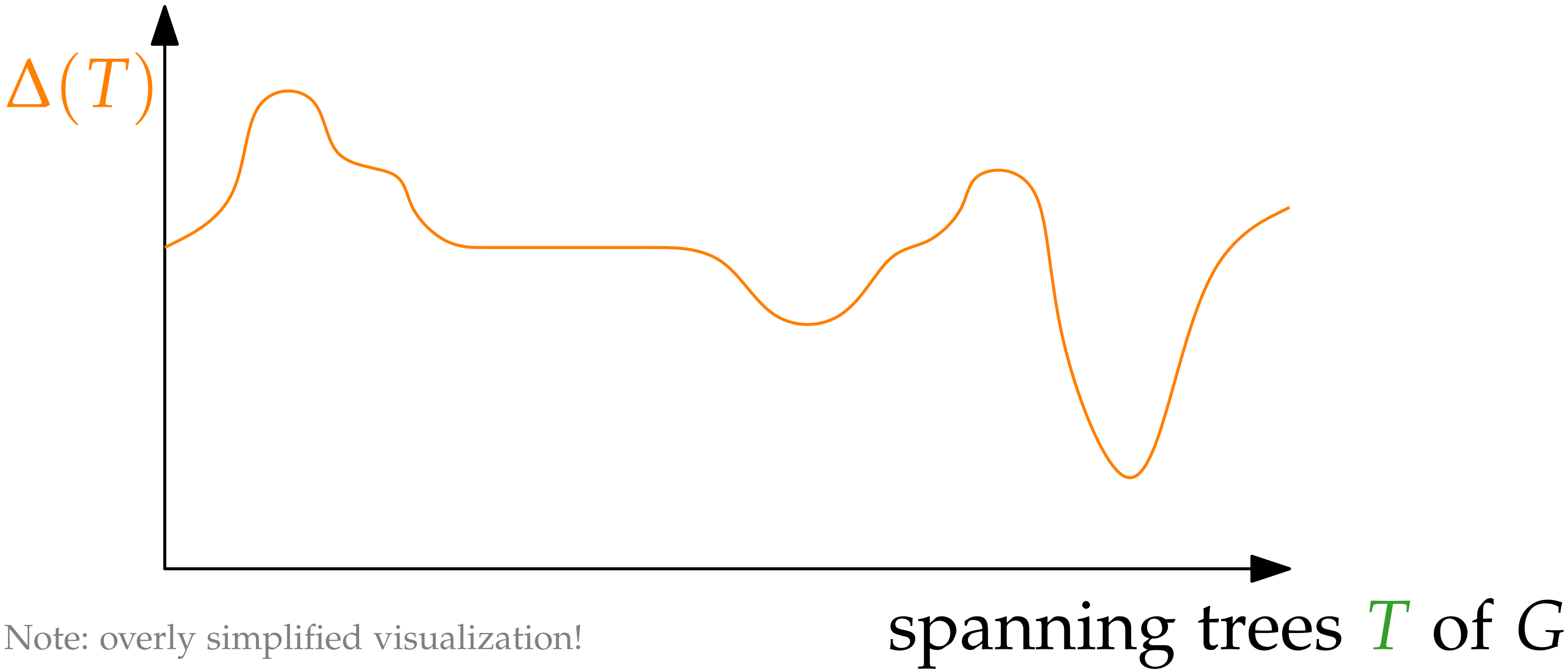

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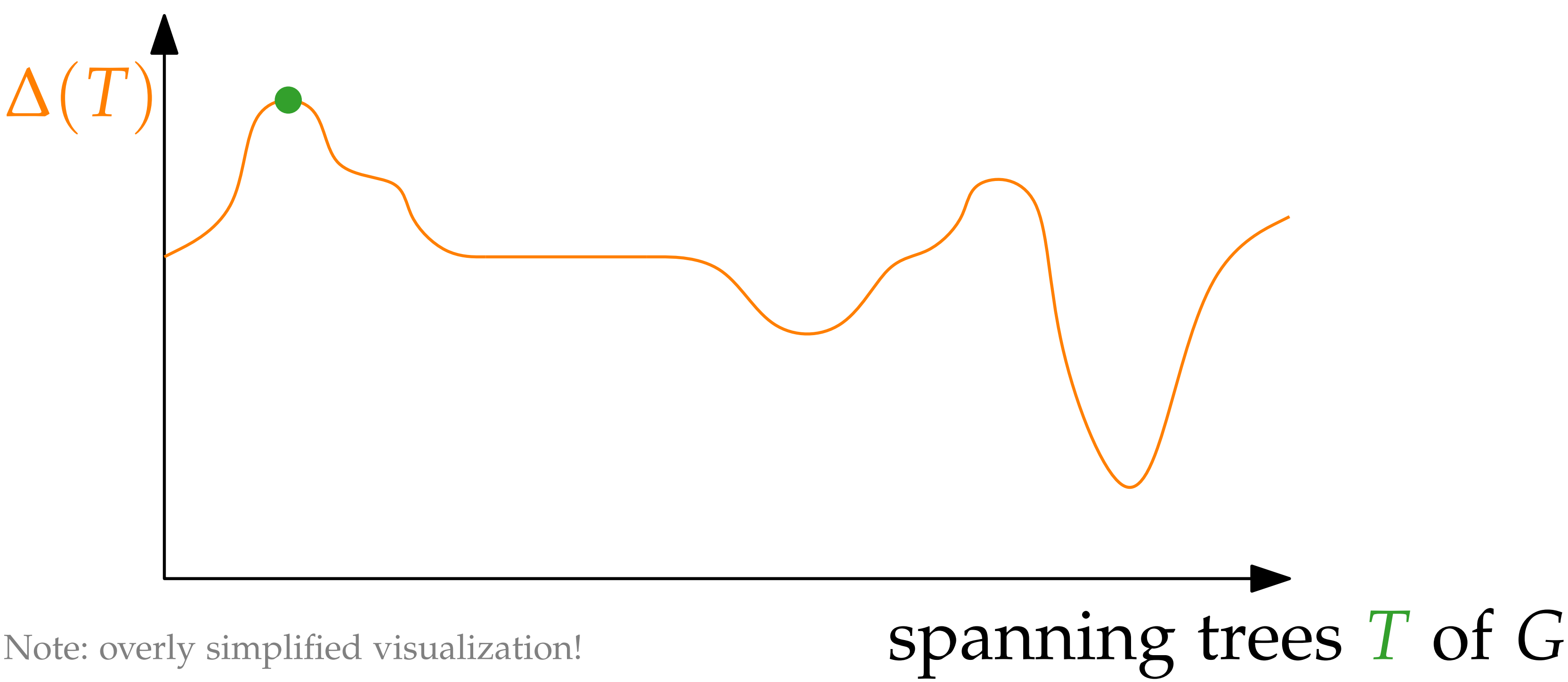

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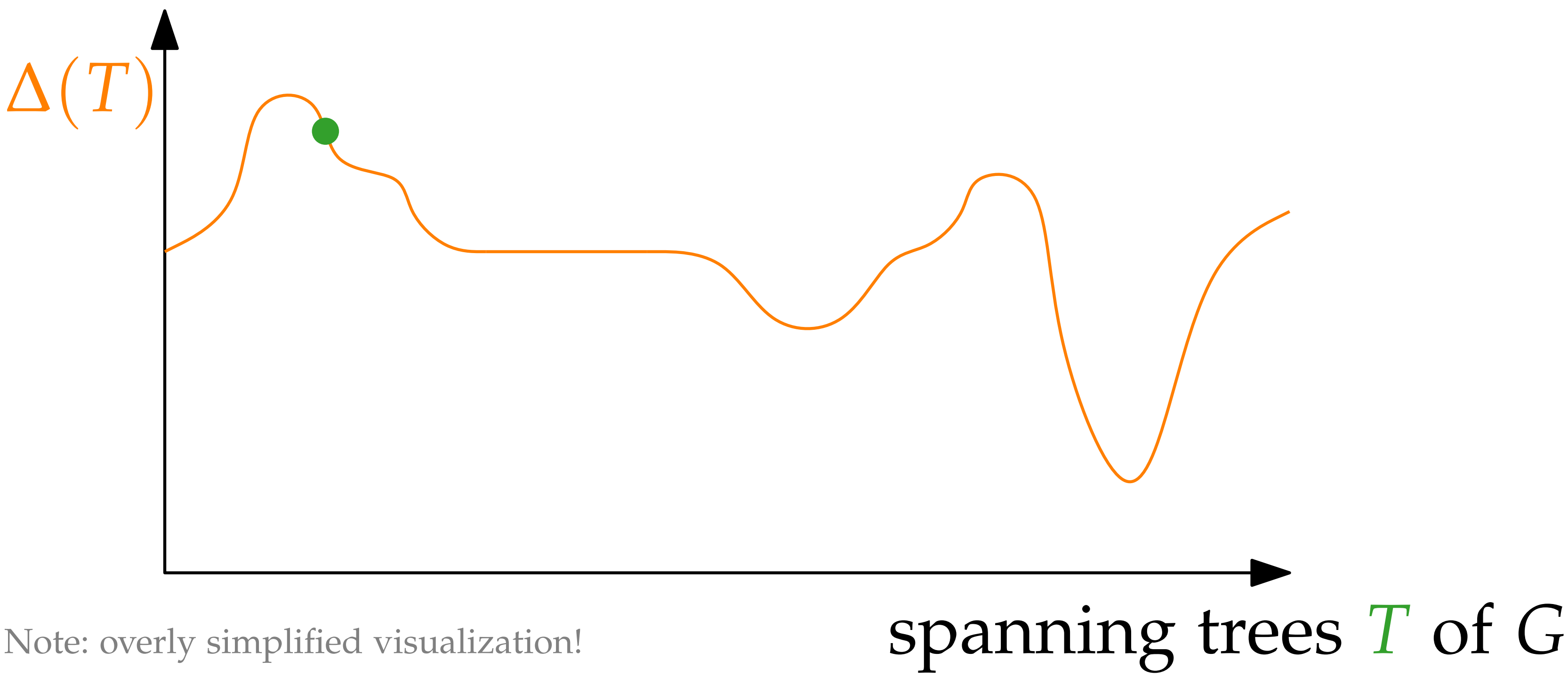

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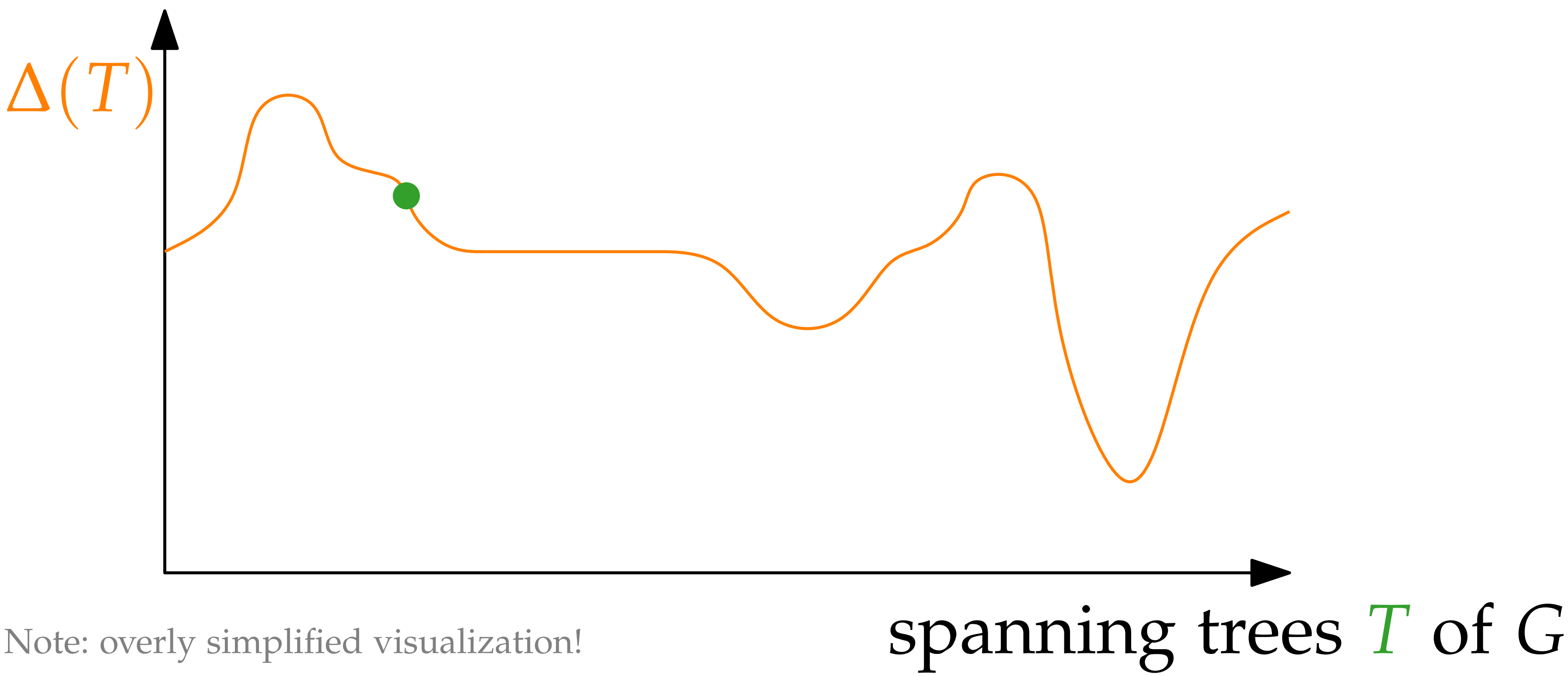

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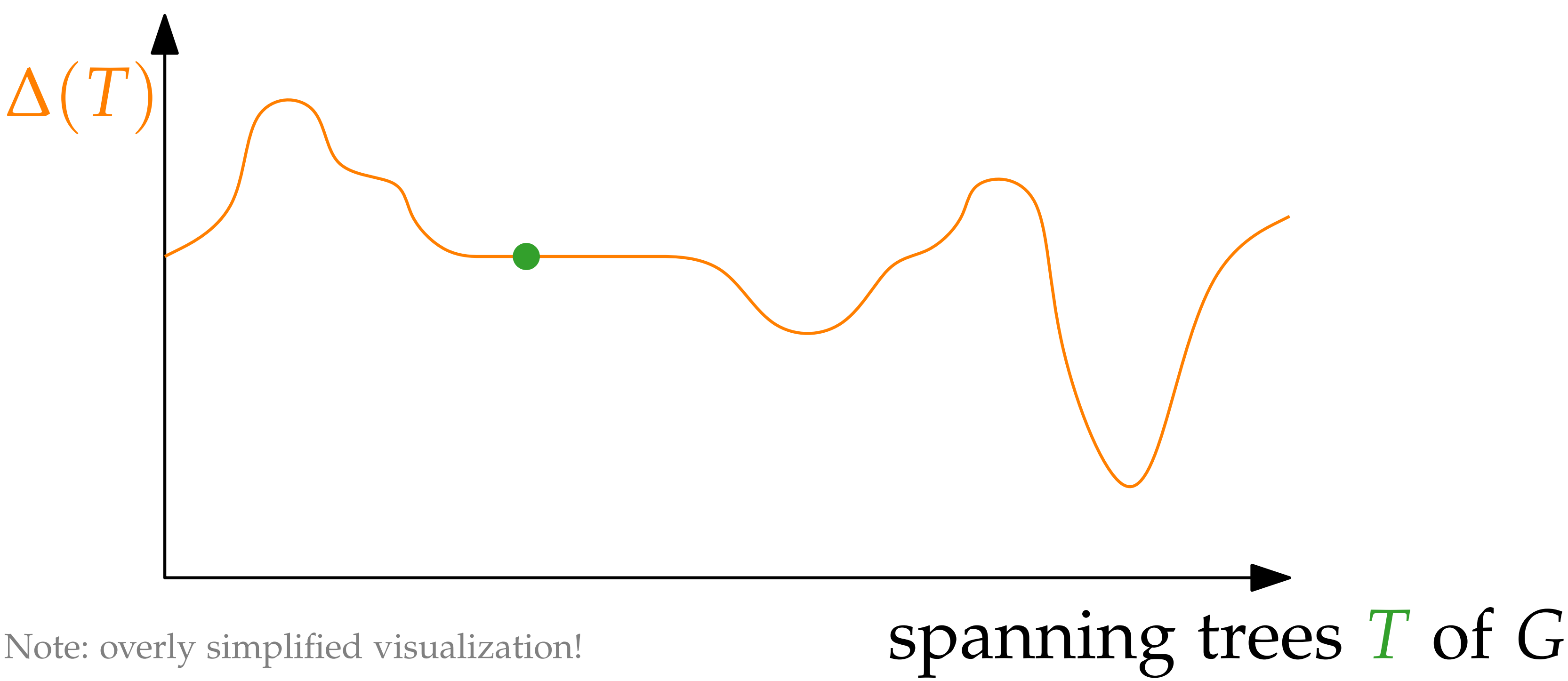

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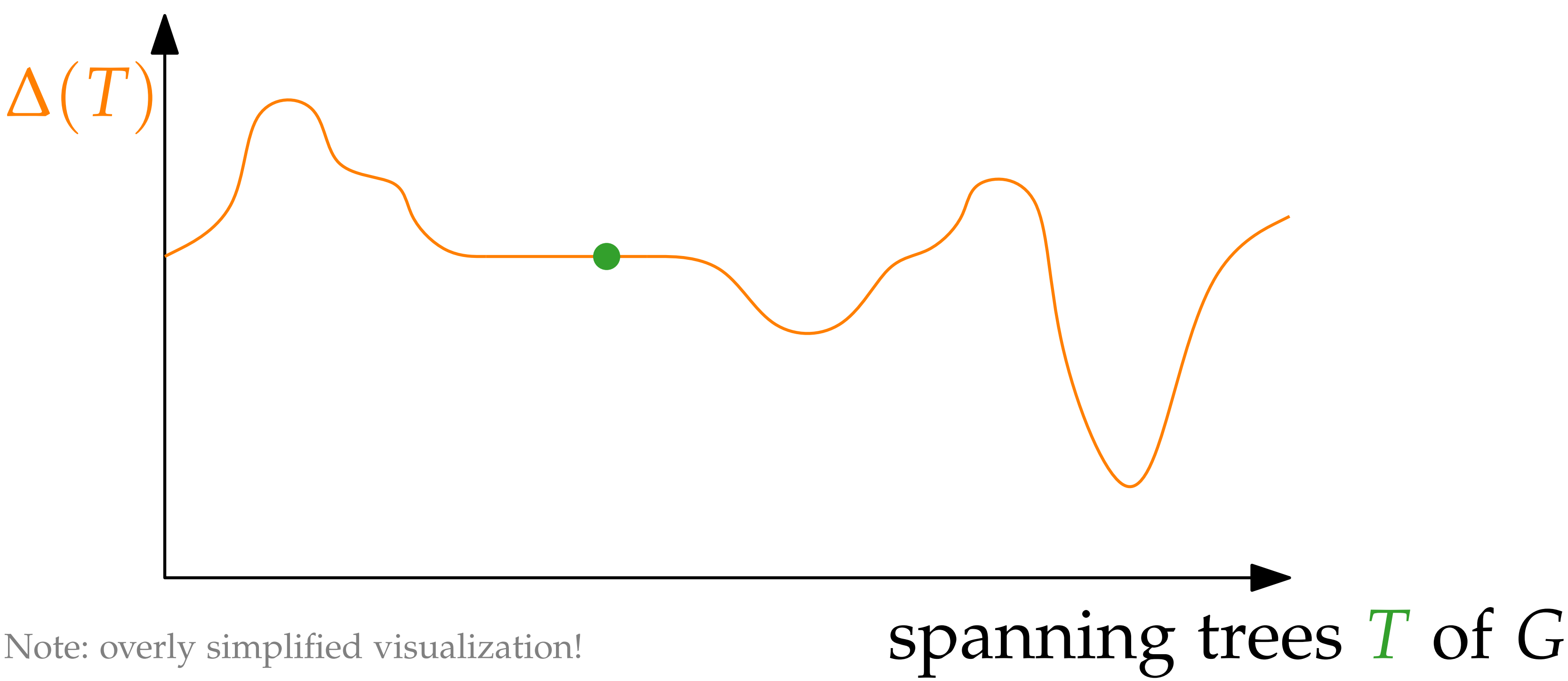

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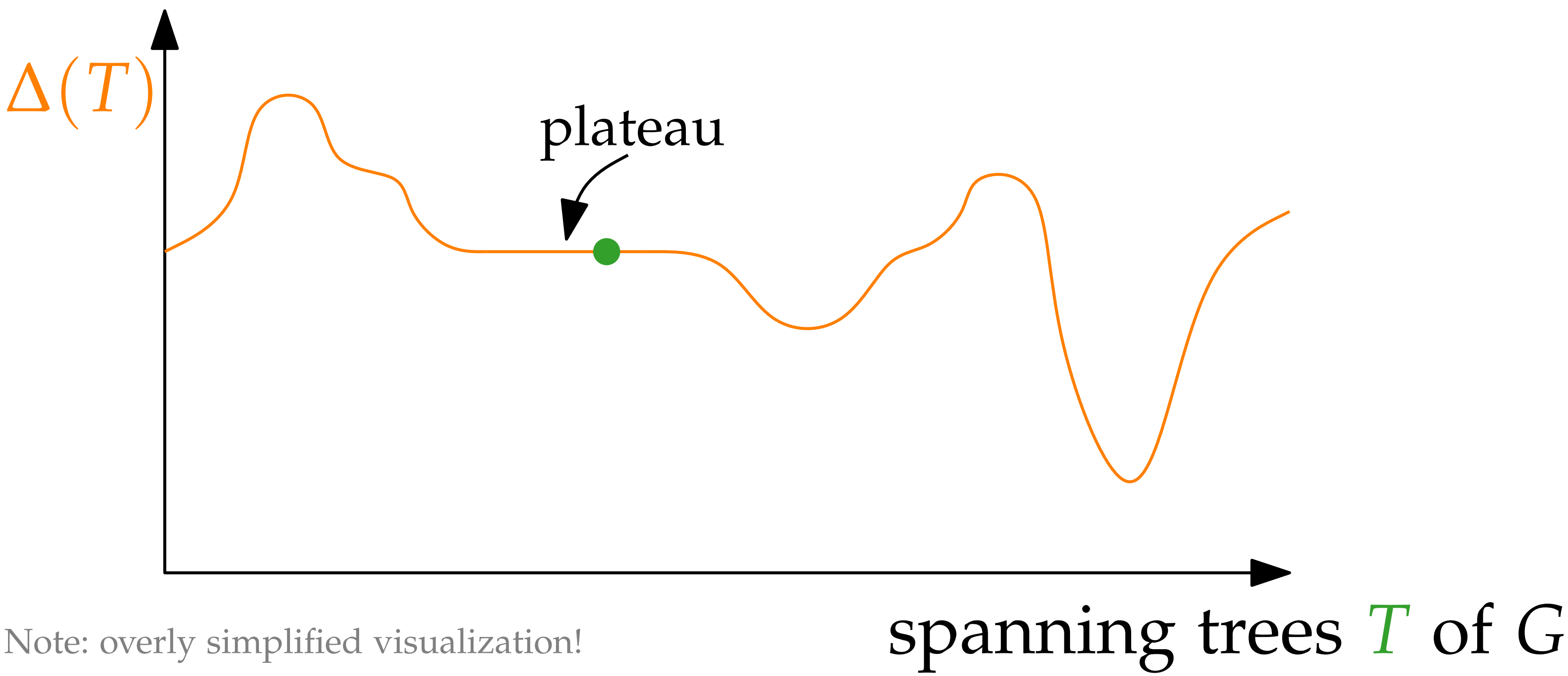

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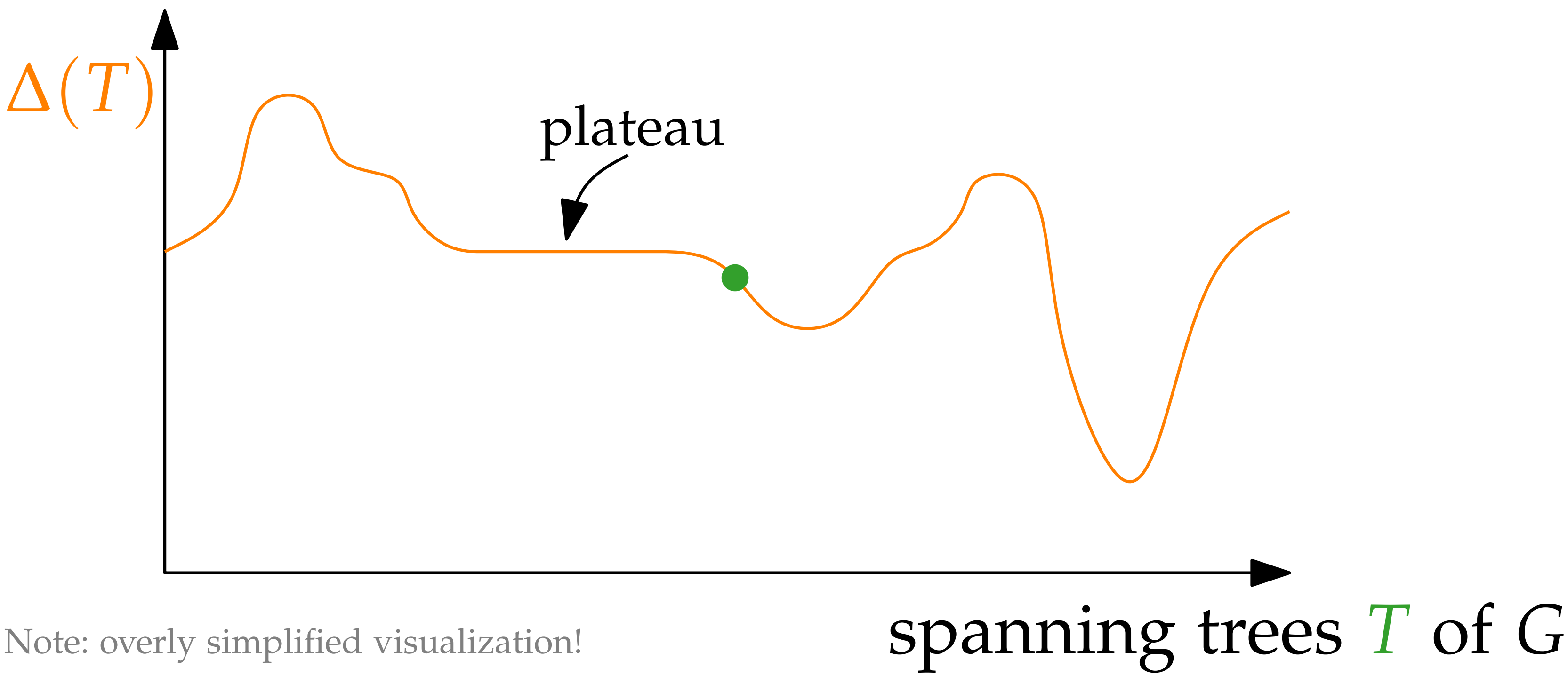

