Approximation Algorithms

Lecture 2:
SetCover and ShortestSuperString

Part I:
SetCover

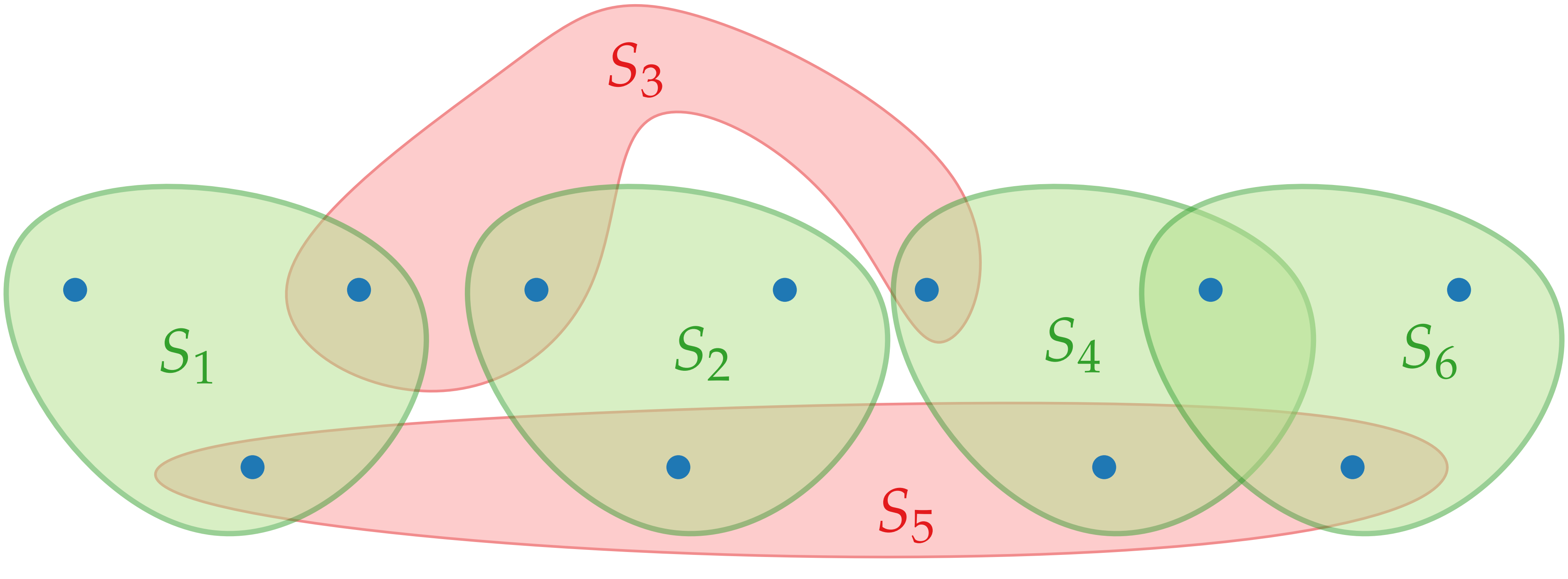
Joachim Spoerhase

Winter 2021/22

SetCover (card.)

Given a **ground set** U and a family $\mathcal{S}$ of **subsets** of U with $\bigcup \mathcal{S} = U$.

Find a **cover** $\mathcal{S}' \subseteq \mathcal{S}$ of U (i.e. with $\bigcup \mathcal{S}' = U$) of minimum cardinality.