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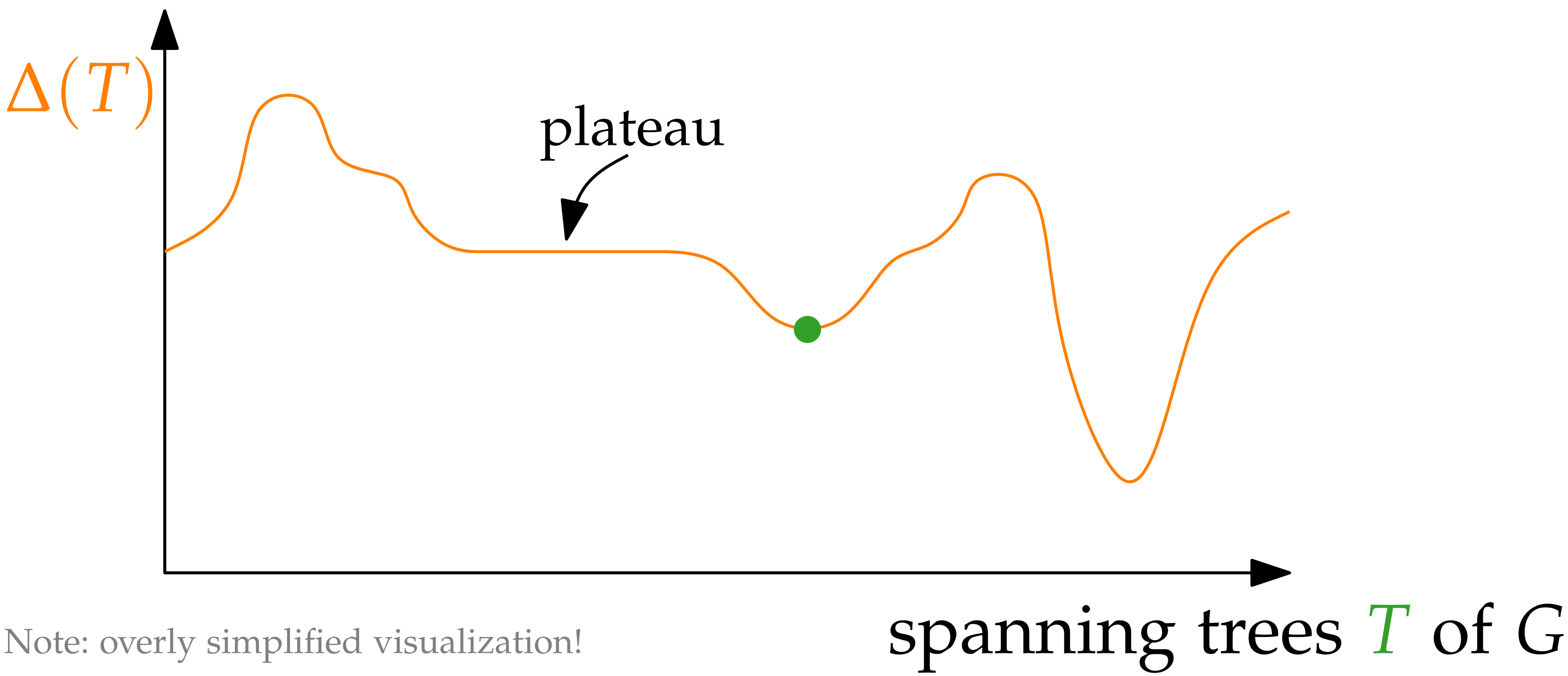

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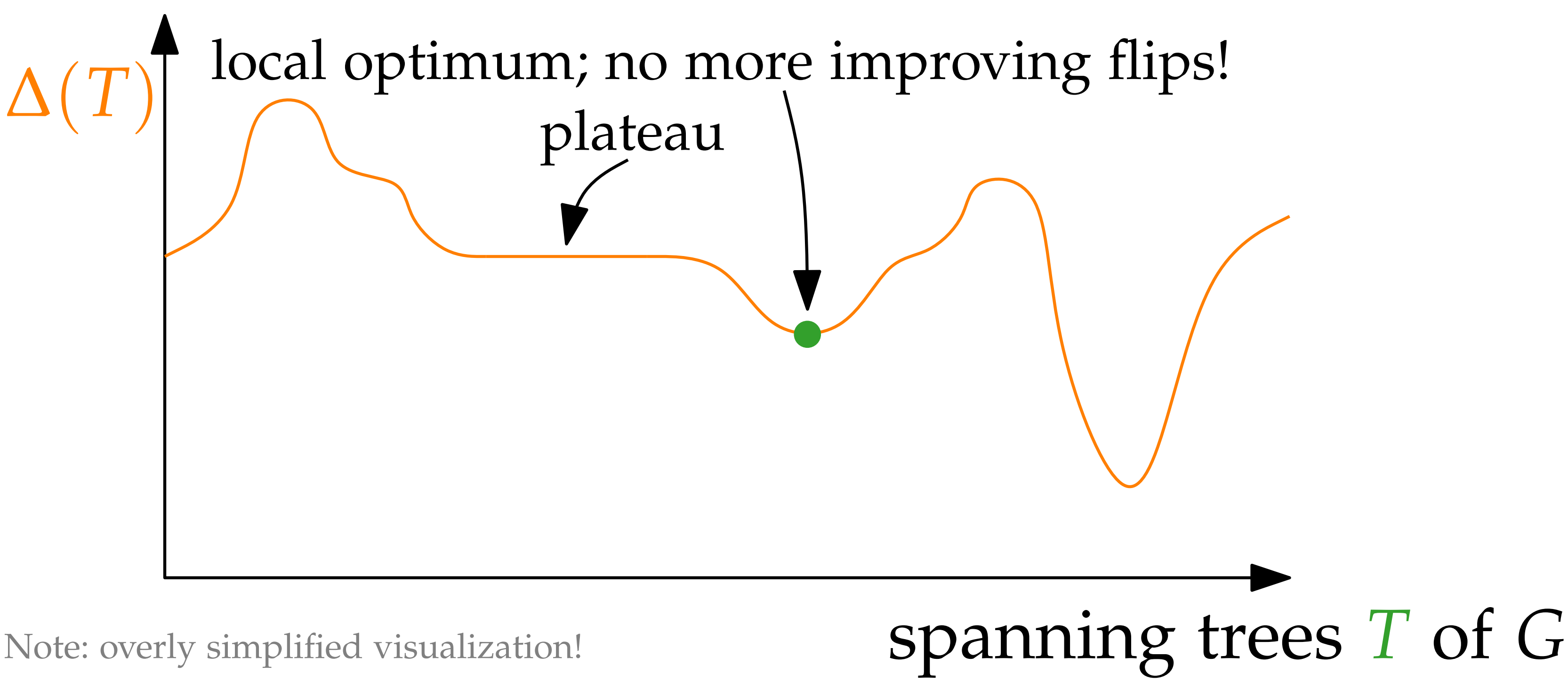

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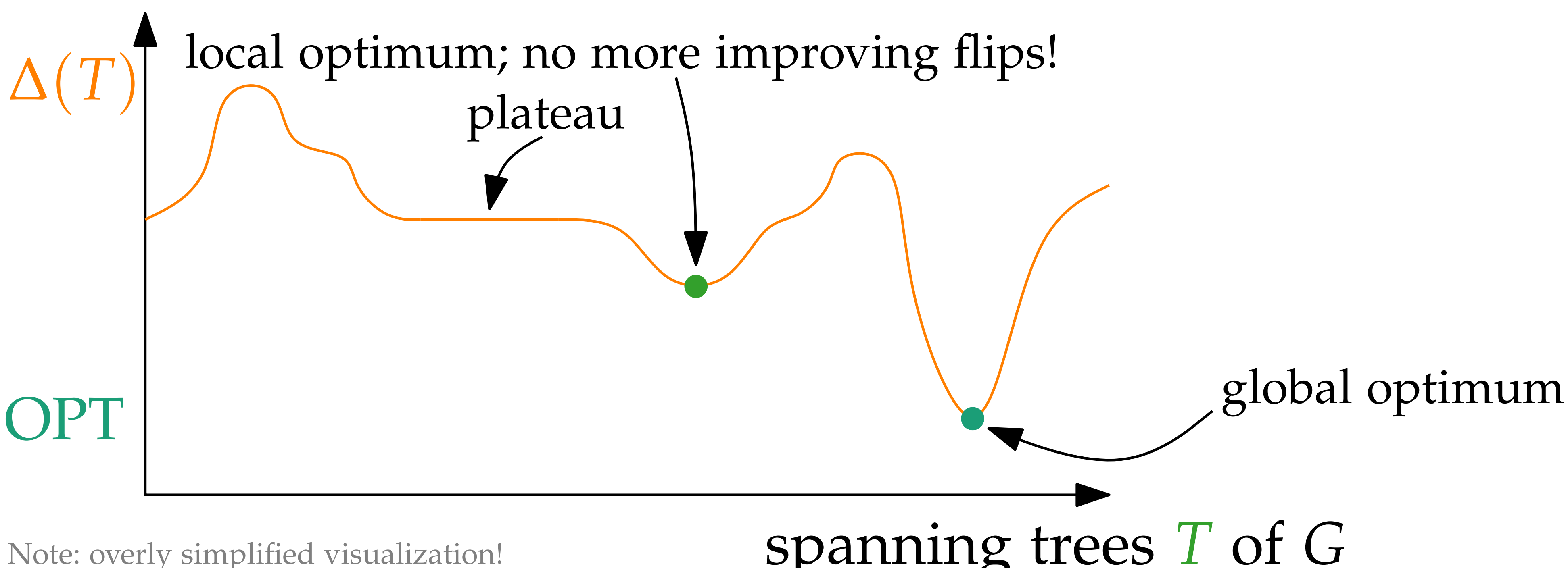

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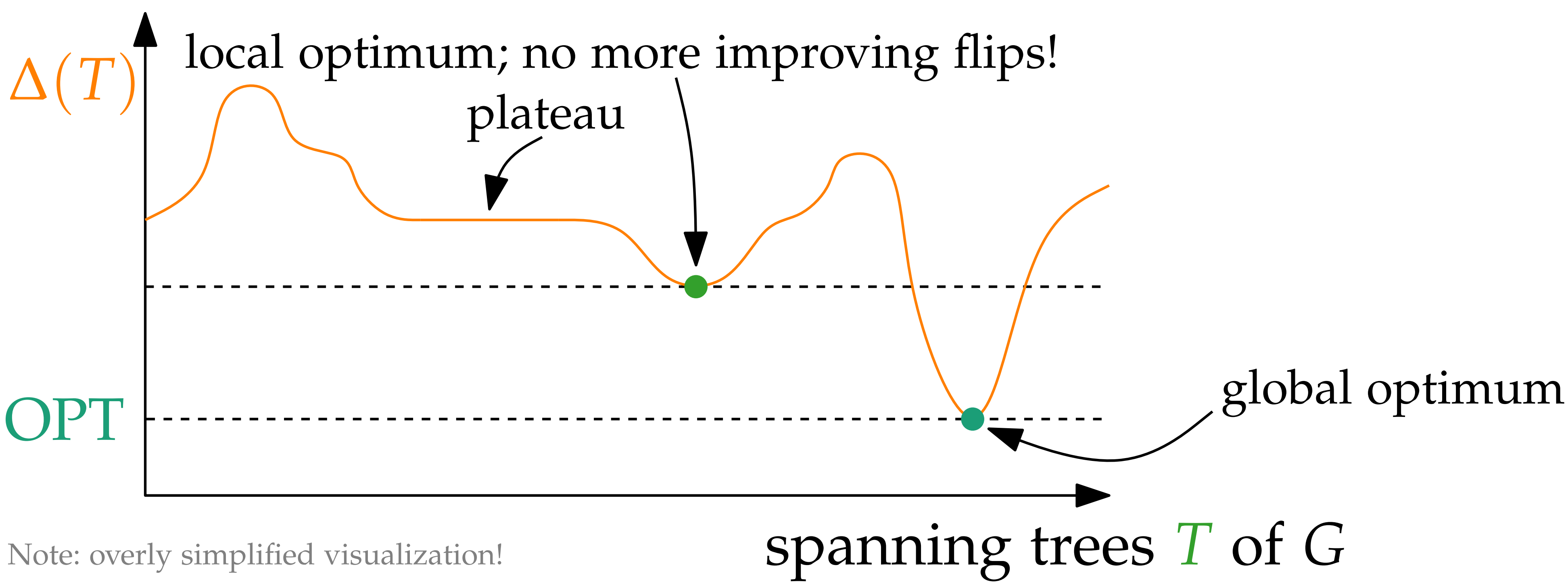

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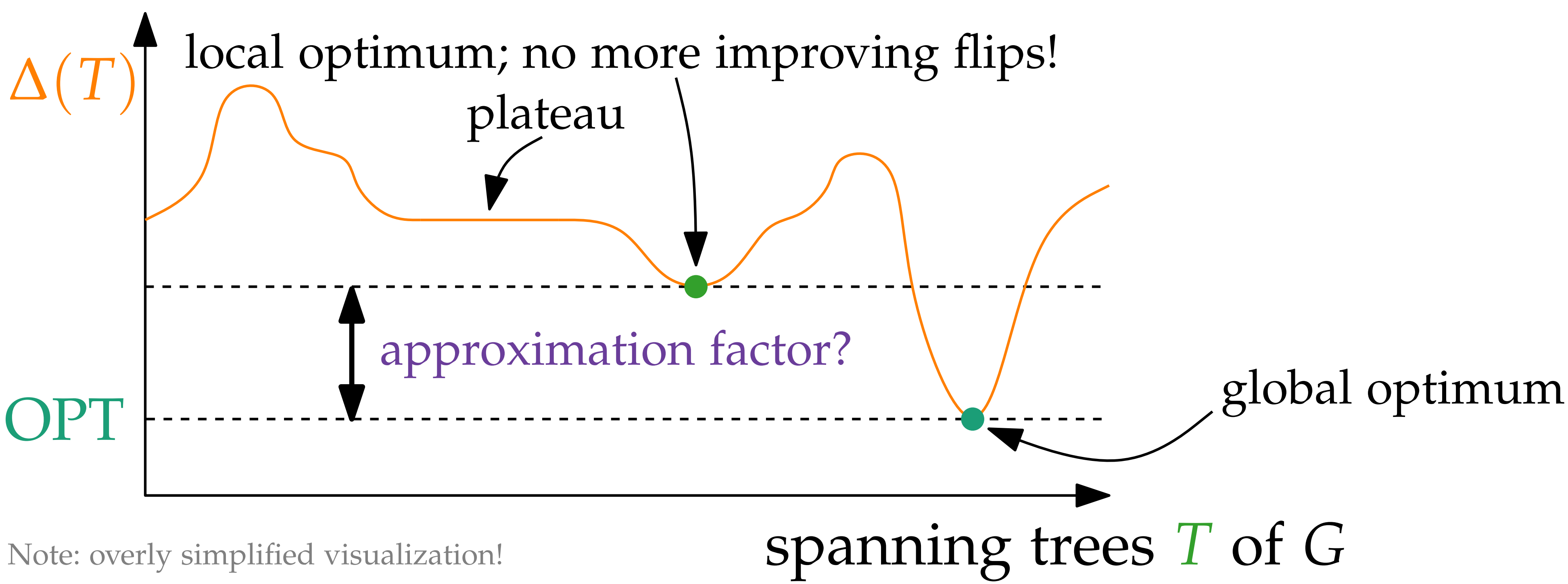

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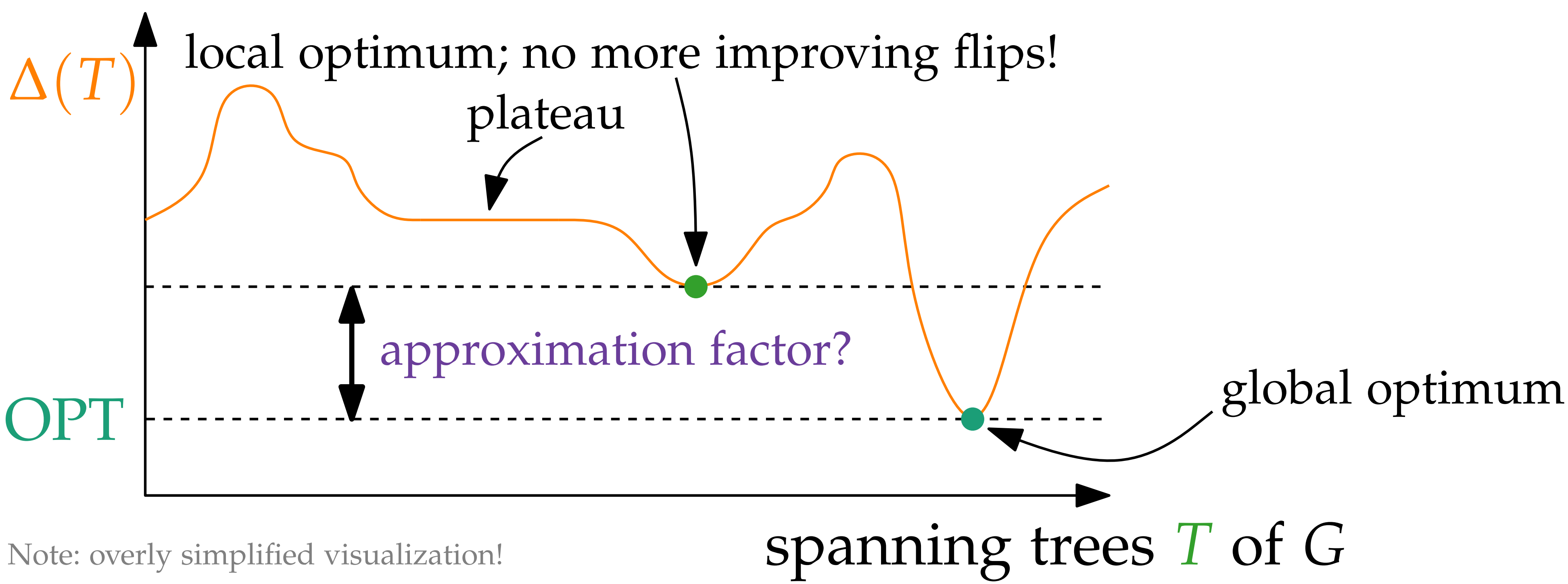
- Termination?



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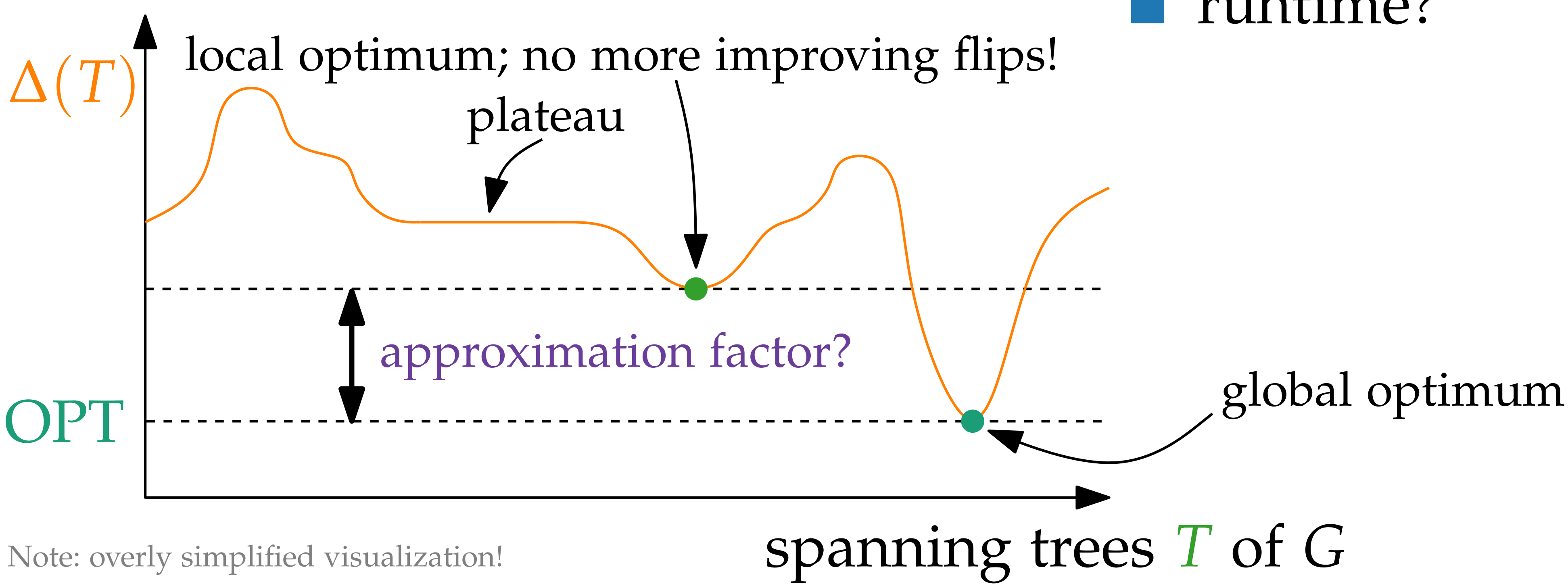
- Termination?
- runtime?



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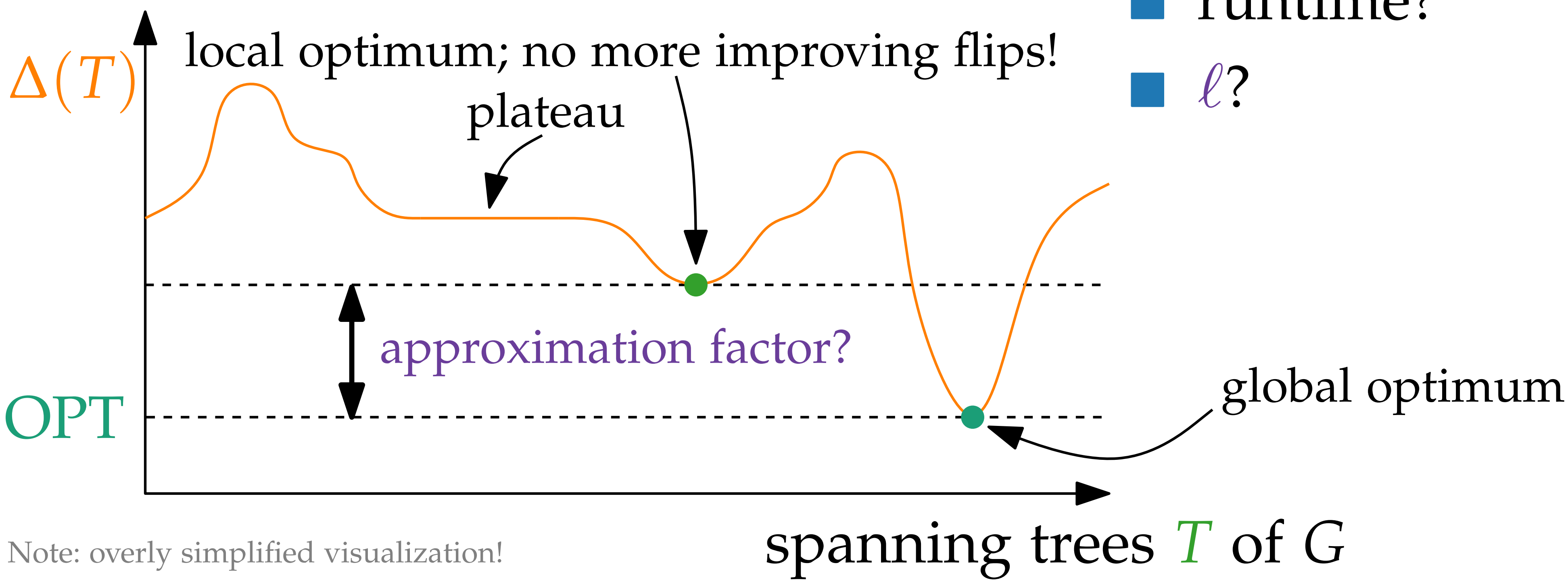
- Termination?
- runtime?
- ℓ?



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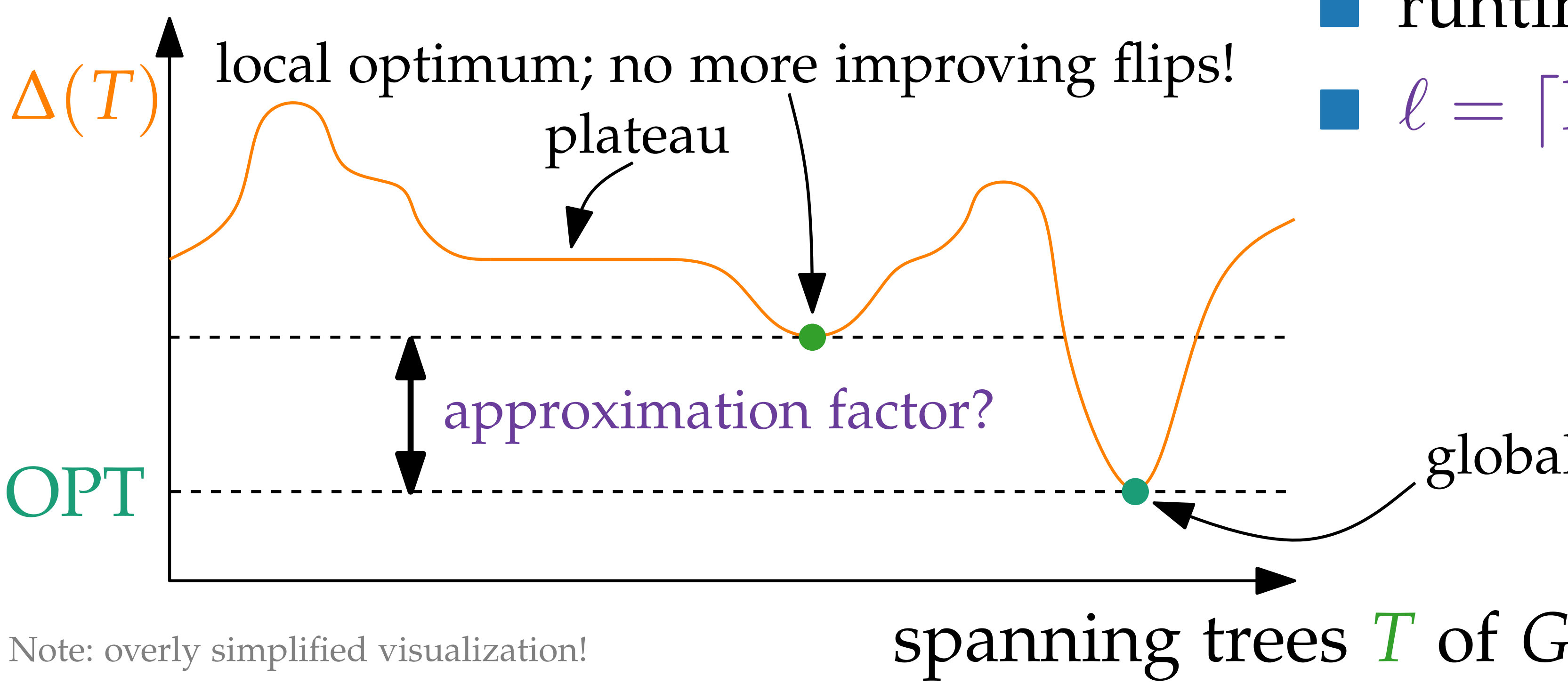
- Termination?
- runtime?
- $\ell = \lceil \log_2 n \rceil$



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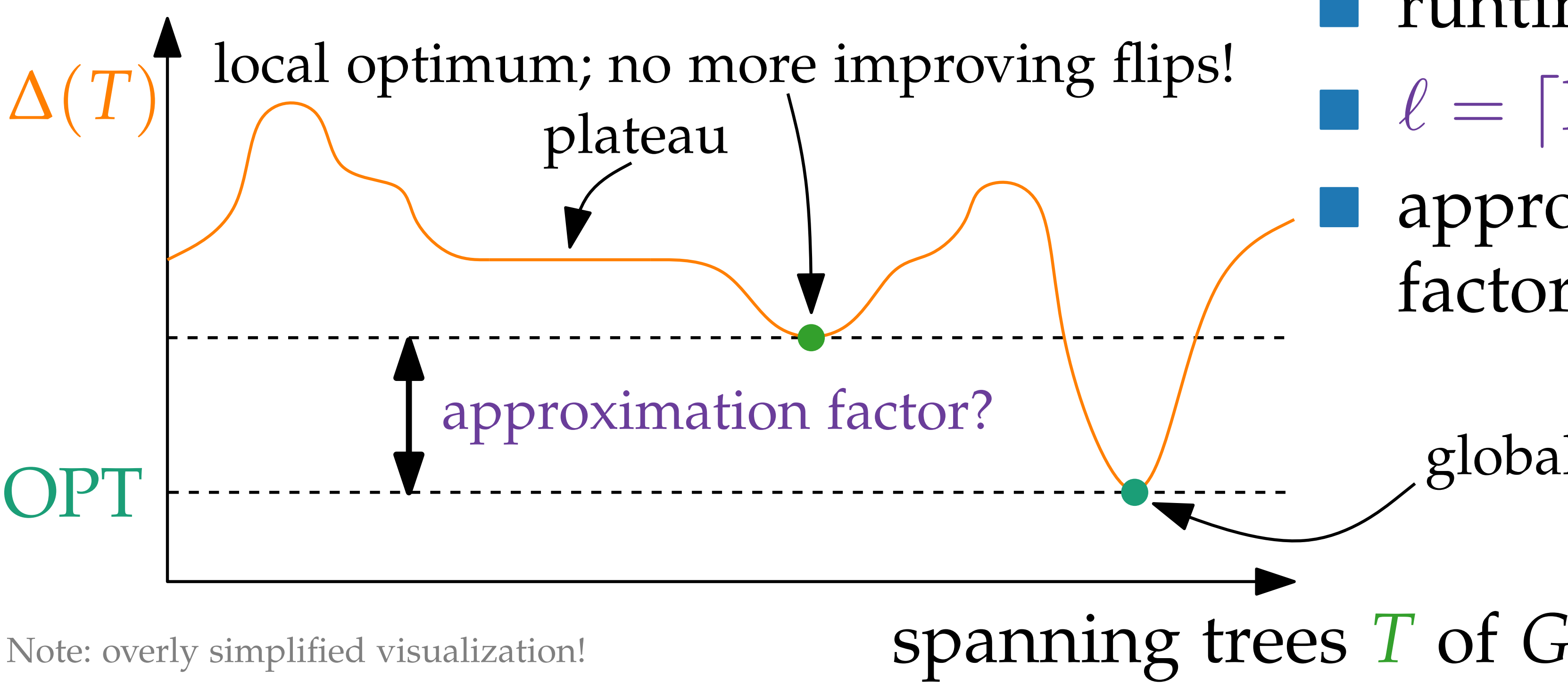
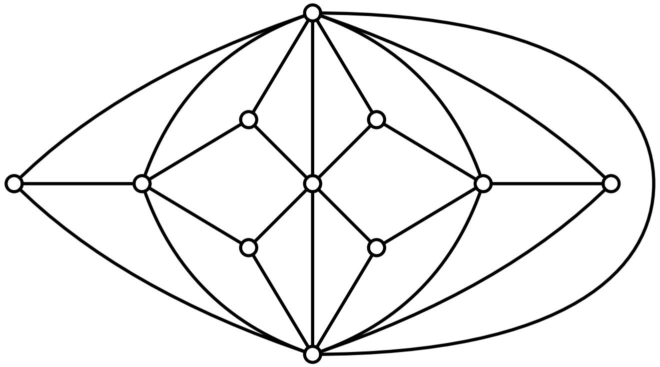
- Termination?
- runtime?
- $\ell = \lceil \log_2 n \rceil$
- approximation factor?