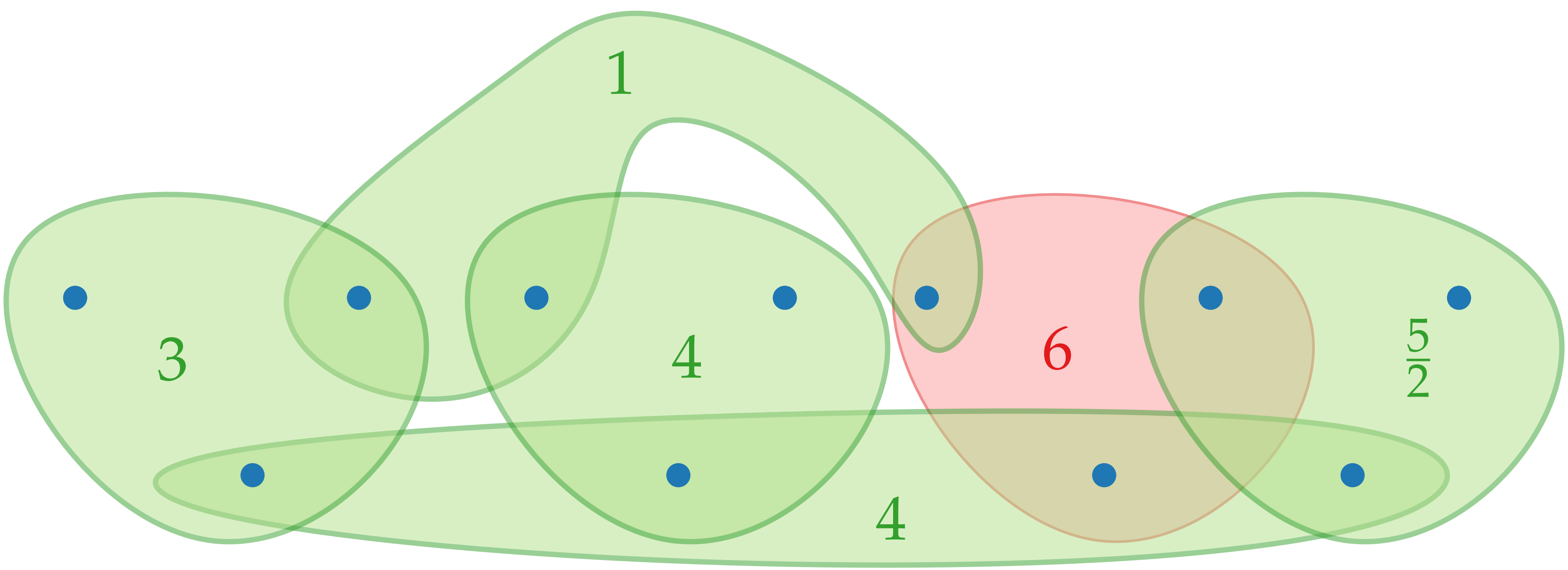


SetCover (general)

Given a **ground set** U and a family $\mathcal{S}$ of **subsets** of U with $\bigcup \mathcal{S} = U$.

Each $S \in \mathcal{S}$ has cost $c(S) > 0$.

Find a **cover** $\mathcal{S}' \subseteq \mathcal{S}$ of U (i.e. with $\bigcup \mathcal{S}' = U$) of minimum ~~cardinality.~~ total cost $c(\mathcal{S}') := \sum_{S \in \mathcal{S}'} c(S)$.



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Part II:
Greedy for SetCover

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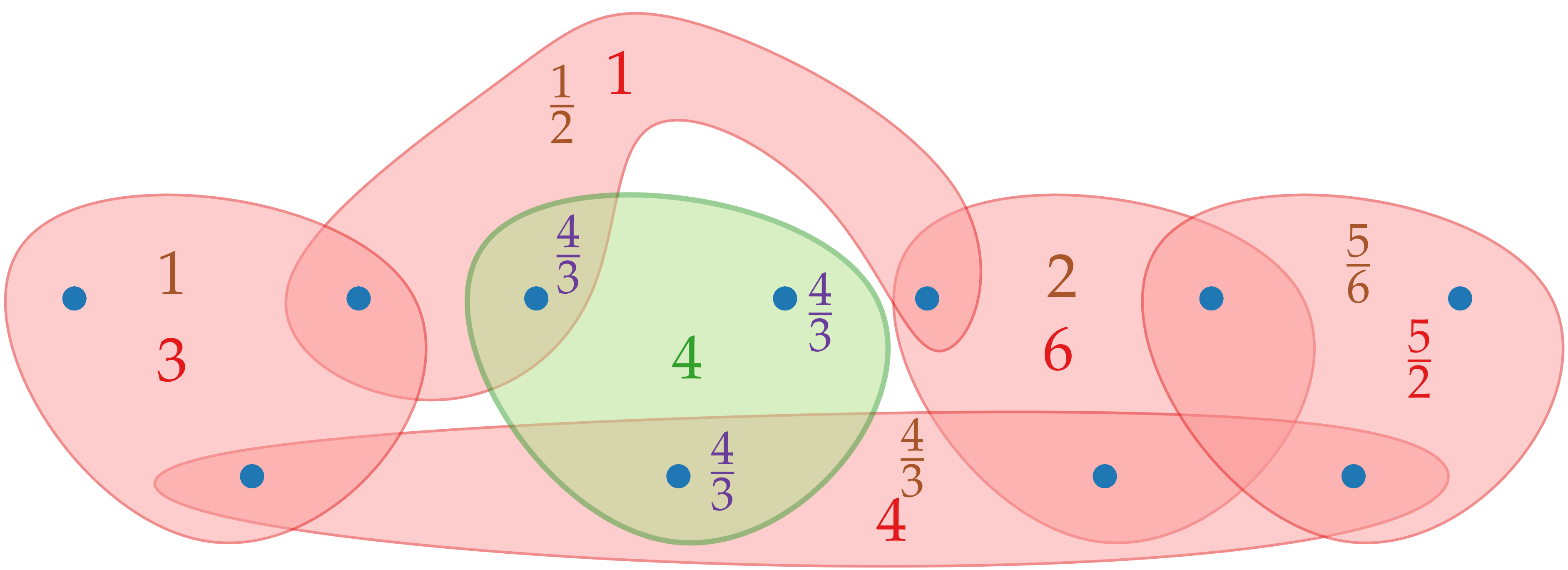
Iterative "Buying" of Elements

What is the real cost of picking a set?

Set with k elements and cost c has per-element cost $\frac{c}{k}$.

What happens if we "buy" a set?

Fix price of elements bought and recompute per-element cost.



Iterative "Buying" of Elements

What is the real cost of picking a set?

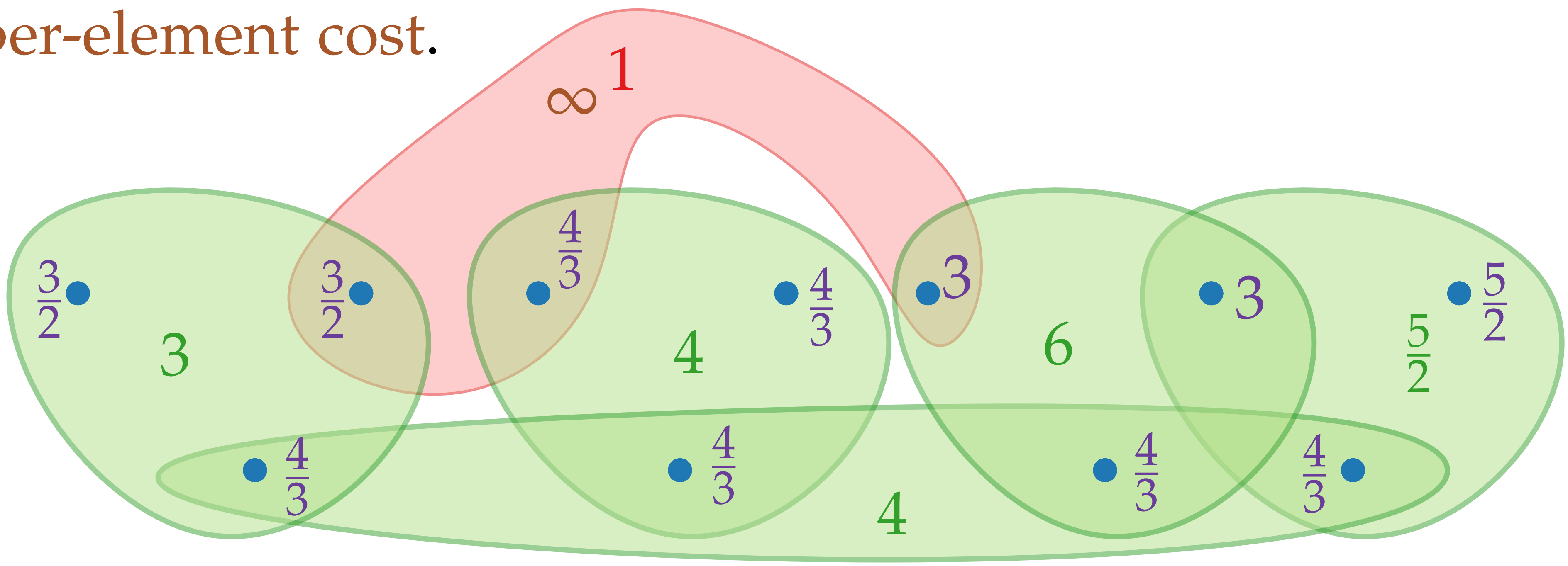
Set with k elements and cost c has per-element cost $\frac{c}{k}$.