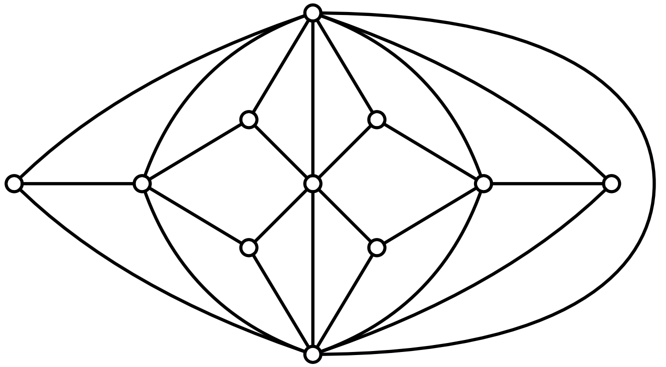


Note: overly simplified visualization!

Example

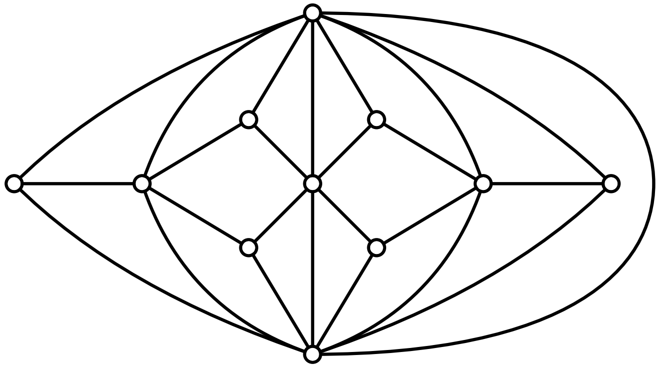


Example



Goldner–Harary graph (minus two edges)

Example

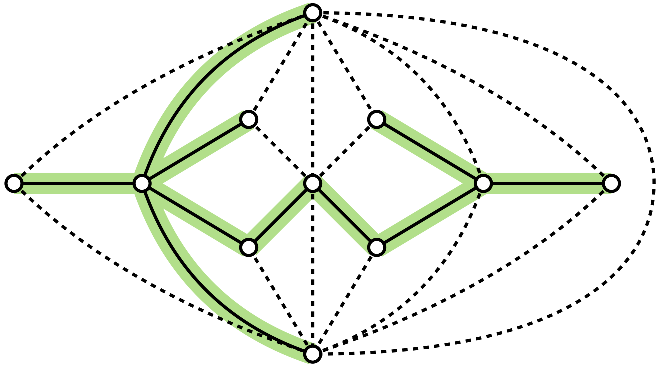

Goldner–Harary graph (minus two edges)

choose any

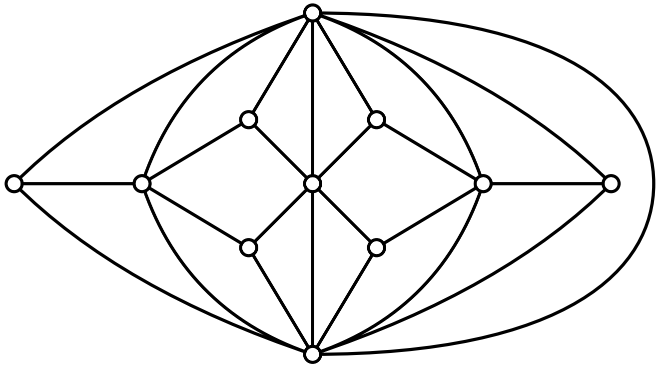
spanning tree

T

$\Delta(T) = 5$



Example



Goldner–Harary graph (minus two edges)

choose any

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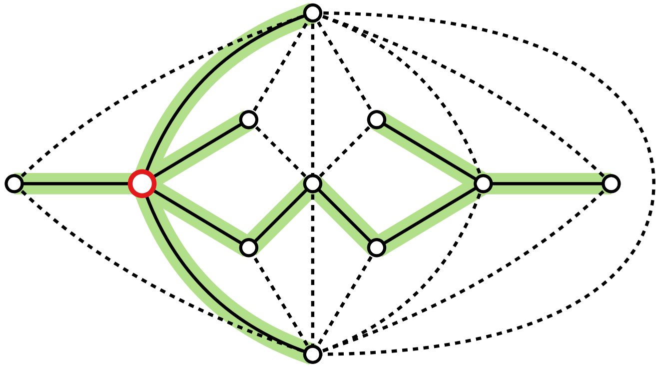
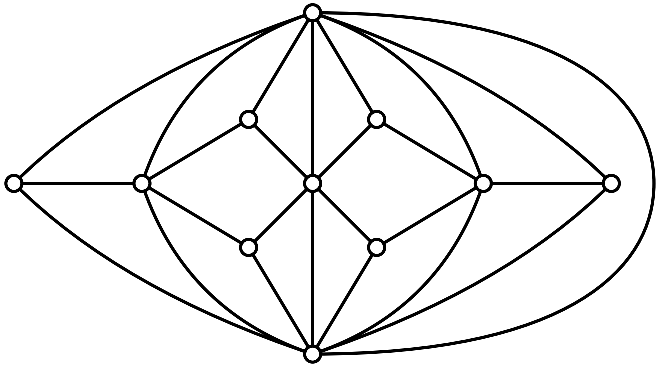


Example



Goldner–Harary graph (minus two edges)

choose any

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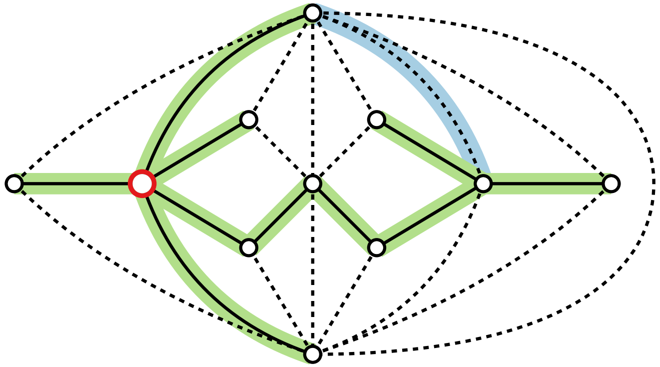
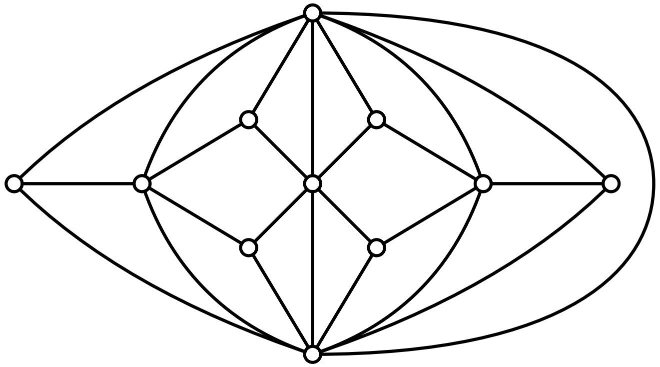


Example



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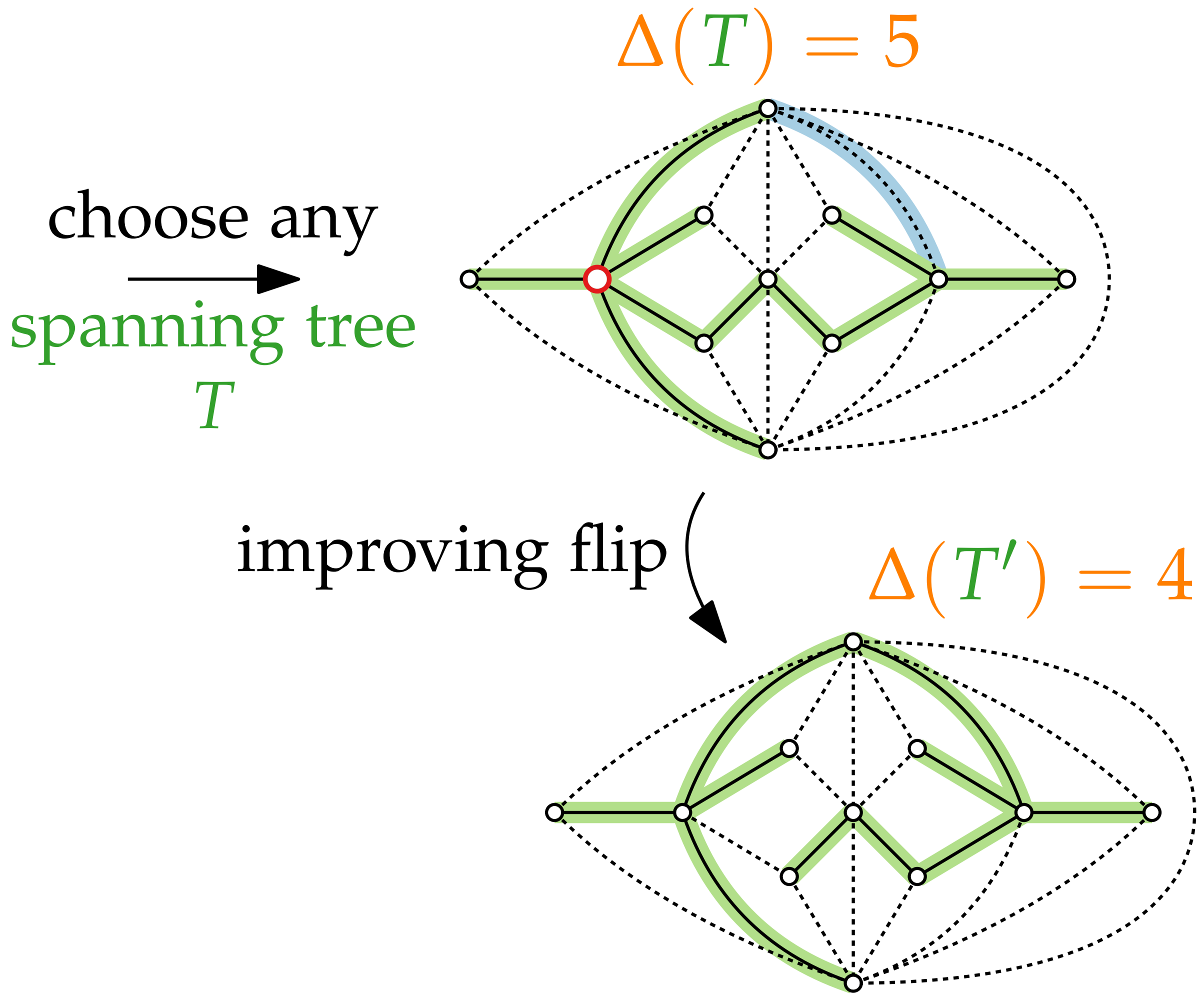
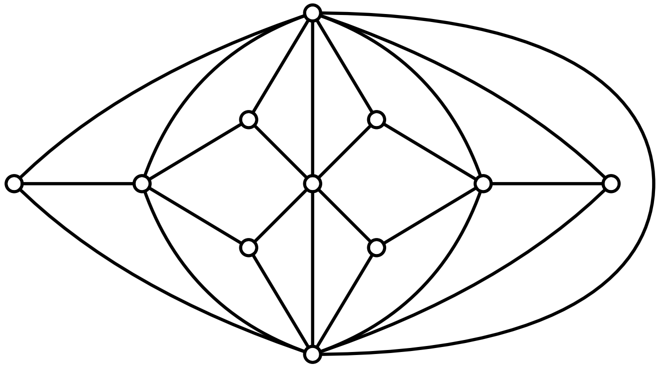


Example



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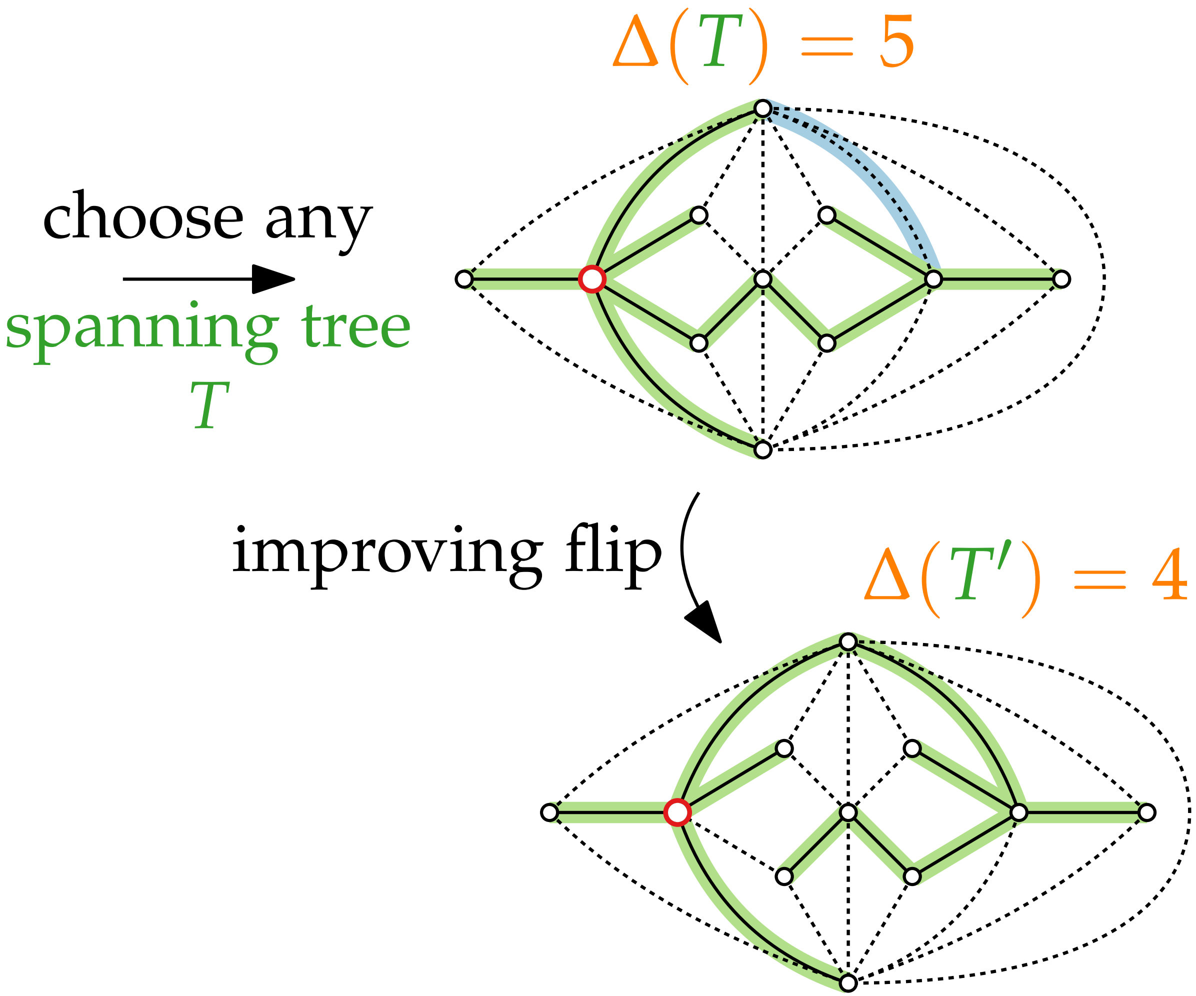
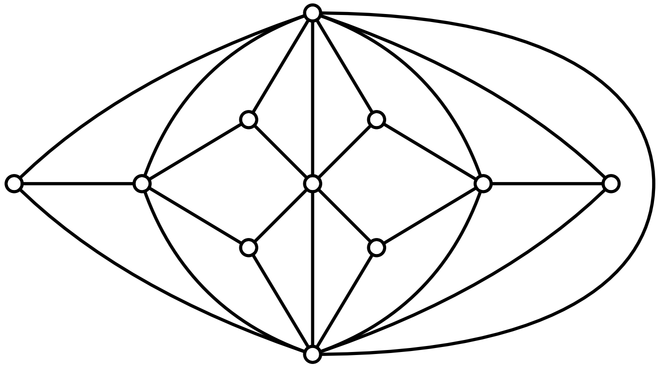


Example

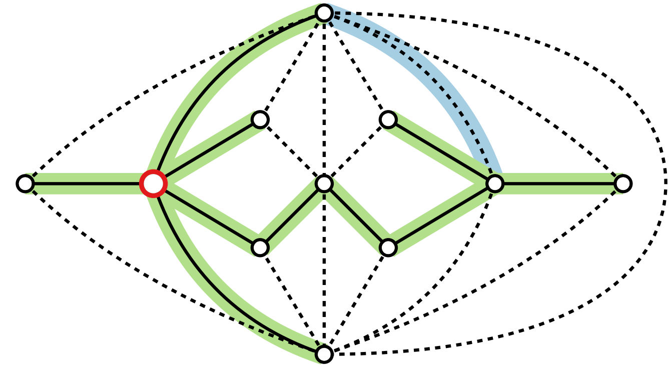


Goldner–Harary graph (minus two edges)

choose any

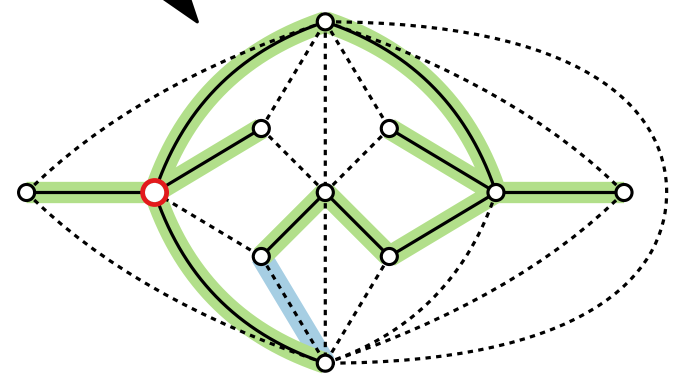
spanning tree T

$\Delta(T) = 5$

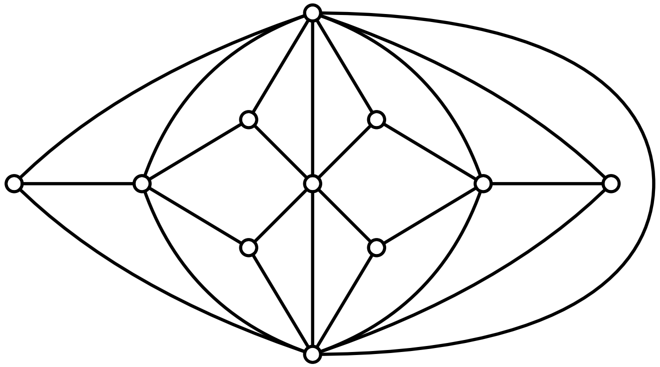


improving flip

$\Delta(T') = 4$



Example

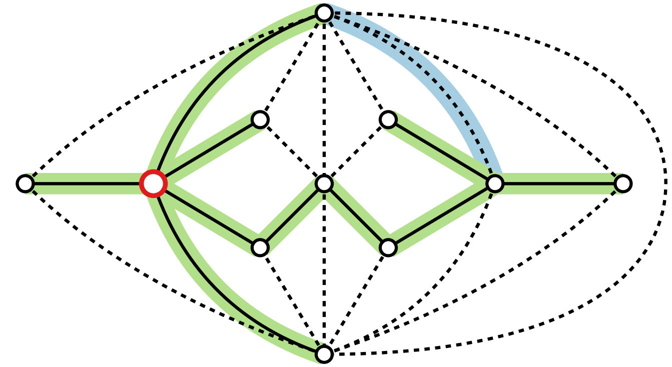


Goldner–Harary graph (minus two edges)

choose any

spanning tree T

$\Delta(T) = 5$



improving flip

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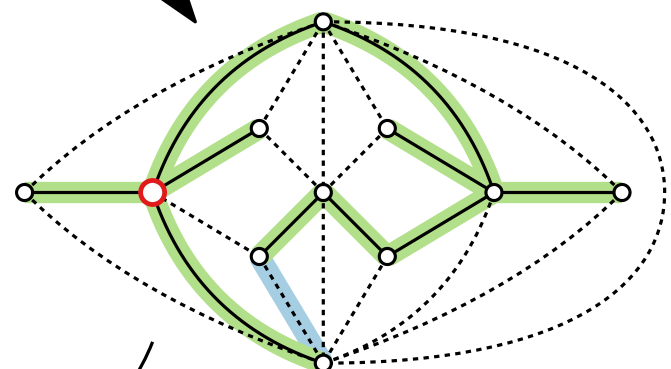
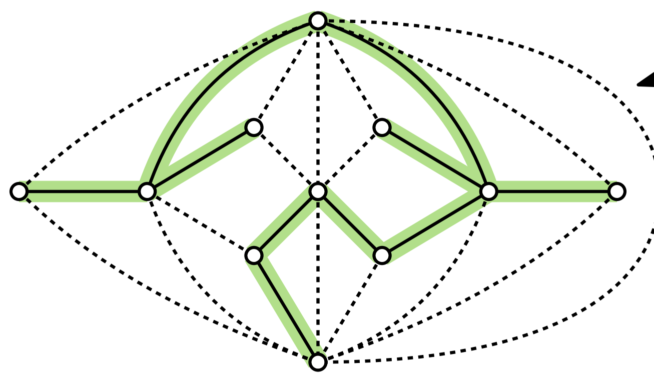
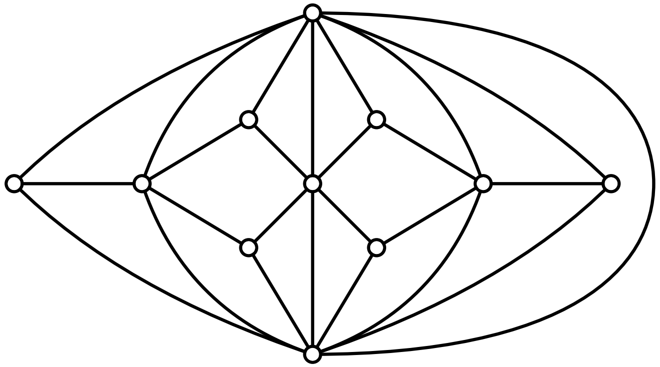


improving flip

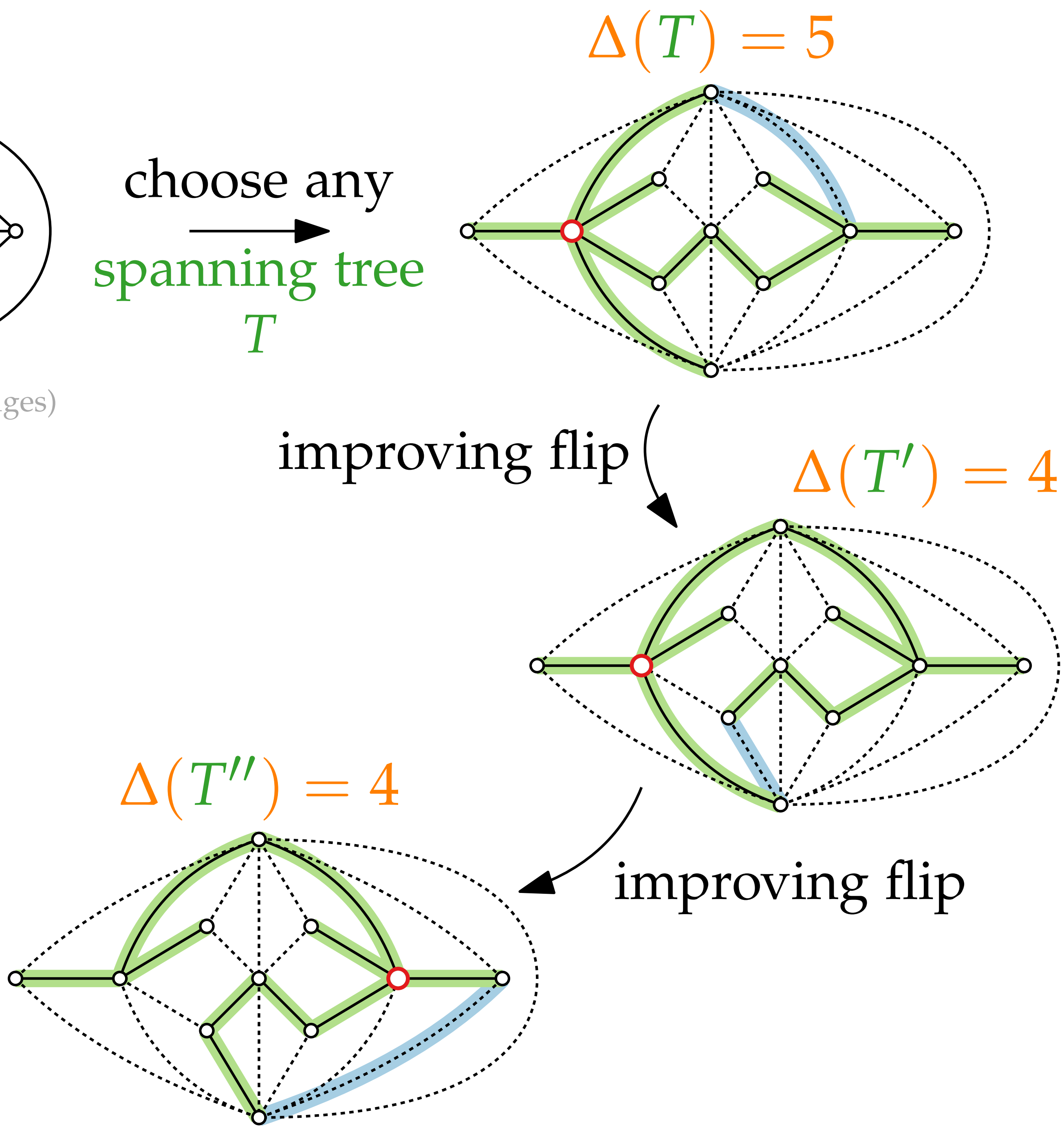
$\Delta(T'') = 4$



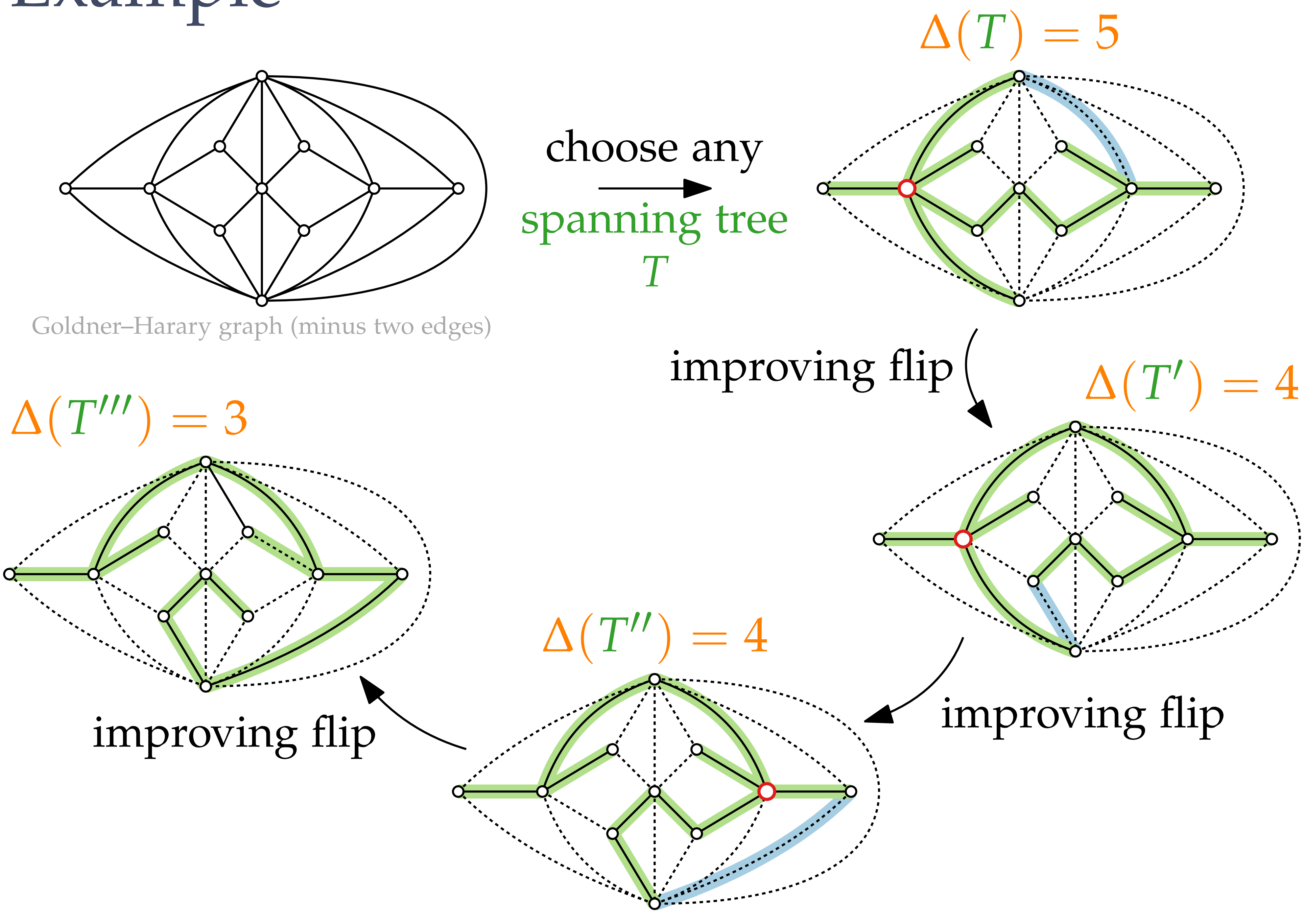
Example



Goldner–Harary graph (minus two edges)

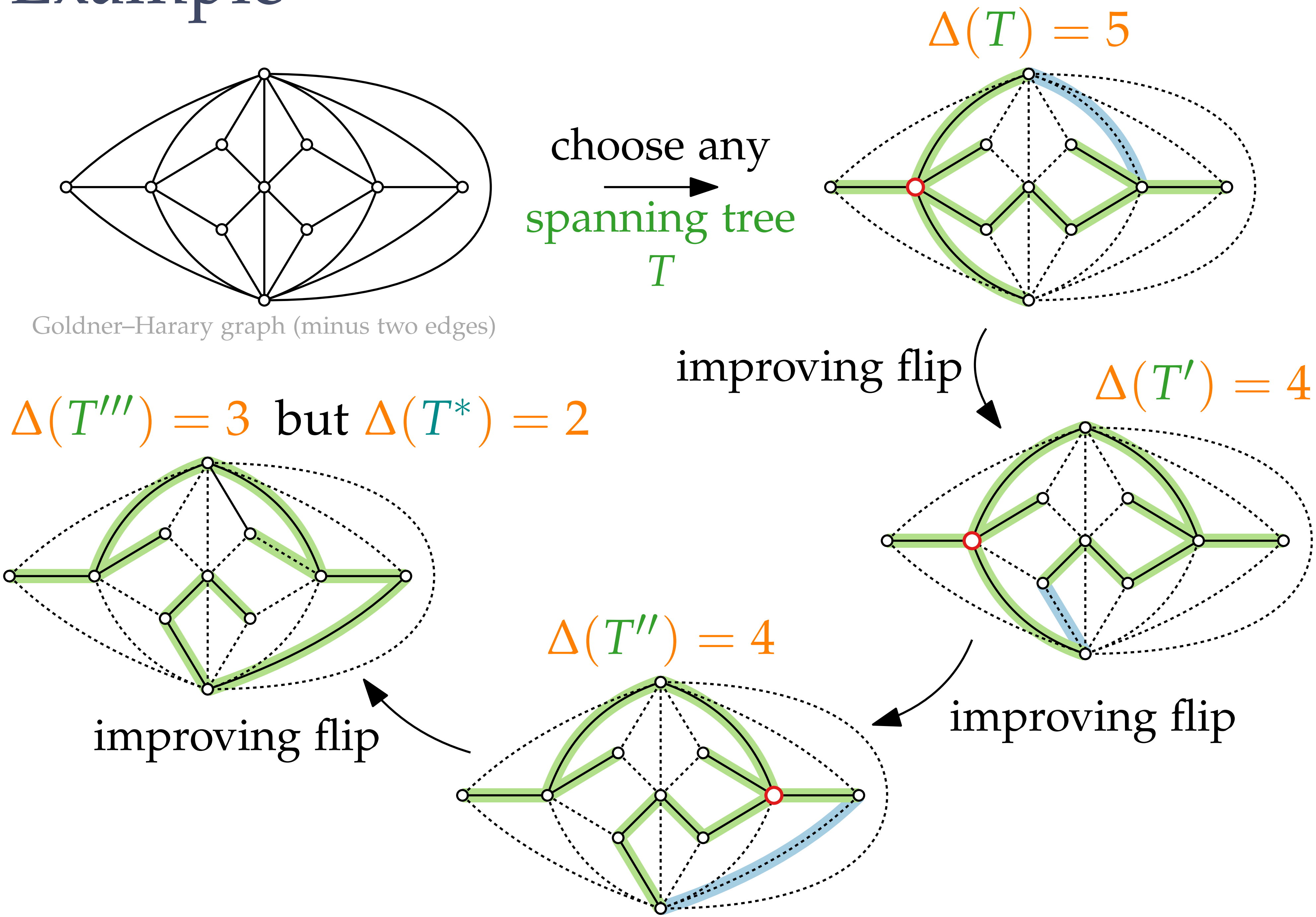


Example



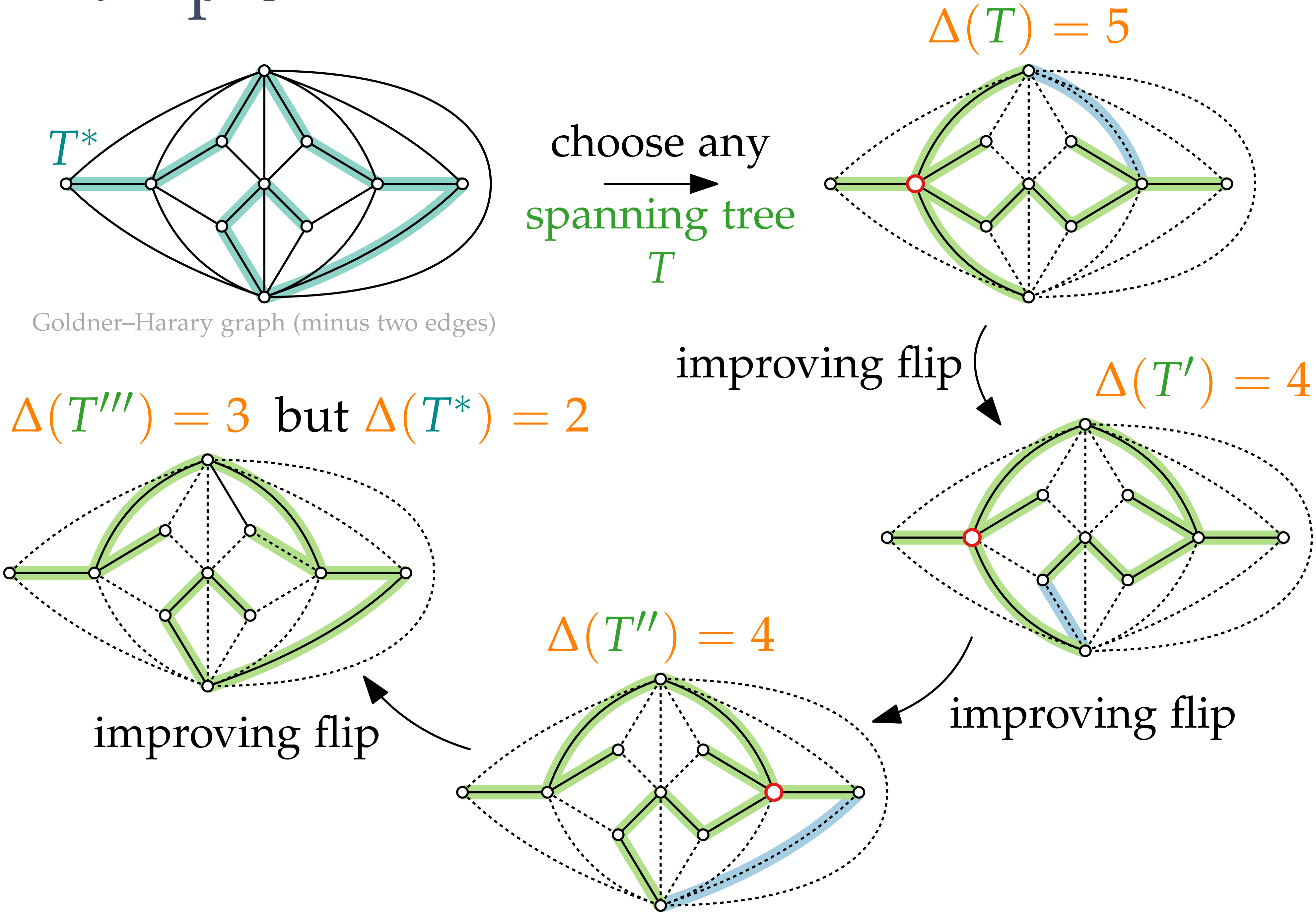
Goldner–Harary graph (minus two edges)
choose any
spanning tree
T
Δ(T) = 5
improving flip
Δ(T′) = 4
improving flip
Δ(T″) = 4
improving flip
Δ(T‴) = 3

Example



Goldner–Harary graph (minus two edges)
choose any
spanning tree
T
$\Delta(T) = 5$
improving flip
$\Delta(T') = 4$
improving flip
$\Delta(T'') = 4$
improving flip
$\Delta(T''') = 3$ but $\Delta(T^*) = 2$

Example

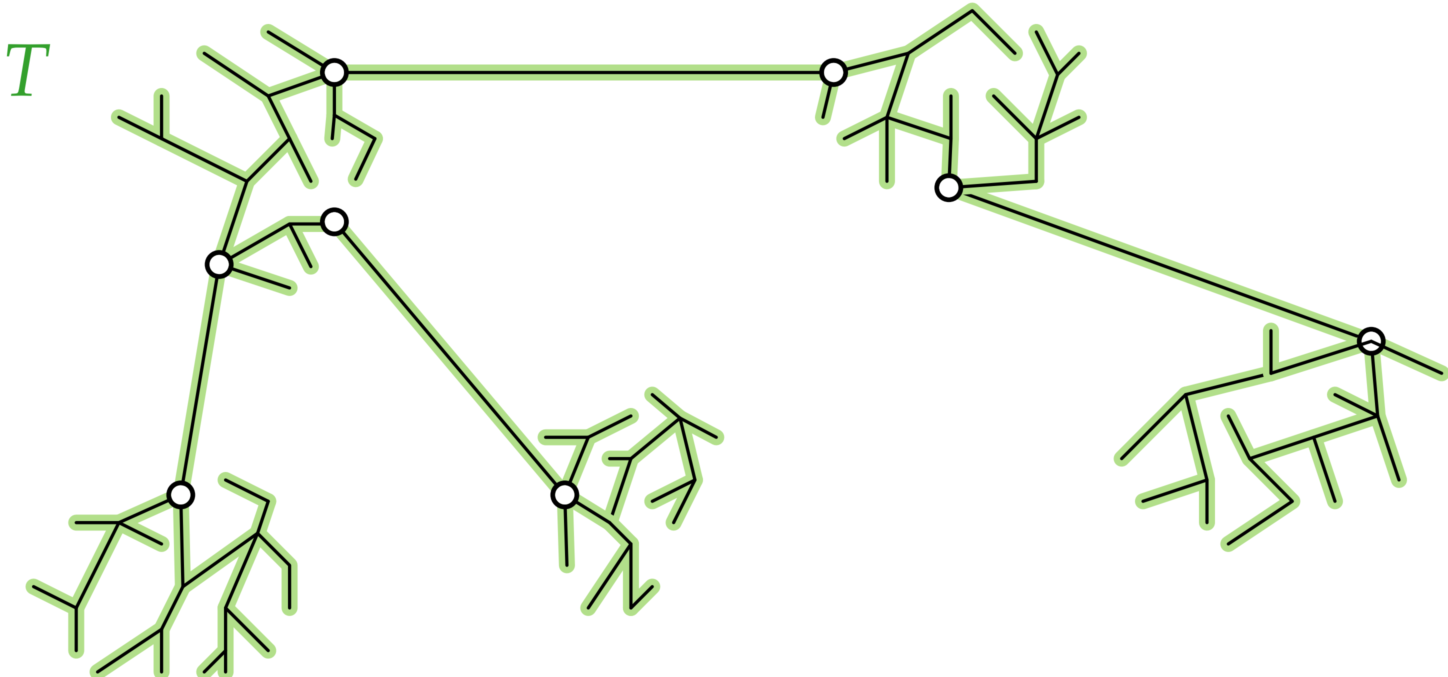

T^*
Goldner–Harary graph (minus two edges)
choose any
spanning tree
T
$\Delta(T) = 5$
improving flip
$\Delta(T') = 4$
improving flip
$\Delta(T'') = 4$
improving flip
$\Delta(T''') = 3$ but $\Delta(T^*) = 2$

Approximation Algorithms

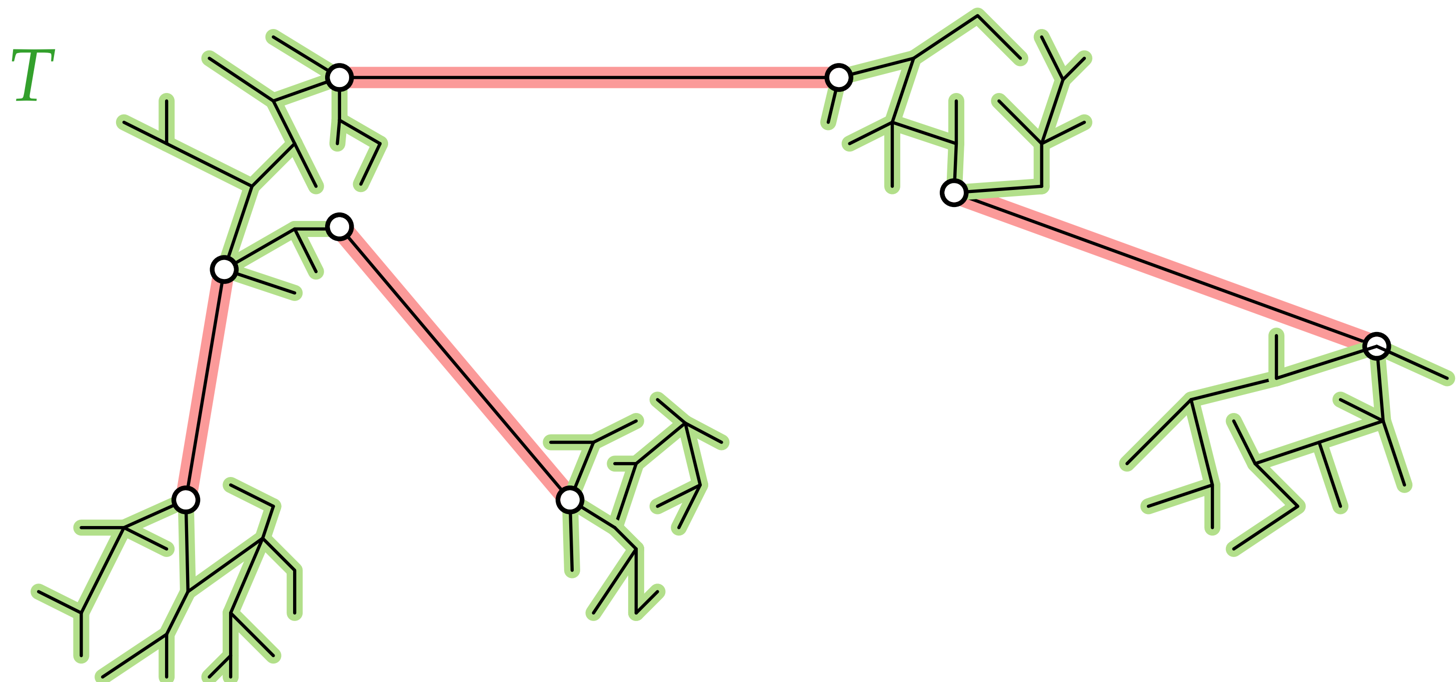
Lecture 10:
Minimum-Degree Spanning Tree
via Local Search

Part III:
Lower Bound

Decomposition

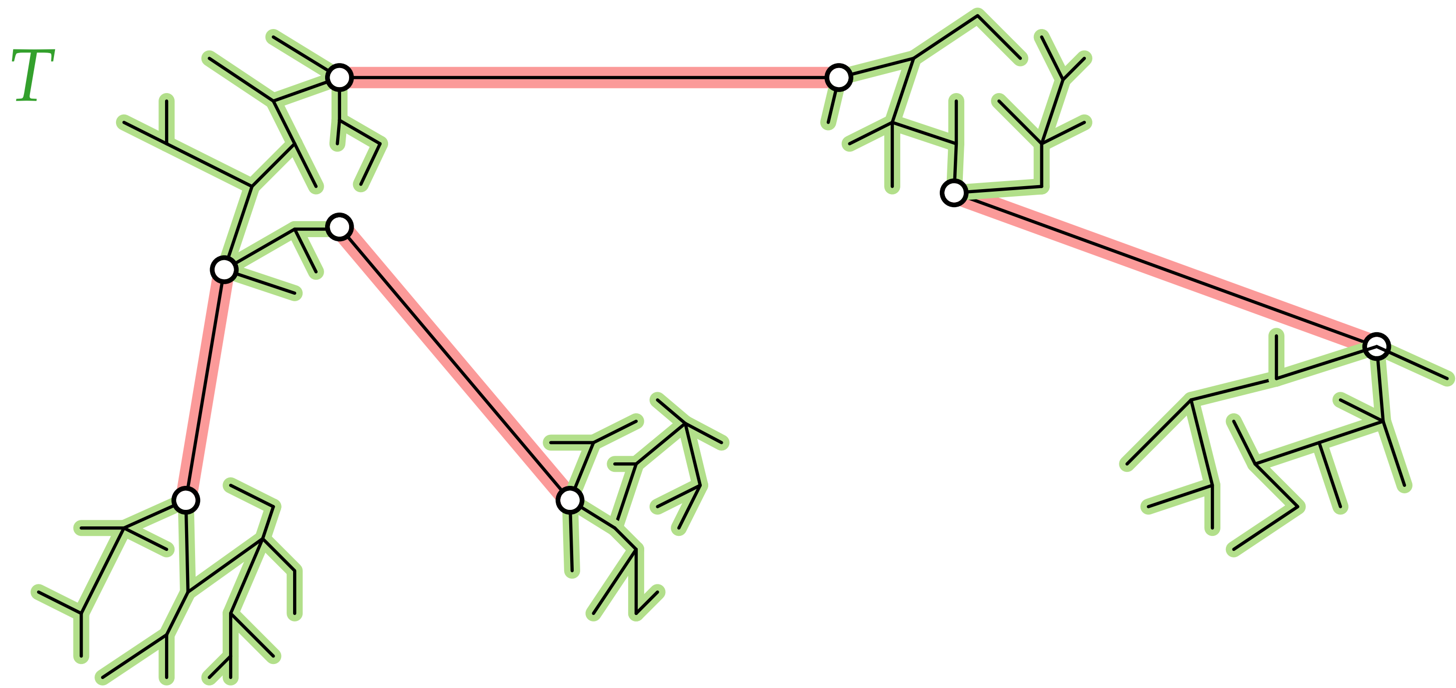


Decomposition



Decomposition

- Removing k edges decomposes T into $k+1$ components



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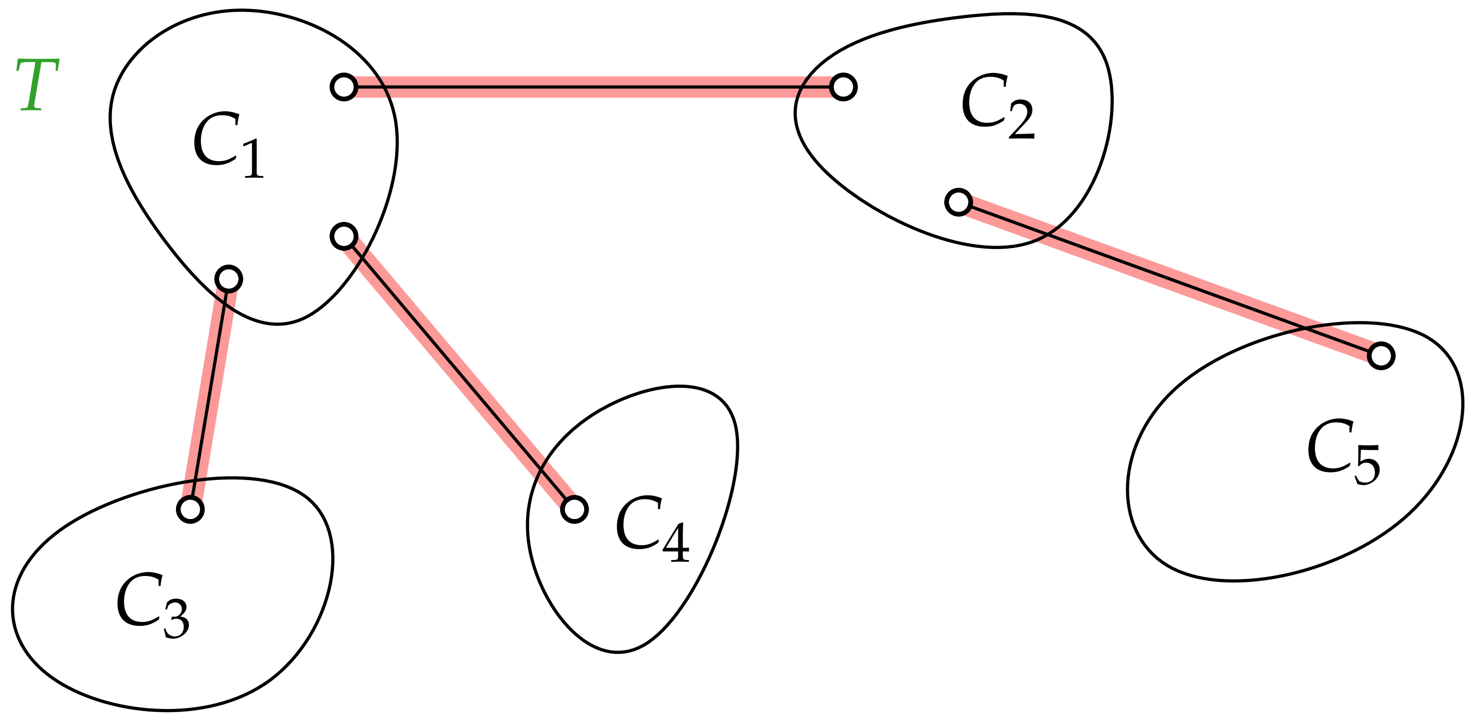
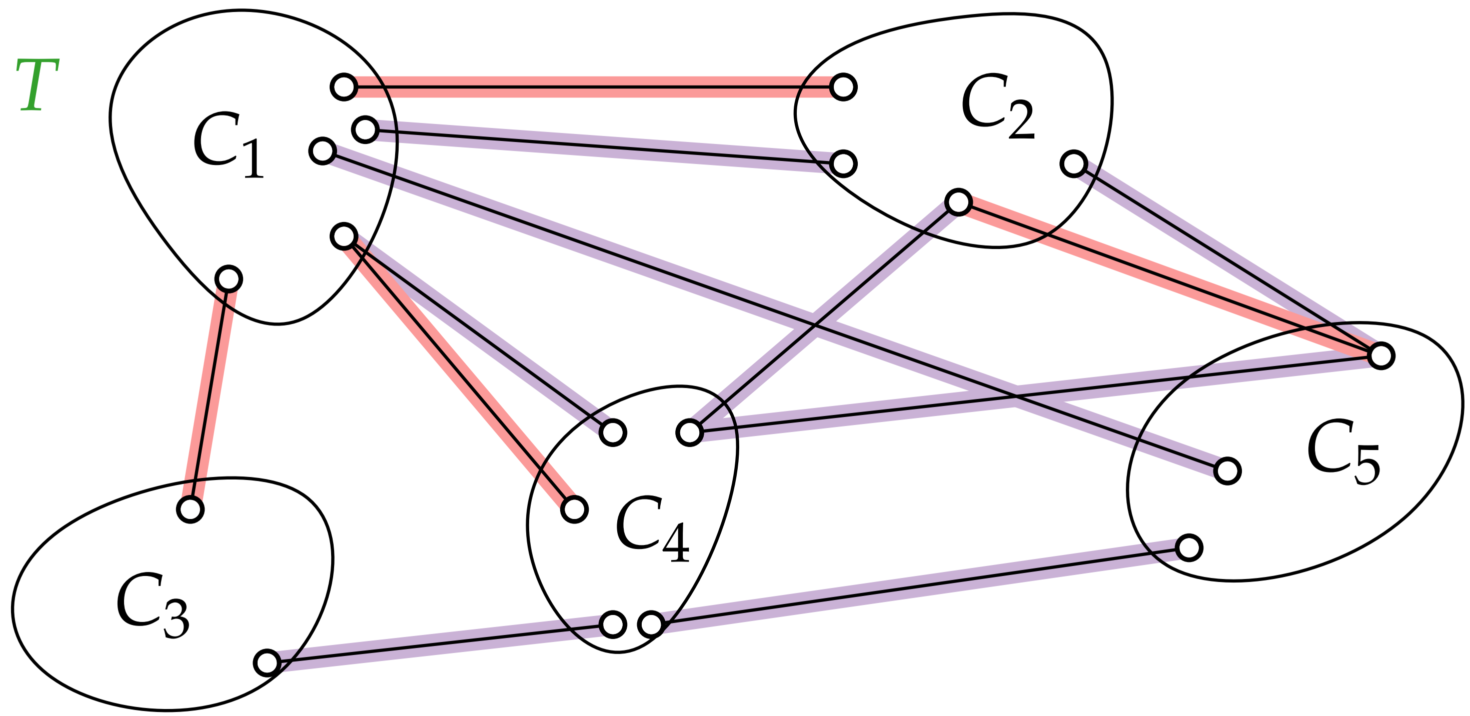