What happens if we "buy" a set?

Fix price of elements bought and recompute per-element cost.

total cost: $\sum_{u \in U} \text{price}(u)$

Greedy: Always choose the set with the minimum per-element cost.



Greedy for SetCover

```
GreedySetCover(U, 𝒮, c)
  C ← ∅
  𝒮' ← ∅
  while C ≠ U do
    S ← Set from 𝒮 that minimizes c(S)/|S\C|
    foreach u ∈ S \ C do
      price(u) ← c(S)/|S\C|
    C ← C ∪ S
    𝒮' ← 𝒮' ∪ {S}
  return 𝒮'                                  // Cover of U
```

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Part III:
Analysis

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Analysis

Theorem. GreedySetCover is a factor-$\mathcal{H}_k$-approximation algorithm for SetCover, where k is the cardinality of the largest set in $\mathcal{S}$ and $\mathcal{H}_k := 1 + \frac{1}{2} + \frac{1}{3} + \ldots + \frac{1}{k} \to 0.5 + \ln k$.

Lemma. Let $S \in \mathcal{S}$ and $u_1, \ldots, u_\ell$ be the elements of S in the order they are covered ("bought") by GreedySetCover. Then $\text{price}(u_j) \leq c(S)/(\ell - j + 1)$.

Proof. Iteration at which alg. buys $u_j \Rightarrow$

- $\leq j - 1$ elements of S already bought
- $\geq \ell - j + 1$ elements of S not yet bought
- per-element cost for S: $\leq c(S)/(\ell - j + 1)$
- price by alg. no larger due to greedy choice

Analysis

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Lemma. $\text{price}(S) := \sum_{i=1}^{\ell} \text{price}(u_i) \le c(S) \cdot \mathcal{H}_\ell$.

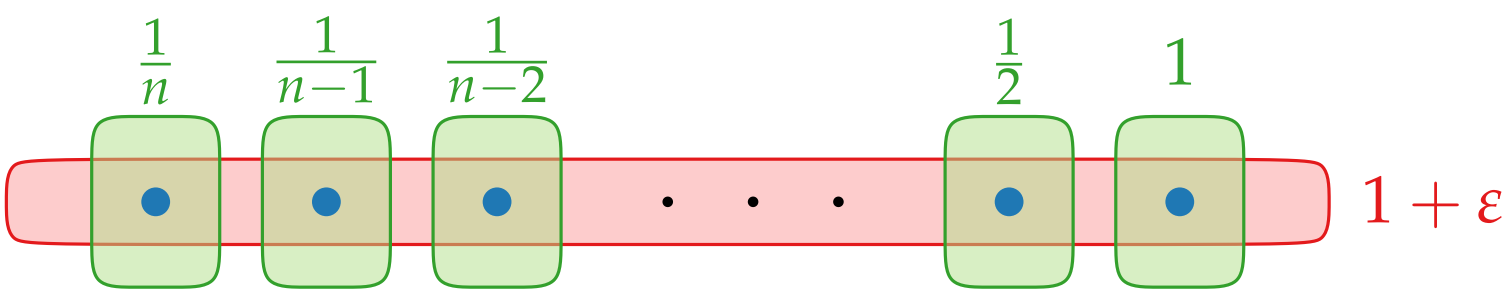
Proof. Let $\{S_1, \ldots, S_m\}$ be opt. sol. $\text{OPT} = \sum_{i=1}^{m} c(S_i)$

$$\text{price}(U) = \sum_{u \in U} \text{price}(u) \le \sum_{i=1}^{m} \text{price}(S_i)$$

$$\le \sum_{i=1}^{m} c(S_i) \cdot \mathcal{H}_k = \text{OPT} \cdot \mathcal{H}_k$$

Analysis tight?

Theorem. GreedySetCover is a factor-$\mathcal{H}_k$-approximation algorithm for SetCover, where k is the cardinality of the largest set in $\mathcal{S}$ and $\mathcal{H}_k := 1 + \frac{1}{2} + \frac{1}{3} + \ldots + \frac{1}{k} \leq 1 + \ln k$.



$\text{price}(U) = \mathcal{H}_n$ $\quad$ $\text{OPT} = 1 + \varepsilon$

better?