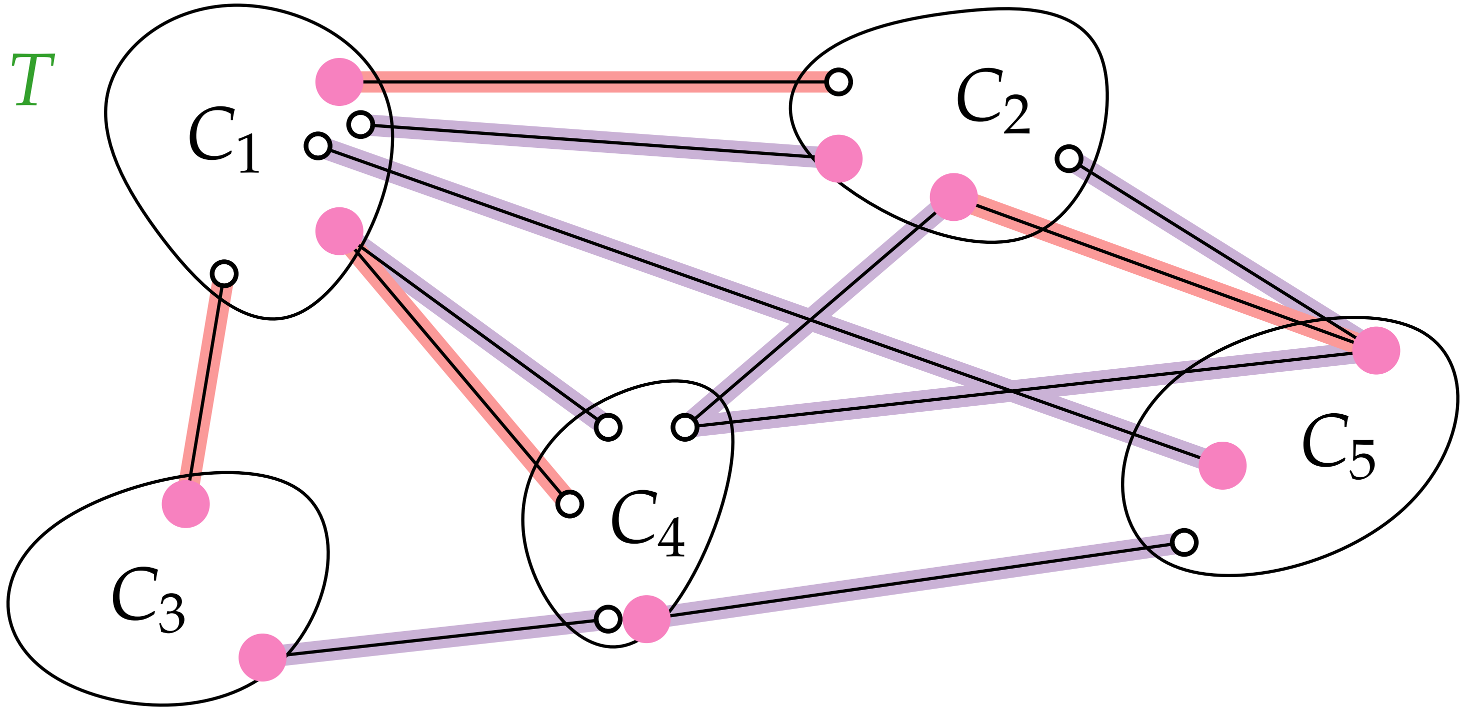
Decomposition

- Removing k edges decomposes T into $k+1$ components
- $E' = \{$edges in G between different components $C_i \neq C_j\}$.



Decomposition

- Removing k edges decomposes T into $k+1$ components
- $E' = \{\text{edges in } G \text{ between different components } C_i \neq C_j\}$.
- $S :=$ vertex cover of E'.



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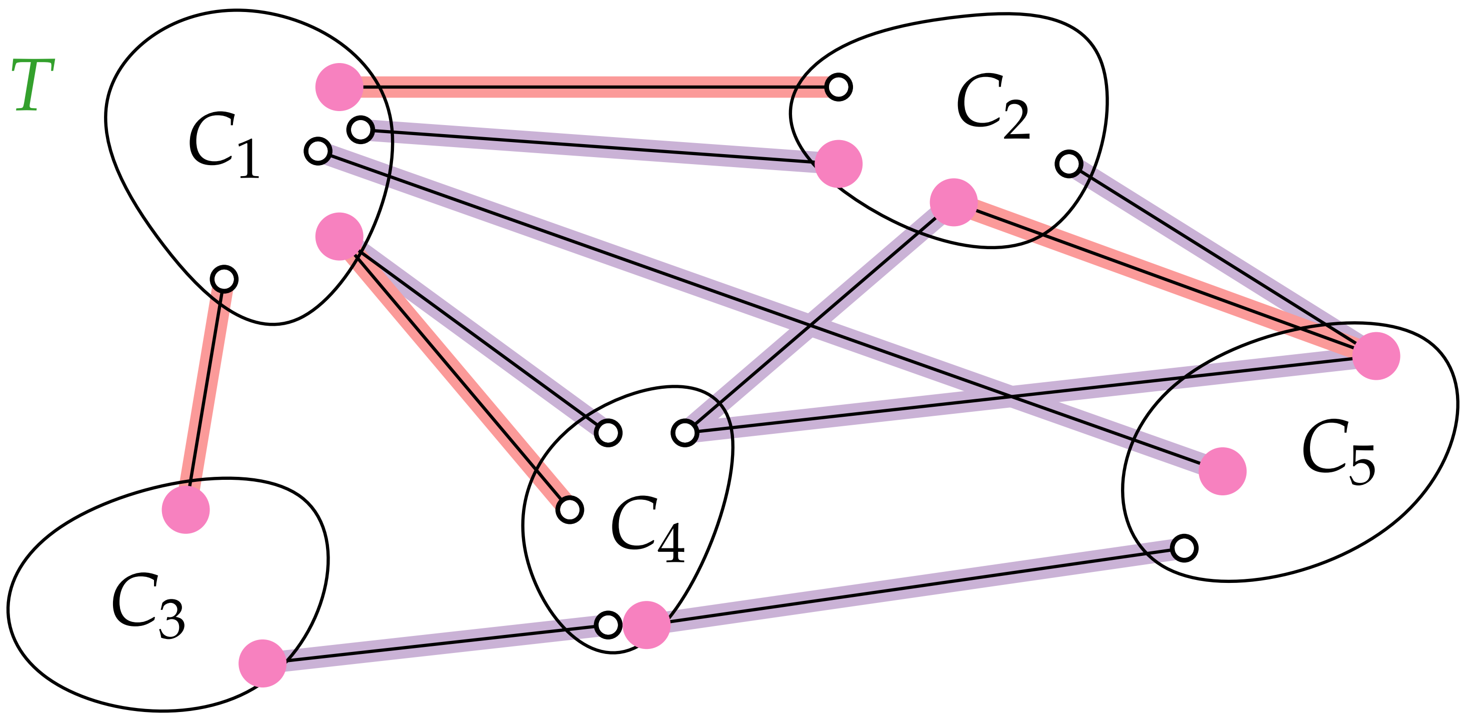

- $|E(T^*) \cap E'| \geq k$ for opt. spanning tree T^*

Decomposition

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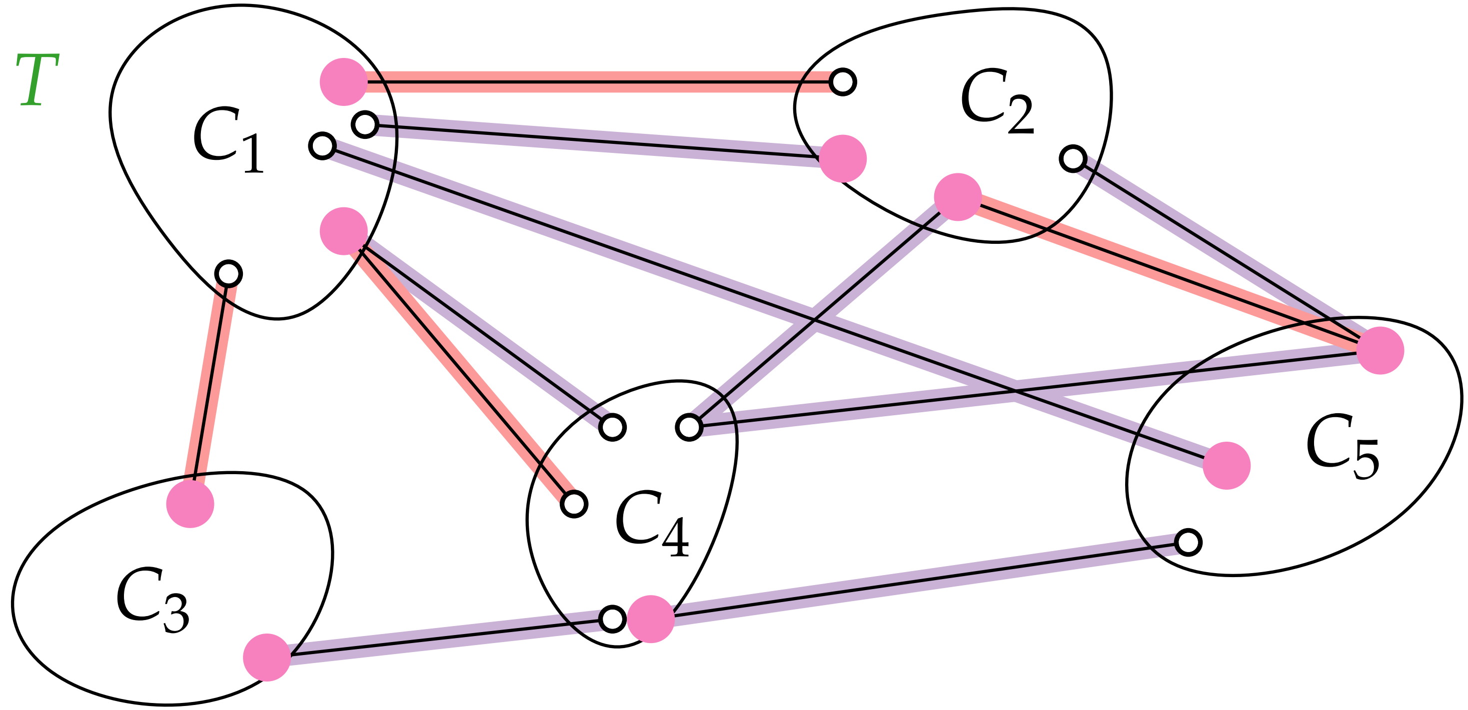

- $|E(T^*) \cap E'| \geq k$ for opt. spanning tree T^*
- $\sum_{v \in S} \deg_{T^*}(v) \geq k$

Decomposition $\Rightarrow$ Lower Bound for OPT

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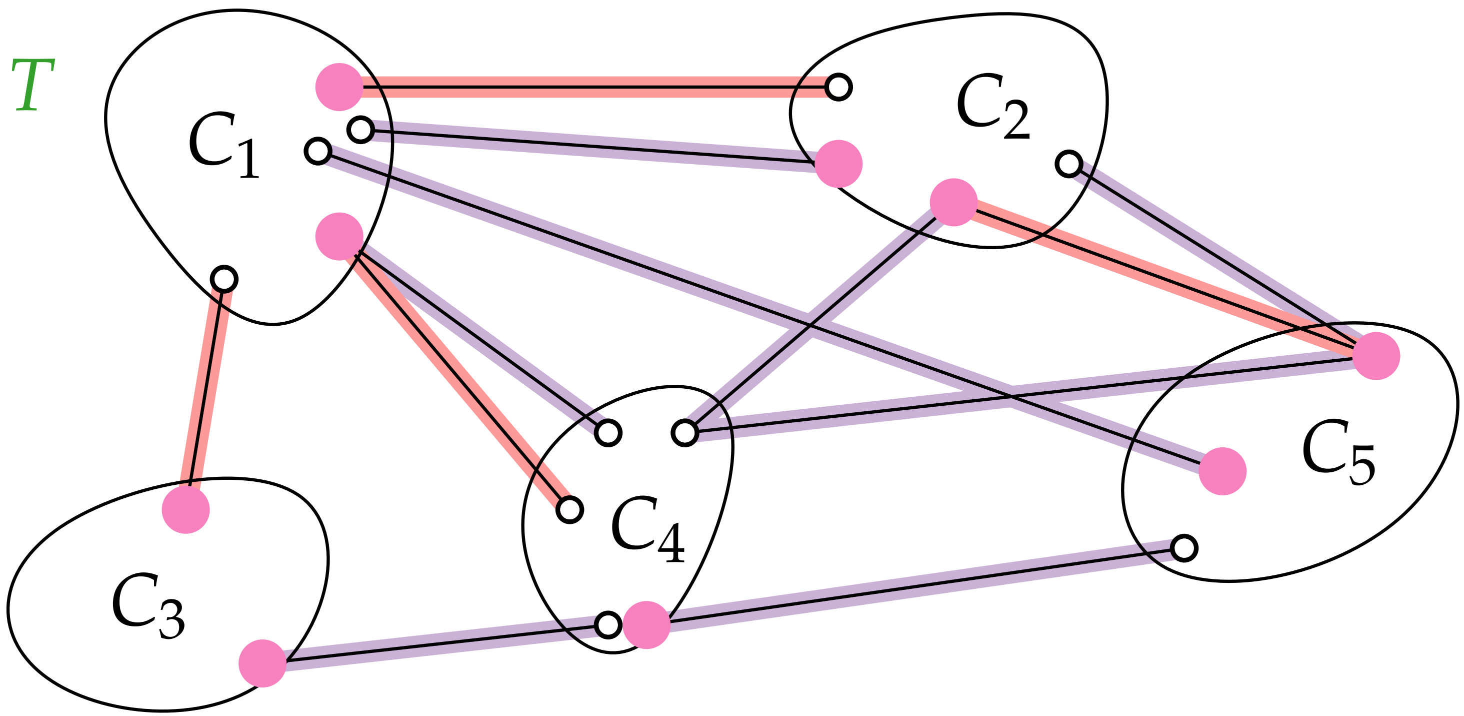

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Lemma 1.
$\Rightarrow$ OPT $\geq$

Decomposition $\Rightarrow$ Lower Bound for OPT

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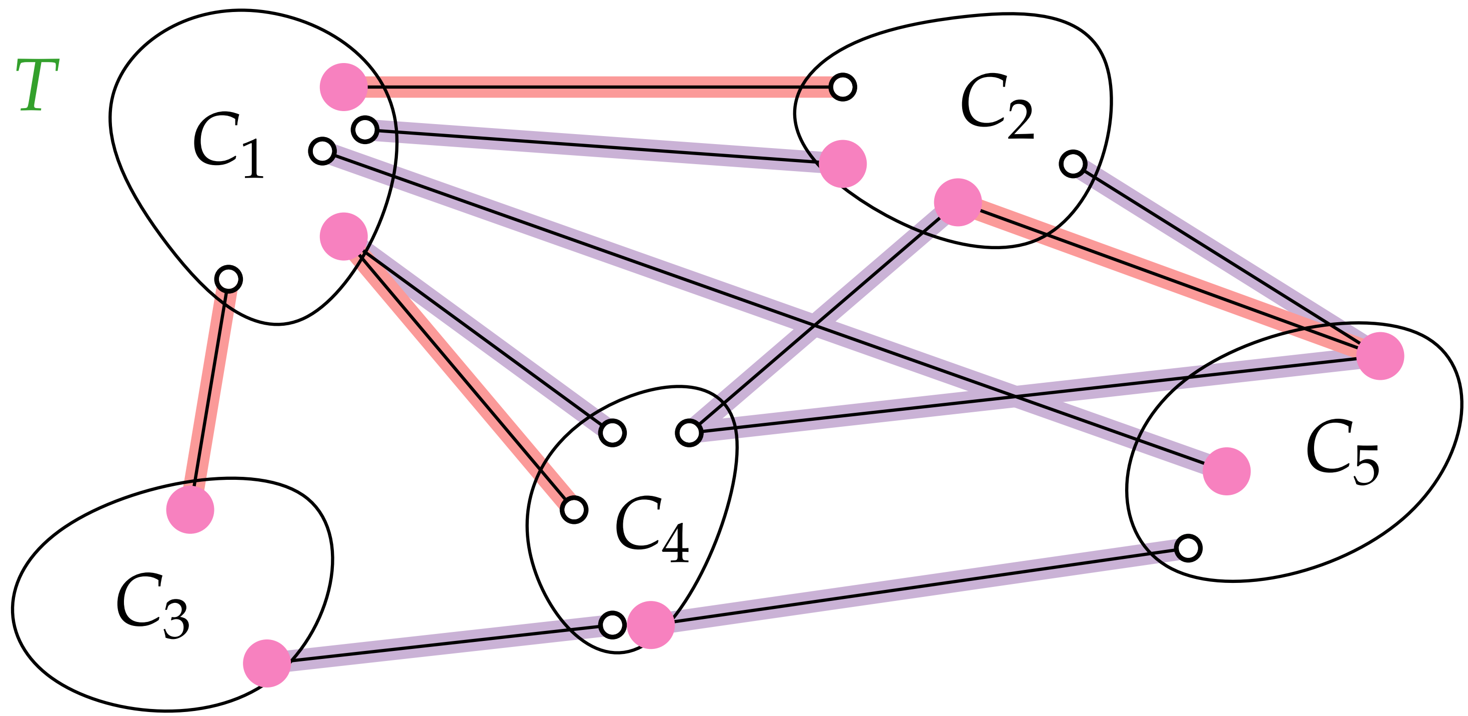

- $|E(T^*) \cap E'| \geq k$ for opt. spanning tree T^*
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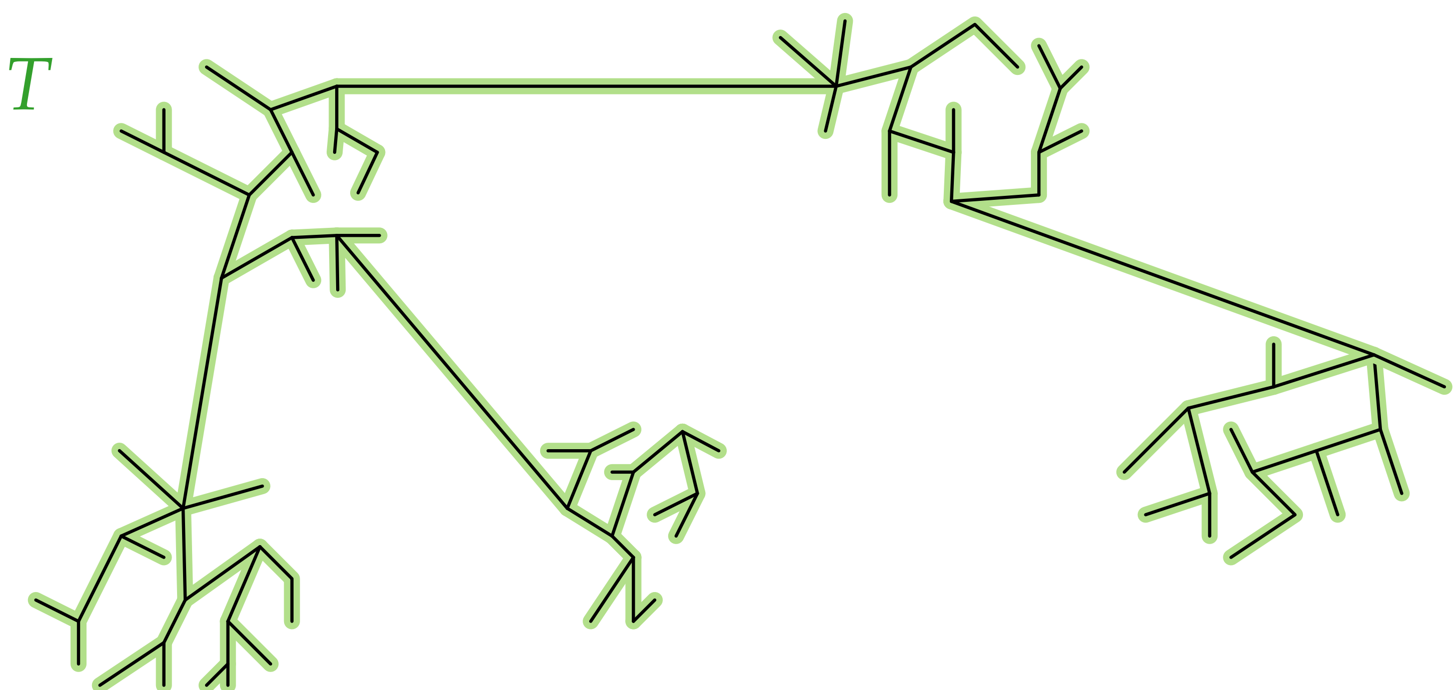
Lemma 1.
$\Rightarrow \mathrm{OPT} \geq k/|S|$

Approximation Algorithms

Lecture 10:
Minimum-Degree Spanning Tree
via Local Search

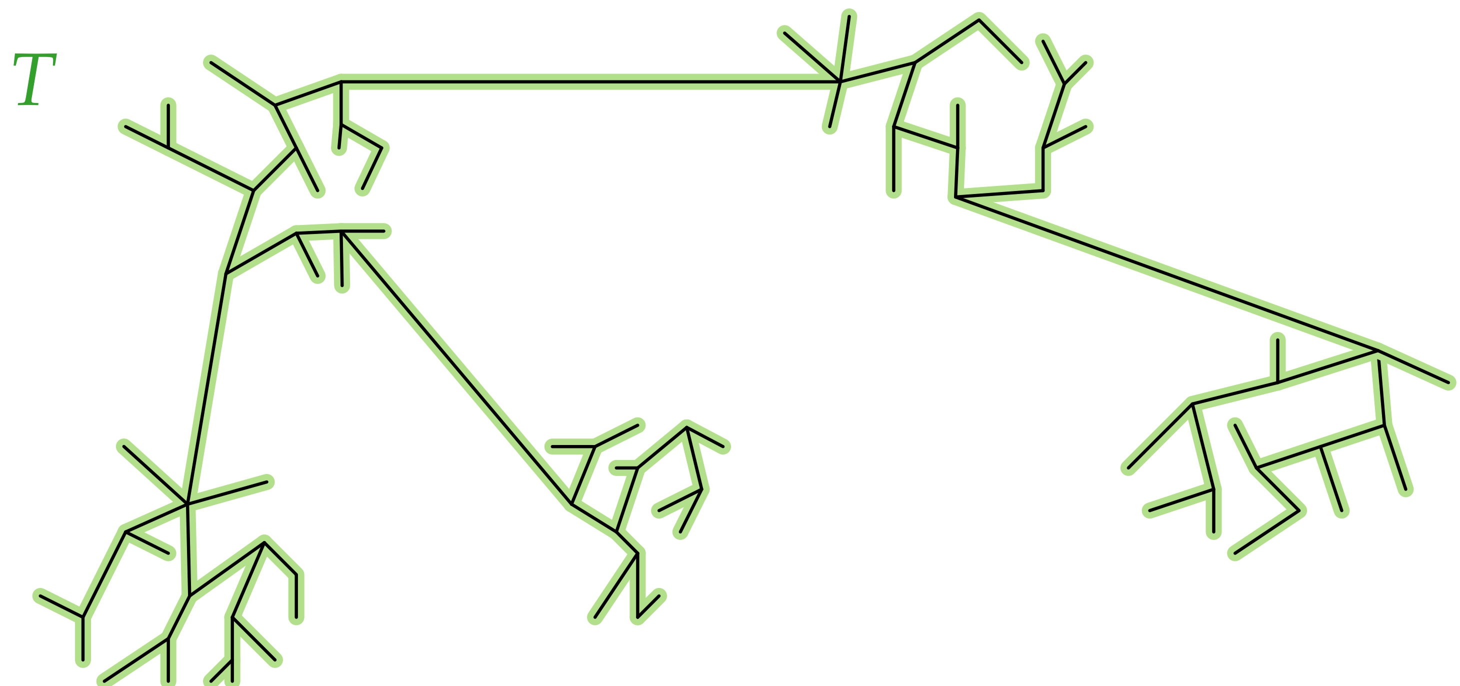
Part IV:
More Lemmas

More Lemmas



More Lemmas

Let S_i be the set of vertices v in T with $\deg_T(v) \geq i$.