SetCover cannot be approximated within factor $(1 - o(1)) \cdot \ln(n)$ (unless P=NP)

Approximation Algorithms

Lecture 2:
SetCover and ShortestSuperString

Part IV:
ShortestSuperString

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ShortestSuperString (SSS)

Given a set $\{s_1, \ldots, s_n\} \subseteq \Sigma^+$ of strings over a finite alphabet Σ.
Find a **shortest string** s (*superstring*) such that each s_i, $i = 1, \ldots, n$ is a *substring* of s.

Example. $U := \{cbaa, abc, bcb\}$ cbaabcb

"covers" all strings in U

W.l.o.g.: No string s_i is a substring of any other string s_j.

$abcbaa$

abc
bcb
$cbaa$

SSS as a SetCover Problem

SetCover Instance: ground set U, set family $\mathcal{S}$, costs c.

ground set $U := \{s_1, \ldots, s_n\}$

Let be σ_{ijk} be the unique string with prefix s_i and suffix s_j where s_i and s_j *overlap* on k characters (for suitable i, j, k)

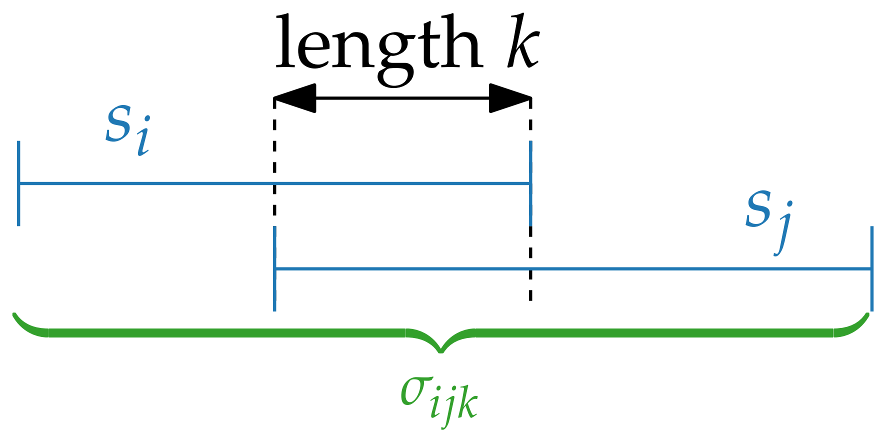
s_i : cabab s_j : ababc

cabab
  ababc

σ_{ij2} : cabababc

cabab
 ababc

σ_{ij4} : cababc



$S(\sigma_{ijk}) = \{s \in U \mid s \text{ substring of } \sigma_{ijk}\}$ contains the elements of the ground set covered by σ_{ijk}.

$c\left(S(\sigma_{ijk})\right) = |\sigma_{ijk}|$ (number of characters in σ_{ijk})

$\mathcal{S} = \{S(\sigma_{ijk}) \mid k > 0\}$ (possibly $i = j$)

Approximation Algorithms

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Part V:
Solving ShortestSuperString via SetCover

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Relating SSS and SetCover

Lemma. Let $\mathrm{OPT}_{\mathrm{SSS}}$ be the length of a shortest superstring of U and $\mathrm{OPT}_{\mathrm{SC}}$ be the minimum cost of the corresponding SetCover instance. Then:

$$\mathrm{OPT}_{\mathrm{SSS}} \leq \mathrm{OPT}_{\mathrm{SC}}$$

Proof.

Consider an optimal set cover $\{S(\pi_1), \dots, S(\pi_k)\}$ of U.

$s := \pi_1 \circ \cdots \circ \pi_k$ is a superstring of U of length $\sum_{i=1}^{k} |\pi_i| = \sum_{i=1}^{k} c(S(\pi_i)) = \mathrm{OPT}_{\mathrm{SC}}$.

Thus, $\mathrm{OPT}_{\mathrm{SSS}} \leq |s| = \mathrm{OPT}_{\mathrm{SC}}$.

Relating SSS and SetCover

Lemma. $\mathrm{OPT}_{\mathrm{SC}} \leq 2 \cdot \mathrm{OPT}_{\mathrm{SSS}}$

Proof. Consider optimal superstring s.
Construct set cover with cost $\leq 2|s| = 2 \cdot \mathrm{OPT}_{\mathrm{SSS}}$.