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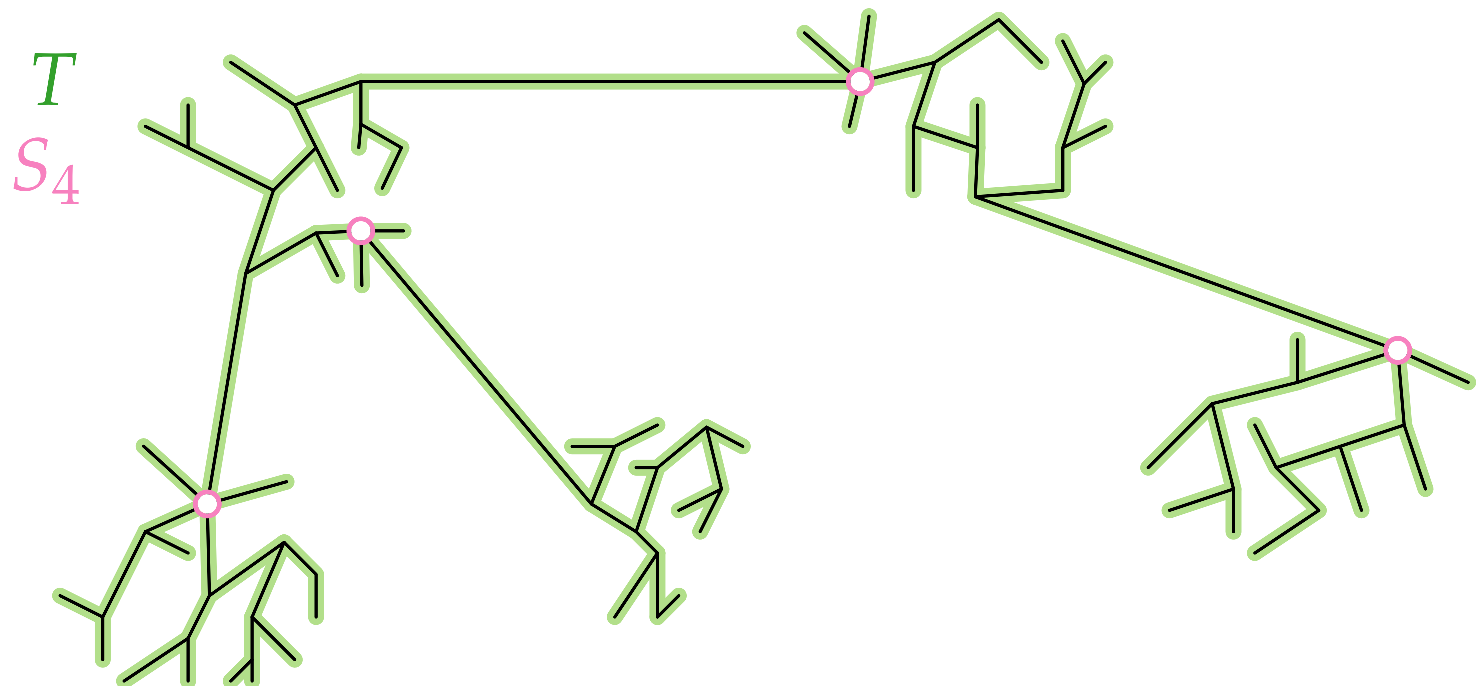

More Lemmas

Let S_i be the set of vertices v in T with $\deg_T(v) \geq i$.
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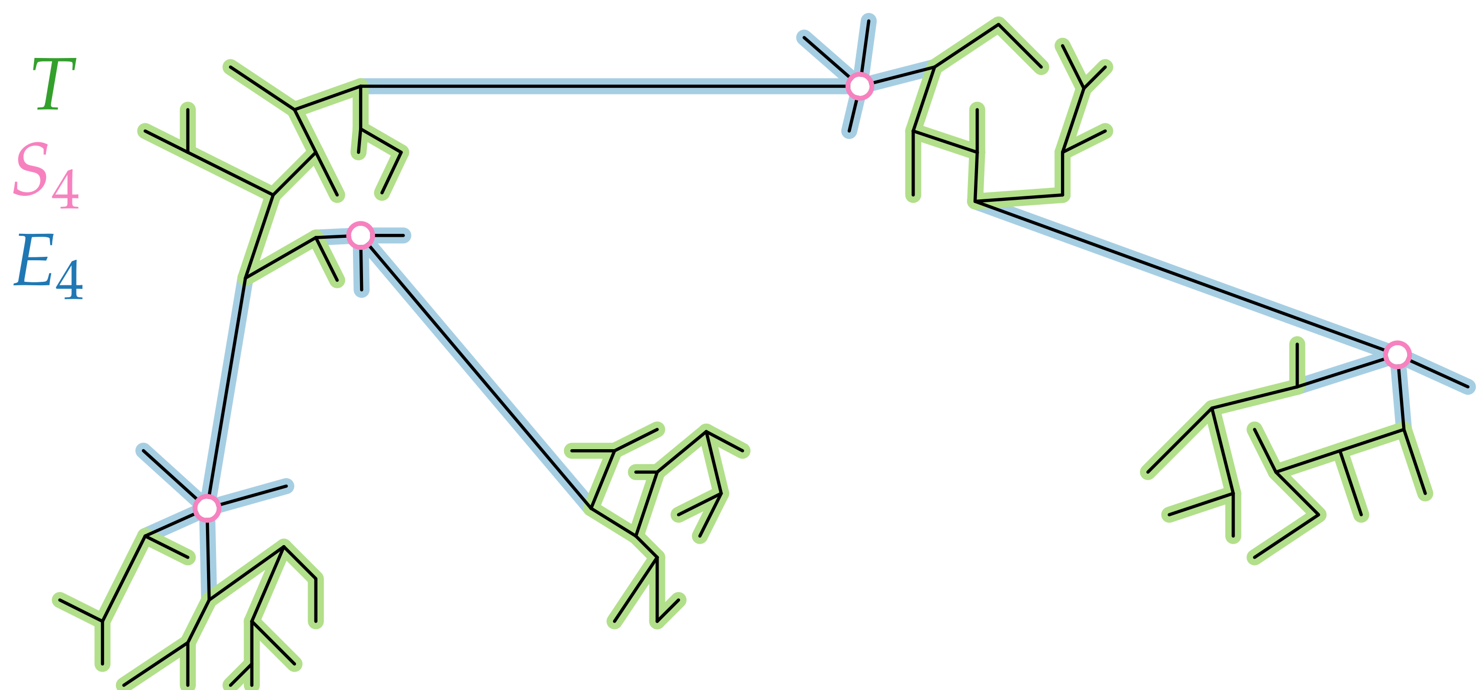

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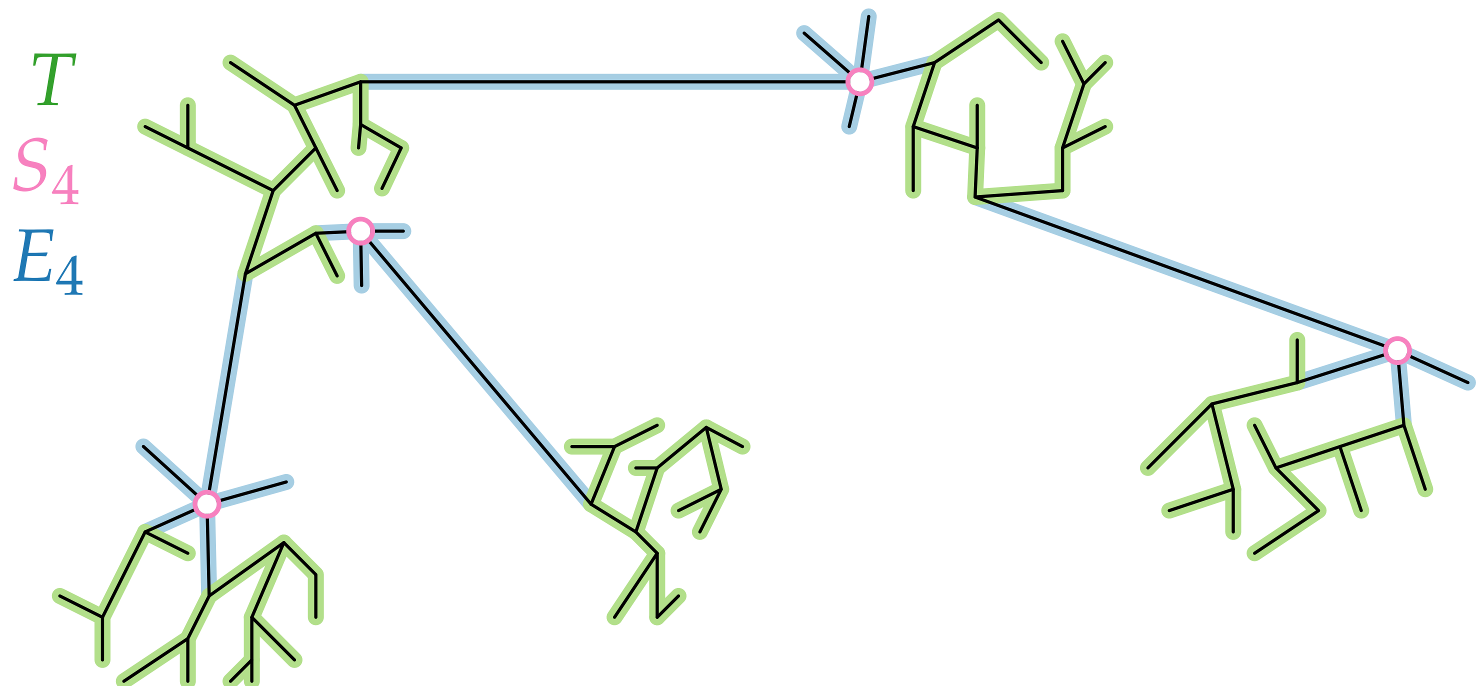

$\Rightarrow S_1 \supseteq S_2 \supseteq \dots$
$\Rightarrow S_1 = V(G)$

More Lemmas

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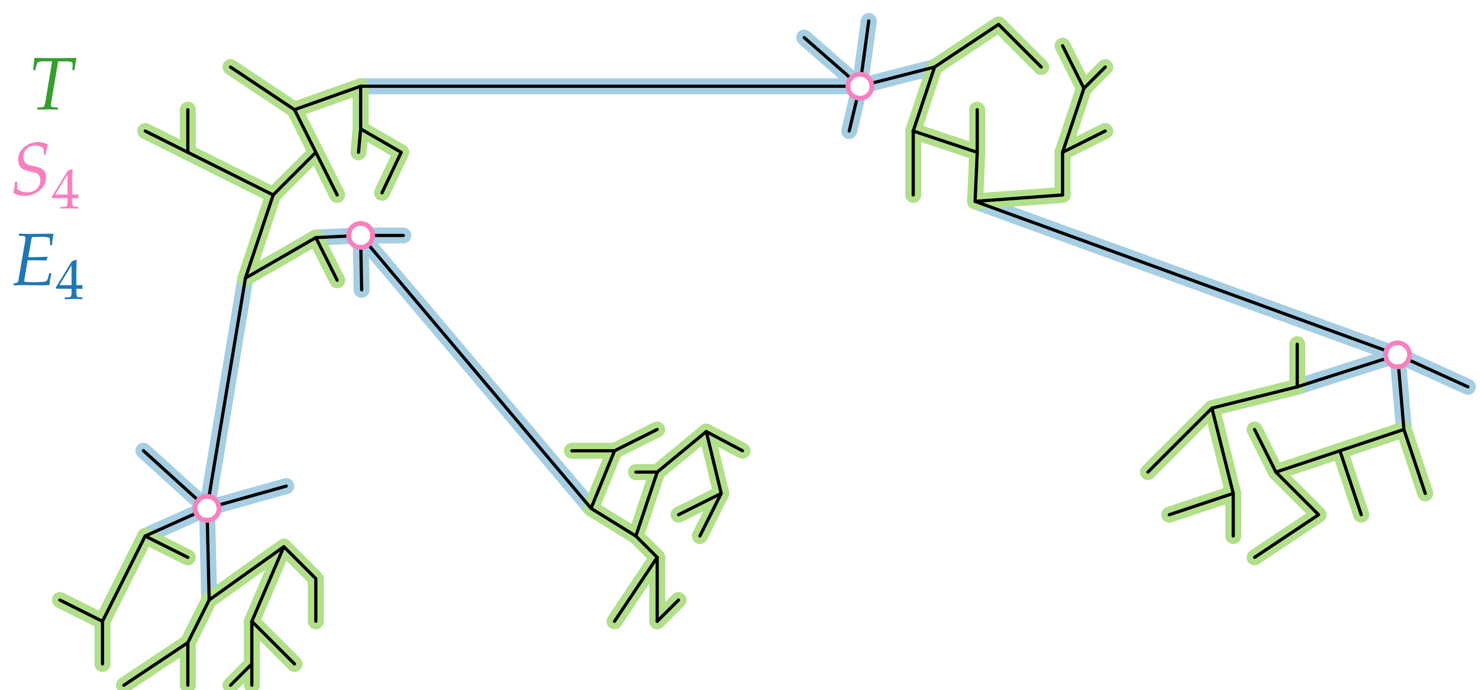

More Lemmas

$\Rightarrow S_1 \supseteq S_2 \supseteq \dots$
$\Rightarrow S_1 = V(G)$
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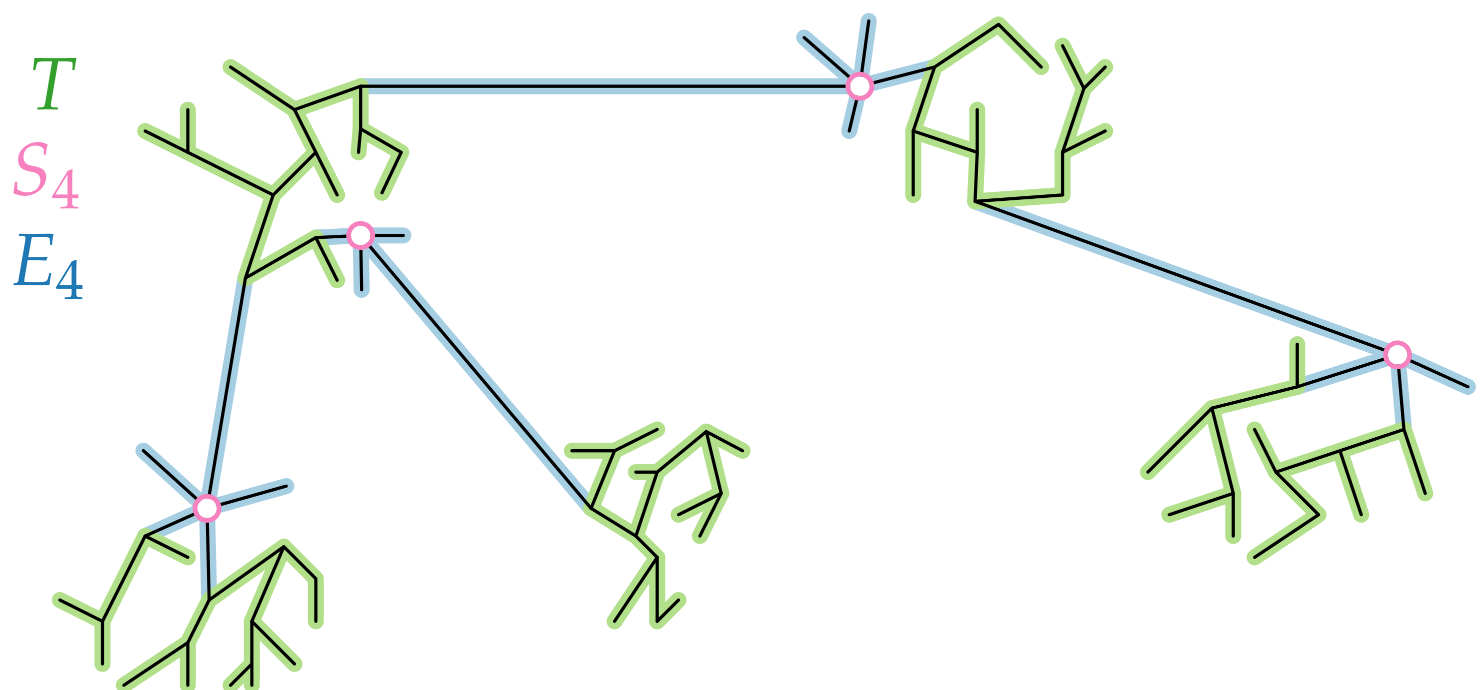
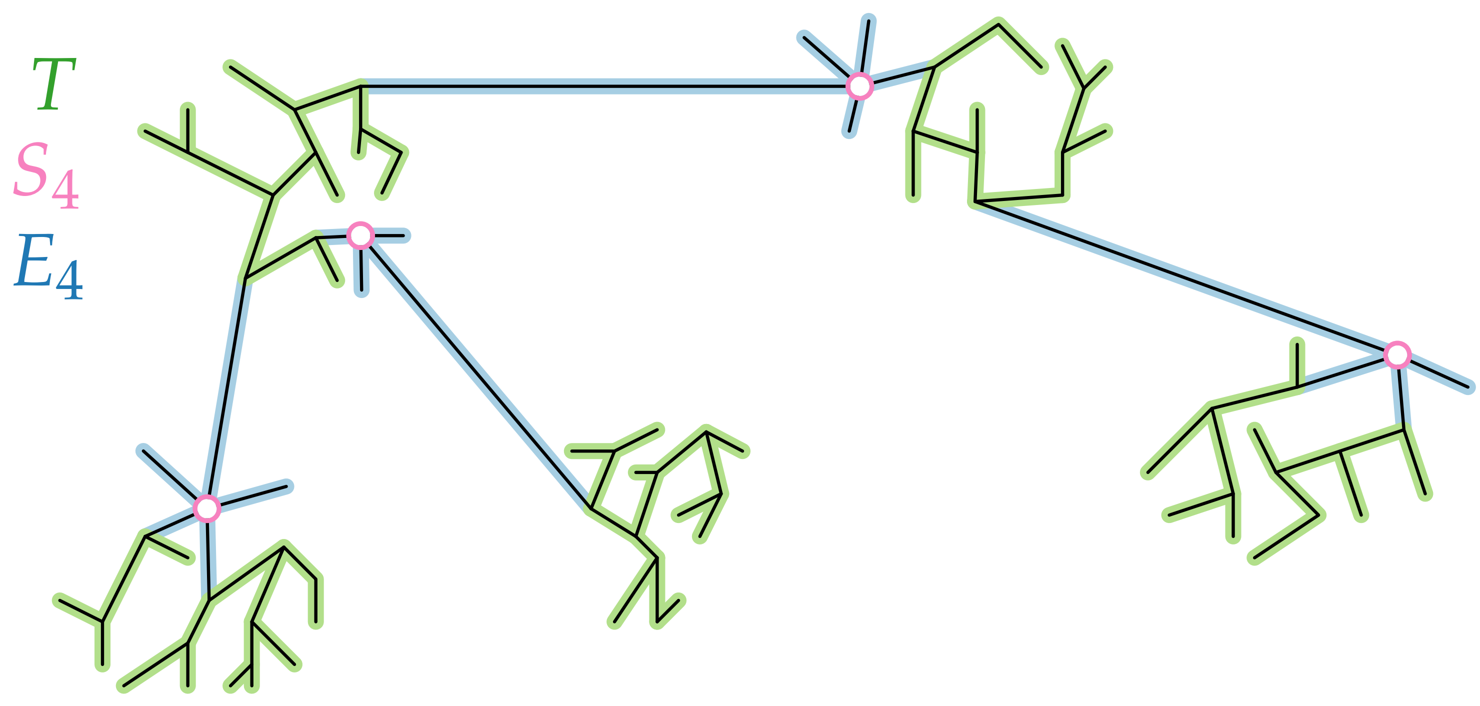


More Lemmas

$\Rightarrow S_1 \supseteq S_2 \supseteq \ldots$
$\Rightarrow S_1 = V(G)$
$\Rightarrow E_1 = E(T)$

Let S_i be the set of vertices v in T with $\deg_T(v) \geq i$.
Let E_i be the set of edges in T incident to S_i.

Lemma 2. There is some $i \geq \Delta(T) - \ell + 1$ with $|S_{i-1}| \leq 2|S_i|$.



More Lemmas

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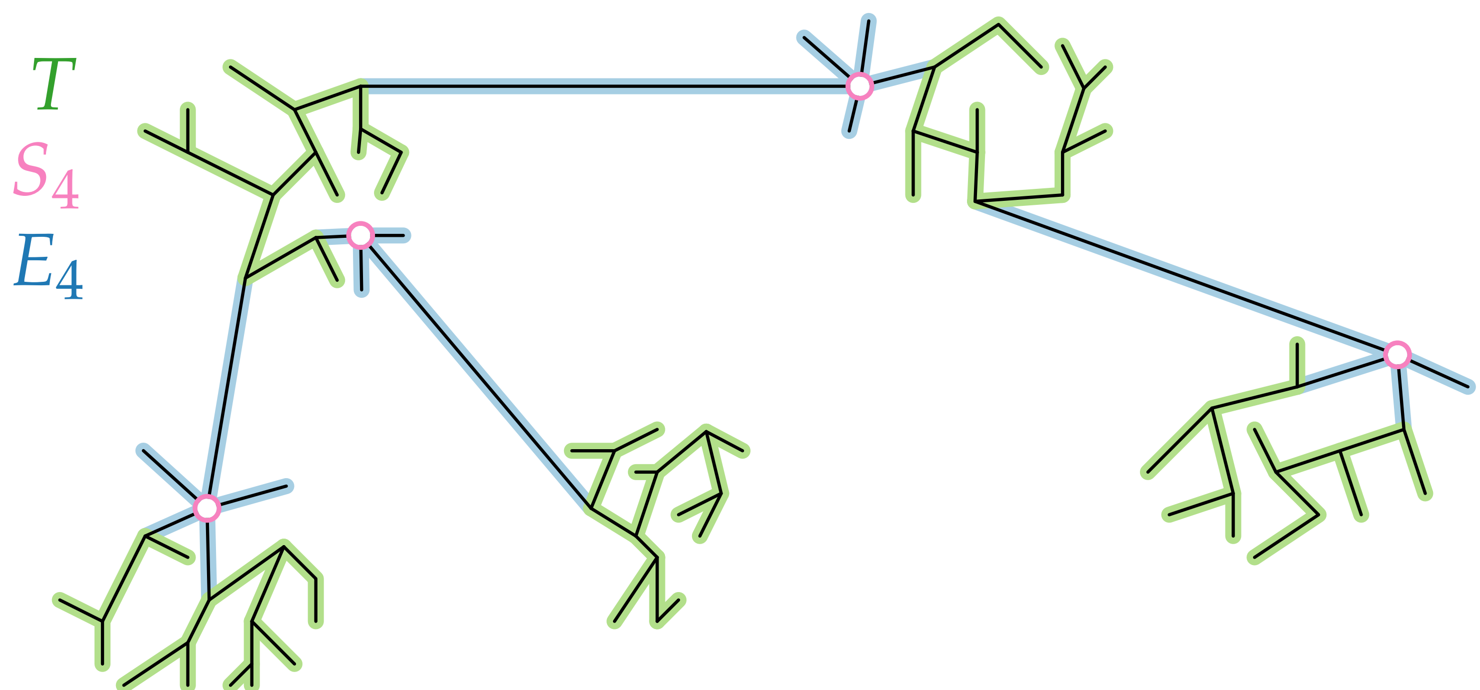
Proof. $|S_{\Delta(T)-\ell}| > 2^{\ell}|S_{\Delta(T)}|$

Otherwise



More Lemmas

$\Rightarrow S_1 \supseteq S_2 \supseteq \dots$
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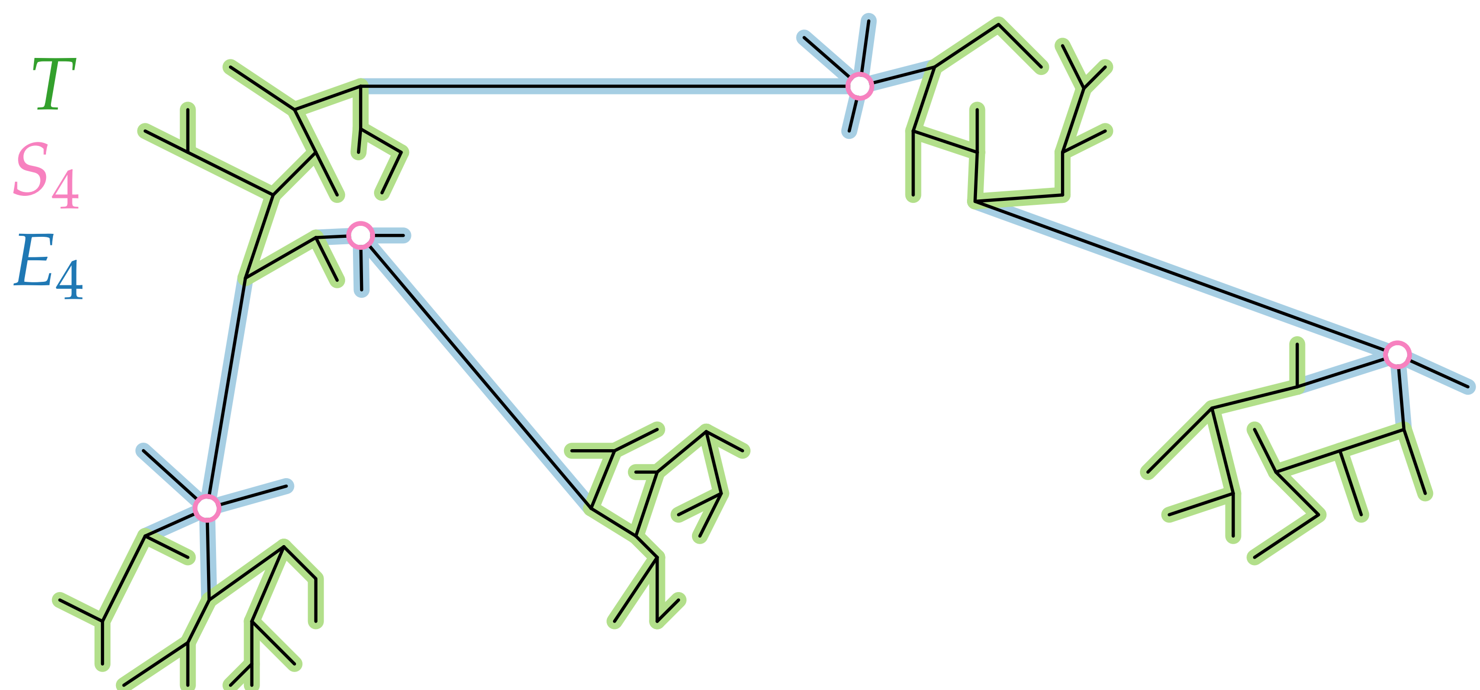
Proof. $|S_{\Delta(T)-\ell}| > 2^{\ell} |S_{\Delta(T)}| = 2^{\lceil \log_2 n \rceil} |S_{\Delta(T)}| \geq$

$\ell = \lceil \log_2 n \rceil$

Otherwise



More Lemmas

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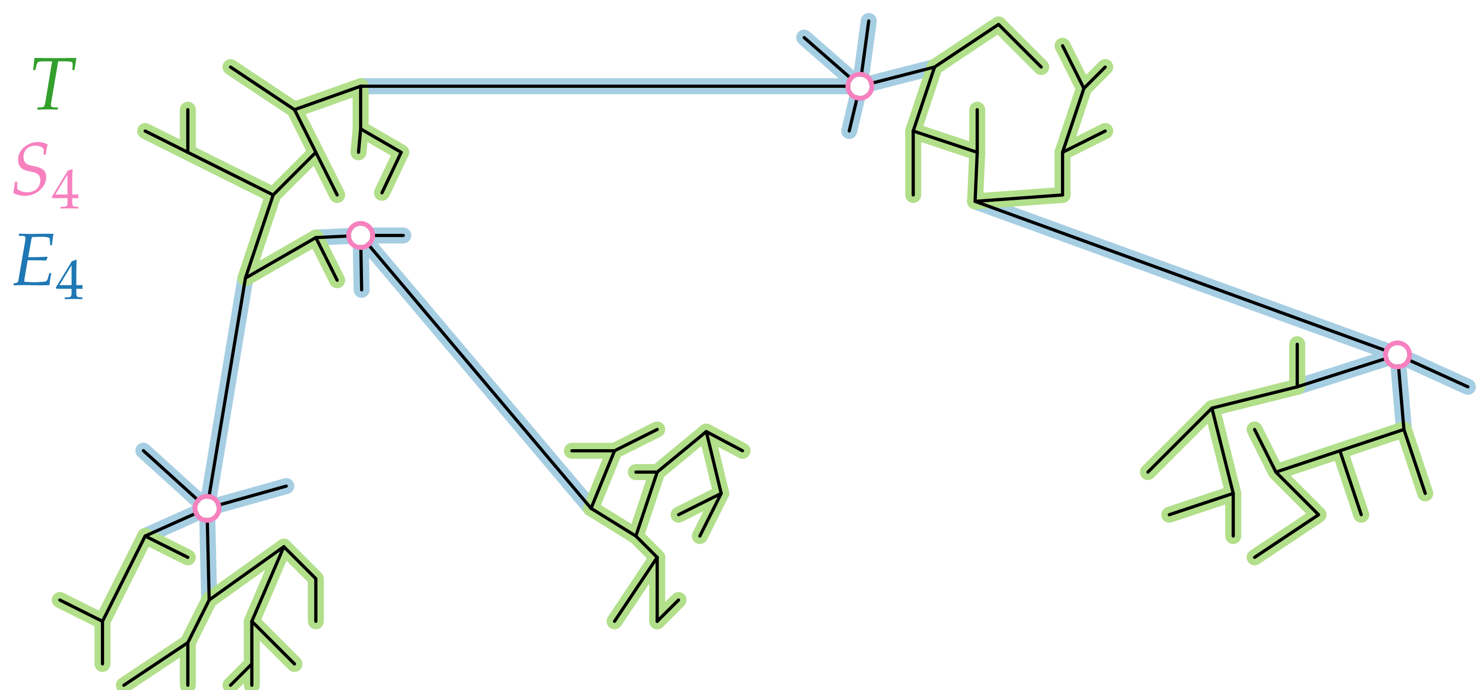
$\ell = \lceil \log_2 n \rceil$