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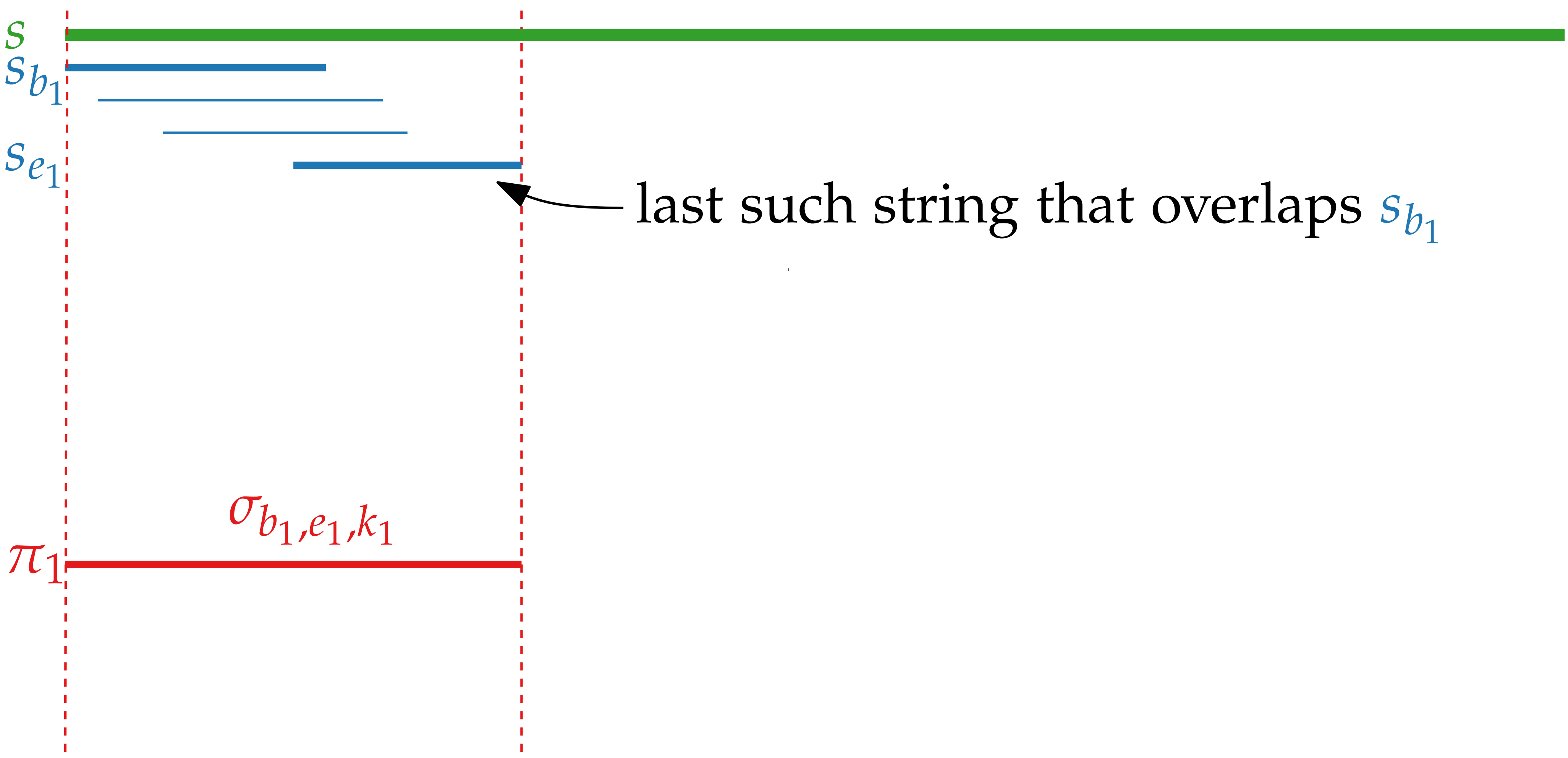

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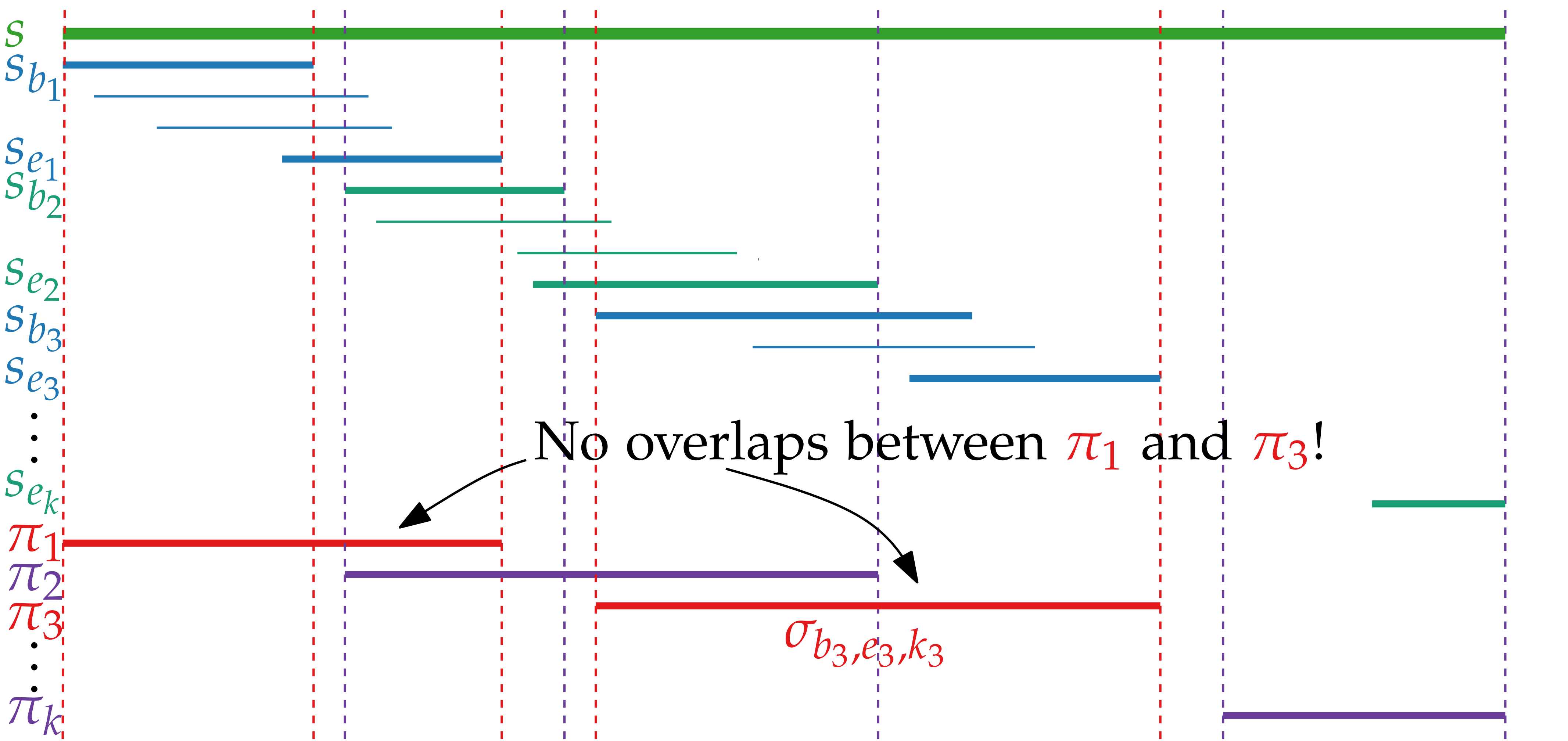

Relating SSS and SetCover

Lemma. $\mathrm{OPT}_{\mathrm{SC}} \leq 2 \cdot \mathrm{OPT}_{\mathrm{SSS}}$

Proof.

Each string $s_i \in U$ is a substring of some π_j.

$\{S(\pi_1), \dots, S(\pi_k)\}$ is a solution for the SetCover instance with cost $\sum_i |\pi_i|$.

Substrings π_j, π_{j+2} do not overlap.

Each character in s lies in at most **two** (subsequent) substrings, namely π_j and π_{j+1}.

$\sum_i |\pi_i| \leq 2|s| = 2 \cdot \mathrm{OPT}_{\mathrm{SSS}}$

Algorithm for SSS

1. Construct SetCover instance U, $\mathcal{S}$, c .
2. Compute a set cover $\{S(\pi_1), \ldots, S(\pi_k)\}$ with algorithm GreedySetCover.
3. Return $\pi_1 \circ \cdots \circ \pi_k$ as the superstring.

Theorem. This algorithms is a factor-$2\mathcal{H}_n$-approximation algorithm for ShortestSuperString.

better?

The best-known approximation factor for ShortestSuperString is $\frac{71}{30} \approx 2.367$.

ShortestSuperString cannot be approximation within factor $\frac{333}{332} \approx 1.003$ (unless P=NP).