Otherwise



More Lemmas

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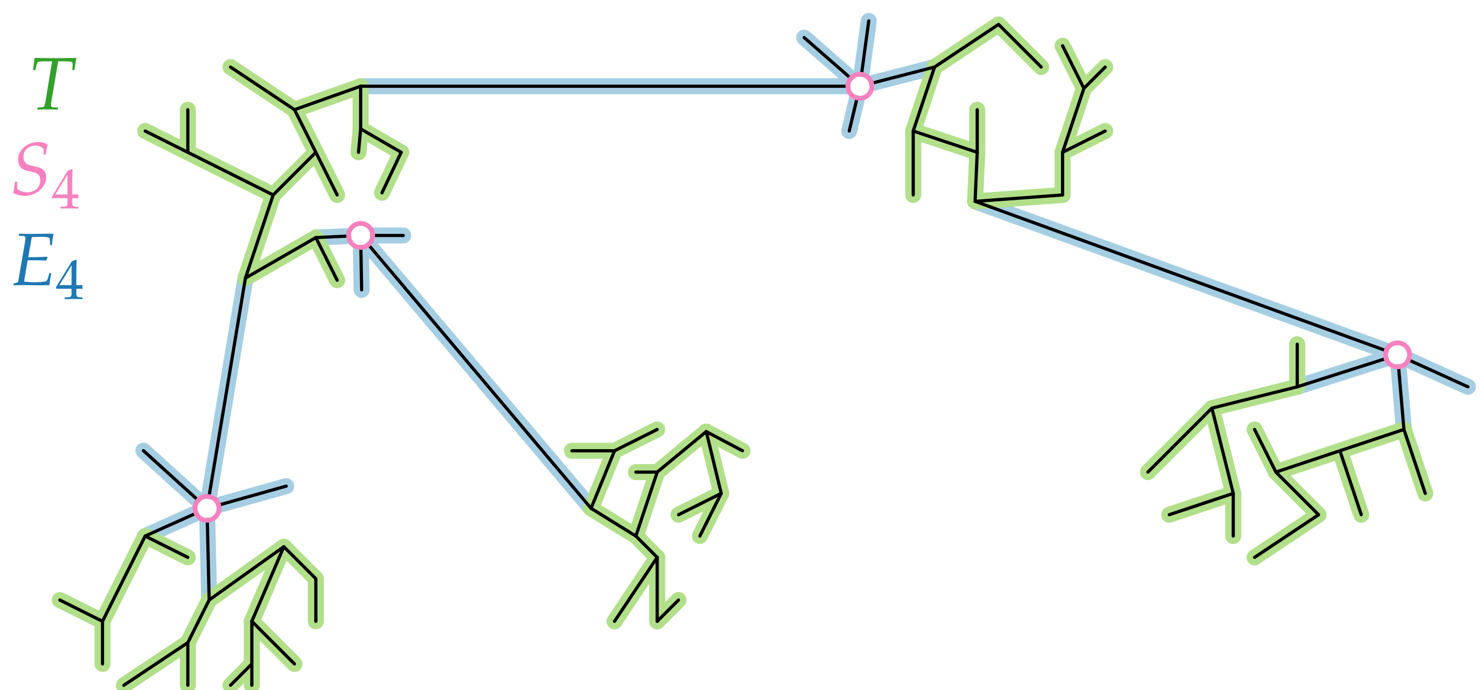
$\ell = \lceil \log_2 n \rceil$

Otherwise



More Lemmas

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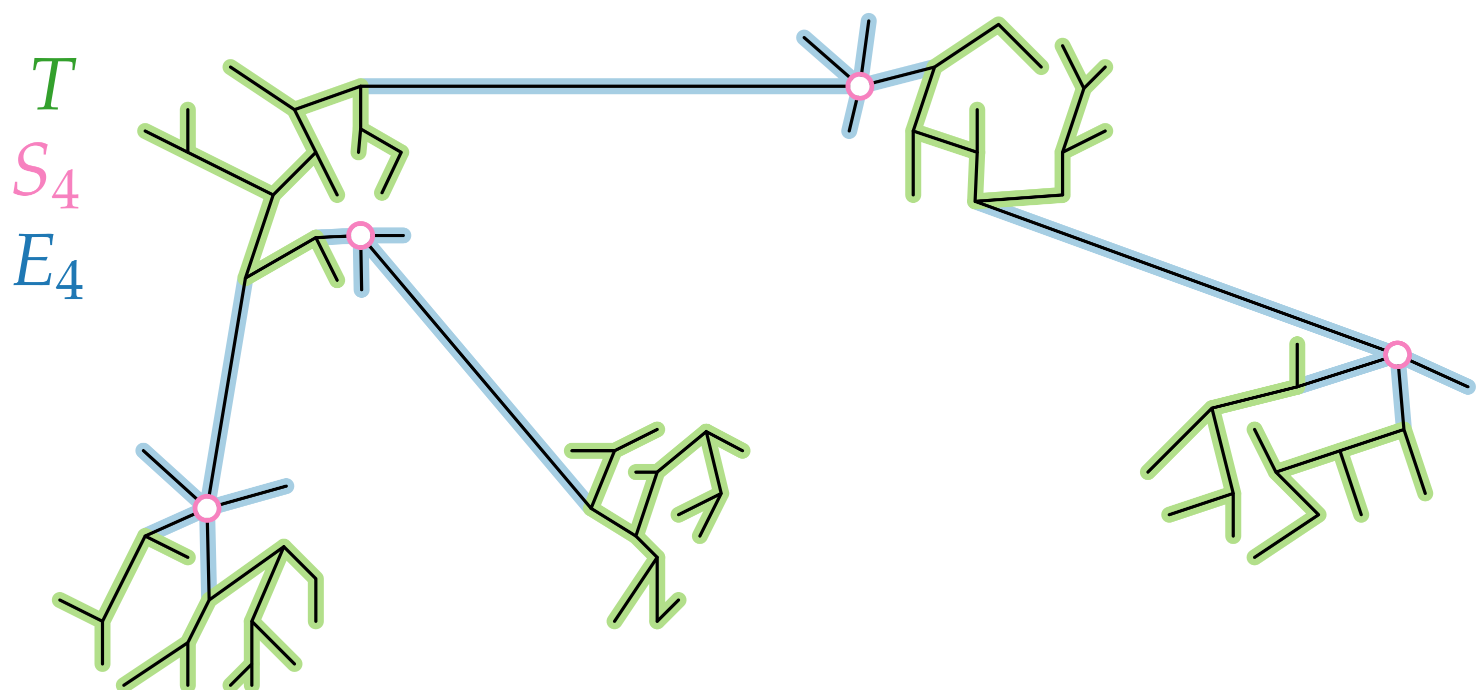
Proof. $|S_{\Delta(T)-\ell}| > 2^\ell |S_{\Delta(T)}| \underset{\ell = \lceil \log_2 n \rceil}{=} 2^{\lceil \log_2 n \rceil} |S_{\Delta(T)}| \geq n \cdot |S_{\Delta(T)}|$ ↯

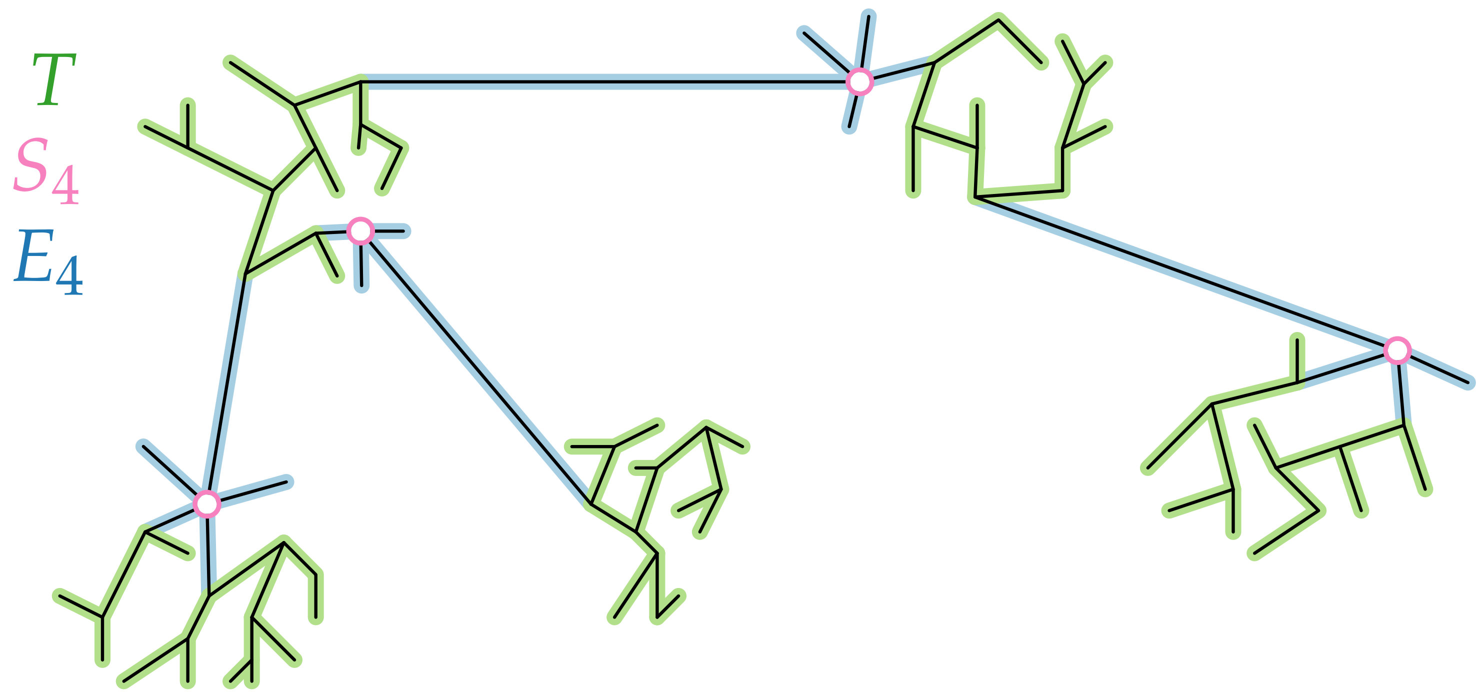
Otherwise

TODO: What if $\ell > \Delta(T)$?



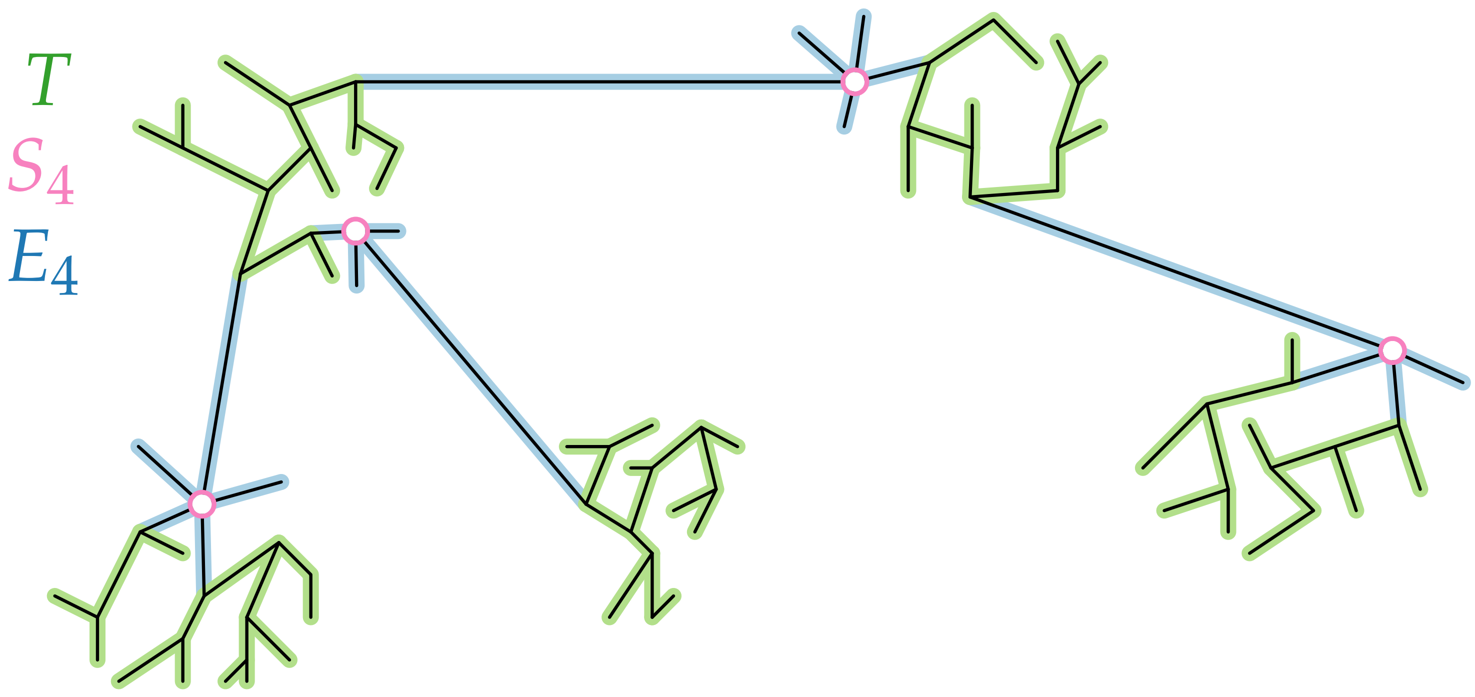
More Lemmas

Lemma 3. For $i \geq \Delta(T) - \ell + 1$,



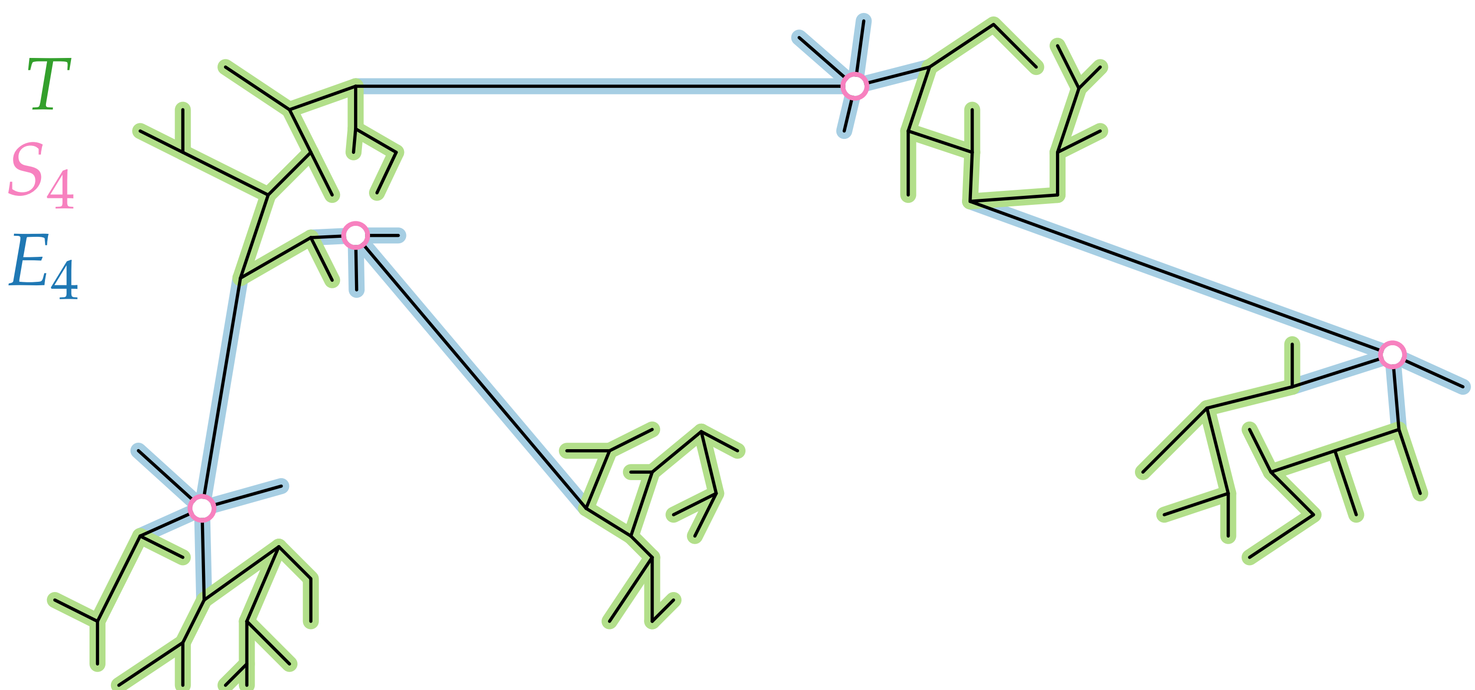
More Lemmas

Lemma 3. For $i \geq \Delta(T) - \ell + 1$,
(i) $|E_i| \geq (i-1)|S_i| + 1$,



More Lemmas

Lemma 3. For $i \geq \Delta(T) - \ell + 1$,
(i) $|E_i| \geq (i-1)|S_i| + 1$,
(ii) Each edge $e \in E(G) \setminus E_i$ connecting distinct components of $T \setminus E_i$ is incident to a node of S_{i-1}.



More Lemmas

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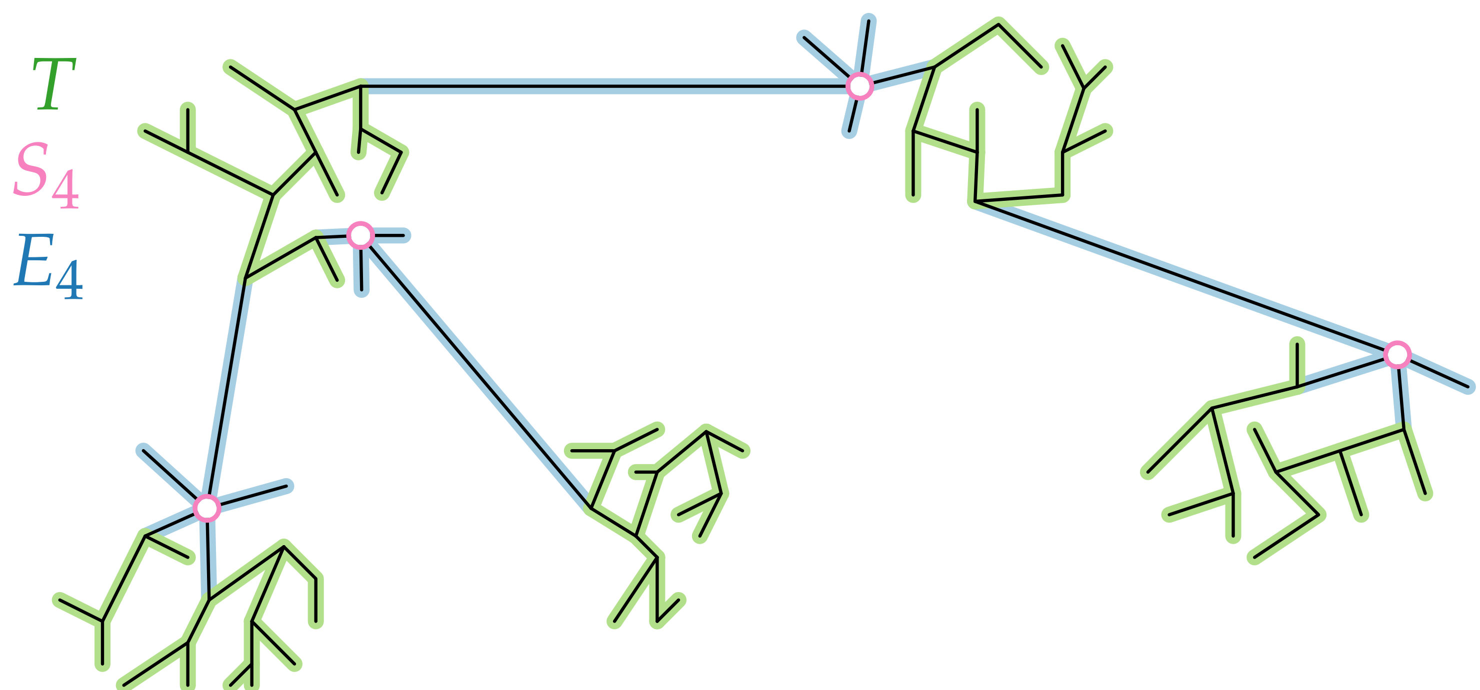

More Lemmas

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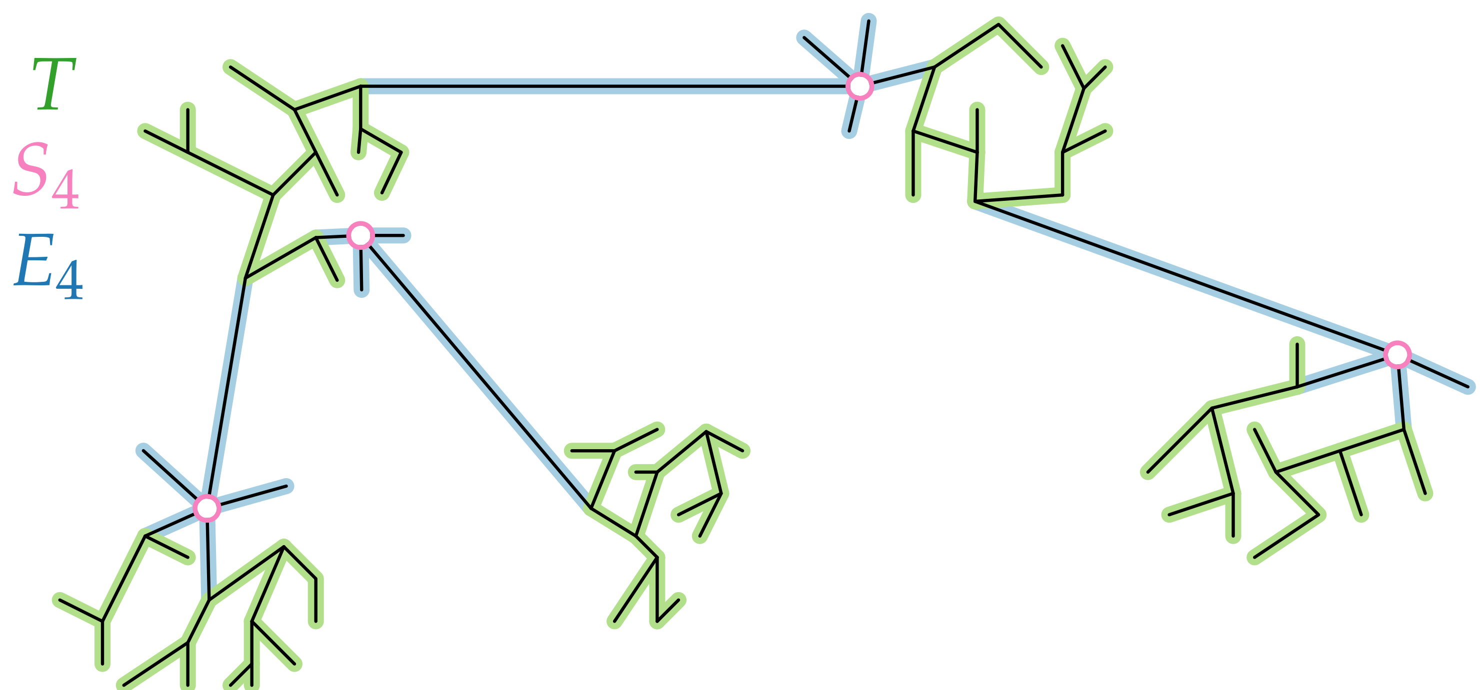

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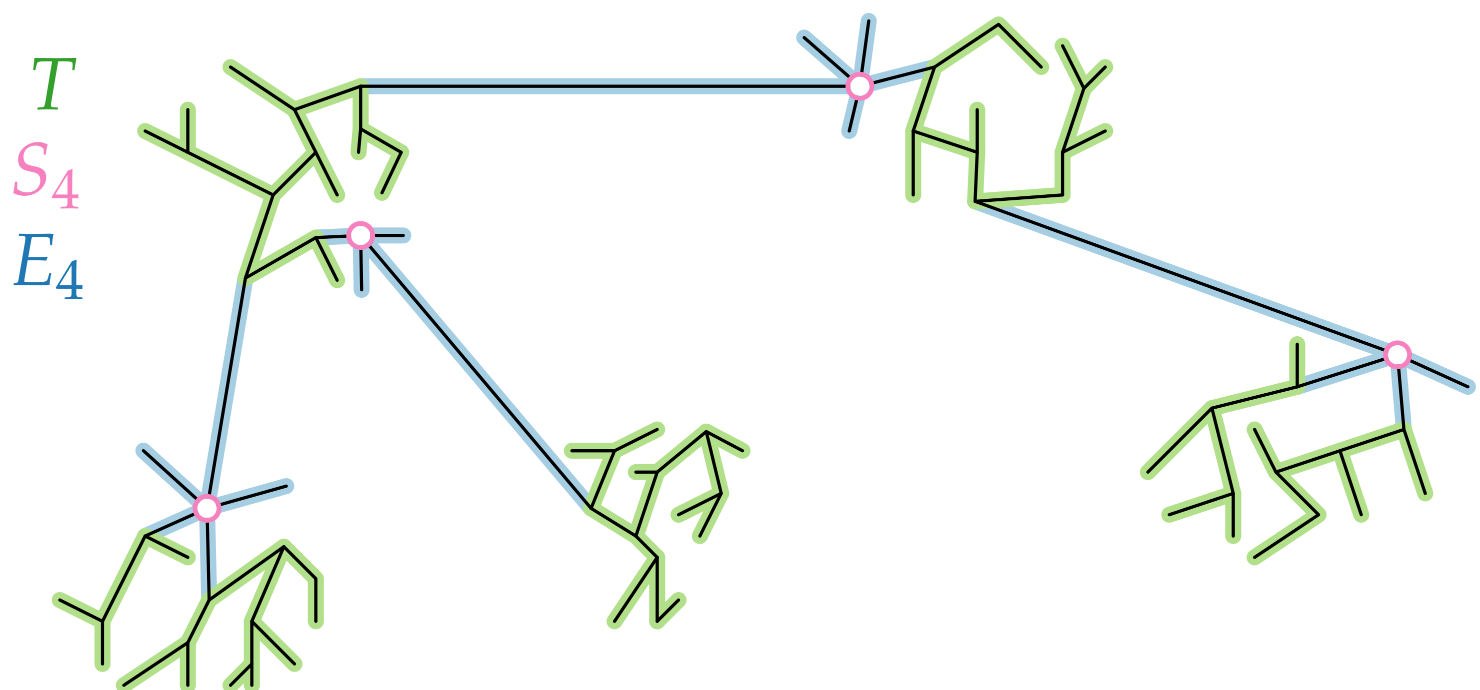

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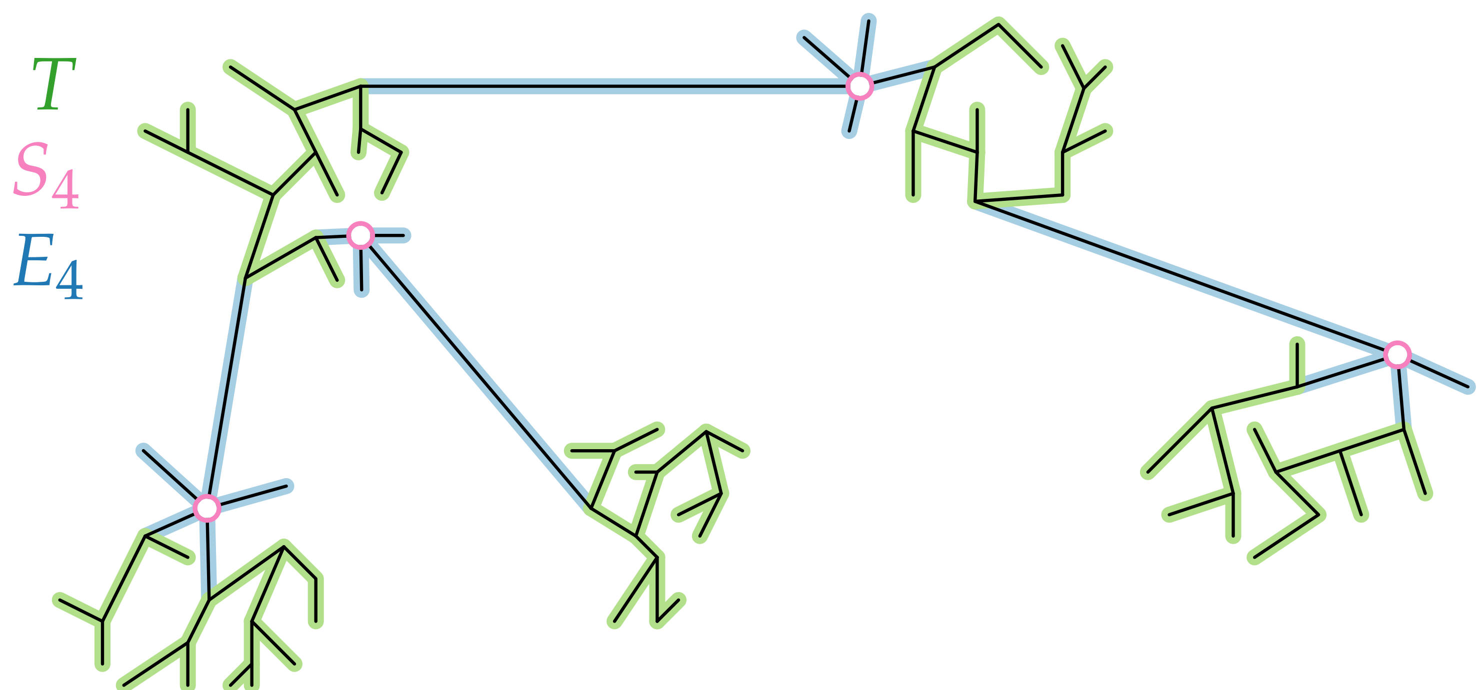

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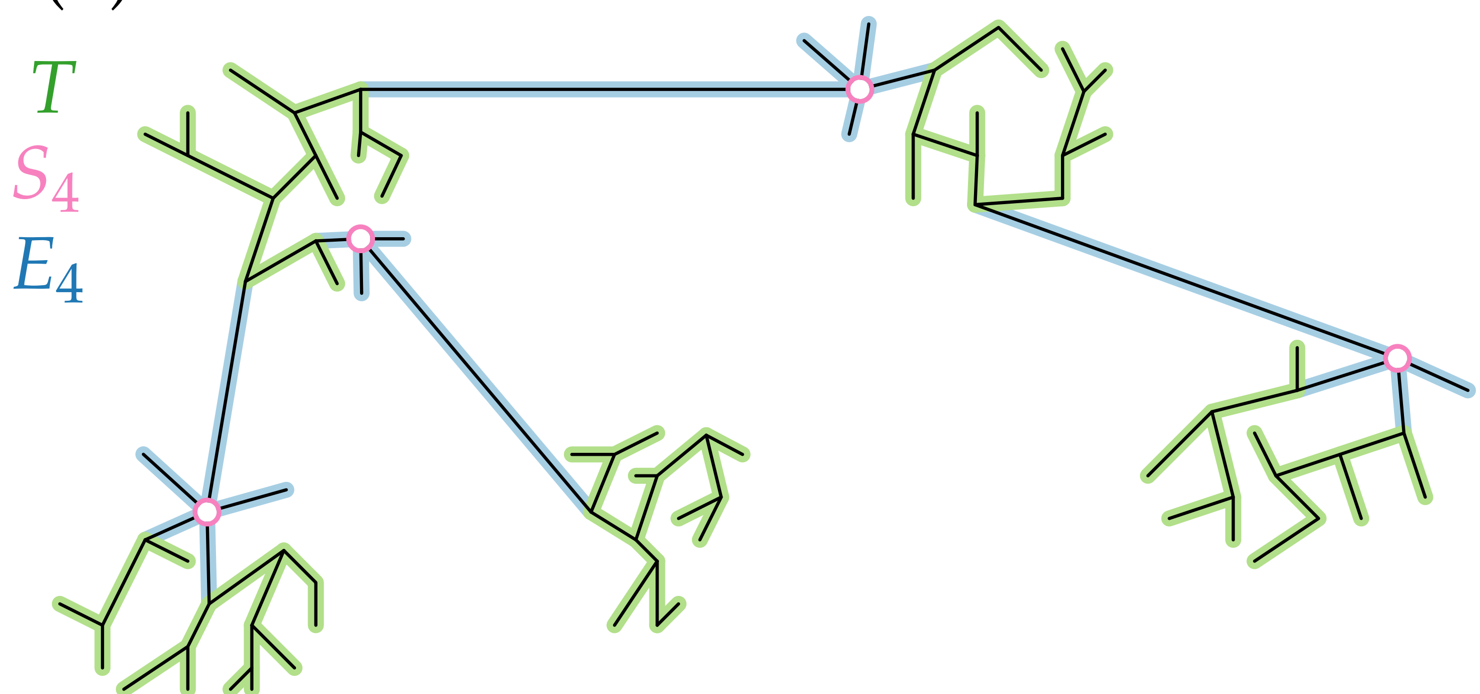

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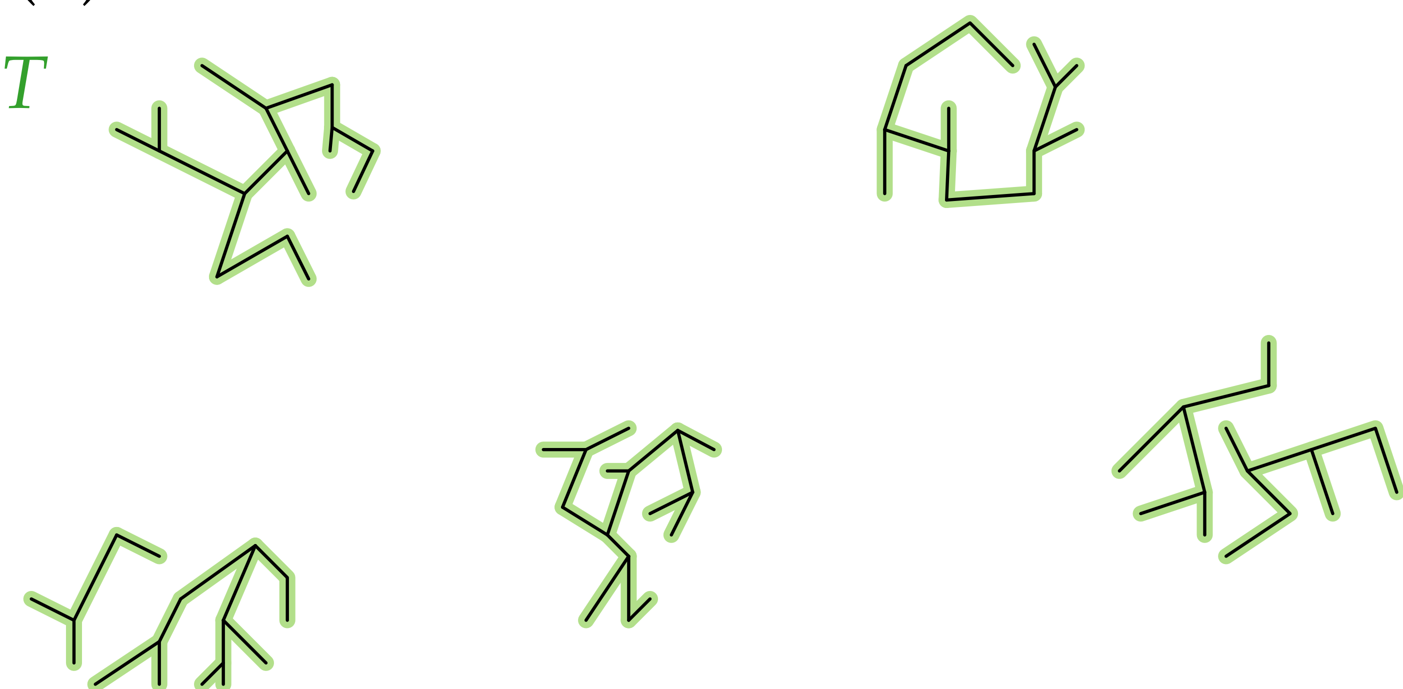

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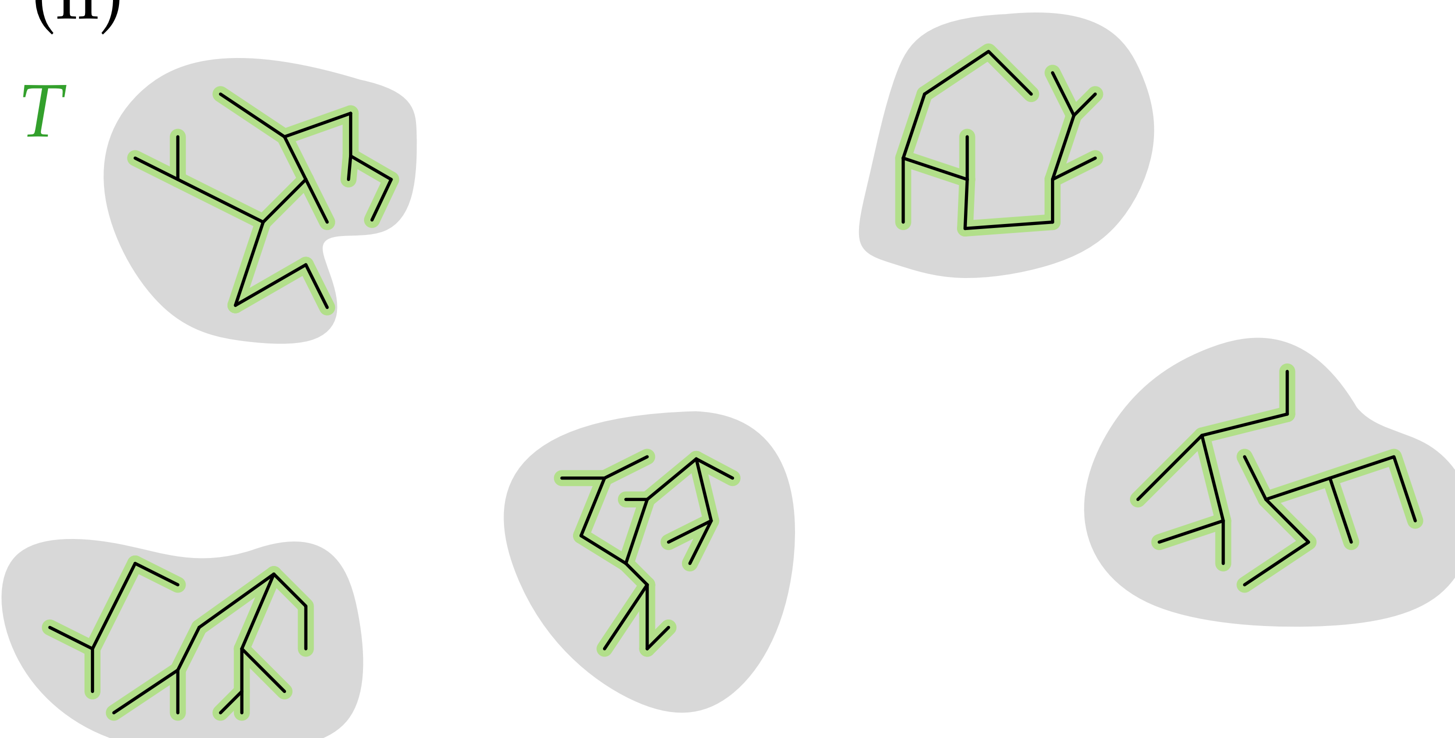

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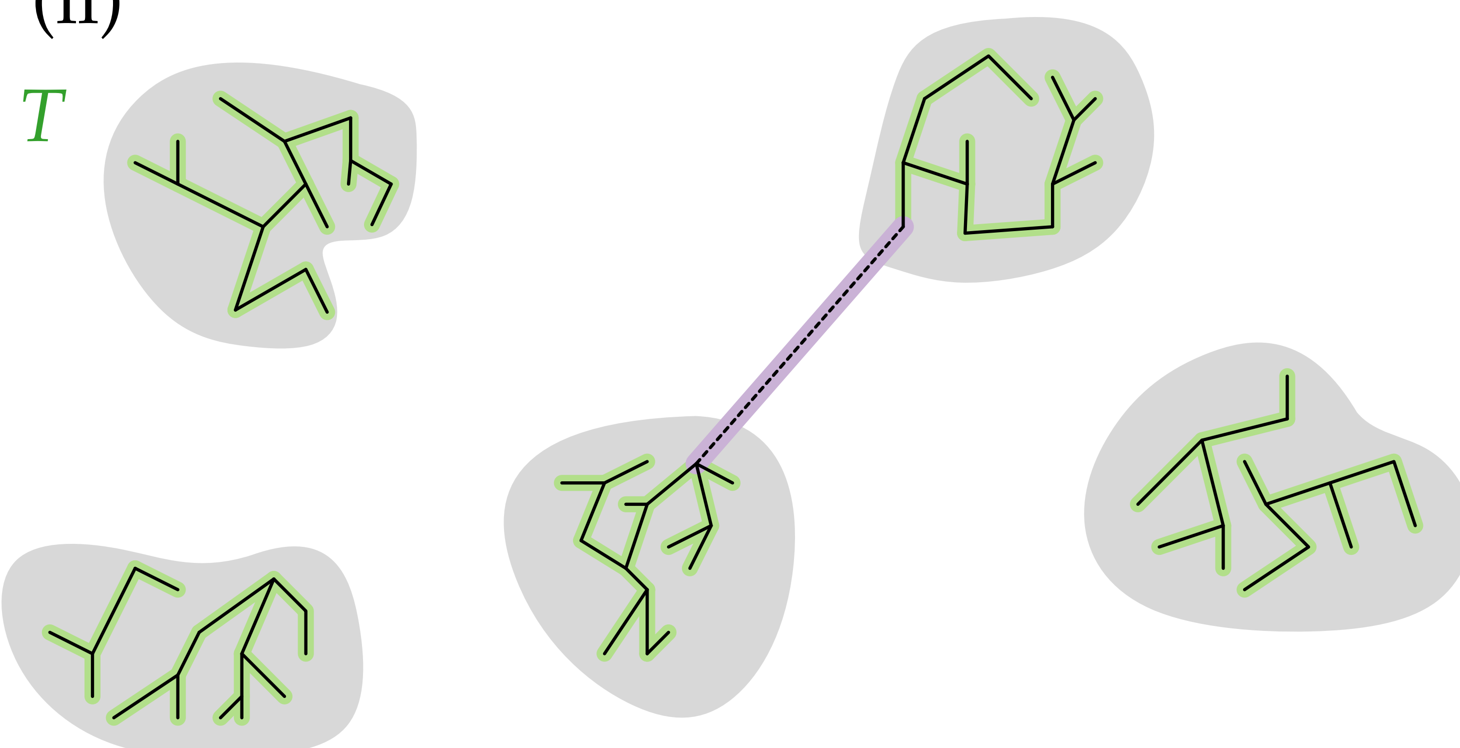

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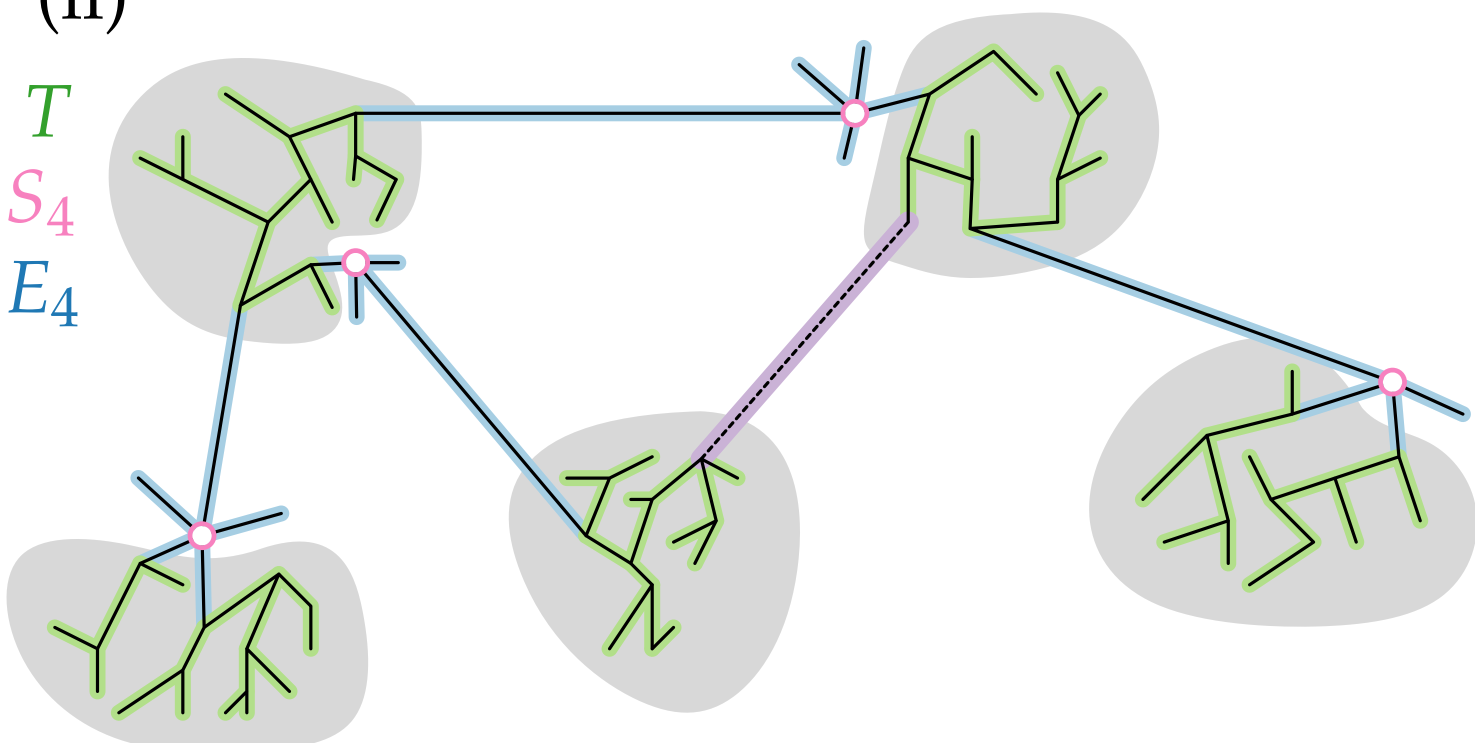

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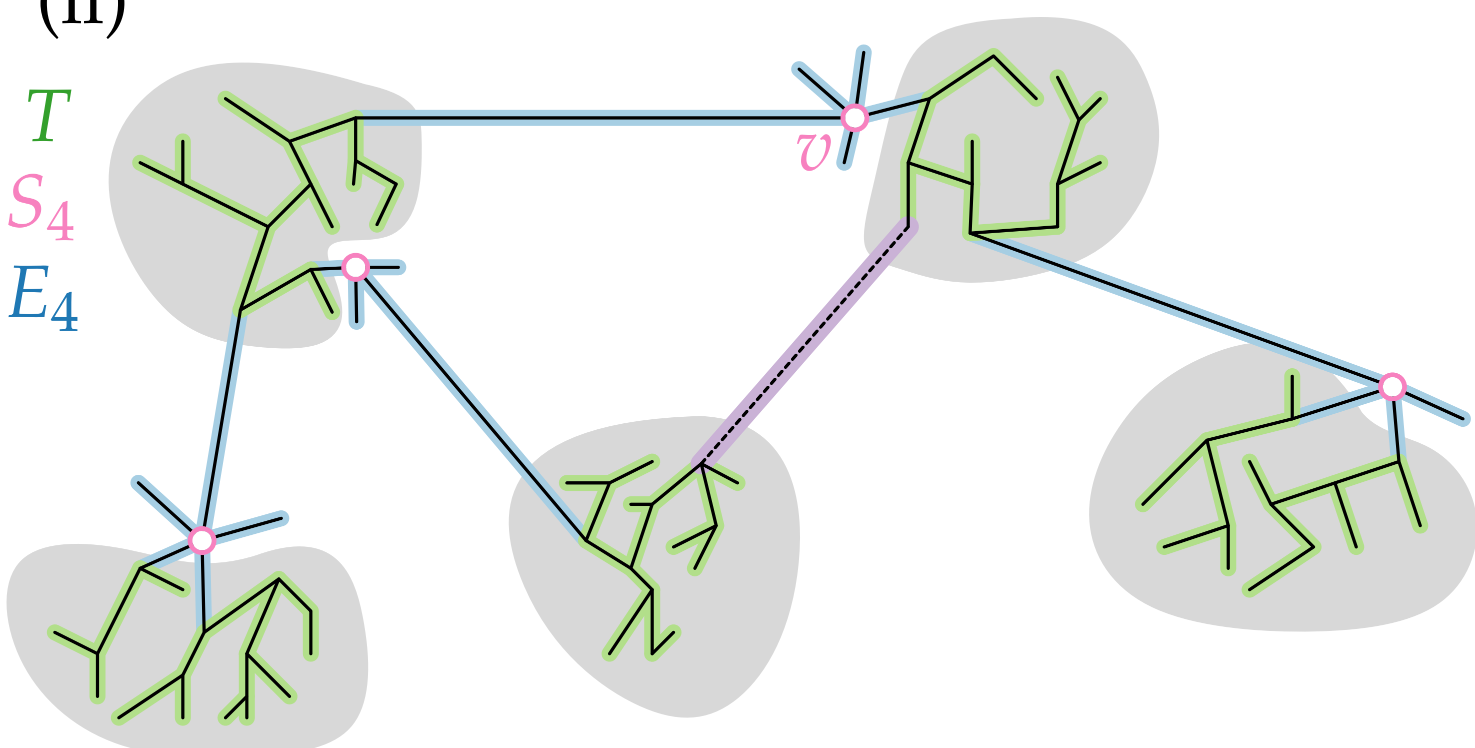

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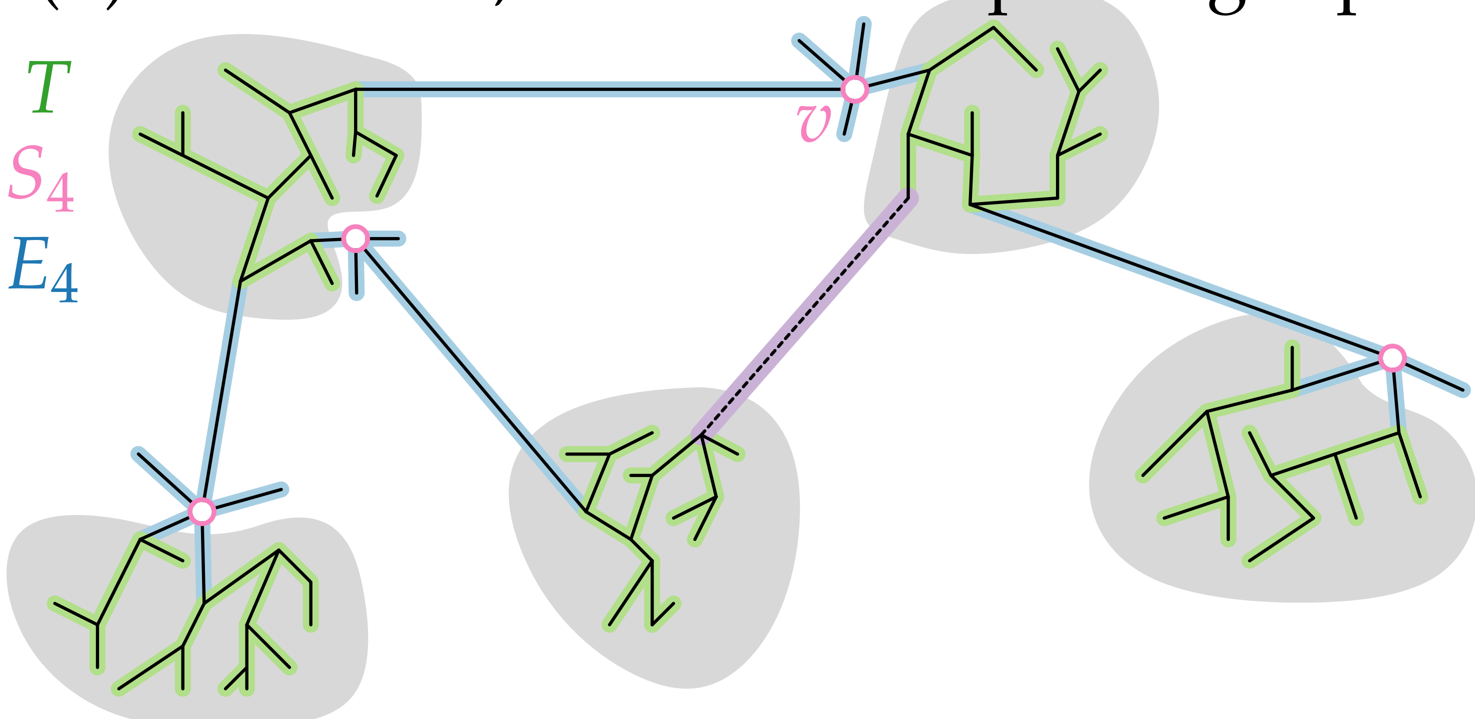

Approximation Algorithms

Lecture 10:
Minimum-Degree Spanning Tree
via Local Search

Part V:
Approximation Factor

Approximation Factor

Approximation Factor

[Fürer & Raghavachari: SODA'92, JA'94]

Theorem. Let T be a locally optimal spanning tree.
Then $\Delta(T) \leq 2 \cdot \mathrm{OPT} + \ell$, where $\ell = \lceil \log_2 n \rceil$.

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Approximation Algorithms

Lecture 10:
Minimum-Degree Spanning Tree
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Part VI:
Termination, Running Time & Extensions

Termination and Running Time

Theorem. The algorithm finds a locally optimal spanning tree efficiently.

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Lemma. After each flip $T \to T'$, $\Phi(T') \leq (1 - \frac{2}{27n^3})\Phi(T)$.

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Lemma. After each flip $T \to T'$, $\Phi(T') \leq (1 - \frac{2}{27n^3})\Phi(T)$.

Homework

- The function is bounded both from above and below.

Lemma. For each spanning tree T, $\Phi(T) \in [3n, n3^n]$.

- Executing $f(n)$ iterations would exceed the lower bound.

How does $\Phi(T)$ change?

$\Phi(T)$ decreases by: $(1 - \frac{2}{27n^3})^{f(n)} \leq (e^{-\frac{2}{27n^3}})^{f(n)} =$

$1 + x \leq e^x$

Termination and Running Time

Theorem. The algorithm finds a locally optimal spanning tree after at most $f(n)$ iterations.

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Termination and Running Time

Theorem. The algorithm finds a locally optimal spanning tree after $O(n^4)$ iterations.

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Extensions

[Fürer & Raghavachari: SODA'92, JA'94]

Corollary. For any constant $b > 1$ and $\ell = \lceil \log_b n \rceil$, the local search algorithm runs in polynomial time and produces a spanning tree T with $\Delta(T) \leq b \cdot \mathrm{OPT} + \lceil \log_b n \rceil$.

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Theorem. There is a local search algorithm that runs in $O(EV\alpha(E,V)\log V)$ time and produces a spanning tree T with $\Delta(T) \leq \mathrm{OPT} + 1$.